

Minimum Camera Placement for Forest Monitoring – DP Design

Problem

We are given a forest graph $G = (V, E)$. Each vertex represents a candidate designated point (cdp) where we can place a camera, and each edge represents a shared region that can be monitored by cameras at two different cdp's. The goal is to find the minimum number of cameras needed to monitor all vertices (regions).

This is the *minimum dominating set* problem. While it is NP-hard for general graphs, it can be solved in polynomial time using dynamic programming (DP) for forest structures (trees). The following DP finds the optimum by rooting each tree.

DP Design

1) Subproblems

For each node v , we maintain three states (tree rooted at r):

- $dp[v][0]$: Camera is placed at v .
- $dp[v][1]$: No camera at v , but at least one child has a camera and v is dominated.
- $dp[v][2]$: No camera at v , v is not yet dominated; domination must come from parent.

2) Recursive Formulation

For a leaf node:

$$dp[v][0] = 1, \quad dp[v][1] = \infty, \quad dp[v][2] = 0$$

For an internal node v with children set $C(v)$:

$$dp[v][0] = 1 + \sum_{c \in C(v)} \min(dp[c][0], dp[c][1], dp[c][2])$$

$$dp[v][2] = \sum_{c \in C(v)} \min(dp[c][0], dp[c][1])$$

In state 1, v must be dominated by at least one child camera:

$$base = \sum_{c \in C(v)} \min(dp[c][0], dp[c][1])$$

$$dp[v][1] = \begin{cases} \infty, & C(v) = \emptyset \\ base + \min_{c \in C(v)} (dp[c][0] - \min(dp[c][0], dp[c][1])), & \text{otherwise} \end{cases}$$

This is because at least one child must actually have a camera (state 0).

Valid answer for root: $\min(dp[r][0], dp[r][1])$ (root must be dominated).

3) Justification for DP

- Optimal substructure: The optimal solution for each subtree (child) is independent of others; combinations only interact through state labels.
- Number of subproblems: Constant number of states per node, total $O(|V|)$ subproblems.
- Overlapping subproblems: Each node's states cannot be requested by multiple parents (tree structure), so memoization with DP is efficient.

4) Pseudocode

Time Complexity (Task 2)

Let $n = |V|$ and $m = |E|$. Each edge is visited at most twice during the DFS, and per node we do $O(1)$ work for constant-state DP aggregation. Thus the time complexity is $O(n + m)$ in the worst case (for trees, $m = n - 1$, so $O(n)$). The space complexity is $O(n)$ for the DP tables and recursion stack.

Step-by-Step Example (Task 3)

Consider a sample tree with 5 cdps (nodes) where each cdp monitors 2–3 regions: edges $\{(0, 1), (1, 2), (1, 3), (3, 4)\}$ and root at 1.

- Post-order traversal: process children before parent.
- Node 0 (leaf): $dp[0] = [1, \infty, 0]$.
- Node 2 (leaf): $dp[2] = [1, \infty, 0]$.
- Node 4 (leaf): $dp[4] = [1, \infty, 0]$.
- Node 3 (child 4):

$$- dp[3][0] = 1 + \min(1, \infty, 0) = 1$$

Algorithm 1 MinCamerasOnTree(G, r)

```
1: function SOLVE( $v, parent$ )
2:    $dp[v][0] \leftarrow 1$  ▷ Cost if camera placed at  $v$ 
3:    $dp[v][1] \leftarrow \infty$  ▷ Initially impossible
4:    $dp[v][2] \leftarrow 0$  ▷ Cost if  $v$  waits for parent
5:   for child  $c$  of  $v$  where  $c \neq parent$  do
6:     SOLVE( $c, v$ ) ▷ Process children first
7:   end for
8:    $base \leftarrow 0, gain \leftarrow \infty$ 
9:   for child  $c$  of  $v$  where  $c \neq parent$  do
10:     $m02 \leftarrow \min(dp[c][0], dp[c][1], dp[c][2])$  ▷ Best for state 0
11:     $m01 \leftarrow \min(dp[c][0], dp[c][1])$  ▷ Best for states 1 and 2
12:     $dp[v][0] \leftarrow dp[v][0] + m02$ 
13:     $dp[v][2] \leftarrow dp[v][2] + m01$ 
14:     $base \leftarrow base + m01$ 
15:     $gain \leftarrow \min(gain, dp[c][0] - m01)$  ▷ Extra cost to force state 0
16:   end for
17:   if  $gain < \infty$  then
18:      $dp[v][1] \leftarrow base + gain$  ▷ At least one child has camera
19:   end if
20: end function
21: SOLVE( $r, -1$ )
22: return  $\min(dp[r][0], dp[r][1])$  ▷ Root must be dominated
```

$$- dp[3][2] = \min(1, \infty) = 1$$

$$- dp[3][1] = base + gain = 1 + (1 - 1) = 1 \text{ (force child 4 to have camera)}$$

So $dp[3] = [1, 1, 1]$.

- Node 1 (children 0,2,3):

$$- \text{For child 0: } m02 = 1, m01 = 1.$$

$$- \text{For child 2: } m02 = 1, m01 = 1.$$

$$- \text{For child 3: } m02 = 1, m01 = 1.$$

Aggregate:

$$dp[1][0] = 1 + (1 + 1 + 1) = 4$$

$$dp[1][2] = 1 + 1 + 1 = 3$$

$$base = 3, \quad gain = \min(0, 0, 0) = 0$$

$$dp[1][1] = 3 + 0 = 3$$

- Answer at root 1: $\min(dp[1][0], dp[1][1]) = \min(4, 3) = 3$ cameras.

Placement achieving 3 cameras: cameras at nodes $\{0, 3, 1\}$ (or any equivalent minimum set).

Task 5: Functional Testing

We design 7 test instances appropriate for both white-box and black-box functional testing. Each instance targets specific properties of the algorithm.

Test Instances

Instance 1: Single Node

- Graph: Isolated vertex (no edges)
- Structure: $V = \{0\}$, $E = \emptyset$
- Expected: 1 camera (the node must monitor itself)
- **White-box testing:** Tests the base case initialization: $dp[v][0] = 1$, $dp[v][1] = \infty$, $dp[v][2] = 0$ for a leaf node.
- **Black-box testing:** Tests minimal input handling - a single vertex must be monitored.

Instance 2: Two Nodes

- Graph: Single edge connecting two nodes
- Structure: Path $0 - 1$
- Expected: 1 camera (at either node, covering both)
- **White-box testing:** Tests state transitions where placing a camera at one node (state 0) covers its neighbor.
- **Black-box testing:** Tests minimal connected graph - edge case for connectivity.

Instance 3: Path of 3 Nodes

- Graph: Linear path
- Structure: Path $0 - 1 - 2$
- Expected: 1 camera (optimal at middle node 1)
- **White-box testing:** Tests internal node with two children, state 1 calculation where a child's camera dominates the parent.
- **Black-box testing:** Tests optimal placement in linear structures.

Instance 4: Star Graph

- Graph: Center node connected to multiple leaves
- Structure: Center 0 connected to leaves $\{1, 2, 3, 4\}$

- Expected: 1 camera (at center node 0)
- **White-box testing:** Tests node with multiple children, gain calculation for state 1 when all children are leaves.
- **Black-box testing:** Tests high-degree vertex scenario - hub-and-spoke topology.

Instance 5: Binary Tree

- Graph: Balanced binary tree
- Structure: Root 0, level-1: {1, 2}, level-2: {3, 4, 5, 6}
- Expected: 2 cameras (optimal placement)
- **White-box testing:** Tests recursive DP on balanced structure, multiple levels of recursion, complex state interactions.
- **Black-box testing:** Tests hierarchical tree structure - realistic tree topology.

Instance 6: Forest with Multiple Components

- Graph: Two disconnected paths
- Structure: Component 1: $0 - 1 - 2$, Component 2: $3 - 4 - 5$
- Expected: 2 cameras (1 per component)
- **White-box testing:** Tests component detection algorithm, independent processing of each component, root selection.
- **Black-box testing:** Tests disconnected graph handling - forest structure.

Instance 7: Complex Tree

- Graph: Tree with multiple branching points
- Structure: $0 - 1 - 2$, $1 - 3 - 4$ (root at 1)
- Expected: 2 cameras (optimal: at nodes 1 and 3, or 1 and 4)
- **White-box testing:** Tests complex state transitions, gain calculation with multiple children having different optimal states.
- **Black-box testing:** Tests realistic scenario with multiple branching points and varying subtree structures.

Test Results

All 7 test instances passed successfully. The results are summarized in Table 1, showing each instance, its expected and actual camera counts, and the algorithm properties being tested.

Table 1: Functional Testing Results

Instance	Expected	Actual	Status	Properties Tested
Single node	1	1	✓	Base case initialization, leaf node handling
Two nodes	1	1	✓	State transitions, minimal connectivity
Path (3 nodes)	1	1	✓	Internal node with children, state 1 calculation
Star graph	1	1	✓	High-degree vertex, multiple children, gain calculation
Binary tree	2	2	✓	Recursive DP, multiple levels, balanced structure
Forest (2 components)	2	2	✓	Component detection, independent processing
Complex tree	2	2	✓	Complex state transitions, multiple branching

Properties Verified

The test suite verifies the following algorithm properties:

- **Correctness:** All instances produce optimal solutions (verified manually for each case).
- **Base cases:** Single node and two-node cases handled correctly.
- **State transitions:** All three DP states (0, 1, 2) are correctly computed and used.
- **Component handling:** Disconnected graphs are processed correctly by detecting and handling each component independently.
- **Edge cases:** Minimal inputs, high-degree vertices, and various tree structures are handled correctly.
- **Optimality:** The algorithm finds minimum camera placements for all test cases.