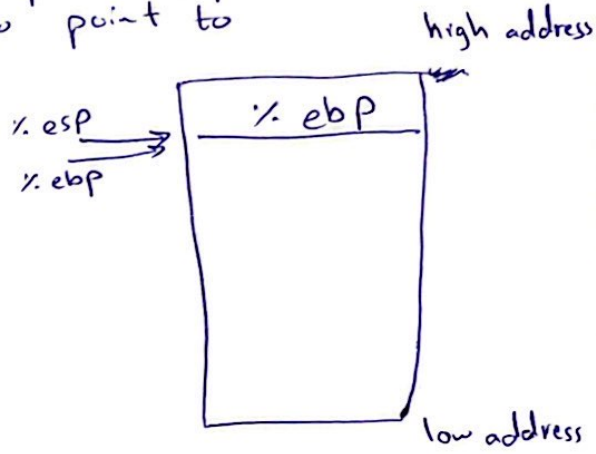
(1)

-main:

LFB13:

```
pushl %.ebp
movl %.esp, %.ebp
```

# pushing the value of the base stack pointer into the stack (decrement %.esp by 4)
# and then move the value of %.esp to %.ebp which sets the base pointer to point to the same location as %.esp

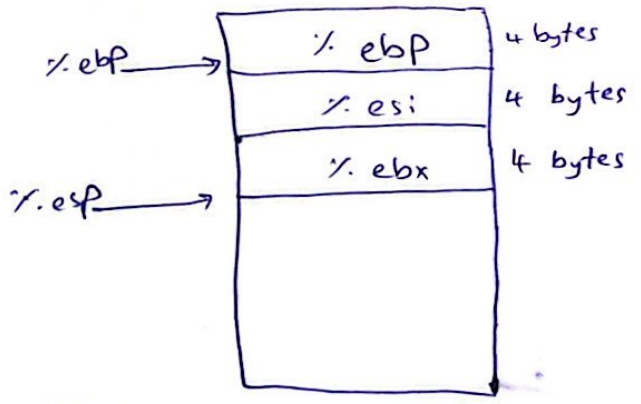



```
pushl %.esi
```

# decrement the stack pointer %.esp by 4 and push the 32-bit %.esi value into the stack

```
pushl %.ebx
```

# decrement %.esp by 4 and push %.ebx value




```
andl $-16, %.esp
```

# 16 ⟶ ...00010000 (32-bits) 2's comp ⟶ ...11101111 + 1

...11110000 (32-bits)

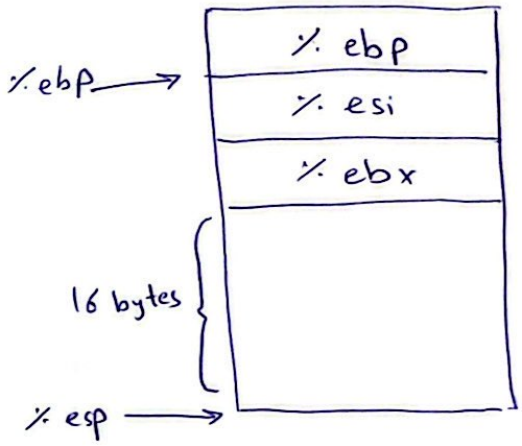
# So -16 in hexa is 0xfffffff0 → %.esp = %.esp &0xfffff

# when anding the %.esp register with 0xfffffff0 we set the %.esp register to the nearest multiple of 16 (stack alignment) which is sometimes required to improve performance for processors that operate more efficiently when data is aligned.

# reference : https://stackoverflow.com/questions/23309863/why-does-gcc-produce-andl-16

```
subl $16, %.esp
```

# %.esp = %.esp -16 , allocate 16 bytes in the stack (for local variable




```
call ___main
```

# this line refers to the ___main in the C library which eventually calls the main() function (internal procedure)

# reference : https://community.st.com/t5/stm32-mcus-products/main-in-startup-assembly/td-p/391658

```
movl $1 , %.ebx
movl $0 , %.esi
```

# %.ebx = 1 , i is now stored in %.ebx

# %.esi = 0 , result in %.esi

```
jmp L7
```

# jump to L7

L7:

```
cmpl $5, %ebx
# i : 5
# compare i with 5
jle L8
```

| result | %esi |
|---|---|
| i | %ebx |

# if less than or equal then jump to L8 which is the loop body ($i \leq 5$) ⟶ L8 explained in page (5)

# we noticed that the compiler didn't save the value of the variable "number" in any register and dealt with it as a constant

# the compiler optimized the code by treating that variable as a compile time constant rather than storing it in a register

```
movl %esi, 4(%esp)
```

# move the content in the register %esi to the memory location of (4+%esp), (4+%esp) = %esi (memory location)

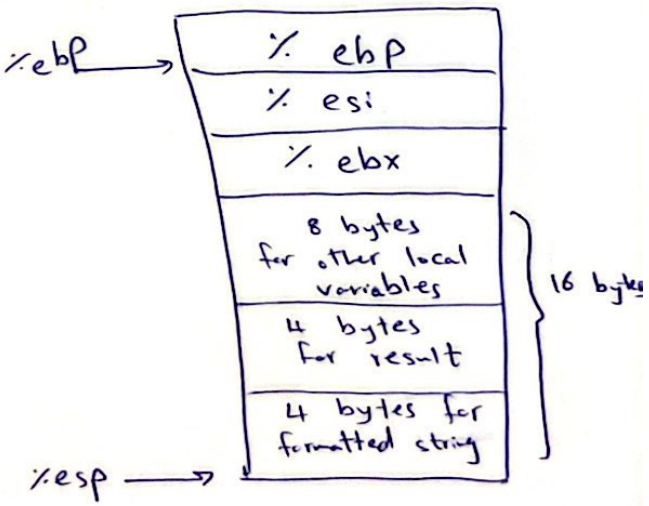
# 4 bytes from the stack pointer, preparing it to be passed to printf

```
movl $LC0, (%esp)
```

# the format string "%d" is stored at the memory location pointed to by %esp to be prepared to be passed to printf, the "%d" in the format string suggest that it is expecting integar argument

#$LC0 is explained in details in page (6)

```
call -printf
# to print result
```



(4)

```
movl   $0, %eax
```

# %eax = 0 , moving the immediate value 0 to %eax
for returning 0 at the end of the main function

```
leal   -8 (%ebp), %esp
```

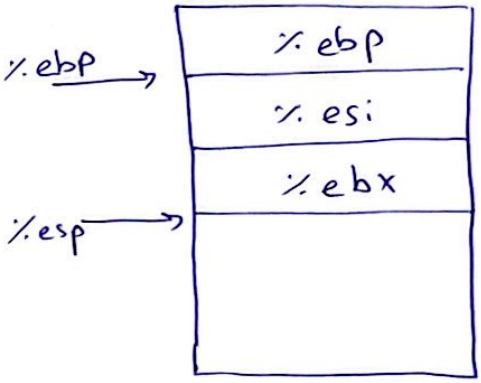
# %esp = %ebp - 8 (cleaning up the stack)

```
popl   %ebx
```

# increment %esp by 4 and
storing the value at %ebx

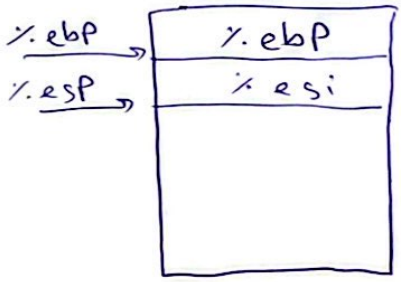
```
popl   %esi

popl   %ebp
```

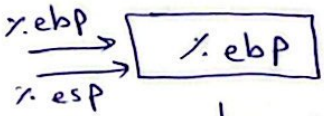



popl %ebx

# same thing incrementing %esp by 4 each step
to free the stack




popl %esi




popl %ebp

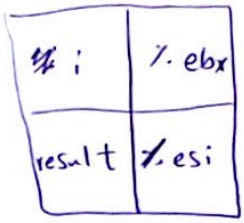
freed

```
ret
```

# returning the value in the register %eax
which is 0 (return 0;)

(5) #if i ≤ 5 then we jump to L8

L8:

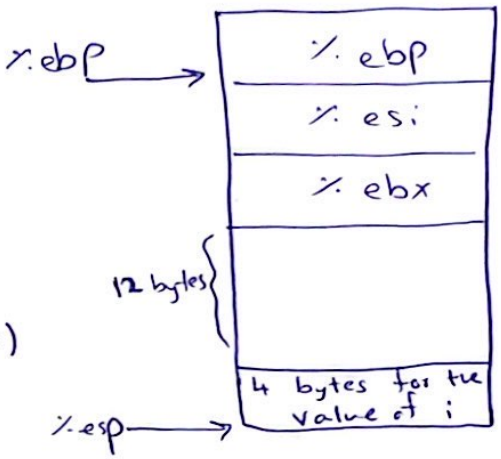



movl %ebx, (%esp)

# move the value of %ebx (i) to the memory location of the stack pointer to use it by factorial function

call _factorial

# calling the function → explained in page (7)




addl %eax, %esi

# %esi = %esi + %eax → result += factorial(i)

~~# i is~~ result

the value returned from the function

~~# addl~~ $1, %ebx

# %ebx = %ebx + 1 → i++

# then L8 body is completed it will go again to L7 and does comparison with 5 (see the code)

if i ≤ 5 goto L8

```
L8:
    movl  %ebx , (%esp)
    call _factorial
    addl  %eax, %esi  #result+= factorial(i)
    addl  $1 , %ebx   # i++
L7:
    cmpl  $5, %ebx  # i:5
    jle  L8
```

(6)

#.L : <5 then we jump to L8

# before we jump to explaing factorial function
# we explained LCO because it is an important part of the program
# this is extra information

.ascii "%d\0"

# include these characters in the data section of the program

.text

# starts a new section, the "text" section is the section in object files that stores code

.global _main

# tells the assembler that this is a global symbol and should be visible to the linker because other object files will use it (declares it as a global symbol)

.def _main ; .scl 2 ; .type 32 ; .endef

# .def defines a symbol _main

# .scl 2 set the storage class of the symbol to 2 (external symbol)
external symbol means that is defined or used in another module

# .type 32 indicates that _main is a function symbol (sets the type of the symbol)
number 32 indicates a function

# .endef : end definition

# reference : https://stackoverflow.com/questions/17794533/what-does-this-assembly-language-code-mean

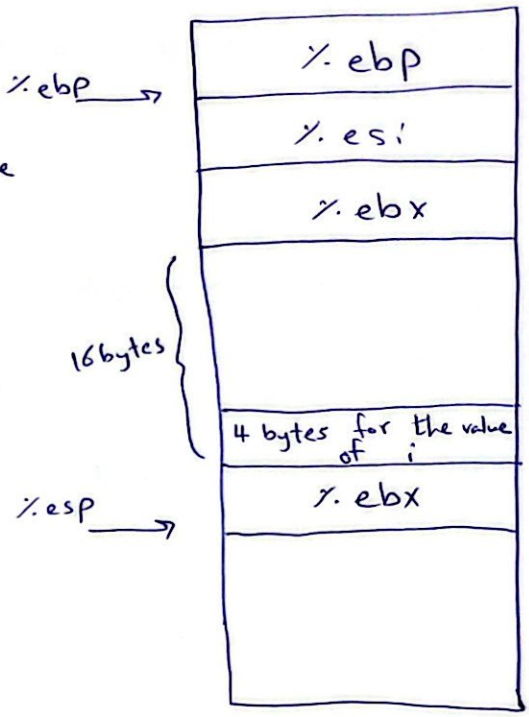
(7)

- factorial:

LFB12:

pushl %ebx

# decrement esp by 4 to push %ebx value

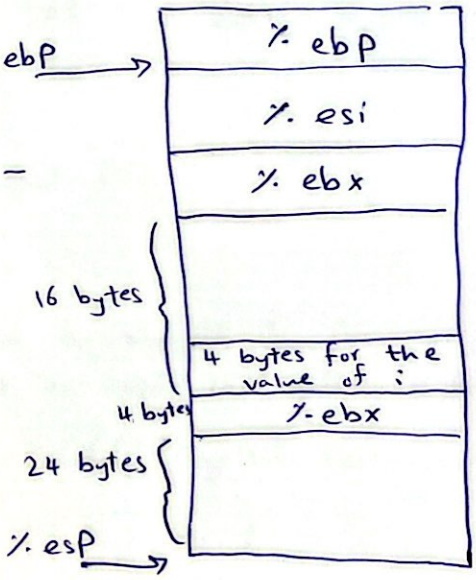



subl $24, %esp

# decrement %esp by 24

# %esp = %esp - 24

# allocating 24 bytes for local variables etc -




movl 32(%esp), %ebx

# ~~32~~ %ebx = (32 + %esp)

# value of the memory stored at (%esp+32) is the value of i

# this line is ~~equivalent to~~ setting the parameter of the function "num" to i (int num = i)

# num in %ebx and it took the value of i which was the argument of the function call

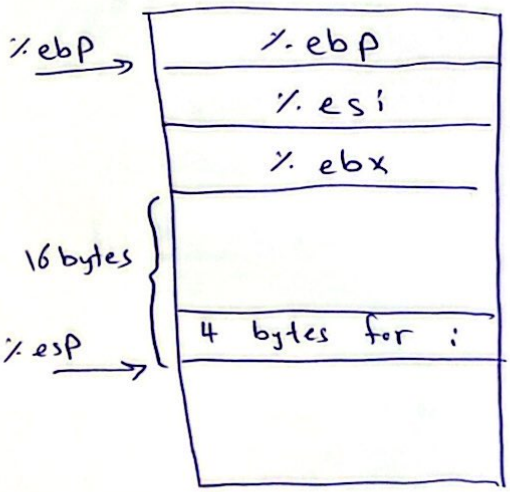
(8)

```
testl   %ebx, %ebx
# testl ands the two arguments and sets the condition codes
  without storing the result
# anding the same argument with itself to see if it equals zero
# result of anding ~~with~~ will only equal zero if %ebx value is zero

  jne  L5
# L5 explained in page (9)
# jne → jnz   ~ZF   jumps if ZF not equal zero
  so if it equals zero:

  movl   $1, %eax
# move immediate value 1 to %eax to be returned
# then we dig into L1
L1:
    addl   $24, %esp
  # increments the stack pointer by 24 which means freeing the
    24 allocated bytes

    popl   %ebx
  # increment %esp by 4
```




```
    ret
  # returns the value stored in %eax which is 1
  # ~~return~~ equivalent to return 1;
```

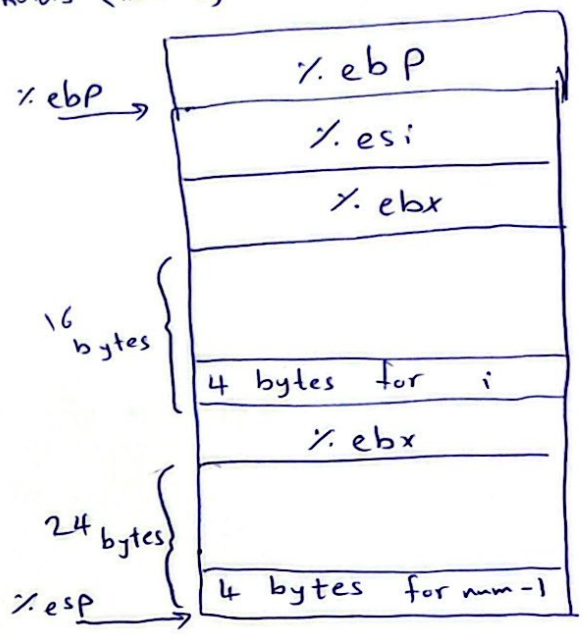
(9)

L5:

```
leal  -1(%ebx), %eax
```

# %eax = %ebx - 1

# o/o ebx represents num so now %eax holds (num-1)




```
movl  %eax, (%esp)
```

# move the value of %eax which is num-1 to the memory location of %esp to be used by the factorial function again

```
call  _factorial
```

# now factorial will use the 4 bytes for num-1 as its new "num"

```
imull  %ebx , %eax
```

# %eax = %eax * %ebx

result of the factorial (num-1) — num

# equivalent to num * factorial (num-1) and putting the value in %eax to be returned

```
jmp L1
```

# goto L1 to free the stack and return %eax value as explained in page 8