

# Chapter 6

## FFT Decimation in Frequency

Spring  
2009

© Ammar Abu-Hudrouss -Islamic  
University Gaza

### Decimation in Frequency

The discrete Fourier transform can be found using

$$X(k) = \sum_{n=0}^{N-1} x(n)W_N^{kn} \quad k = 0, 1, 2, \dots, N-1$$

Where  $N = 2, 4, 8, 16, \dots$  and

$$W_N = e^{-j2\pi/N}$$

$X(k)$  can be expressed as

$$X(k) = \sum_{n=0}^{N/2-1} x(n)W_N^{kn} + \sum_{n=N/2}^{N-1} x(n)W_N^{kn}$$

$$X(k) = \sum_{n=0}^{N/2-1} x(n)W_N^{kn} + \sum_{n=0}^{N/2-1} x(n+N/2)W_N^{k(n+N/2)}$$

$$X(k) = \sum_{n=0}^{N/2-1} x(n)W_N^{kn} + W_N^{kN/2} \sum_{n=0}^{N/2-1} x(n+N/2)W_N^{kn}$$

## Decimation in Frequency

But

$$W_N^{(N/2)k} = e^{2\pi k(N/2)/N} = e^{\pi k} = (-1)^k$$

Then

$$X(k) = \sum_{n=0}^{N/2-1} x(n)W_N^{kn} + (-1)^k \sum_{n=0}^{N/2-1} x(n+N/2)W_N^{kn}$$

If  $k = 2m$  or an even number

$$X(2m) = \sum_{n=0}^{N/2-1} (x(n) + x(n+N/2))W_N^{2mn}$$

$$X(2m) = \sum_{n=0}^{N/2-1} a(n)W_{N/2}^{mn}$$

## Decimation in Frequency

Noting That

$$W_N^{2mn} = e^{-j2mn \times 2\pi/N} = e^{-j2\pi mn/(N/2)} = W_{N/2}^{mn}$$

$$a(n) = x(n) + x(n+N/2)$$

Then  $X(2m)$  is  $N/2$ -point DFT for  $a(n)$

If  $k = 2m+1$  (odd number) and using the same method

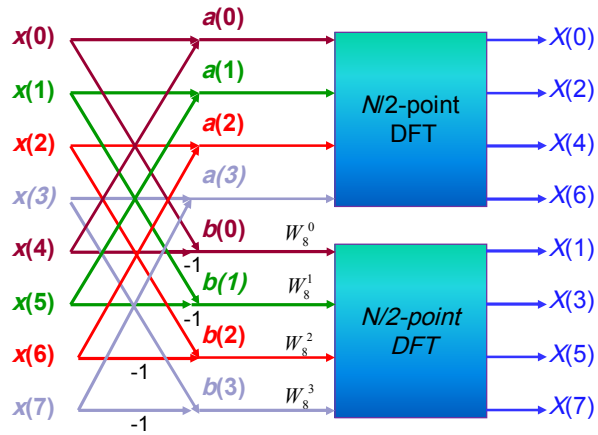
$$X(2m+1) = \sum_{n=0}^{N/2-1} (x(n) - x(n+N/2))W_N^n W_N^{2mn}$$

$$X(2m+1) = \sum_{n=0}^{N/2-1} (b(n)W_N^n)W_{N/2}^{mn}$$

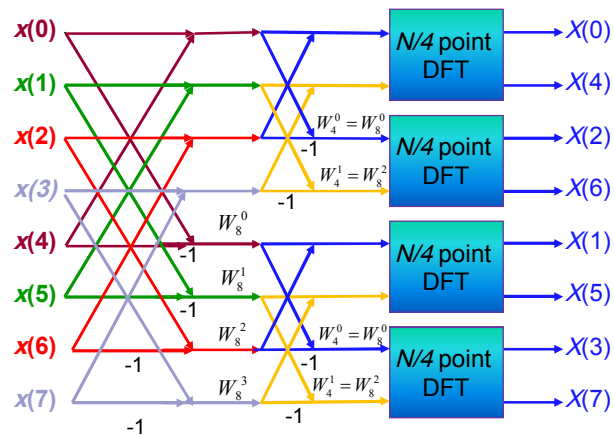
$$b(n) = x(n) - x(n+N/2)$$

$X(2m+1)$  is  $N/2$ -point DFT for  $b(n)W_N^n$

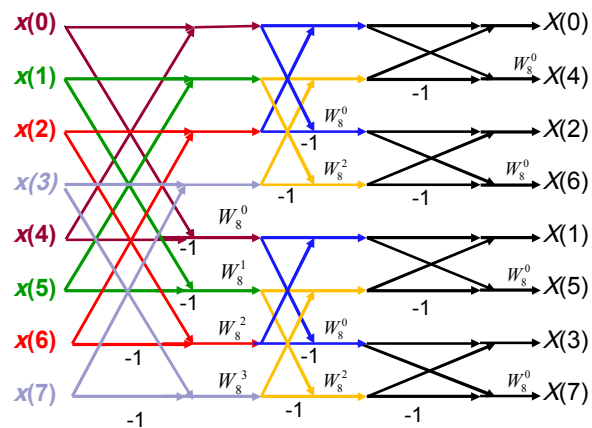
## Decimation in Frequency



## Decimation in Frequency



## Decimation in Frequency



Digital Signal Processing

Slide 9

## Decimation in Frequency

Using the previous algorithm, the complex multiplications needed is only 12. While using the normal DFT would require 64 complex multiplications

In general

Complex multiplication of DFT is:  $N^2$

Complex multiplication of FFT is  $(N/2) \log_2(N)$

If  $N = 1024$

Complex multiplication of DFT is: 1,048,576

Complex multiplication of FFT is: 5,120

Digital Signal Processing

Slide 10

## Decimation in Frequency

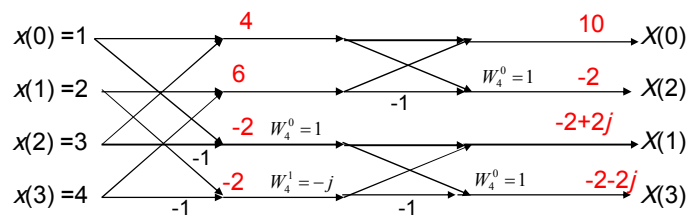
Index mapping for Fast Fourier Transform

Input Data index $n$	Index Bits	Reversal Bits	Output data index $k$
0	000	000	0
1	001	100	4
2	010	010	2
3	011	110	6
4	100	001	1
5	101	101	5
6	110	011	3
7	111	111	7

## Decimation in Frequency

**Example** Given a sequence  $x(n)$  where  $x(0) = 1$ ,  $x(1) = 2$ ,  $x(2) = 3$ ,  $x(3) = 4$  and  $x(n) = 0$  elsewhere, find DFT for the first four points

**solution**



## Inverse Fourier Transform

The inverse discrete Fourier transform can be found using

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) W_N^{-kn} \quad n = 0, 1, 2, \dots, N-1$$

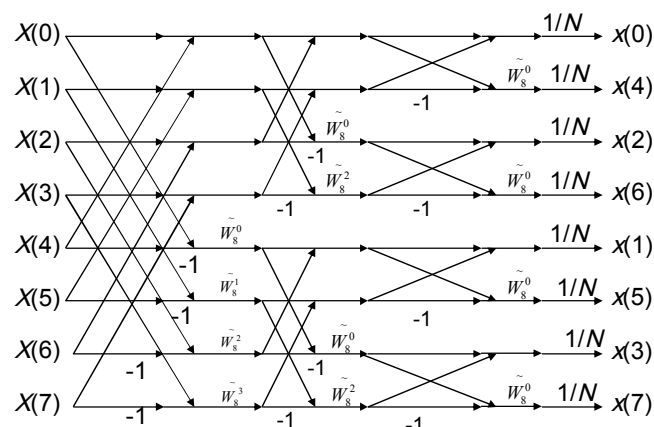
Which can be expressed as

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) \tilde{W}_N^{kn} \quad n = 0, 1, 2, \dots, N-1$$

where  $\tilde{W}_N^{kn} = W_N^{-kn}$  is called the twiddled factor

We can see that the difference between the inverse discrete Fourier and forward Fourier transform is the twiddled factor and the division by  $1/N$

## Inverse Fourier Transform



## Inverse Fourier Transform

**Example** Given a sequence  $X(n)$  given in the previous example.  
Find the IFFT using decimation in frequency method

**solution**

