

Presentations

- EE295

Why read your paper?

- With support
 - What is the paper topic
 - The benefit of reading the paper
 - How will it change your life


Hierarchy of understanding

- Images
- Diagrams
- Charts/graphs
- Text
- Tables

Simple rules to follow

- 4 x 4 (four words, four lines)
- Phrases ★
 - Not speaker notes
- Single message per slide
- Builds to a conclusion

Make it readable

- Minimum 36 point font
- 50 point font
- Good contrast
 - Projectors have **bad** color
- Highlight important points 

Appearance

- Body language
 - Panic ?
 - Actions ?
- Dress
 - Casual professional



Tone and emotion

- Most positive == most important
- Finish strong and positive

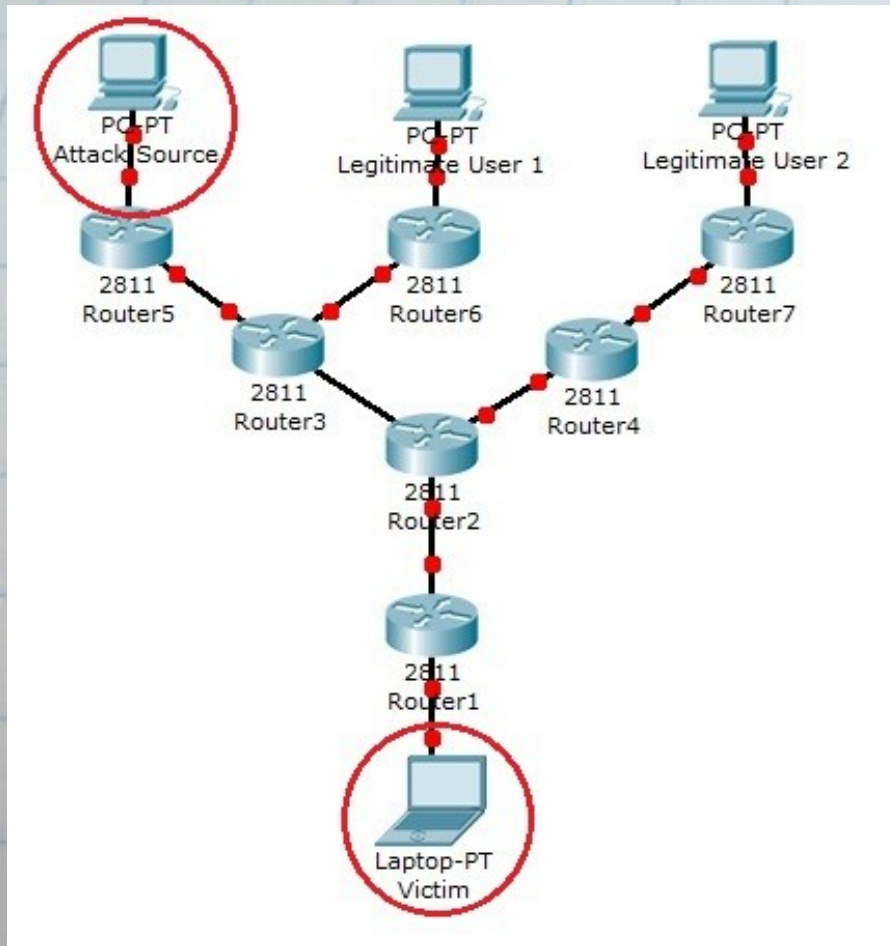
Rehearse

- Not natural
- Time factors
- Message factors
- Reduces panic

Slides from actual Presentations

Some Critical Review

DoS Attack



- ✓ Malicious attempt by an attacker to flood a network with packets.
- ✓ Overwhelm its data handling capacity such that the system becomes inaccessible.
- ✓ DoS attack is initiated by a single host.

Algorithms

For each packet 'p' from Source 'x' to Destination 'y':

if (logging entry not present)

 write address of router's interface through which the packet is going to be forwarded;

 set flag = 0;

else

 if (flag == 1)

 block the packet going through this router

 else

 do nothing

Logging Algorithm at Router

Receive the incoming packets;

if (rate of incoming packets goes beyond acceptable limit)

 send a blocking message with flag = 1 to the router from where the packet was received;

Traceback algorithm at the Victim computer

if (message is a blocking message)

 for the specified route (source-destination), set flag bit = 1;

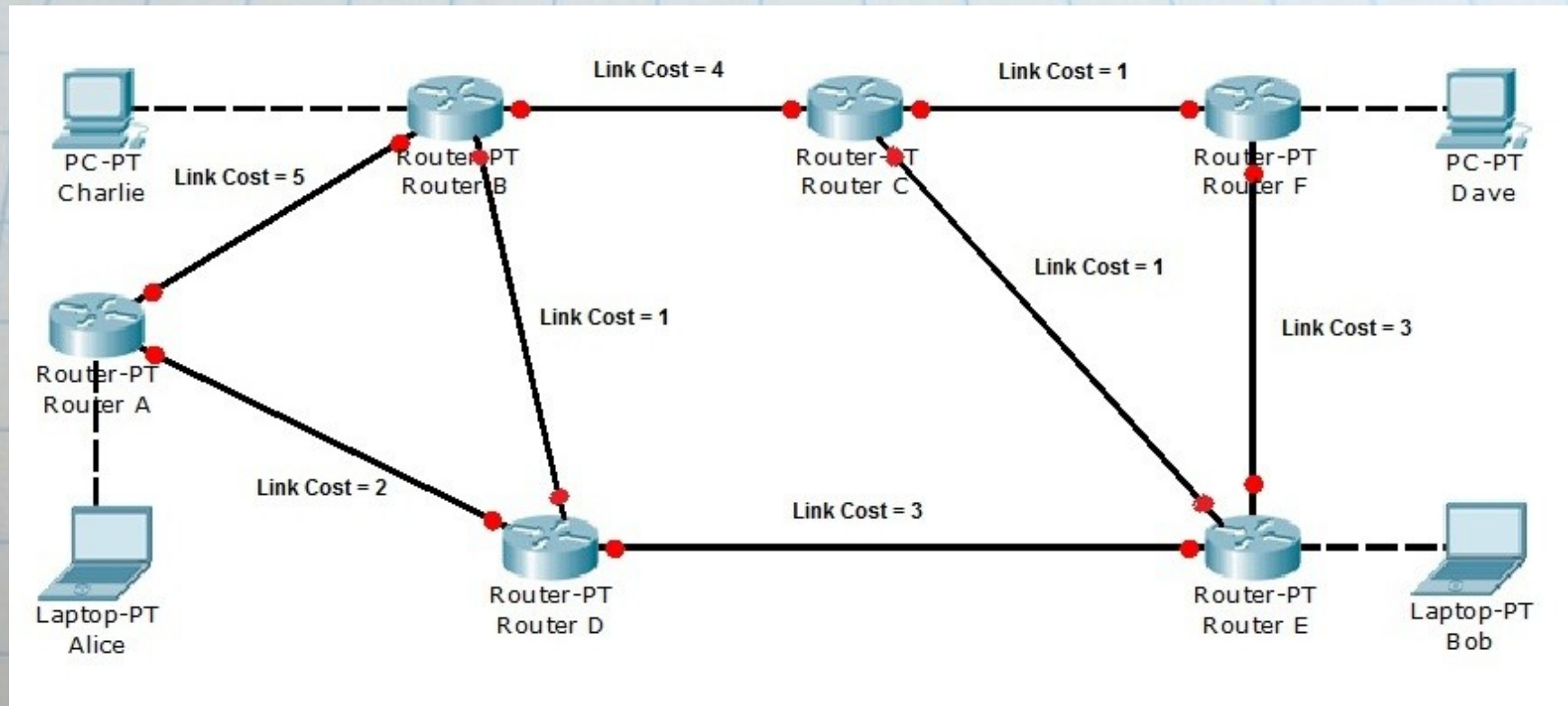
 forward the blocking message to the next router (referring to the log table for the address)

Traceback algorithm at the Router

Proposed IP Traceback Mechanism

- Efficiently traces the attack source without any false positives.
- Simpler, faster and effective way to tackle DoS attacks.
- Blocks all the routers along the attack path.
- Low overhead on routers as most the information used is already available with the routers.
- Only 33 bits of additional data is required to be logged into router.
- Allows tracing of multiple sources of attack.
- Low detection delay.

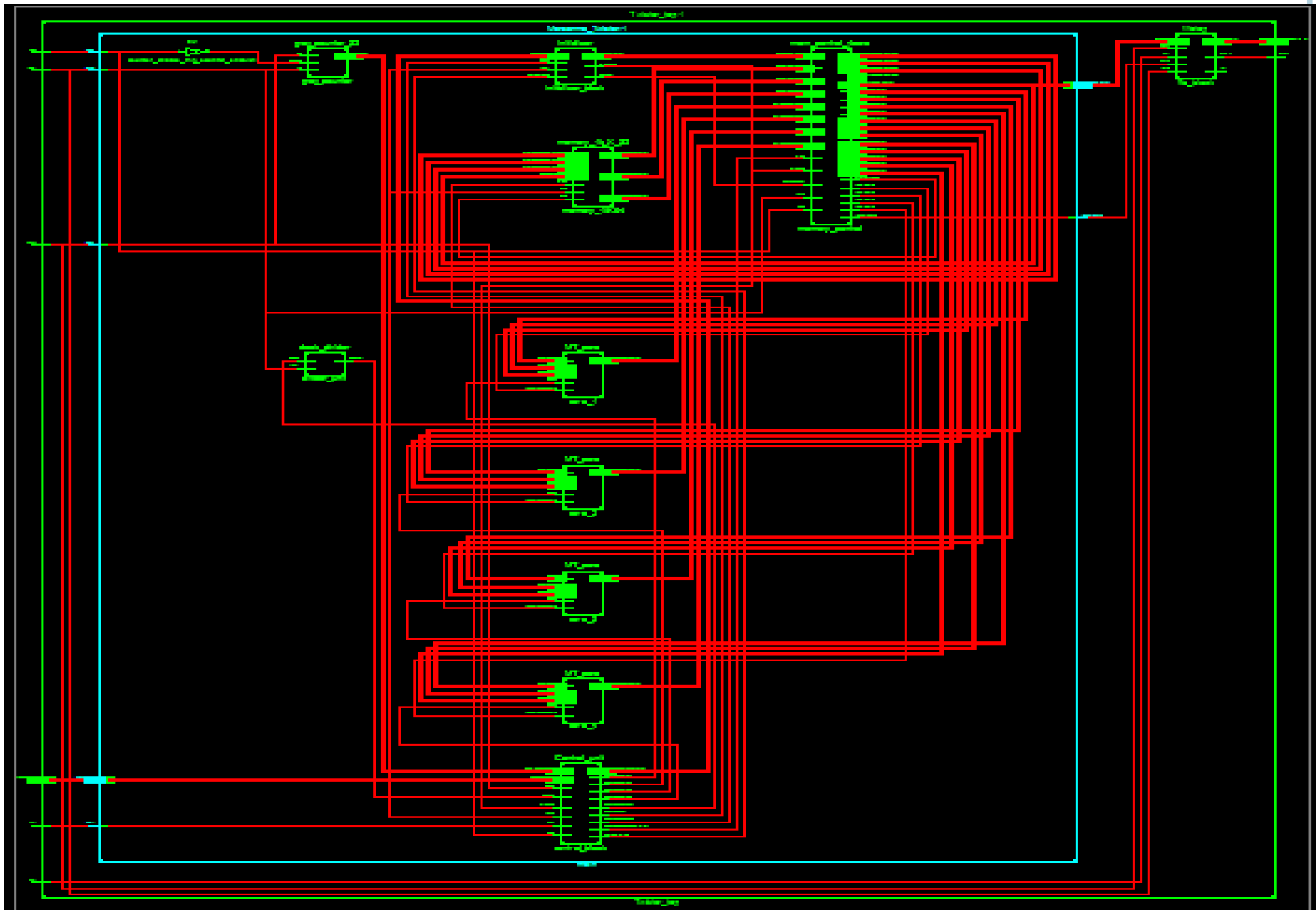
Attack Network



Randomness Results

Randomness Test	Measure	Ideal value	Software	Hardware
Chi-Square Test	Percentage	20% to 80%	72.44%	29.77%
Arithmetic mean	Mean value	127.500	127.5066	127.4789
Monte-carlo value of Pi	Pi value	3.141574059	3.141574059	3.141806341
Serial Correlation coefficient	Correlation coefficient	0.0 = uncorrelated	-0.000433	0.000304
Birthday Spacing test	p-value	[0,1]	0.083175	0.440985
The overlapping-5 permutation test	p-value	[0,1]	0.366844	0.852751
The Bit stream Test	p-value	[0,1]	0.58389	0.88988
OPSO, OQSO and DNA	p-value	[0,1]	0.3087	0.714
Parking lot test	p-value	[0,1]	0.620989	0.725731
Minimum distance test	p-value	[0,1]	0.003322	0.861518
The 3D Spheres test	p-value	[0,1]	0.622752	0.899760
The overlapping sums test	p-value	[0,1]	0.645970	0.322378
Runs test	p-value	[0,1]	0.550942	0.499031
Craps Test	p-value	[0,1]	0.386168	0.630164

RTL Schematic



```

/**
 * The Packet that is currently being transmitted on the channel (or last
 * packet to have been transmitted on the channel if the channel is
 * free.)
 */
Ptr<Packet> m_currentPkt;

/**
 * Device Id of the source that is currently transmitting on the
 * channel. Or last source to have transmitted a packet on the
 * channel, if the channel is currently not busy.
 */
uint32_t      m_currentSrc;

/**
 * Current state of the channel
 */
WireState      m_state;
bool           m_collision;
};

} // namespace ns3

namespace ns3 {

class Queue;
class RsmaChannel;
class ErrorModel;

/**
 * \defgroup csma RsmaNetDevice
 *
 * This section documents the API of the ns-3 csma module. For a generic functional description, please refer to the ns-3 manual.
 */

/**
 * \ingroup csma
 * \class RsmaNetDevice
 * \brief A Device for a Csma Network Link.
 *
 * The Csma net device class is analogous to layer 1 and 2 of the
 * TCP stack. The NetDevice takes a raw packet of bytes and creates a
 * protocol specific packet from them.
 */
class RsmaNetDevice : public NetDevice

```


Comparison with other Diversities.

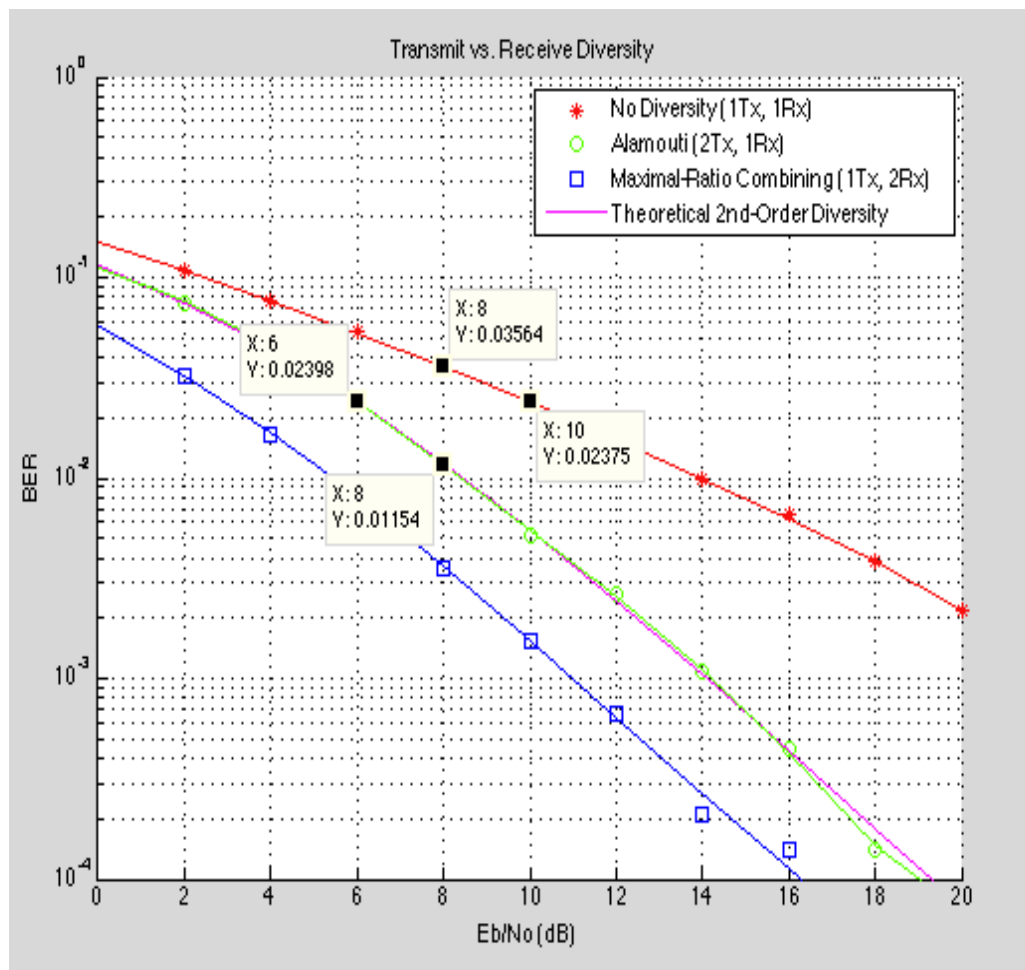


Figure 6[4]

- No Diversity Slope of BER curve is -1.
- Maximal Ratio diversity BER curve slope -2.
- Alamouti's Diversity curve BER slope of -2 but at a loss of 3dB.

Slides from prior proposal presentations

Overview of current protocols

- 1) HTTP - Header size of Hyper Text Transfer Protocol (HTTP) is 400 Bytes to 8KB^[3]. Fig^[4]
- 2) Continuous transmission of HTTP messages

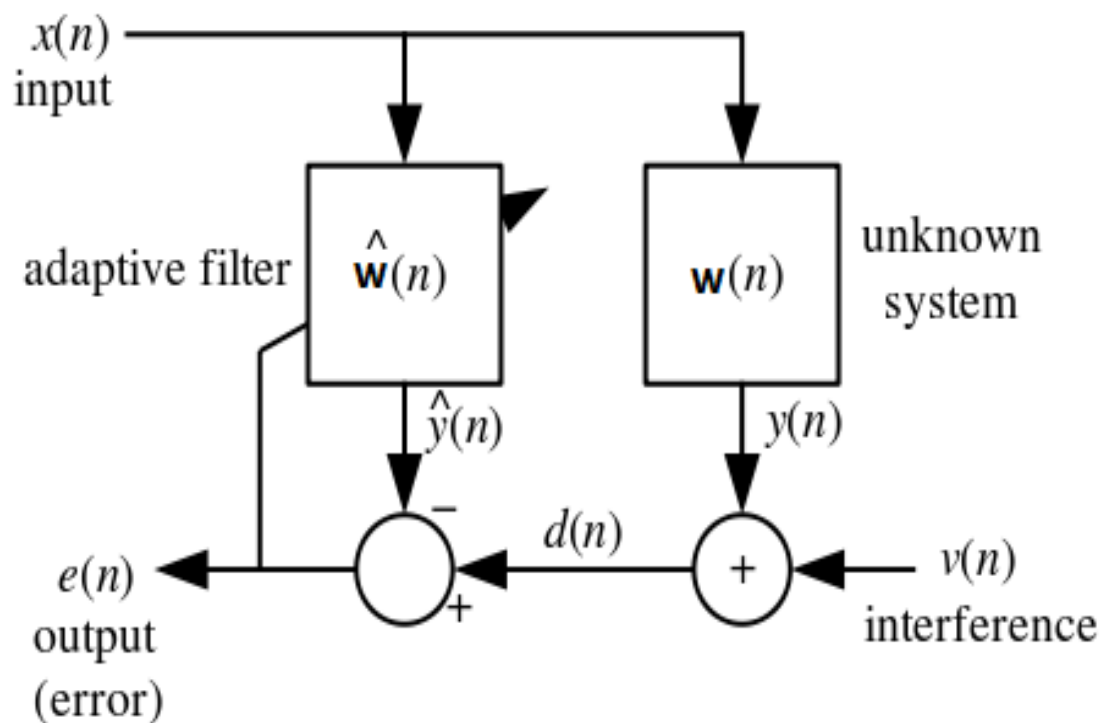
Example 2: HTTP request header

```
GET /PollingStock//PollingStock HTTP/1.1
Host: localhost:8080
User-Agent: Mozilla/5.0 (Windows; U; Windows NT 5.1; en-US; rv:1.9.1.5)
Gecko/20091102 Firefox/3.5.5
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8
Accept-Language: en-us
Accept-Encoding: gzip,deflate
Accept-Charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7
Keep-Alive: 300
Connection: keep-alive
Referer: http://www.example.com/PollingStock/
Cookie: showInheritedConstant=false;
showInheritedProtectedConstant=false; showInheritedProperty=false;
showInheritedProtectedProperty=false; showInheritedMethod=false;
showInheritedProtectedMethod=false; showInheritedEvent=false;
showInheritedStyle=false; showInheritedEffect=false
```

Example 3: HTTP response header

```
HTTP/1.x 200 OK
X-Powered-By: Servlet/2.5
Server: Sun Java System Application Server 9.1_02
Content-Type: text/html; charset=UTF-8
Content-Length: 21
Date: Sat, 07 Nov 2009 00:32:46 GMT
```

Mathematics



$$y(n) = \sum_{i=0}^{N-1} w(n)x(n-i) = \mathbf{w}^T(n)\mathbf{x}(n)$$

$$e(n) = d(n) - y(n)$$

$$\mu(n) = \frac{1}{\mathbf{x}^T(n)\mathbf{x}(n)}$$

$$\mathbf{w}(n+1) = \mathbf{w}(n) + \mu(n)e(n)\mathbf{x}(n)$$

RLS Noise Cancellation Circuit

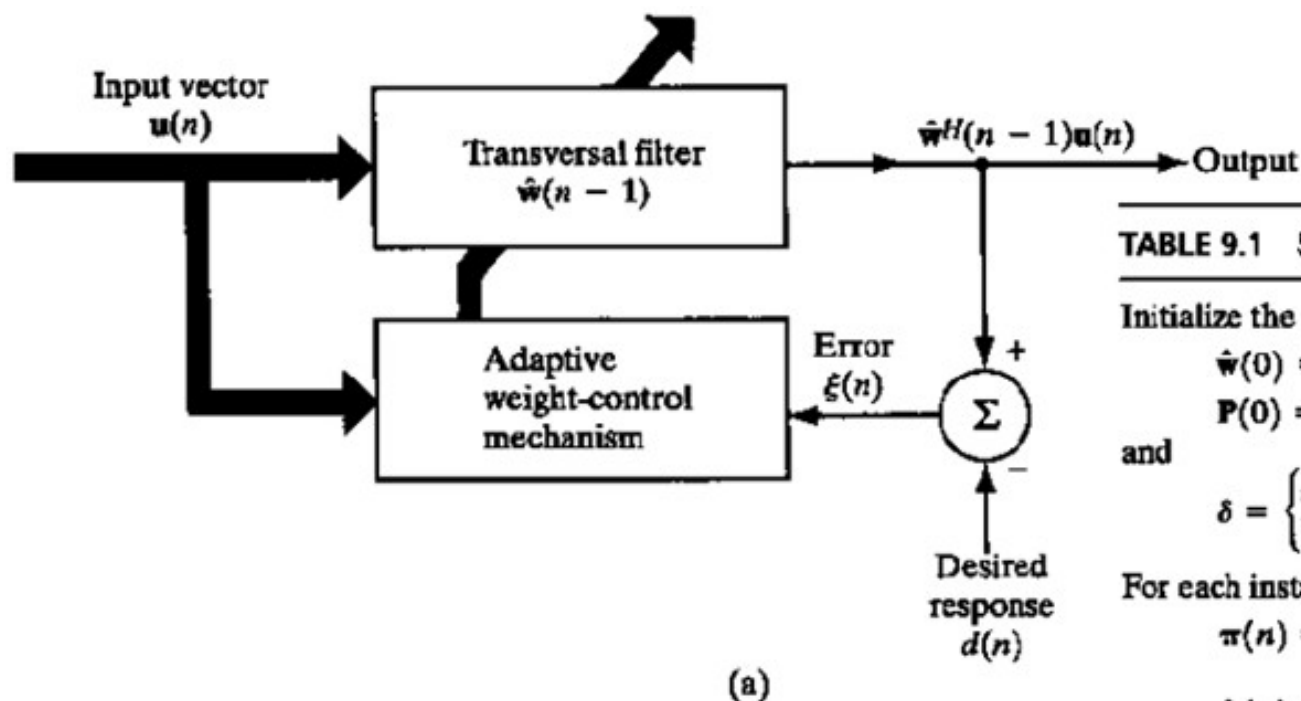


TABLE 9.1 Summary of the RLS Algorithm

Initialize the algorithm by setting

$$\hat{\mathbf{w}}(0) = \mathbf{0},$$

$$\mathbf{P}(0) = \delta^{-1} \mathbf{I},$$

and

$$\delta = \begin{cases} \text{small positive constant for high SNR} \\ \text{large positive constant for low SNR} \end{cases}$$

For each instant of time, $n = 1, 2, \dots$, compute

$$\boldsymbol{\pi}(n) = \mathbf{P}(n-1)\mathbf{u}(n),$$

$$\mathbf{k}(n) = \frac{\boldsymbol{\pi}(n)}{\lambda + \mathbf{u}^H(n)\boldsymbol{\pi}(n)},$$

$$\xi(n) = d(n) - \hat{\mathbf{w}}^H(n-1)\mathbf{u}(n),$$

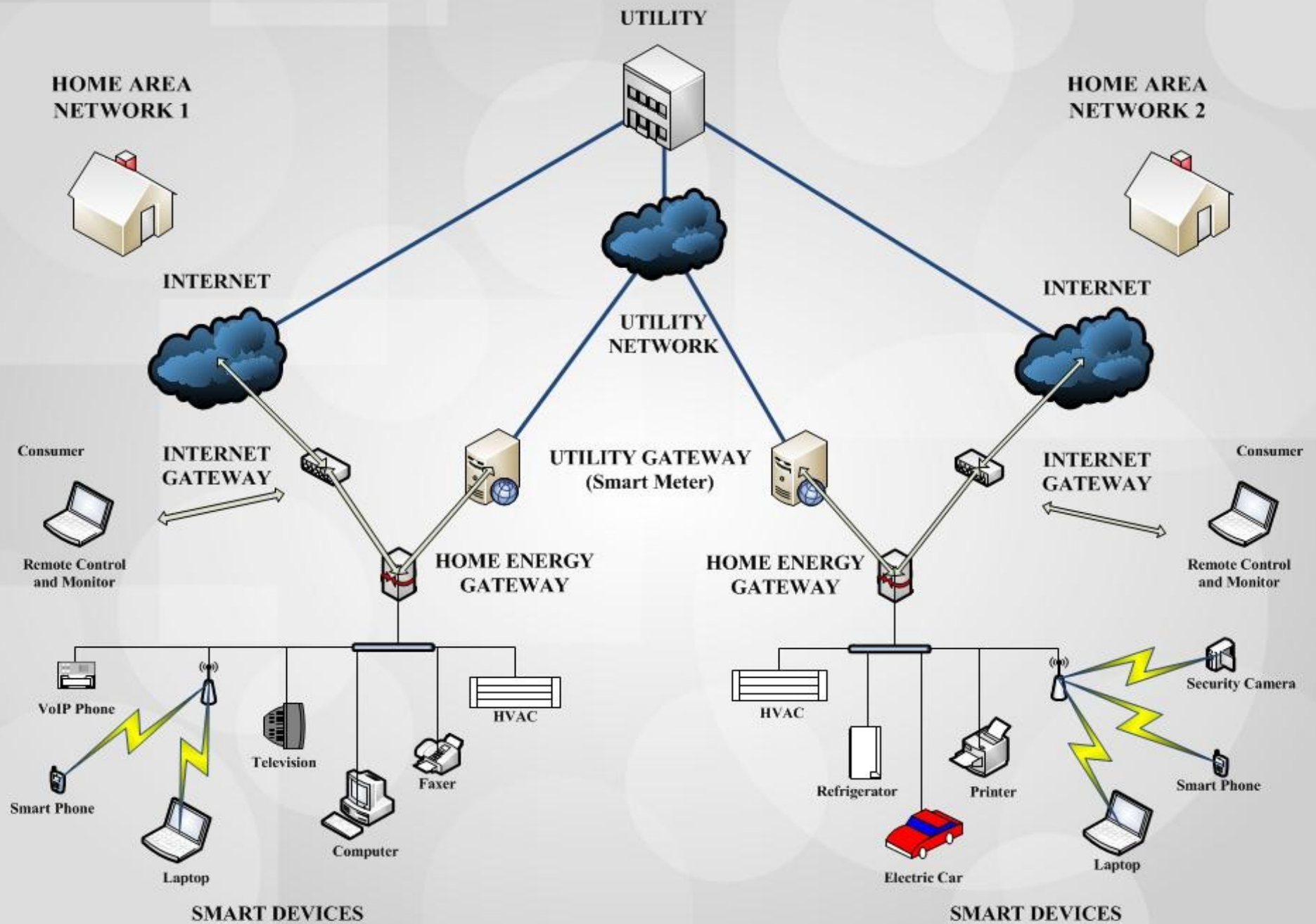
$$\hat{\mathbf{w}}(n) = \hat{\mathbf{w}}(n-1) + \mathbf{k}(n)\xi^*(n),$$

and

$$\mathbf{P}(n) = \lambda^{-1}\mathbf{P}(n-1) - \lambda^{-1}\mathbf{k}(n)\mathbf{u}^H(n)\mathbf{P}(n-1).$$

Traditional Method - Low pass Filters

- In image or video processing, one of the most prominent problems is degradation by noise and to overcome this issue, the images are passed through a low pass filter that helps reduce visible noise.
- Traditionally all the Low Pass filter uses the principle of Averaging the neighborhood of pixels and applying this averaged value as input.
- The problem occurs when large numbers of iterations are performed and **edges regions** are also included in the averaging process.
- This results in degradation of quality of image and makes it look very blurred.



SCALABILITY ON A LOAD BALANCER

- Achieved in applications by creating a server cluster or a server farm.
- These are collections of servers that distribute the workload amongst themselves
- The solution is designed to make adding more servers easy and efficient in order to handle increasing workloads
- This is far more preferable to adding more resources to a single server for reasons of cost, reliability and efficiency