

V L Mahesh (CS16M046)  
Muhammad Arsath K F (CS16S035)



# **Journey of MAD**



# What does it stand for?

- Is it the creator's names?
- Is it the names of the processes involved in it?
- Is it a reference in pop culture?



# What does it stand for?

Currently Nothing

(But, we want it to mean something)



# Presentation Agenda

**M**

**A**

**D**



# Presentation Agenda

**M**istakes

**A**pproach

**D**eployment



# Mistakes



# Mistakes

## Goals

- Complex (imitating/inspired by nature)
- Fast
- Solid (no compromise)
- <10 rounds
- Grand, Unique, Fresh



# Mistakes

- Can we do it without Sboxes?  
(like SHACAL-2)



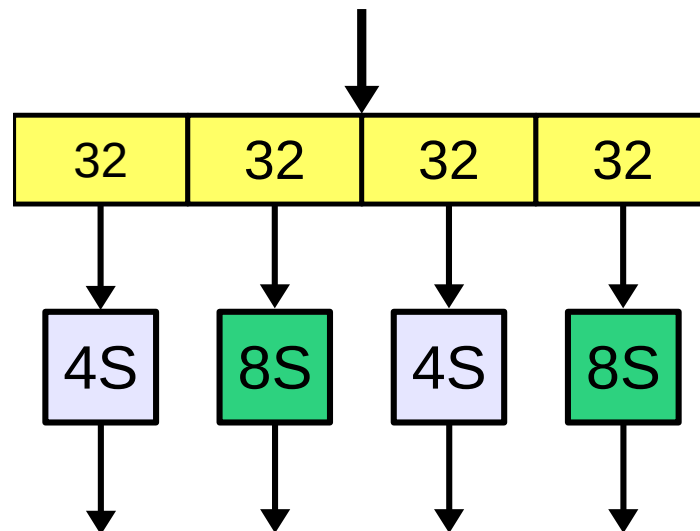


# Mistakes

- Can we do it without Sboxes? (Obviously Not - MtoQ)  
(like SHACAL-2)
- Complicated S-layer, Simple P-layer
- Sbox Selection

# Mistakes | Complicated S-layer

- Different patterns and sizes of Sboxes
- 4x4, 8x8 and a combination of both
- 4-8-4 blocks
- Groups of 4x4 sboxes and 8x8 sboxes

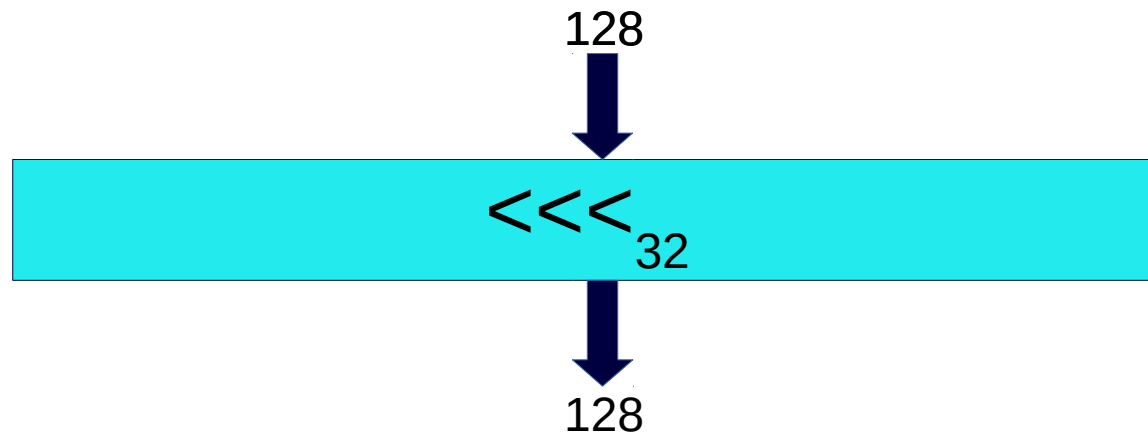


# **Mistakes** | Complicated S-layer

- Different patterns and sizes of Sboxes
- 4x4, 8x8 and a combination of both
- 4-8-4 blocks
- Groups of 4x4 sboxes and 8x8 sboxes
- Made analysis very difficult and error prone

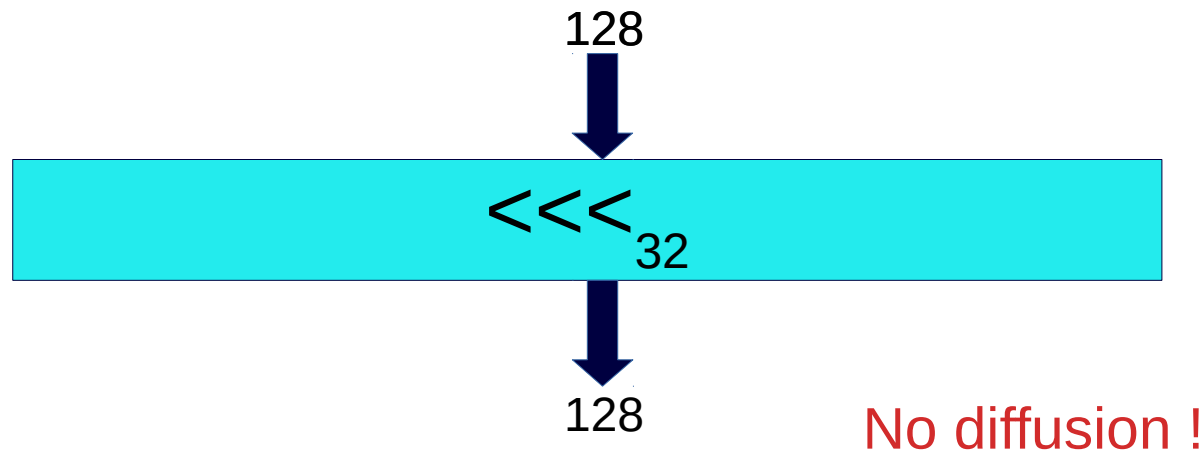
# Mistakes | Simple P-layer

- What is wrong with this? (In conjunction with previous layer)



# Mistakes | Simple P-layer

- What is wrong with this? (In conjunction with previous layer)



# Mistakes | Sbox Selection

- First, we used the **PRESENT** Sbox

$x$	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
$S[x]$	C	5	6	B	9	0	A	D	3	E	F	8	4	7	1	2

# Mistakes | Sbox Selection

- First, we used the **PRESENT** Sbox
- We got into a problem while doing LC
  - It's output bit 0 had a linear dependence with input bits {0,3}

Algorithm & Source	S-Box 0123456789ABCDEF	Bit 0		Bit 1		Bit 2		Bit 3	
		LS	deg	LS	deg	LS	deg	LS	deg
Lucifer S0 [37]	CF7AEDB026319458	{}	3	{}	3	{}	3	{}	3
Lucifer S1 [37]	72E93B04CD1A6F85	{}	3	{}	3	{}	3	{}	3
Present [9]	C56B90AD3EF84712	{0,3}	2	{}	3	{}	3	{}	3



# Approach





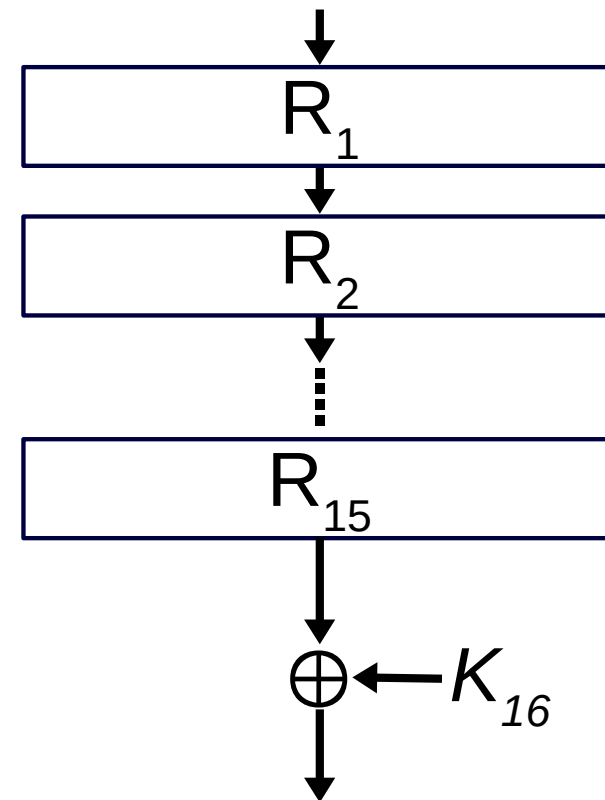
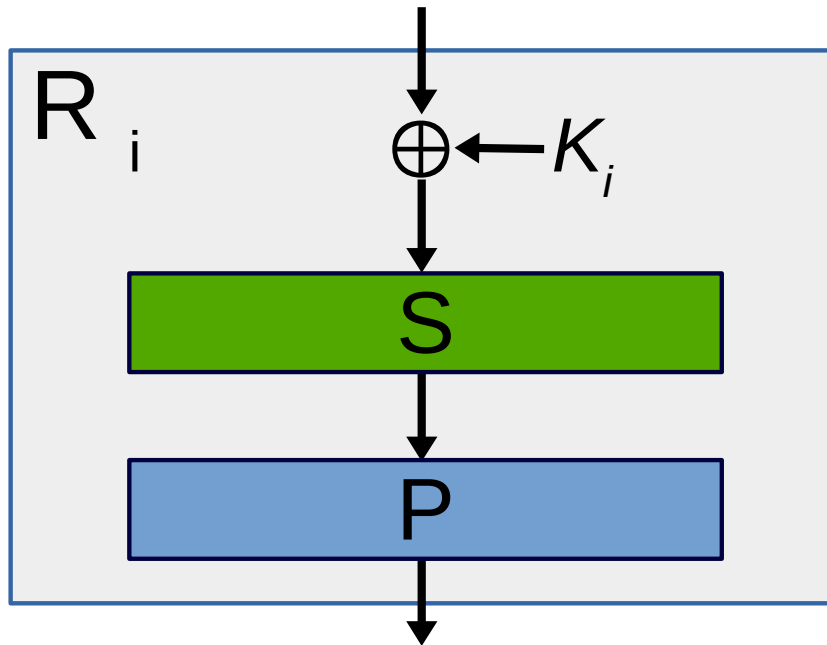
# Approach

## New Goals

- Simple
- Efficient
- Scalable
- $<20$  rounds
- Learn something

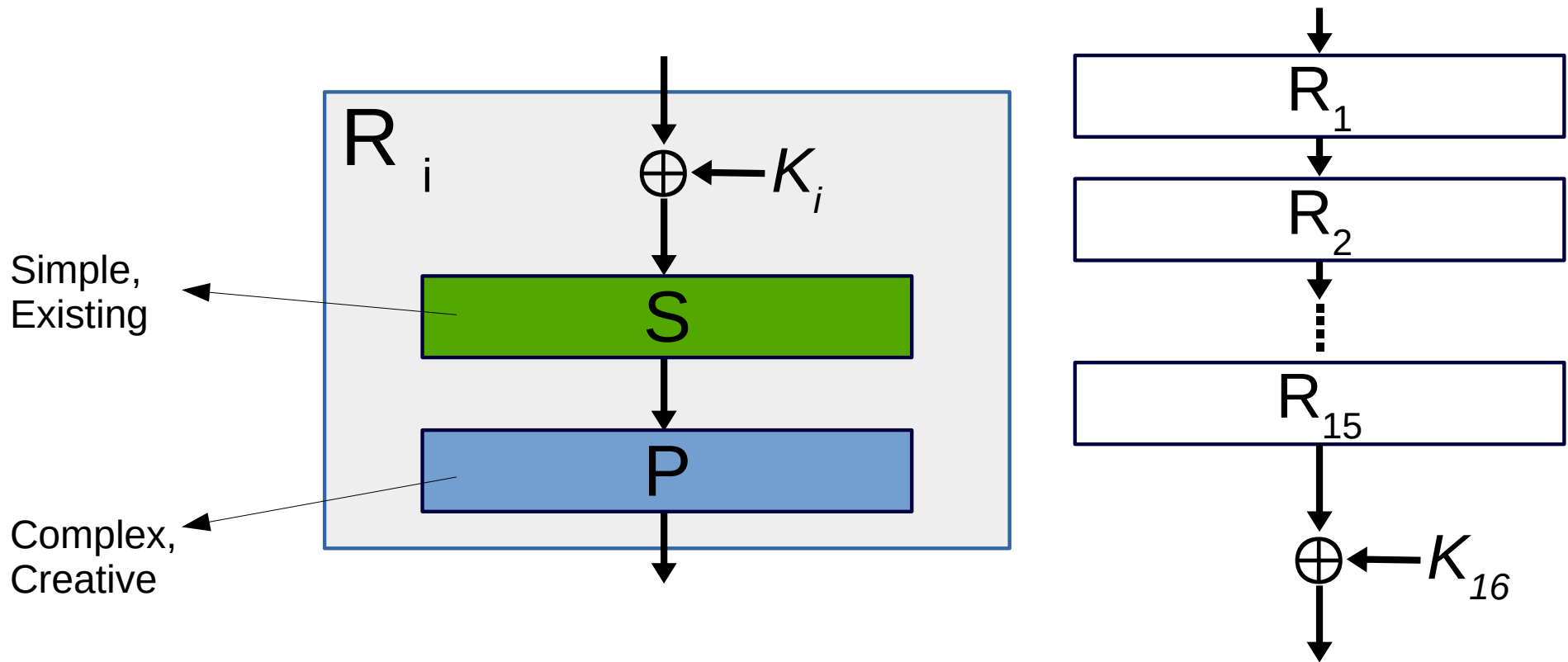
# Approach | General Structure

- Straightforward SPN



# Approach | General Structure

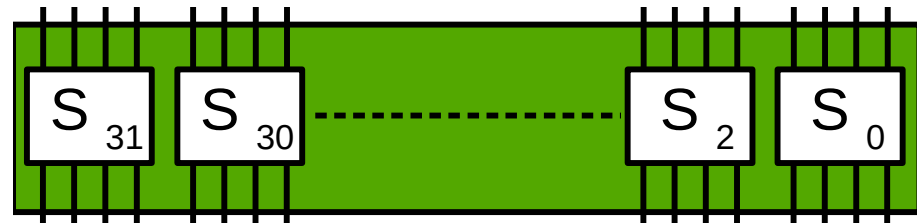
- Straightforward SPN



# Approach | S-layer

- Extremely Simple
- Uses the **SERPENT S3** Sbox

- 4x4
- Said to be Golden
- Very nice properties



Algorithm & Source	S-Box 0123456789ABCDEF	BN #	DC		LC		Bit 0		Bit 1		Bit 2		Bit 3	
			$p$	$n_d$	$\epsilon$	$n_l$	LS	deg	LS	deg	LS	deg	LS	deg
Serpent S0 [1]	38F1A65BED42709C	3	$1/4$	24	$1/4$	36	{}	3	{}	3	{}	3	{1,2}	2
Serpent S1 [1]	FC27905A1BE86D34	3	$1/4$	24	$1/4$	36	{}	3	{}	3	{2,3}	2	{}	3
Serpent S2 [1]	86793CAFD1E40B52	3	$1/4$	24	$1/4$	36	{1,3}	2	{}	3	{}	3	{}	3
Serpent S3 [1]	0FB8C963D124A75E	3	$1/4$	18	$1/4$	32	{}	3	{}	3	{}	3	{}	3

(We did calculate the above properties)

# Approach | P-layer

- Concept of EO Shuffle

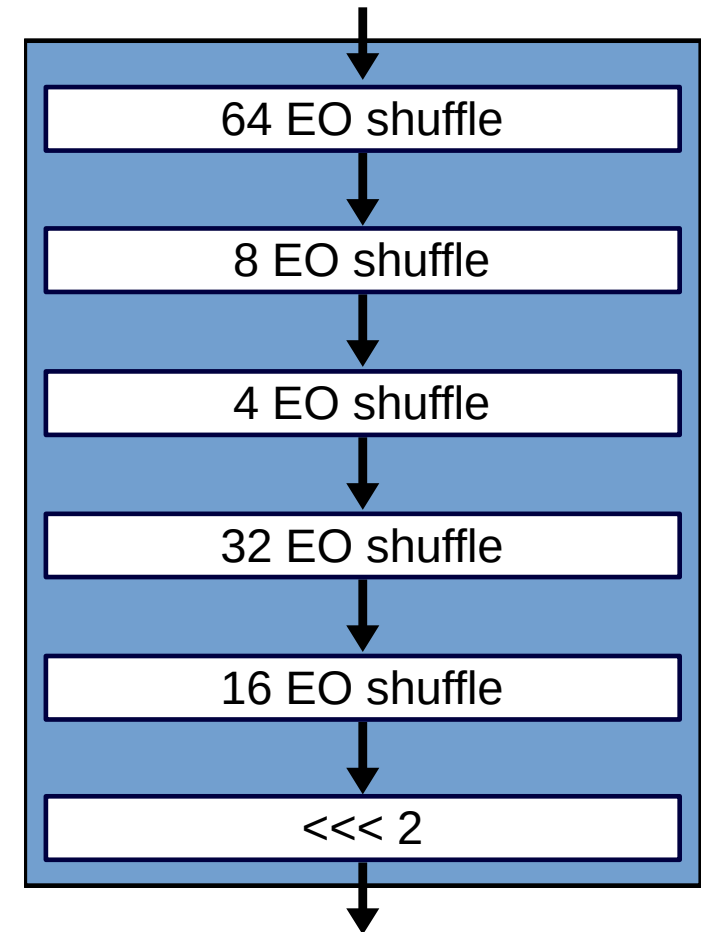
EO

$$\begin{aligned} c &= (a_{\text{odd}} \gg 1) \mid b_{\text{odd}} \\ d &= (b_{\text{even}} \ll 1) \mid a_{\text{even}} \end{aligned}$$

EO<sup>-1</sup>

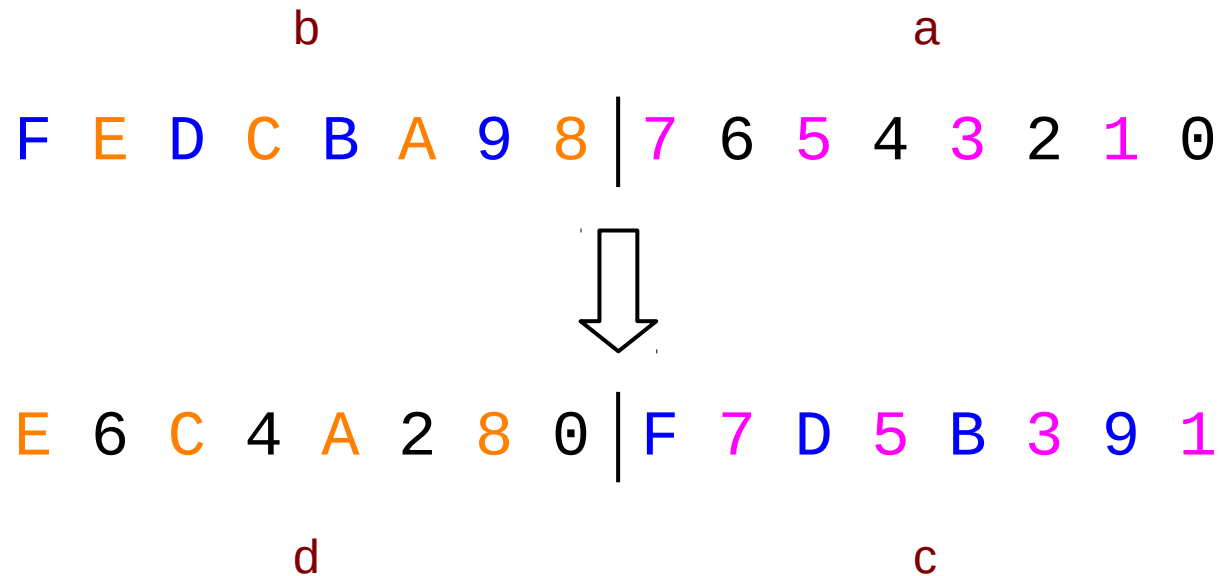
$$\begin{aligned} a &= (c_{\text{even}} \ll 1) \mid d_{\text{even}} \\ b &= (d_{\text{odd}} \gg 1) \mid c_{\text{odd}} \end{aligned}$$

- We tried different sequences
- More inversions → Better randomness



# Approach | EO Shuffle

8 EO Shuffle on a 16-bit number



- Odd & Even bits are extracted using bit masks

# Approach | Analysis

- Achieves Resistance to both Linear & Differential Cryptanalysis in 10 rounds
- Calculated Cumulative Bias & Propagation Ratios respectively (just like in the Stinson Textbook)
- Since Serpent S3 only had  $\frac{1}{4}$  &  $\frac{1}{2}$  in LAT/DDT, we could easily calculate them w/o Multiprecision Libs.
- But, we've used 15 rounds to be resistant to even advanced cryptanalysis techniques, just in case!



# Deployment





# Deployment

## Implementation Guidelines

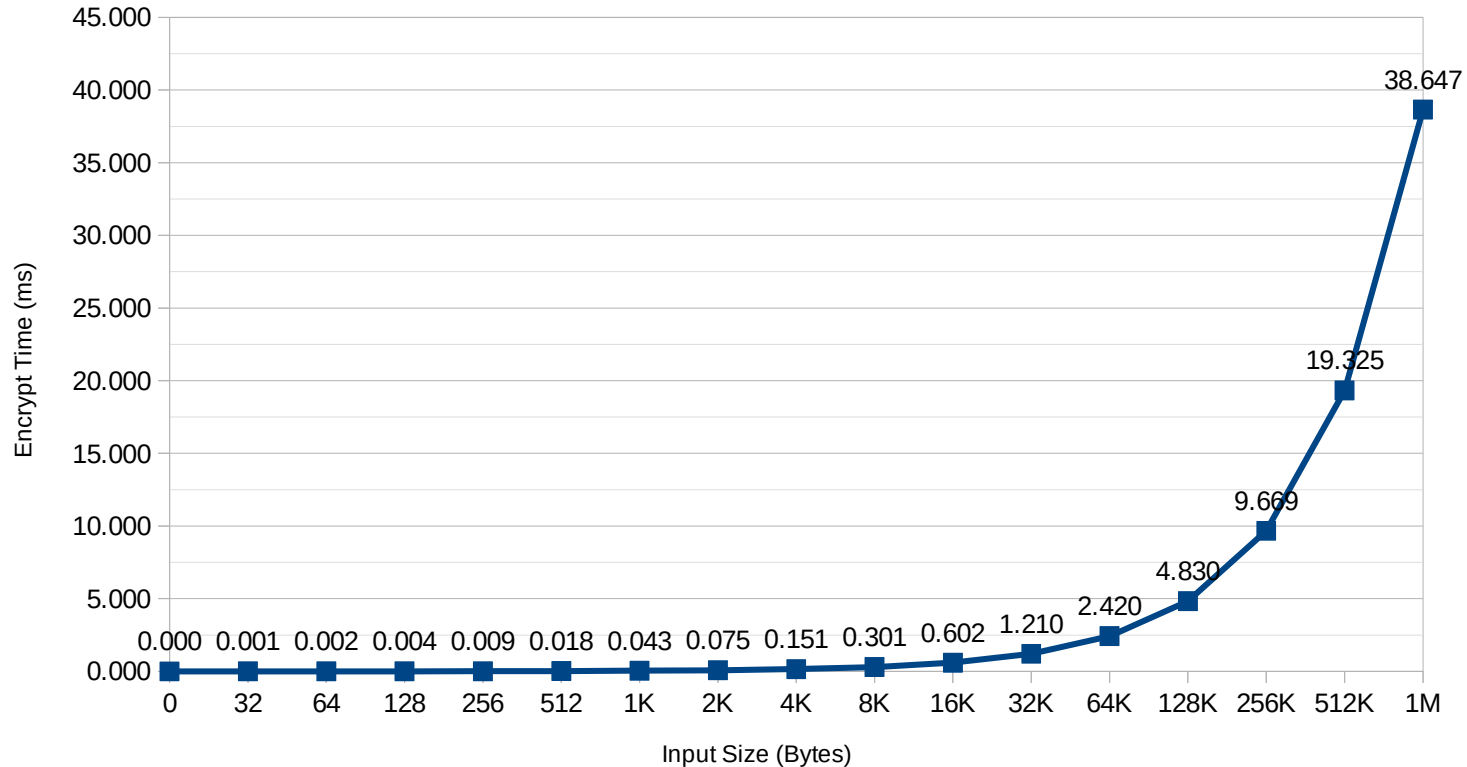
- Use previously learned concepts from different areas (Cryptography, Computer Architecture, Compilers)
- Readable
- Portable
- Auto-Optimized
- Efficient Hardware Implementation

# Deployment | At-a-glance

- Stores the 128-bit block as two `uint64_t`
- Relies on a Compiler's matured toolset, rather than our own experience in assembly code optimization
- Separate set of functions for encrypt and decrypt to avoid branches (we need them because SPN)
- Cache Efficient  
(only 1 4x4 Sbox/R-Sbox = 32 Bytes)

# Deployment | Performance

Encryption Time with Size (on a 7th Gen i7 Desktop CPU w/ -O3)



- As can be seen, 38.6 ms » ~26 MB/s
- Pretty good for a first implementation



# Deployment | Details

How did we achieve this?

- No expensive operations (only shifts, and, or)



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- Auto-vectorized (SSE, AVX, etc...)



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- No expensive operations (only shifts, and, or)
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- Constant size loops (Loop Unrolling)
- Auto-vectorized (SSE, AVX, etc...)

All of these without losing portability, readability