

P - NP PROBLEMS: MOTIVATION

ANALYSIS OF ALGORITHMS AND COMPLEXITY

THE \$1M QUESTION

**The Clay Mathematics Institute
Millenium Prize Problems**

- 1. Birch and Swinnerton-Dyer Conjecture**
- 2. Hodge Conjecture**
- 3. Navier-Stokes Equations**
- 4. P vs NP**
- 5. Poincaré Conjecture**
- 6. Riemann Hypothesis**
- 7. Yang-Mills Theory**

THE P VERSUS NP PROBLEM

**Is perhaps one of the biggest open problems
in computer science (and mathematics!)
today.**

(Even featured in one of TV shows)

But what is the P-NP problem?

Sudoku

2			3		8		5	
		3		4	5	9	8	
		8			9	7	3	4
6		7		9				
9	8						1	7
				5		6		9
3	1	9	7			2		
	4	6	5	2		8		
	2		9		3			1

3x3x3

Sudoku

2	9	4	3	7	8	1	5	6
1	7	3	6	4	5	9	8	2
5	6	8	2	1	9	7	3	4
6	5	7	1	9	2	3	4	8
9	8	2	4	3	6	5	1	7
4	3	1	8	5	7	6	2	9
3	1	9	7	8	4	2	6	5
7	4	6	5	2	1	8	9	3
8	2	5	9	6	3	4	7	1

3x3x3

Sudoku

	F		2						6			C	B	3	
	C				4	8	E	A			0		D		
D	A	8			3		2	7	F			6		5	
6			E	D	F		C		8						7
	9	3		7					A						2
E						6	F	5		8	4		3		1
C	8		1	3	9	D		0	2		E				
	D		6		5	E	B		1					0	4
9	6					1		F	3	2		0		A	
				4		A	8		D	0	9	B		2	5
2		A		0	D		5	6	C						F
5						2					A		4	8	
B						4		1		A	2	F			0
	0		7			F	3	C		D			2	9	B
		5		1			A	9	0	B				D	
	2	D	A			9						1		4	

4x4x4

Sudoku

0	F	9	2	A	7	5	1	4	6	E	D	C	B	3	8
7	C	1	3	6	4	8	E	A	B	5	0	2	D	F	9
D	A	8	4	9	3	B	2	7	F	C	1	6	0	5	E
6	5	B	E	D	F	0	C	2	8	9	3	4	A	1	7
4	9	3	5	7	1	C	0	D	A	F	B	8	E	6	2
E	B	7	0	2	A	6	F	5	9	8	4	D	3	C	1
C	8	F	1	3	9	D	4	0	2	6	E	5	7	B	A
A	D	2	6	8	5	E	B	3	1	7	C	9	F	0	4
9	6	4	8	E	B	1	7	F	3	2	5	0	C	A	D
3	7	C	F	4	6	A	8	E	D	0	9	B	1	2	5
2	1	A	B	0	D	3	5	6	C	4	8	7	9	E	F
5	E	0	D	F	C	2	9	B	7	1	A	3	4	8	6
B	3	6	9	C	E	4	D	1	5	A	2	F	8	7	0
1	0	E	7	5	8	F	3	C	4	D	6	A	2	9	B
8	4	5	C	1	2	7	A	9	0	B	F	E	6	D	3
F	2	D	A	B	0	9	6	8	E	3	7	1	5	4	C

4x4x4

2			3		8		5	
		3		4	5	9	8	
		8			9	7	3	4
6		7		9				
9	8						1	7
				5		6		9
3	1	9	7			2		
	4	6	5	2		8		
	2		9		3			1

Sudoku

Suppose it takes you $S(n)$ to solve $n \times n \times n$

	F		2					6		C	B	3	
	C				4	8	E	A			0		D
D	A	8			3		2	7	F			6	5
6			E	D	F		C		8				7
	9	3		7				A					2
E					6	F		5	8	4		3	1
C	8		1	3	9	D		0	2	E			
	D		6		5	E	B		1				0
9	6				1			F	3	2		0	A
			4		A	8		D	0	9	B	2	5
2		A		0	D		5	6	C				F
5					2					A		4	8
B					4			1	A	2	F		0
	0	7			F	3	C	D				2	9
		5		1		A	9	0	B				D
2	D	A			9						1		4

$V(n)$ time to verify the solution

Fact: $V(n) = O(n^2 \times n^2)$

Question: is there some constant such that

$S(n) \leq n^{\text{constant}}$?

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$n \times n \times n$

2			3		8		5	
		3		4	5	9	8	
		8			9	7	3	4
6		7		9				
9	8						1	7
				5		6		9
3	1	9	7			2		
	4	6	5	2		8		
	2		9		3			1

Sudoku

P vs NP problem

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	F		2					6			C	B	3	
	C				4	8	E	A		0			D	
D	A	8			3		2	7	F			6		5
6			E	D	F		C		8					7
	9	3		7				A						2
E						6	F	5	8	4			3	1
C	8		1	3	9	D		0	2		E			
	D		6		5	E	B		1					0
9	6					1		F	3	2		0	A	
				4		A	8		D	0	9	B		2
2		A		0	D		5	6	C					F
5					2						A		4	8
B					4			1	A	2	F			0
	0		7			F	3	C		D			2	9
		5		1		A	9	0	B				D	
2	D	A			9						1		4	

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$n \times n \times n$

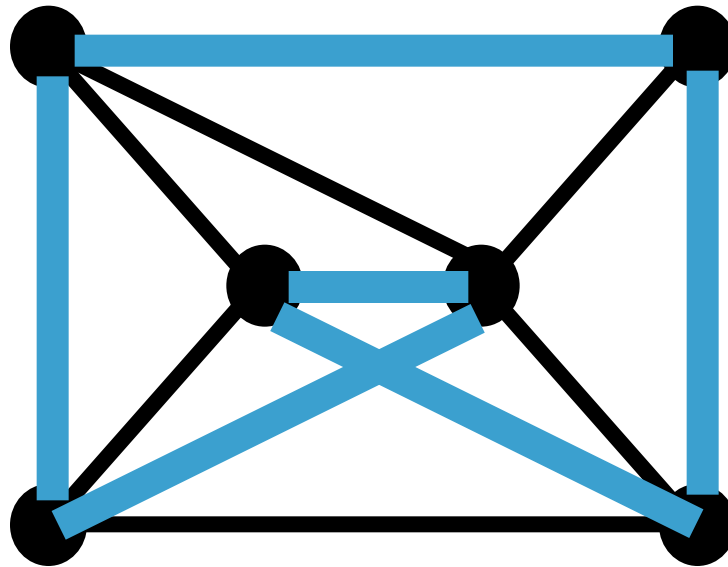
Does there exist an algorithm for $n \times n \times n$ Sudoku that runs in time $p(n)$ for some polynomial $p()$?

THE P VERSUS NP PROBLEM (INFORMALLY)

Is proving a theorem **much** more difficult
than checking the proof of a theorem?

Hamilton Cycle

Given a graph $G = (V, E)$, a cycle that visits all the nodes exactly once



The Problem “HAM”

Input: Graph $G = (V, E)$

Output: YES if G has a Hamilton cycle

NO if G has no Hamilton cycle

The Set “HAM”

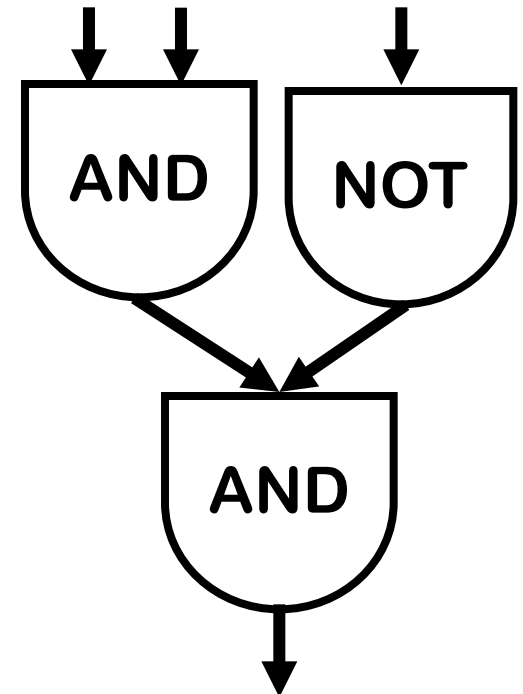
$\text{HAM} = \{ \text{graph } G \mid G \text{ has a Hamilton cycle} \}$

Circuit-Satisfiability

Input: A circuit C with one output

Output: YES if C is satisfiable

NO if C is not satisfiable



The Set “SAT”

$\text{SAT} = \{ \text{all satisfiable circuits } C \}$

Sudoku

Input: $n \times n \times n$ sudoku instance

Output: YES if this sudoku has a solution
NO if it does not

The Set “SUDOKU”

SUDOKU = { All solvable sudoku instances }

Decision Versus Search Problems

Decision Problem

YES/NO answers

**Does G have a
Hamilton cycle?**

**Can G be
3-colored ?**

Search Problem

**Find a Hamilton cycle
in G if one exists,
else return NO**

**Find a 3-coloring of
 G if one exists, else
return NO**

Reducing Search to Decision

Given an algorithm for decision Sudoku, devise an algorithm to find a solution

Idea:

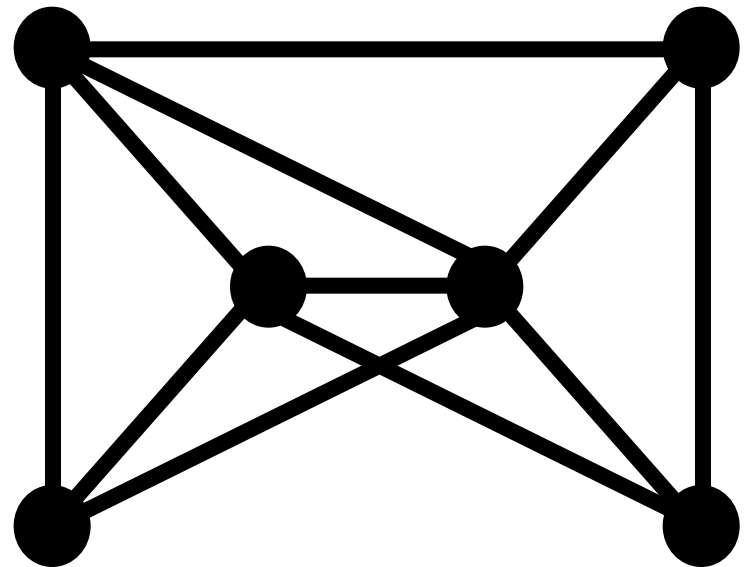
**Fill in one-by-one and
use decision algorithm**

2			3		8		5	
		3		4	5	9	8	
		8			9	7	3	4
6		7		9				
9	8						1	7
				5		6		9
3	1	9	7			2		
	4	6	5	2		8		
	2		9		3			1

Reducing Search to Decision

Given an algorithm for decision HAM,
devise an algorithm to find a solution

Idea:
Find the edges of the
cycle one by one



Decision/Search Problems

We' ll study decision problems because they are almost the same (asymptotically) as their search counterparts

Polynomial Time and The Class “P” of Decision Problems

WHAT IS AN EFFICIENT ALGORITHM?

Is an $O(n)$ algorithm efficient?

How about $O(n \log n)$?

$O(n^2)$?

$O(n^{10})$?

$O(n^{\log n})$?

$O(2^n)$?

$O(n!)$?

polynomial time

$O(n^c)$ for some
constant c

non-polynomial
time

Does an algorithm
running in $O(n^{100})$ time
count as efficient?

We consider **non-polynomial** time
algorithms to be **inefficient**.

And hence a **necessary** condition for an
algorithm to be efficient is that it should
run in poly-time.

Asking for a poly-time algorithm for a problem sets a (very) low bar when asking for efficient algorithms.

The question is: can we achieve even this for 3-coloring?

SAT?

Sudoku?