



# Chapter2: Boolean Algebra and Logic Gates

## Lecture3- Canonical Forms

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# Objectives

- Study Canonical Forms
- Standard and non-standard Forms
- Conversion of Canonical Forms

# Canonical Forms

- A **canonical form** is a standard method for representing Boolean functions.
- The two canonical forms that are used are:
  - Sum of Minterms
  - Product of Maxterms
- These forms are sometimes considered the “brute force” method of representing functions as they seldom represent a function in a minimized form.
- Examples of these two forms are:

$$F_1 = xyz' + xy'z + x'y'z'$$

$$F_2 = (x+y+z')(x+y'+z)(x'+y'+z')$$

# Minterms

- Any given binary variable can be represented in two forms:
  - $x$ , its normal form, and
  - $x'$ , its complement
- If we consider two binary variables and the AND operation, there are four combinations of the variables:
  - $xy$
  - $xy'$
  - $x'y$
  - $x'y'$
- Each of the above four AND terms is called a **minterm** or a **standard product**.
- $n$  variables can be combined to form  $2^n$  minterms.

# Minterms Expressed

## Minterms for Three Variables

<b>X</b>	<b>Y</b>	<b>Z</b>	<b>Product Term</b>	<b>Symbol</b>
0	0	0	$\overline{X}\overline{Y}\overline{Z}$	$m_0$
0	0	1	$\overline{X}\overline{Y}Z$	$m_1$
0	1	0	$\overline{X}Y\overline{Z}$	$m_2$
0	1	1	$\overline{X}YZ$	$m_3$
1	0	0	$X\overline{Y}\overline{Z}$	$m_4$
1	0	1	$X\overline{Y}Z$	$m_5$
1	1	0	$XY\overline{Z}$	$m_6$
1	1	1	$XYZ$	$m_7$

# Maxterms Expressed

## Maxterms for Three Variables

X	Y	Z	Sum Term	Symbol
0	0	0	$X + Y + Z$	$M_0$
0	0	1	$X + Y + \bar{Z}$	$M_1$
0	1	0	$X + \bar{Y} + Z$	$M_2$
0	1	1	$X + \bar{Y} + \bar{Z}$	$M_3$
1	0	0	$\bar{X} + Y + Z$	$M_4$
1	0	1	$\bar{X} + Y + \bar{Z}$	$M_5$
1	1	0	$\bar{X} + \bar{Y} + Z$	$M_6$
1	1	1	$\bar{X} + \bar{Y} + \bar{Z}$	$M_7$

# Truth Table to Expression (Sum of Minterms)

- Any Boolean function can be expressed as a **sum of minterms** or **sum of products** (i.e. the ORing of terms).
  - You can form the function algebraically by forming a **minterm** for each combination of the variables that produces a **1** in the function. (Each row with output of **1** becomes a **product term**)
  - **Sum (OR)** product terms together.

x	y	z	G
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1


$$xyz + xyz' + x'yz$$

# Minterms and Maxterms Expressed

**Table 2.3**

*Minterms and Maxterms for Three Binary Variables*

<b>x</b>	<b>y</b>	<b>z</b>	<b>Minterms</b>		<b>Maxterms</b>	
			<b>Term</b>	<b>Designation</b>	<b>Term</b>	<b>Designation</b>
0	0	0	$x'y'z'$	$m_0$	$x + y + z$	$M_0$
0	0	1	$x'y'z$	$m_1$	$x + y + z'$	$M_1$
0	1	0	$x'yz'$	$m_2$	$x + y' + z$	$M_2$
0	1	1	$x'yz$	$m_3$	$x + y' + z'$	$M_3$
1	0	0	$xy'z'$	$m_4$	$x' + y + z$	$M_4$
1	0	1	$xy'z$	$m_5$	$x' + y + z'$	$M_5$
1	1	0	$xyz'$	$m_6$	$x' + y' + z$	$M_6$
1	1	1	$xyz$	$m_7$	$x' + y' + z'$	$M_7$

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# Sum of Minterms Example

x	y	z		Function $F_1$	Required Minterms
0	0	0		1	$x'y'z'$
0	0	1		0	
0	1	0		0	
0	1	1		1	$x'yz$
1	0	0		1	$xy'z'$
1	0	1		0	
1	1	0		0	
1	1	1		0	

$$F_1 = x'y'z' + x'yz + xy'z'$$

$$= m_0 + m_3 + m_4$$

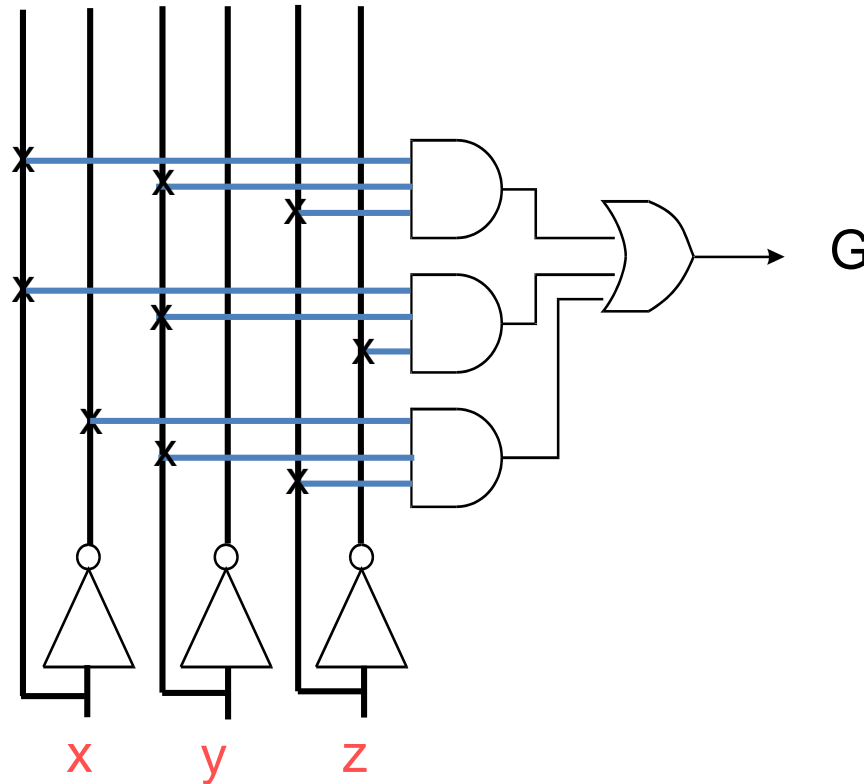
$$= \sum(0,3,4)$$

# Equivalent Representations of Circuits

- All three formats are equivalent
- Number of 1's in truth table output column equals AND terms for Sum-of-Products (SOP)

x	y	z	G
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

$$G = xyz + xyz' + x'yz$$



# Truth Table to Expression (Product of Maxterms)

- Any Boolean function can be expressed as a **product of maxterms** or **product of sums** (i.e. the ANDing of terms).
  - You can form the function algebraically by forming a **maxterm** for each combination of the variables that produces a **0** in the function. (Each row with output of **0** becomes a **standard sums**)
  - **AND** these maxterms together.

# Product of Maxterms Example

x	y	z		Function $F_1$	Required Maxterms
0	0	0		1	
0	0	1		0	$x + y + z'$
0	1	0		0	$x + y' + z$
0	1	1		1	
1	0	0		1	
1	0	1		0	$x' + y + z'$
1	1	0		0	$x' + y' + z$
1	1	1		0	$x' + y' + z'$

$$\begin{aligned}
 F_1 &= (x + y + z')(x + y' + z)(x' + y + z')(x' + y' + z)(x' + y' + z') \\
 &= M_1 M_2 M_5 M_6 M_7 \\
 &= \pi(1, 2, 5, 6, 7)
 \end{aligned}$$

# Minterms and Maxterms

- Each variable in a Boolean expression is a **literal**
- Boolean variables can appear in normal (x) or complement form (x')
- Each AND combination of terms is a **minterm**
- Each OR combination of terms is a **maxterm**
- Example:

## Minterms

x	y	z	Minterm	
0	0	0	$x'y'z'$	$m_0$
0	0	1	$x'y'z$	$m_1$
...				
1	0	0	$xy'z'$	$m_4$
...				
1	1	1	$xyz$	$m_7$

## Maxterms

x	y	z	Maxterm	
0	0	0	$x+y+z$	$M_0$
0	0	1	$x+y+z'$	$M_1$
...				
1	0	0	$x'+y+z$	$M_4$
...				
1	1	1	$x'+y'+z'$	$M_7$

# Obtaining Sum of Minterms Form

A	B	C	$F = A'B + B' + C$	Required Minterms	Required Designations
0	0	0	1	$A'B'C'$	$m_0$
0	0	1	1	$A'B'C$	$m_1$
0	1	0	1	$A'BC'$	$m_2$
0	1	1	1	$A'BC$	$m_3$
1	0	0	1	$AB'C'$	$m_4$
1	0	1	1	$AB'C$	$m_5$
1	1	0	0		
1	1	1	1	$ABC$	$m_7$

$$F = A'B'C' + A'B'C + A'BC' + A'BC + AB'C' + AB'C + ABC$$

$$= m_0 + m_1 + m_2 + m_3 + m_4 + m_5 + m_7$$

$$F(A, B, C) = \sum(0, 1, 2, 3, 4, 5, 7)$$

# Obtaining Product of Maxterms

A	B	C		$F = A'B + B'C$	Required Maxterms	Required Designations
0	0	0		0	$A + B + C$	$M_0$
0	0	1		1		
0	1	0		1		
0	1	1		1		
1	0	0		0	$A' + B + C$	$M_4$
1	0	1		1		
1	1	0		0	$A' + B' + C$	$M_6$
1	1	1		0	$A' + B' + C'$	$M_7$

$$F = (A+B+C)(A'+B+C)(A'+B'+C)(A'+B'+C')$$

$$= M_0 \cdot M_4 \cdot M_6 \cdot M_7$$

$$F(A, B, C) = \pi(0, 4, 6, 7)$$

# Canonical Form Conversion

- A function represented as Sum of **minterms** can be represented as the Product of **maxterms** of the remaining terms.
- The complement of a function expressed in sum of minterms equals the sum of minterms missing from the original function
  - $F(A, B, C) = \sum(0, 3, 4) = m_0 + m_3 + m_4$
  - $F'(A, B, C) = \sum(1, 2, 5, 6, 7) = m_1 + m_2 + m_5 + m_6 + m_7$
- Now if we take the complement of  $F'$  using DeMorgan's theorem, we obtain  $F$  in the product of maxterms form:
  - $(F')' = (m_1 + m_2 + m_5 + m_6 + m_7)'$
  - $F = m_1' \cdot m_2' \cdot m_5' \cdot m_6' \cdot m_7'$  [Complement of minterms]
  - $= M_1 M_2 M_5 M_6 M_7$  [maxterms]
  - $= \pi(1, 2, 5, 6, 7)$
- This implies the following relation:
$$m_i' = M_j$$
- So sum of minterms:  $\sum(0, 3, 4) =$  product of maxterms:  $\pi(1, 2, 5, 6, 7)$



# Table A: Conversion of Forms

Desired Form					
Given Form		Minterm Expansion of F	Maxterm Expansion of F	Minterm Expansion of F'	Maxterm Expansion of F'
	Minterm Expansion of F	-	maxterm nos are those nos, not on the minterm list of F	List minterms not present in F	Maxterm nos are same as minterm nos of F
	Maxterm Expansion of F	minterm nos are those nos, not on the maxterm list of F	-	minterm nos are same as maxterm nos of F	List maxterms not present in F

# Table B: Application of Table A

Desired Form					
Given Form		Minterm Expansion of F	Maxterm Expansion of F	Minterm Expansion of F'	Maxterm Expansion of F'
	$F = \sum(3,4,5,6,7)$	-	$F = \pi(0,1,2)$	$\sum(0,1,2)$	$\pi(3,4,5,6,7)$
	$F = \pi(0,1,2)$	$\sum(3,4,5,6,7)$	-	$\sum(0,1,2)$	$\pi(3,4,5,6,7)$

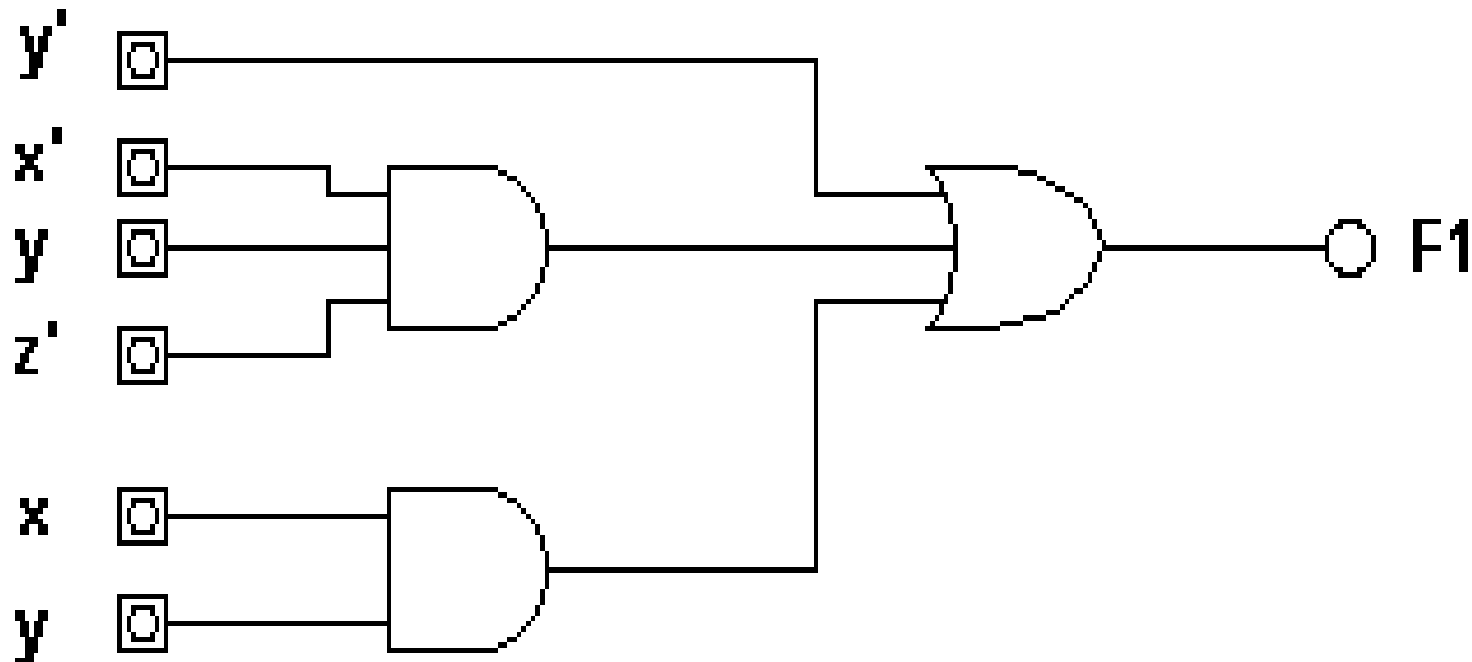
# Standard Forms

- **Standard forms** are those forms that allow the terms forming the function to consist of any number of the variables.
- There are two standard forms:
  - sum of products (SOP)
  - product of sums (POS)
- Examples of these two forms are:  
 $F = xy' + x'yz$                       SOP  
 $G = (X + Y')(X' + Y + Z)$                       POS

# Sum of Products

- The **Sum of Products (SOP)** is a Boolean expression containing AND terms, called **product terms**, of one or more literals each.

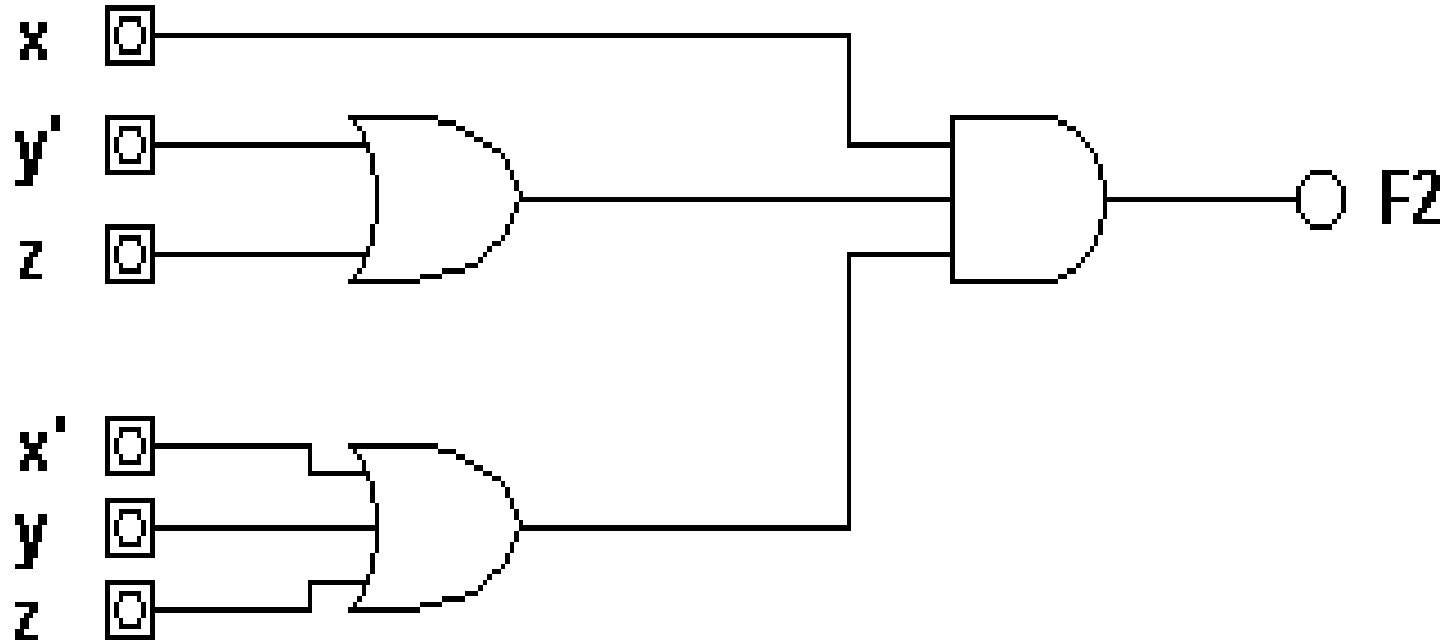
$$- F_1 = y' + xy + x'yz'$$



# Product of Sums

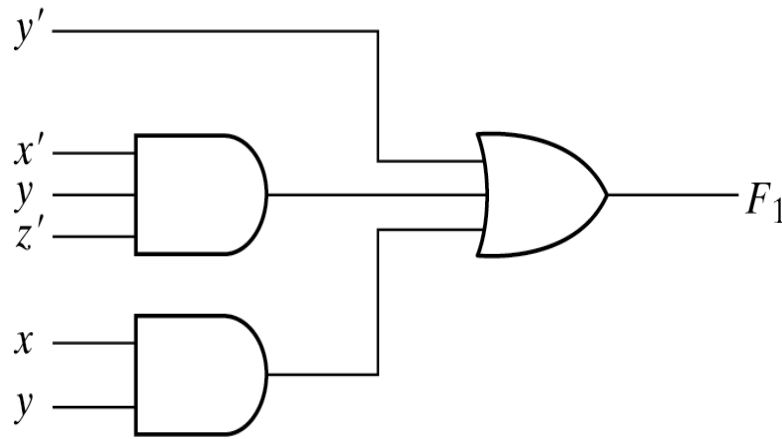
- The **Product of Sums (POS)** is a Boolean expression containing OR terms, called **sum terms**, of one or more literals each.

➤  $F_2 = x(y' + z)(x' + y + z')$

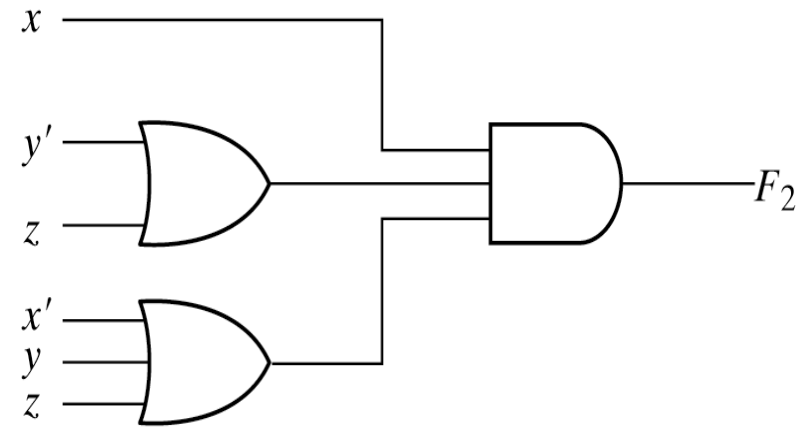


# Two Level Implementations

- The **standard type** of expression results in a **two-level** gating structure



(a) Sum of Products

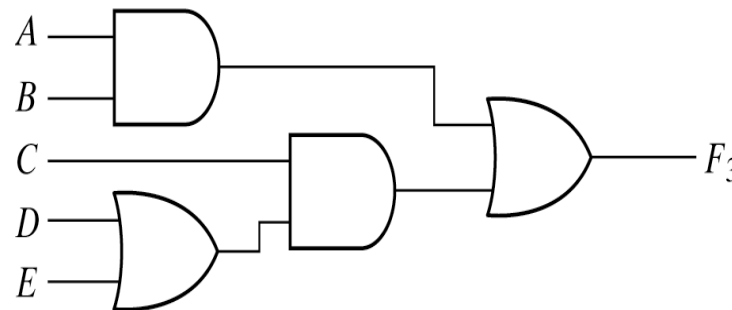


(b) Product of Sums

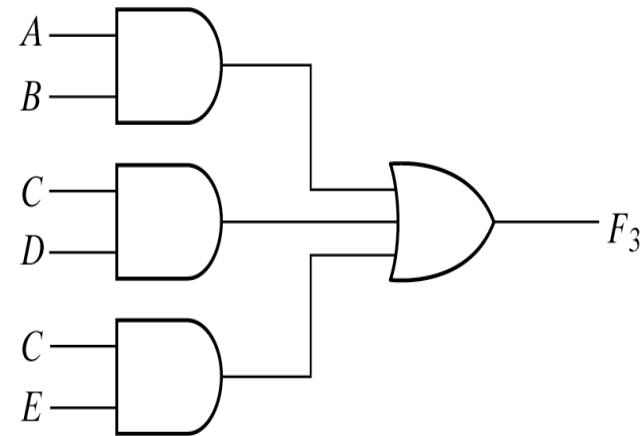
Fig. 2-3 Two-level implementation

# Conversion from Nonstandard to Standard Form

- A Boolean function may be expressed in a **nonstandard** form (fig 2.4a shows a function that is neither in sum of products nor in product of sums). It has three levels of gating
- It can be converted to a **standard form** (Sum of product) by using distributive law to remove parenthesis
- **Two-level** implementation is preferred as it produces the **least amount of delay**



(a)  $AB + C(D + E)$



(b)  $AB + CD + CE$

# The End