Data Structures & Algorithms

Linked Lists Variants

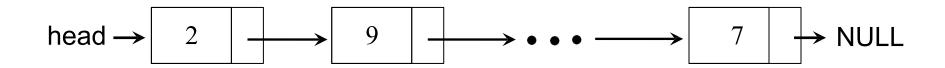
Today's lecture



- Linked structures
 - Singly Linked Lists
 - ▶ Doubly Linked Lists
 - ▶ Circular Linked Lists

Singly Linked List Problems





Problem:

- Cannot get backwards or to the beginning since no information related to previous node is available
 - Need a loop
 while (nodePtr->next) // Find the last node in the list
 nodePtr = nodePtr->next;

Other list structures



- Doubly-linked list
 - ► Each node has a pointer to its successor and its predecessor
 - ► Faster insert/delete, but more space
- Circular list
 - The last node points back to the head
- Sorted list
 - Items stored in sorted order
- Skip list
 - Skip certain nodes to avoid sequential processing
- Self Organizing list
 - Dynamically organizes the list in a certain manner

Doubly Linked List



- Contains two references to other nodes within each node: one to the next node on the list, and one to the previous node
 - ▶ I.e., every node (except the last node) contains the address of the next node, and every node (except the first node) contains the address of the previous node
 - Allows traversal in either direction

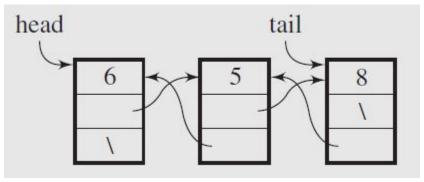
 Implementations: ALT+TAB and ALT+SHIFT+TAB (Window Browsing)

Doubly Linked List



- Add a prev pointer to Node class
 - Allows backward iteration
- When adding or removing a node, we must fix the prev and next pointers to have the correct value!

 Can make it easier to implement some methods such as remove



Declaration – Singly vs Doubly



Singly Linked List

```
class Node {
public:
    Node() {
         next = 0;
    Node(int i, Node *in = 0) {
         info = i; next = in;
    int info;
    Node *next;
};
```

Doubly Linked List

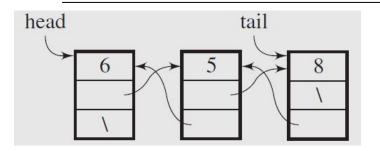
```
class Node {
public:
    Node() {
        next = prev = 0;
    Node(int i, Node *n = 0 , Node *p = 0)
        info = i; next = n; prev = p;
    int info;
    Node *next, *prev;
};
```



Singly Linked List

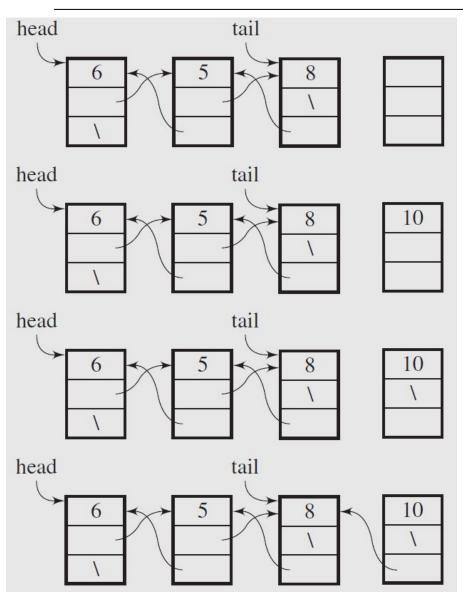
Doubly Linked List





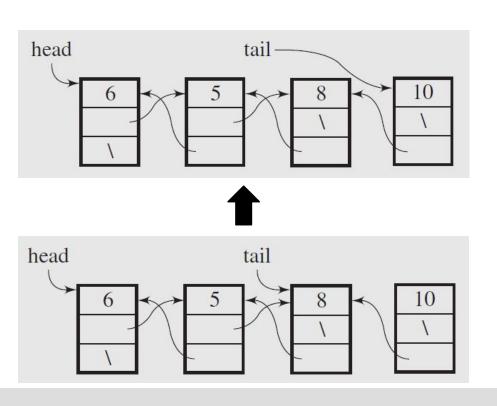
Doubly Linked List





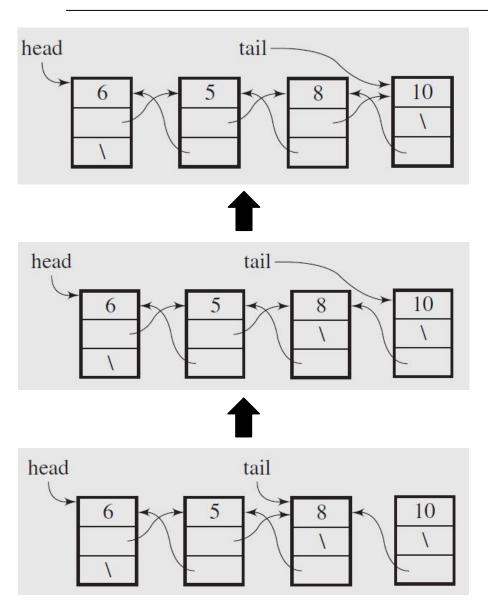
```
Doubly Linked List
void AccessNode::addToTail(int el) {
if (tail != 0) { // if list not empty;
         tail = new Node(el,0,tail);
         tail->prev->next = tail;
else
         head = tail = new Node(el);
```





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Doubly Linked List
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Adding a node in middle

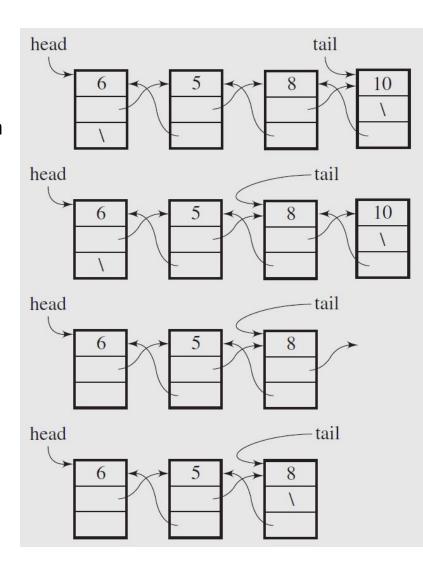


- When adding a node to the list at a given index, the following steps must be taken:
 - Advance through the list to the node just before the one with the proper index
 - Create a new node, and attach it to the nodes that should precede and follow it

Deleting a node



```
int AccessNode::deleteFromTail() {
  if (!tail) {
     // The list is empty, nothing to delete
     throw std::logic_error("Cannot delete from an
                              empty list.");
  el = tail->info;
  if (head == tail) { // If only one node in the list
     delete head;
     head = tail = nullptr;
  else { // If more than one node in the list
     tail = tail->prev;
     delete tail->next;
     tail->next = nullptr;
  return el;
```



Deleting a node in middle



- When deleting a node from the list at a given index, the following steps must be taken:
 - Advance through the list to the node with the proper index
 - Detach it from the nodes that used to precede and follow it

Today's lecture

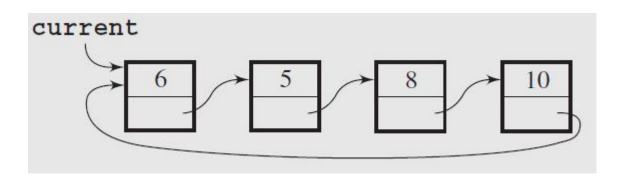


- Linked structures
 - Singly Linked Lists
 - Doubly Linked Lists
 - ▶ Circular Linked Lists

Circular Linked Lists



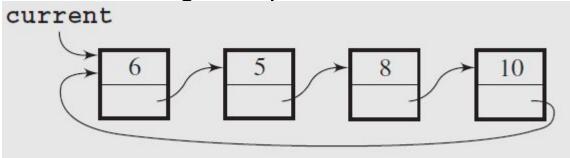
- Variant of Singly Linked List
 - Last node references the first node
 - Every node has a successor
 - ▶ No node in a circular linked list contains NULL



Advantage



- Allows searching from any node of the linked list and get to any other node
- Application example:
 - ► To ensure several processors sharing the same resource
 - ► **USAGE**: All processes are put on a circular list accessible through the pointer current



Implementation



 In an implementation of a circular singly linked list, we can use only one permanent pointer, i.e., tail, as successor of tail node is head node

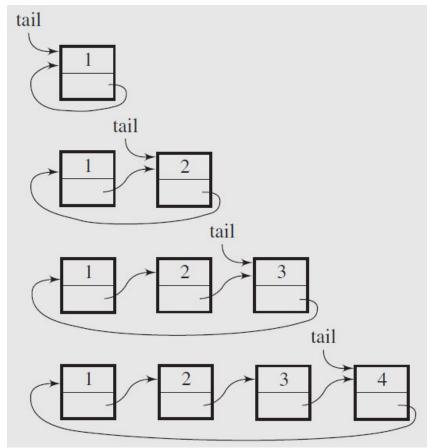
```
void addToTail(int el) {
if (isEmpty()) { // empty list
        tail = new Node(el);
        tail->next = tail;
     }
else { // not empty list
        tail->next = new Node(el,tail->next);
     tail = tail->next;
     }
}Call from main() with el =1,2,3,4
```

Implementation



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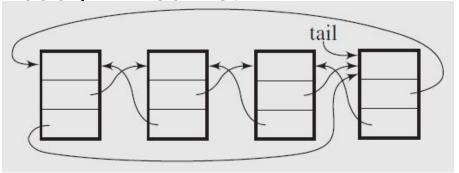


Problem with Circular Singly Linked List



- In deletion, still requires loop to search so that tail can be set to its predecessor after deleting the node
- Also, processing data in the reverse order (printing, searching, etc.) is not very efficient

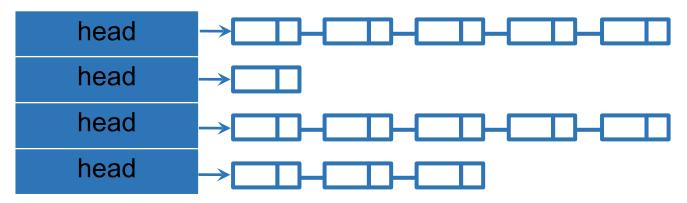
- One possible solution:
 - Circular Doubly Linked List



Arrays of Linked Lists



- Interesting data structure that combines static structure (an array) with dynamic structure (linked lists)
- Example: Array of head pointers
 - Appropriate for applications where number of categories is known in advance, but how many nodes in each category is NOT known



Arrays of Linked Lists



- Example application
 - A job portal website maintain categories and one member falls only in one category

where

- ► Each category index of array
- ► Member within category item of linked list

