

# DFA to Code

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# What is DFA?

## Deterministic Finite Automaton (DFA)

In DFA, for each input symbol, one can determine the state to which the machine will move. Hence, it is called **Deterministic Automaton**. As it has a finite number of states, the machine is called **Deterministic Finite Machine** or **Deterministic Finite Automaton**.

## Formal Definition of a DFA

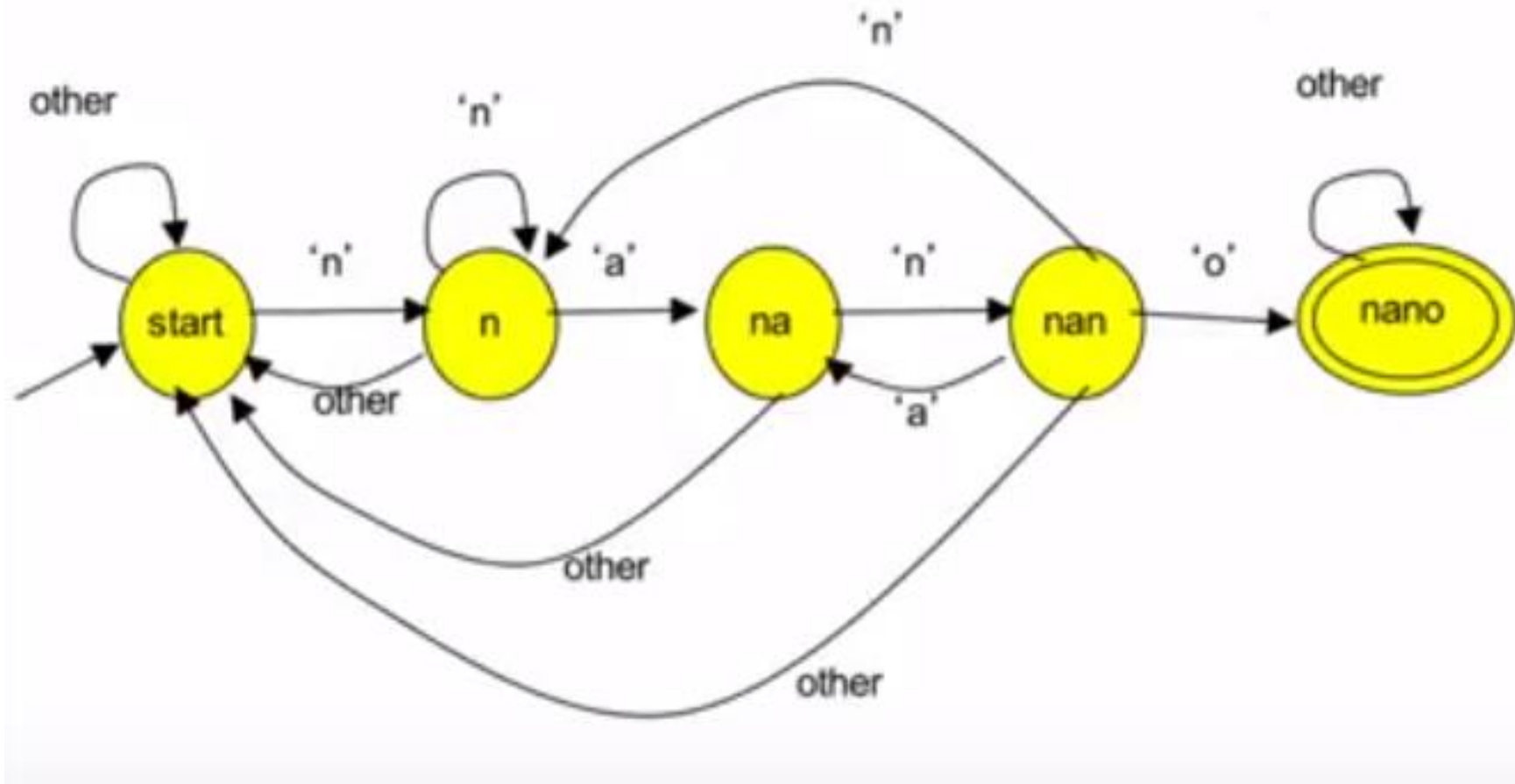
A DFA can be represented by a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$  where:

- $Q$  is a finite set of states.
- $\Sigma$  is a finite set of symbols called the alphabet.
- $\delta$  is the transition function where  $\delta: Q \times \Sigma \rightarrow Q$
- $q_0$  is the initial state from where any input is processed ( $q_0 \in Q$ ).
- $F$  is a set of final state/states of  $Q$  ( $F \subseteq Q$ ).

# DFA that can recognize string pattern “nano”



# DFA-State Diagram



# DFA-Program Logic

```
START (reject) <-- the start state
```

```
  n: N
```

```
  other: START
```

```
N (reject)
```

```
  n: N
```

```
  a: NA
```

```
  other: START
```

```
NA (reject)
```

```
  n: NAN
```

```
  other: START
```

```
NAN (reject)
```

```
  n: N
```

```
  a: NA
```

```
  o: NANO
```

```
  other: START
```

```
NANO (accept)
```

```
  other: NANO
```

## Hint!

- States as Function
- Symbol/Alphabet as Input Character
- Calling functions from a function
- Print that nano is present in the final state

# DFA-Regular Expression & Languages

	REGULAR EXPRESSION	REGULAR LANGUAGES
set of vowels	$(a \cup e \cup i \cup o \cup u)$	$\{a, e, i, o, u\}$

$\{w | w \in \{a, e, i, o, u\}^*\} \leftarrow$  Regular Language written formally