

# Combinators for Algebraic Structures (CAS)

## Version 1

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### Abstract

The document presents a small language of combinators for construction semigroups and semirings.

## 1 Motivation

We assume  $2 \leq |S|$  for all carrier sets  $S$ .

## 2 Semigroup Combinators

### 2.1 Adding an identity

$$\text{AddId}(\alpha, (S, \bullet)) \equiv (S \uplus \{\alpha\}, \bullet_{\alpha}^{\text{id}})$$

where

$$a \bullet_{\alpha}^{\text{id}} b \equiv \begin{cases} a & (\text{if } b = \text{inr}(\alpha)) \\ b & (\text{if } a = \text{inr}(\alpha)) \\ \text{inl}(x \bullet y) & (\text{if } a = \text{inl}(x), b = \text{inl}(y)) \end{cases}$$

### 2.2 Adding an annihilator

$$\text{AddAn}(\omega, (S, \bullet)) \equiv (S \uplus \{\omega\}, \bullet_{\omega}^{\text{an}})$$

where

$$a \bullet_{\omega}^{\text{an}} b \equiv \begin{cases} \text{inr}(\omega) & (\text{if } b = \text{inr}(\omega)) \\ \text{inr}(\omega) & (\text{if } a = \text{inr}(\omega)) \\ \text{inl}(x \bullet y) & (\text{if } a = \text{inl}(x), b = \text{inl}(y)) \end{cases}$$

### 2.3 Direct Product

Let  $(S, \bullet)$  and  $(T, \diamond)$  be semigroups.

The direct product is denoted

$$(S, \bullet) \times (T, \diamond) \equiv (S \times T, \star)$$

where

$$\star = \bullet \times \diamond$$

is defined as

$$(s_1, t_1) \star (s_2, t_2) = (s_1 \bullet s_2, t_1 \diamond t_2).$$

## 2.4 Lexicographic Product

Suppose that semigroup  $(S, \bullet)$  is commutative, idempotent, and selective and that  $(T, \diamond)$  is a semigroup.

$$(S, \bullet) \vec{\times} (T, \diamond) \equiv (S \times T, \star)$$

where  $\star \equiv \bullet \vec{\times} \diamond$  is defined as

$$(s_1, t_1) \star (s_2, t_2) = \begin{cases} (s_1 \bullet s_2, t_1 \diamond t_2) & s_1 = s_1 \bullet s_2 = s_2 \\ (s_1 \bullet s_2, t_1) & s_1 = s_1 \bullet s_2 \neq s_2 \\ (s_1 \bullet s_2, t_2) & s_1 \neq s_1 \bullet s_2 = s_2 \end{cases}$$

Examples for  $(\mathbb{N}, \min) \vec{\times} (\mathbb{N}, \min)$ .

$$\begin{aligned} (1, 17) \star (2, 3) &= (1, 17) \\ (2, 17) \star (2, 3) &= (2, 3) \\ (2, 3) \star (2, 3) &= (2, 3) \end{aligned}$$

Examples for  $(\mathbb{N}, \min) \vec{\times} (\mathbb{N}, \max)$ .

$$\begin{aligned} (1, 17) \star (2, 3) &= (1, 17) \\ (2, 17) \star (2, 3) &= (2, 17) \\ (2, 3) \star (2, 3) &= (2, 3) \end{aligned}$$

Examples for  $(\mathbb{N}, \max) \vec{\times} (\mathbb{N}, \min)$ .

$$\begin{aligned} (1, 17) \star (2, 3) &= (2, 3) \\ (2, 17) \star (2, 3) &= (2, 3) \\ (2, 3) \star (2, 3) &= (2, 3) \end{aligned}$$

## 2.5 Lifted Product

Assume  $(S, \bullet)$  is a semigroup. Let  $\text{lift}(S, \bullet) \equiv (\text{fin}(2^S), \hat{\bullet})$  where

$$X \hat{\bullet} Y = \{x \bullet y \mid x \in X, y \in Y\}.$$

Example.

$$\{1, 3, 17\} \hat{+} \{1, 3, 17\} = \{2, 4, 6, 18, 20, 34\}$$

## 2.6 Minimal set union

Let  $\leq$  be a partial order on  $S$ . For  $X \subseteq S$ , define

$$\min_{\leq}(X) \equiv \{x \in X \mid \forall y \in X, \neg(y < x)\}.$$

Define

$$\mathcal{P}_{\text{fin}}(S, \leq) \equiv \{X \subseteq S \mid X \text{ finite and } \min_{\leq}(X) = X\}$$

and

$$A \cup_{\min}^{\leq} B \equiv \min_{\leq}(A \cup B)$$

$$\text{MinUnion}(S, \leq) \equiv (\mathcal{P}_{\text{fin}}(S, \leq), \cup_{\min}^{\leq})$$

NOTE: perhaps we can always define  $\leq$  in terms of a  $(S, \text{oplus})$ , and so get rid of the class of structures  $(S, \leq)$ .

## 2.7 Minimal set lift

$$A \otimes_{\min}^{\leq} B \equiv \min_{\leq}(\{a \otimes b \mid a \in A, b \in B\})$$

$$\text{MinLift}(S, \leq, \otimes) \equiv (\mathcal{P}_{\text{fin}}(S, \leq), \otimes_{\min}^{\leq})$$

NOTE: perhaps we can always define  $\leq$  in terms of a  $(S, \text{oplus})$ , and so get rid of the class of structures  $(S, \leq)$ .

## 3 Pre-order Combinators

## 4 Pre-order Semigroup Combinators

## 5 Bi-Semigroup Combinators

### 5.1 Adding a zero

$$\text{AddZero}(\bar{0}, (S, \oplus, \otimes)) \equiv (S \uplus \{\bar{0}\}, \oplus_0^{\text{id}}, \otimes_0^{\text{an}})$$

### 5.2 Adding a one

$$\text{AddOne}(\bar{1}, (S, \oplus, \otimes)) \equiv (S \uplus \{\bar{1}\}, \oplus_1^{\text{an}}, \otimes_1^{\text{id}})$$

### 5.3 Direct Product

$$(S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T) \equiv (S \times T, \oplus_S \times \oplus_T, \otimes_S \times \otimes_T)$$

## 5.4 Lexicographic Product

$$(S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T) \equiv (S \times T, \oplus_S \vec{\times} \oplus_T, \otimes_S \times \otimes_T)$$

## 5.5 $\mathcal{M}$ and $\mathcal{N}$

Suppose  $D \in \{L, R\}$ . Let  $\leq \equiv \leq_{\oplus}^D$  and define

$$\mathcal{M}(D, (S, \oplus, \otimes)) \equiv (\mathcal{P}_{\text{fin}}(S, \leq), \otimes_{\min}^{\leq}, \cup_{\min}^{\leq})$$

$$\mathcal{N}(D, (S, \oplus, \otimes)) \equiv (\mathcal{P}_{\text{fin}}(S, \leq), \cup_{\min}^{\leq}, \otimes_{\min}^{\leq})$$

Recall:  $a \leq_{\oplus}^L b \equiv a = a \oplus b$  and  $a \leq_{\oplus}^R b \equiv b = a \oplus b$ .

## 6 Initial properties

The following are the properties of interest.

### 6.1 Semigroup

$(S, \bullet)$

associative	AS	$\equiv \forall a, b, c \in S, a \bullet (b \bullet c) = (a \bullet b) \bullet c$
identity	ID	$\equiv \exists \alpha \in S, \forall a \in S, a = \alpha \bullet a = a \bullet \alpha$
annihilator	AN	$\equiv \exists \omega \in S, \forall a \in S, \omega = \omega \bullet a = a \bullet \omega$
commutative	CM	$\equiv \forall a, b \in S, a \bullet b = b \bullet a$
selective	SL	$\equiv \forall a, b \in S, a \bullet b \in \{a, b\}$
idempotent	IIP	$\equiv \forall a \in S, a \bullet a = a$

### 6.2 Orders

$(S, \leq)$

reflexive	RX	$\equiv \forall a \in S, a \leq a$
transitive	TR	$\equiv \forall a, b, c \in S, a \leq b \wedge b \leq c \implies a \leq c$
anti-symmetric	AY	$\equiv \forall a, b \in S, a \leq b \wedge b \leq a \implies a = b$
total	TO	$\equiv \forall a, b \in S, a \leq b \vee b \leq a$
bottom	BM	$\equiv \exists b \in S, \forall a \in S, b \leq a$
top	TP	$\equiv \exists t \in S, \forall a \in S, a \leq t$

### 6.3 Bi-semigroups

$(S, \oplus, \otimes)$

left distributivity	$\mathbb{LD} \equiv \forall a, b, c \in S, a \otimes (b \oplus c) = (a \otimes b) \oplus (a \otimes c)$
right distributivity	$\mathbb{RD} \equiv \forall a, b, c \in S, (a \oplus b) \otimes c = (a \otimes c) \oplus (b \otimes c)$
zero annihilates	$\mathbb{ZA} \equiv \exists \bar{0} \in S, \mathbb{IID}(S, \oplus, \bar{0}) \wedge \mathbb{IAN}(S, \otimes, \bar{0})$
one annihilates	$\mathbb{OA} \equiv \exists \bar{1} \in S, \mathbb{IID}(S, \otimes, \bar{1}) \wedge \mathbb{IAN}(S, \oplus, \bar{1})$

where

$$\begin{aligned} \mathbb{IID}(S, \bullet, \alpha) &\equiv \forall a \in S, a = \alpha \bullet a = a \bullet \alpha \\ \mathbb{IAN}(S, \bullet, \omega) &\equiv \forall a \in S, \omega = \omega \bullet a = a \bullet \omega \end{aligned}$$

## 6.4 Ordered semigroups

$(S, \leq, \otimes)$

left monotonicity	$\mathbb{LM} \equiv \forall a, b, c \in S, a \leq b \implies (a \otimes c) \leq (b \otimes c)$
right monotonicity	$\mathbb{RM} \equiv \forall a, b, c \in S, a \leq b \implies (c \otimes a) \leq (c \otimes b)$

## 7 Closure

To fully close the theorems about our initial properties with necessary and sufficient conditions, we also need to track the following extra properties.

### 7.1 Semigroups

$(S, \bullet)$

		required for
is right	$\mathbb{IR} \equiv \forall s, t \in S, s \bullet t = t$	$\mathbb{SL}(\times)$
is left	$\mathbb{IL} \equiv \forall s, t \in S, s \bullet t = s$	$\mathbb{SL}(\times)$
left cancellative	$\mathbb{LC} \equiv \forall a, b, c \in S, c \bullet a = c \bullet b \implies a = b$	$\mathbb{LD}(\vec{\times})$
right cancellative	$\mathbb{RC} \equiv \forall a, b, c \in S, a \bullet c = b \bullet c \implies a = b$	$\mathbb{RD}(\vec{\times})$
left constant	$\mathbb{LK} \equiv \forall a, b, c \in S, c \bullet a = c \bullet b$	$\mathbb{LD}(\vec{\times})$
right constant	$\mathbb{RK} \equiv \forall a, b, c \in S, a \bullet c = b \bullet c$	$\mathbb{RD}(\vec{\times})$
anti-left	$\mathbb{AL} \equiv \forall a, b \in S, a \bullet b \neq a$	$\mathbb{LC}(\text{AddId}), \mathbb{LAB}(\vec{\times}, \times)$
anti-right	$\mathbb{AR} \equiv \forall a, b \in S, b \bullet a \neq a$	$\mathbb{RC}(\text{AddId}), \mathbb{RAB}(\vec{\times}, \times)$

### 7.2 Bi-semigroups

$(S, \oplus, \otimes)$

		required for
left absorption	$\mathbb{LAB} \equiv \forall a, b \in S, a = (b \otimes a) \oplus a = a \oplus (b \otimes a)$	$\mathbb{LD}(\text{AddOne})$
right absorption	$\mathbb{RAB} \equiv \forall a, b \in S, a = (a \otimes b) \oplus a = a \oplus (a \otimes b)$	$\mathbb{LD}(\text{AddOne})$

## 8 Closure progress

### 8.1 A few useful implications

$$\begin{aligned}
& \text{SL} \Rightarrow \text{IP} \\
& \text{RAB}(S, \oplus, \otimes) \wedge \text{ID}(S, \oplus) \Rightarrow \text{IP}(S, \otimes) \\
& \text{LAB}(S, \oplus, \otimes) \wedge \text{ID}(S, \oplus) \Rightarrow \text{IP}(S, \otimes) \\
& \text{LD}(S, \oplus, \otimes) \wedge \text{OA}(S, \oplus, \otimes) \Rightarrow \text{RAB}(S, \oplus, \otimes) \\
& \text{RD}(S, \oplus, \otimes) \wedge \text{OA}(S, \oplus, \otimes) \Rightarrow \text{LAB}(S, \oplus, \otimes)
\end{aligned}$$

### 8.2 AddId

$$\begin{aligned}
& \text{AS}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{AS}(S, \bullet) \\
& \text{ID}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{TRUE} \\
& \text{AN}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{AN}(S, \bullet) \\
& \text{CM}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{CM}(S, \bullet) \\
& \text{IP}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{IP}(S, \bullet) \\
& \text{SL}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{SL}(S, \bullet) \\
& \text{IR}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{IL}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{LC}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{LC}(S, \bullet) \wedge \text{AL}(S, \bullet) \\
& \text{RC}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{RC}(S, \bullet) \wedge \text{AR}(S, \bullet) \\
& \text{LK}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{RK}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{AL}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{AR}(\text{AddId}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE}
\end{aligned}$$

### 8.3 AddAn

$$\begin{aligned}
& \text{AS}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{AS}(S, \bullet) \\
& \text{ID}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{ID}(S, \bullet) \\
& \text{AN}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{TRUE} \\
& \text{CM}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{CM}(S, \bullet) \\
& \text{IP}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{IP}(S, \bullet) \\
& \text{SL}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{SL}(S, \bullet) \\
& \text{IR}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{IL}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{LC}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{RC}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{LK}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{RK}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{AL}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE} \\
& \text{AR}(\text{AddAn}(\alpha, (S, \bullet))) \Leftrightarrow \text{FALSE}
\end{aligned}$$

#### 8.4 $\times$

$$\begin{aligned}
\text{AS}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{AS}(S, \bullet) \wedge \text{AS}(T, \diamond) \\
\text{ID}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{ID}(S, \bullet) \wedge \text{ID}(T, \diamond) \\
\text{AN}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{AN}(S, \bullet) \wedge \text{AN}(T, \diamond) \\
\text{CM}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{CM}(S, \bullet) \wedge \text{CM}(T, \diamond) \\
\text{IP}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{IP}(S, \bullet) \wedge \text{IP}(T, \diamond) \\
\text{SL}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow (\text{IR}(S, \bullet) \wedge \text{IR}(T, \diamond)) \vee (\text{IL}(S, \bullet) \wedge \text{IL}(T, \diamond)) \\
\text{IR}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{IR}(S, \bullet) \wedge \text{IR}(T, \diamond) \\
\text{IL}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{IL}(S, \bullet) \wedge \text{IL}(T, \diamond) \\
\text{LC}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{LC}(S, \bullet) \wedge \text{LC}(T, \diamond) \\
\text{RC}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{RC}(S, \bullet) \wedge \text{RC}(T, \diamond) \\
\text{LK}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{LK}(S, \bullet) \wedge \text{LK}(T, \diamond) \\
\text{RK}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{RK}(S, \bullet) \wedge \text{RK}(T, \diamond) \\
\text{AL}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{AL}(S, \bullet) \vee \text{AL}(T, \diamond) \\
\text{AR}((S, \bullet) \times (T, \diamond)) &\Leftrightarrow \text{AR}(S, \bullet) \vee \text{AR}(T, \diamond)
\end{aligned}$$

#### 8.5 $\vec{\times}$ (semigroup)

Assuming  $\text{AS}(S, \bullet)$ ,  $\text{CM}(S, \bullet)$ , and  $\text{SL}(S, \bullet)$ .

$$\begin{aligned}
\text{AS}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{AS}(T, \diamond) \\
\text{ID}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{ID}(S, \bullet) \wedge \text{ID}(T, \diamond) \\
\text{AN}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{AN}(S, \bullet) \wedge \text{AN}(T, \diamond) \\
\text{CM}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{CM}(T, \diamond) \\
\text{IP}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{IP}(T, \diamond) \\
\text{SL}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{SL}(T, \diamond) \\
\text{IR}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{IL}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{LC}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{RC}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{LK}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{RK}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{AL}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE} \\
\text{AR}((S, \bullet) \vec{\times} (T, \diamond)) &\Leftrightarrow \text{FALSE}
\end{aligned}$$

*MinUnion*

## 8.6 Lift

$\text{AS}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{AS}(S, \bullet)$
$\text{ID}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{ID}(S, \bullet) \ (\hat{\alpha} = \{\alpha\})$
$\text{AN}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{TRUE} \ (\omega = \{\})$
$\text{CM}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{CM}(S, \bullet)$
$\text{SL}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{IL}(S, \bullet) \vee \text{IR}(S, \bullet) \vee (\text{IP}(S, \bullet) \wedge  S  = 2)$
$\text{IP}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{SL}((S, \bullet))$
$\text{IL}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{IR}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{LC}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{RC}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{LK}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{RK}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{AL}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{AR}(\text{lift}(S, \bullet))$	$\Leftrightarrow$	$\text{FALSE}$

## 8.7 MinUnion

Assume  $\text{RX}(S, \leq)$  (needed for sensible semantics),  $\text{TR}(S, \leq)$  (needed for associativity),  $\text{AY}(S, \leq)$  (needed for annihilator) and  $\text{BM}(S, \leq)$  (needed for annihilator).

$\text{AS}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TRUE}$
$\text{ID}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TRUE} \ (\alpha = \emptyset)$
$\text{AN}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TRUE} \ (\omega = \{\perp\})$
$\text{CM}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TRUE}$
$\text{SL}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TO}(S, \leq)$
$\text{IP}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{TRUE}$
$\text{IL}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{IR}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{LC}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{RC}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{LK}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{RK}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{AL}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$
$\text{AR}(\text{MinUnion}(S, \leq))$	$\Leftrightarrow$	$\text{FALSE}$

## 8.8 MinLift

Assume  $\text{LM}(S, \leq, \otimes)$ ,  $\text{RM}(S, \leq, \otimes)$ , and  $\text{AY}(S, \leq)$ . (Needed for associativity)  
**Perhaps others constraints are needed to make progress on SL and IP. Lattice Theory to the rescue?**



$$\begin{aligned}
\text{AS}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{AS}(S, \otimes) \\
\text{ID}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{ID}(S, \otimes) \\
\text{AN}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{TRUE} \quad (\omega = \{\}) \\
\text{CM}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{CM}(S, \otimes) \\
\text{SL}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow ??? \quad (\text{SOMEDAY !!!}) \\
\text{IP}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow ??? \quad (\text{SOMEDAY !!!}) \\
\text{IL}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{IR}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{LC}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{RC}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{LK}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{RK}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{AL}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE} \\
\text{AR}(\text{MinLift}(S, \leq, \otimes)) &\Leftrightarrow \text{FALSE}
\end{aligned}$$

### 8.9 AddZero

$$\begin{aligned}
\text{LD}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{LD}(S, \oplus, \otimes) \\
\text{RD}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{RD}(S, \oplus, \otimes) \\
\text{ZA}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{TRUE} \\
\text{OA}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{OA}(S, \oplus, \otimes) \\
\text{RAB}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{RAB}(S, \oplus, \otimes) \\
\text{LAB}(\text{AddZero}(\bar{0}, (S, \oplus, \otimes))) &\Leftrightarrow \text{LAB}(S, \oplus, \otimes)
\end{aligned}$$

### 8.10 AddOne

$$\begin{aligned}
\text{LD}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{LD}(S, \oplus, \otimes) \wedge \text{RAB}(S, \oplus, \otimes) \wedge \text{IP}(S, \oplus) \\
\text{RD}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{RD}(S, \oplus, \otimes) \wedge \text{LAB}(S, \oplus, \otimes) \wedge \text{IP}(S, \oplus) \\
\text{ZA}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{ZA}(S, \oplus, \otimes) \\
\text{OA}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{TRUE} \\
\text{RAB}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{RAB}(S, \oplus, \otimes) \wedge \text{IP}(S, \oplus) \\
\text{LAB}(\text{AddOne}(\bar{1}, (S, \oplus, \otimes))) &\Leftrightarrow \text{LAB}(S, \oplus, \otimes) \wedge \text{IP}(S, \oplus)
\end{aligned}$$

### 8.11 Direct Product, $(\times, \times)$

$$\begin{aligned}
\text{LD}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{LD}(S, \oplus_S, \otimes_S) \wedge \text{LD}(T, \oplus_T, \otimes_T) \\
\text{RD}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{RD}(S, \oplus_S, \otimes_S) \wedge \text{RD}(T, \oplus_T, \otimes_T) \\
\text{ZA}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{ZA}(S, \oplus_S, \otimes_S) \wedge \text{ZA}(T, \oplus_T, \otimes_T) \\
\text{OA}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{OA}(S, \oplus_S, \otimes_S) \wedge \text{OA}(T, \oplus_T, \otimes_T) \\
\text{RAB}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{RAB}(S, \oplus_S, \otimes_S) \wedge \text{RAB}(T, \oplus_T, \otimes_T) \\
\text{LAB}((S, \oplus_S, \otimes_S) \times (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{LAB}(S, \oplus_S, \otimes_S) \wedge \text{LAB}(T, \oplus_T, \otimes_T)
\end{aligned}$$

### 8.12 Lexicographic Product, $(\vec{\times}, \times)$

Assume  $\text{AS}(S, \oplus_S)$ ,  $\text{CM}(S, \oplus_S)$ , and  $\text{SL}(S, \oplus_S)$ .

$$\begin{aligned}
\text{LD}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{LD}(S, \oplus_S, \otimes_S) \wedge \text{LD}(T, \oplus_T, \otimes_T) \\
&\quad \wedge (\text{LC}(S, \otimes_S) \vee \text{LK}(T, \otimes_T)) \\
\text{RD}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{RD}(S, \oplus_S, \otimes_S) \wedge \text{RD}(T, \oplus_T, \otimes_T) \\
&\quad \wedge (\text{RC}(S, \otimes_S) \vee \text{RK}(T, \otimes_T)) \\
\text{ZA}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{ZA}(S, \oplus_S, \otimes_S) \wedge \text{ZA}(T, \oplus_T, \otimes_T) \\
\text{OA}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{OA}(S, \oplus_S, \otimes_S) \wedge \text{OA}(T, \oplus_T, \otimes_T) \\
\text{RAB}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{RAB}(S, \oplus_S, \otimes_S) \wedge (\text{RAB}(T, \oplus_T, \otimes_T) \vee \text{AR}(S, \otimes_S)) \\
\text{LAB}((S, \oplus_S, \otimes_S) \vec{\times} (T, \oplus_T, \otimes_T)) &\Leftrightarrow \text{LAB}(S, \oplus_S, \otimes_S) \wedge (\text{LAB}(T, \oplus_T, \otimes_T) \vee \text{AL}(S, \otimes_S))
\end{aligned}$$

### 8.13 $\mathcal{N}$

Which constraints should we impose on  $(S, \oplus, \otimes)$  in order to make this easier, cleaner? And how can we do this without giving up too much in terms of expressive power?

$$\begin{aligned}
\text{LD}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RD}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{ZA}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{OA}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RAB}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{LAB}(\mathcal{N}(L, (S, \oplus, \otimes))) &\Leftrightarrow ???
\end{aligned}$$

$$\begin{aligned}
\text{LD}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RD}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{ZA}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{OA}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RAB}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{LAB}(\mathcal{N}(R, (S, \oplus, \otimes))) &\Leftrightarrow ???
\end{aligned}$$

### 8.14 $\mathcal{M}$

$$\begin{aligned}
\text{LD}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RD}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{ZA}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{OA}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RAB}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{LAB}(\mathcal{M}(L, (S, \oplus, \otimes))) &\Leftrightarrow ???
\end{aligned}$$

$$\begin{aligned}
\text{LD}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RD}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{ZA}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{OA}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{RAB}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ??? \\
\text{LAB}(\mathcal{M}(R, (S, \oplus, \otimes))) &\Leftrightarrow ???
\end{aligned}$$

## 9 Other Problems

### 9.1 Matrix algebras

### 9.2 swap

We might want to define

$$\text{swap}(S, \oplus, \otimes) \equiv (S, \otimes, \oplus).$$

The disadvantage of this is that we would then have to introduce (and close with respect to) these “dual” properties:

$$\begin{aligned} \text{LD}^\delta(S, \oplus, \otimes) &\equiv \forall a, b, c \in S, a \oplus (b \otimes c) = (a \oplus b) \otimes (a \oplus c) \\ \text{RD}^\delta(S, \oplus, \otimes) &\equiv \forall a, b, c \in S, (a \otimes b) \oplus c = (a \oplus c) \otimes (b \oplus c) \\ \text{RAB}^\delta(S, \oplus, \otimes) &\equiv \forall a, b \in S, a = (a \oplus b) \otimes a = a \otimes (a \oplus b) \\ \text{LAB}^\delta(S, \oplus, \otimes) &\equiv \forall a, b \in S, a = (b \oplus a) \otimes a = a \otimes (b \oplus a) \end{aligned}$$

### 9.3 $k$ -best paths

### 9.4 Semi-direct product

### 9.5 General reductions