### A Formally Verified Schulze Method

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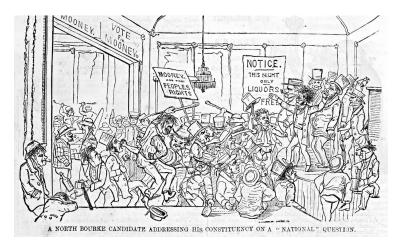
Research School of Computer Science

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### Talk Outline

- Motivation
- Schulze Voting as Evidence Carrying Computation
- Scaling it to count millions ballot
- Verifiable Homomorphic Tallying for the Schulze Vote Counting Scheme
- Machine Checked Properties
- Certificate: Verifying Elections Formally
- Future Work

Election held in 1855 in Victoria, Australia was conducted in pub!



### 2019 Election, NSW, Australia



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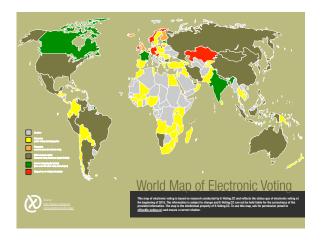
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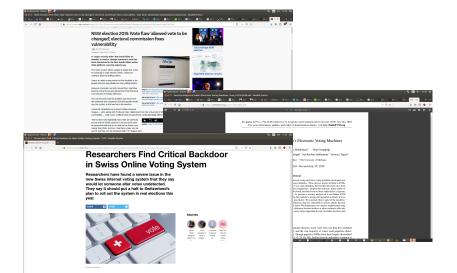


• What about Privacy and Verifiability in electronic voting ?









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- We assume first two part of "End to End verifiability" and work on third part.

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- Trusting on Java code in realized extracted code

# Machine Check Properties [On Going CPP 2020 ?]

- Condercet Winner (Finished)
- Reversal symmetry (Almost There)
- Monotonicity
- Smith Set

## Certificate: Verifying Elections Formally [CCS 2019]

- We developed a general framework for Sigma Protocol
- Instantiate with IACR 2018 election data
- We proved some computationally intense result inside Coq

### Certificate

```
V: [A3 B1 C2 D4,..], I: [], M: [AB:O AC:O AD:O BC:O BD:O CD:O]
V: [A1 B0 C4 D3...], I: [], M: [AB:-1 AC:-1 AD:1 BC:1 BD:1 CD:1]
______
V: [A3 B1 C2 D4,..], I: [A1 B0 C4 D3], M: [AB:-1 AC:-1 AD:1 BC:1 BD:1 CD:1]
V: [A1 B3 C2 D4], I: [A1 B0 C4 D3], M: [AB:2 AC:2 AD:8 BC:5 BD:8 CD:8]
V: [], I: [A1 B0 C4 D3], M: [AB:3 AC:3 AD:9 BC:4 BD:9 CD:9]
winning: A
 for B: path A --> B of strength 3, 4-coclosed set:
    [(B,A),(C,A),(C,B),(D,A),(D,B),(D,C)]
  for C: path A --> C of strength 3, 4-coclosed set:
    [(B,A),(C,A),(C,B),(D,A),(D,B),(D,C)]
  for D: path A --> D of strength 9, 10-coclosed set:
    \lceil (D,A), (D,B), (D,C) \rceil
losing: B
  exists A: path A --> B of strength 3, 3-coclosed set:
    [(A,A),(B,A),(B,B),(C,A),(C,B),(C,C),(D,A),(D,B),(D,C),(D,D)]
losing: C
  exists A: path A --> C of strength 3, 3-coclosed set:
    [(A,A),(B,A),(B,B),(C,A),(C,B),(C,C),(D,A),(D,B),(D,C),(D,D)]
losing: D
  exists A: path A --> D of strength 9, 9-coclosed set:
    [(A,A),(A,B),(A,C),(B,A),(B,B),(B,C),(C,A),(C,B),(C,C),(D,A),(D,B),
     (D,C),(D,D)1
```

### Future Work

- See all of you again in December of final presentation [December 2019]
- Risk limiting audit
- Verifying the cryptographic code

# Thank You!

