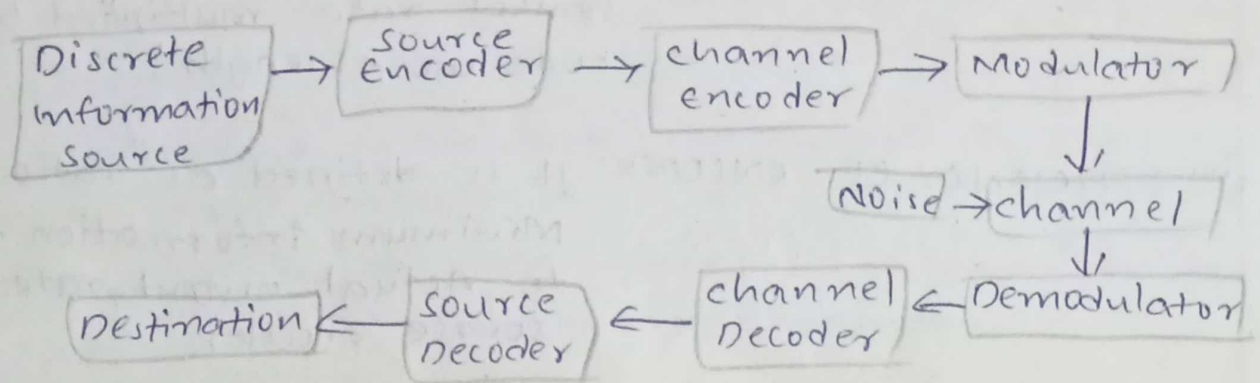


#### 4. DIGITAL CODING OF ANALOG WAVEFORMS

##### Block Diagram of Digital Communication system:



① INFORMATION SOURCE: - Information source may be classified based on nature of output i.e., Analog, (Digital) Discrete information source.

→ Discrete information source can be characterized by  
i) source Alphabet      ii) Symbol rate      iii) source Alphabet probabilities

iv) probabilistic dependence of symbols in a sequence.

→ Information rate =  $\frac{\text{Symbol rate} \times \text{source Entropy}}{(\text{bits/sec}) \quad (\text{symbols/sec}) \quad (\text{bits/symbol})}$

$$R = S \times H$$

② SOURCE ENCODER & DECODER: source encoder assigns code words to symbols. For each distinct symbol there is an unique code word. Code word can be of 4, 8, 16/32 bits length.

→ If it is 8-bit code word, it can generate  $2^8$  distinct code words.

→ source encoders must have the following parameters.

i) Block size: It describes max no. of distinct code words which can be represented by source encoder.

ii) CODE WORD LENGTH: It is no. of bits used to represent a code word.

(iii) AVERAGE DATA RATE: This is o/p bits/sec from source encoder. The data rate is equal to "product of symbol rate multiplied by code word length".

(iv) EFFICIENCY OF ENCODER: It is defined as "ratio of Minimum information rate to Actual output rate of source encoder."

③ CHANNEL ENCODER & DECODER: - The communication channel adds noise to signals. Hence, errors are introduced in binary sequence. To overcome that channel encoding is required.

→ A channel encoder must have:

i. CODING RATE: It depends on redundant bits added by channel encoder.

ii, Coding Method Used

iii, CODING EFFICIENCY: It is defined as "ratio of data rate at i/p to data rate at o/p of encoder".

iv, Error Control Capability

v, Feasibility of Encoder & Decoder

④ DIGITAL MODULATORS & DEMODULATORS:

→ A DM must have following parameters:

i, Bandwidth is needed to transmit signal.

ii, Probability of symbol/Bit error.

iii, Synchronous/Asynchronous method of detection

iv, Complexity of implementation.

→ DMs are classified as:

\* PCM \* Delta Modulator (DM) \* FSK \* PSK ...

⑥ COMMUNICATION CHANNEL: Communication channel can be characterized by following:

- i) SIGNAL ATTENUATION: it occurs due to internal resistance of channel and fading of signal.
- ii) Amplitude & Phase distortion: which is due to non-linear characteristics of communication channel.
- iii) Additive Noise interference: it is produced due to internal solid state devices & resistors.
- iv) Multi Part Distortion: which occurs mostly in wireless communication.

→ Communication channel can be classified as:

\* GUIDED CHANNELS:

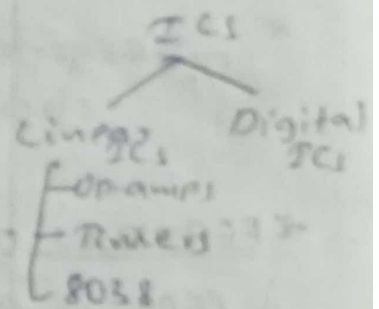
- i) Telephone channels
- ii) coaxial cables
- iii) Optical fibres

\* WIRELESS CHANNELS:

- i) wireless Broadcast channels
- ii) satellite channels
- iii) Mobile channels

\* ADVANTAGES OF DC:

- Simple and cheaper than AC.
- speech, video and other data may be merged and transmitted over a common channel used in multiplexing.
- Using data encryption, only permitted receiver may be allowed to detect transmitted data. This property is of importance in military communication.
- Since, transmission is digital and channel encoding is used, noise doesn't accumulate from repeater to repeater in long distance communication.
- Since, transmitted signal is digital in nature, large amount of noise





interference may be tolerated.

- Since, channel encoding is used, errors may be detected and corrected in receivers.
- It is adaptive to other advanced branches of digital signal processing, Image processing, Image compression etc.

### \* DISADVANTAGES / LIMITATIONS OF DC:

- Due to A-D conversion, data rate becomes high. Therefore, more transmission Bandwidth is required for digital communication.
- DC needs synchronization in case of synchronous modulation.

### \* PULSE CODE MODULATION:

\* Analog signal  $\rightarrow$  LPF  $\rightarrow$  sample  $\rightarrow$  Quantize  $\rightarrow$  Encoder  $\rightarrow$  PCM signal  
(TRANSMITTER BLOCK)

\* Distorted PCM signal  $\rightarrow$  Regenerative Repeater  $\rightarrow$  (No modulation)  $\rightarrow$  Regenerative Repeater  $\rightarrow$  Regenerated PCM signal  
For long distances (TRANSMISSION BLOCK)

\*  $\xrightarrow{i/p}$  Regenerative circuit  $\rightarrow$  Decoder  $\rightarrow$  Reconstruction filter  $\rightarrow$  Destination  
(REGENERATIVE REPEATER)  
(RECEIVER)

### \* FEATURES OF PCM:

- PCM is a type of pulse modulation like PAM, PWM/PPM. But it is digital pulse modulation scheme.
- PCM signal is in coded digital form. It is in the form of digital pulses of constant amplitude, width & position.
- The information is transmitted in the form of code.
- The PCM system consists of: PCM Encoder & Receiver.
- The essential operations in PCM are sampling, quantising and encoding. All these 3 operations are

performed by analog-digital converter.

vi, In other modulation schemes, By varying analog signal, some parameter of modulation schemes will vary, But in PCM it will not happen.

\* Nyquist rate must be maintained to avoid aliasing in PCM systems.

→ uniform quantization ← midtread  
← midrise (Derivation)

→ Non-uniform quantizers

→ Error signal: i/p signal - quantized signal

→ Step size  $\Delta = \frac{2m_{max}}{L}$  ;  $L = 2^R \Rightarrow R = \log_2 L$   
 $\Delta$  : No. of levels

→ quantization noise,  $f_Q(v) = \begin{cases} \frac{1}{\Delta} & -\frac{\Delta}{2} \leq v \leq \frac{\Delta}{2} \\ 0 & \text{elsewhere} \end{cases}$   
 (uniformly distributed)

→ Variance,  $\sigma_Q^2 = E[Q^2]$

$$= \int_{-\Delta/2}^{\Delta/2} v^2 \cdot f_Q(v) dv$$

$$= \frac{1}{\Delta} \int_{-\Delta/2}^{\Delta/2} v^2 dv = \frac{1}{\Delta} \left[ \frac{v^3}{3} \right]_{-\Delta/2}^{\Delta/2} = \frac{1}{\Delta} \left( \frac{\Delta^3}{24} - \left(-\frac{\Delta^3}{24}\right) \right) = \frac{\Delta^2}{12}$$

$$\sigma_Q^2 = \left[ \frac{2m_{max}^2}{2^R} \right] \frac{1}{12} = \frac{4m_{max}^2}{12 \cdot 2^{2R}} = \frac{m_{max}^2}{3 \cdot 2^{2R}}$$

\* 'P' is average power of msg signal "m(t)",

$$* (SNR)_{o/p} = \frac{P}{\sigma_Q^2} = \frac{P}{\frac{m_{max}^2}{3 \cdot 2^{2R}}} = \frac{3 \cdot P \cdot 2^{2R}}{m_{max}^2} = \left[ \frac{3P}{m_{max}^2} \right] 2^{2R}$$

→ From this eqn., the o/p SNR of uniform quantizer increases exponentially with increasing number of bits / sample. (R)

→ For a sinusoidal modulating signal,  $P = \frac{A_m^2}{2}$   
 $\sigma_Q^2 = \frac{1}{3} A_m^2 \cdot 2^{-2R}$

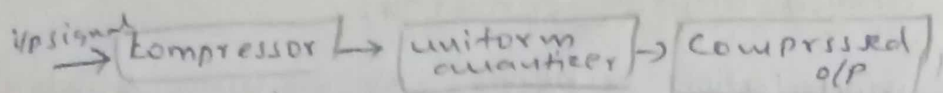
$$* [SNR = 4.8 + 6R]$$

$$(SNR)_{o/p} = \frac{3}{2} \cdot 2^{2R} //$$

$$(SNR)_b = 3 \cdot 2^{(2R-1)}$$

$$\rightarrow \text{In 'db', } (SNR)_b = 10 \log_{10} \frac{3}{2} + 10 \log_{10} 2^{2R} = 1.76 + 6R$$



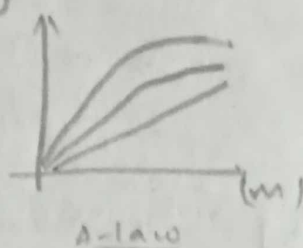
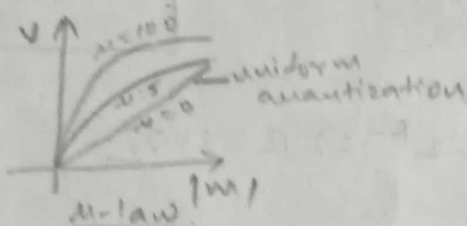


→ standard law of compression ① " $\mu$ -law" ② "A-law".

$$|v| = \frac{\ln(1 + \mu|m|)}{\ln(1 + \mu)} ; \quad m, v = \text{i/p, o/p voltages}$$

$\mu = \text{constant}$

→ It is assumed that  $m$  &  $v$  are scaled so that both lie inside the interval  $[-1, 1]$



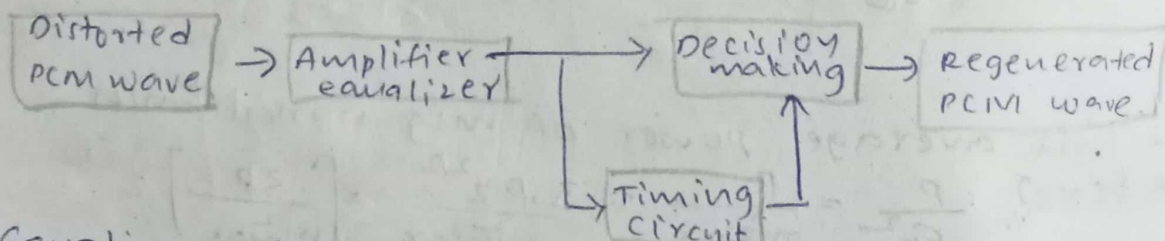
→ For a given value of  $\mu$ , reciprocal slope defined by quantum steps  $\frac{d|m|}{d|v|} = \frac{\ln(1 + \mu)}{\mu} (1 + \mu|m|)$  ( $\mu$ -law)

→

$$|v| = \begin{cases} \frac{A|m|}{1 + \ln A} & , 0 \leq |m| \leq \frac{1}{A} \\ \frac{1 + \ln(A|m|)}{1 + \ln A} & , \frac{1}{A} \leq |m| \leq 1 \end{cases} ; \text{ "A=constant" }$$

\* PCM uses Companding [Compression + Expansion]  
 → ISI is limitation of DC. To overcome, we use equalizers.

\* Analog-code pulse conversion: (REGENERATIVE REPEATED



→ Equalizer shapes received pulses, to compensate the effects of amplitude and phase distortions produced by non-ideal transmission characteristics of channel.

→ Timing circuit provides periodic train of pulses for sampling purpose.

→ Decision making circuit will have predefined threshold to make out a zero/one.

→ Regenerated signal ideally except for delay is same as signal which is originally transmitted. However,

it deviates in practice from original signal for 2 main reasons:

- ① "Bit errors" will appear due to channel noise and interference.
- ② "Jitter" will occur, where by spacing b/w received pulses deviates from their assigned value

If transmitted SNR ratio is high, then regenerated PCM data is same as transmitted PCM data except for a small BER. In other words under these operating conditions, performance degradation in PCM system is essentially confined to transmission Noise i.e., Quantization Noise.

### PROBABILITY OF ERROR:

→ First signal be: (Symbol 1)

$$* s_1(t) = \sqrt{\frac{E_{max}}{T_b}} \quad \text{--- (1)} \quad 0 \leq t \leq T_b$$

$\downarrow$   
 peak signal energy  
 $\downarrow$   
 duration

→ When symbol '0' is sent:

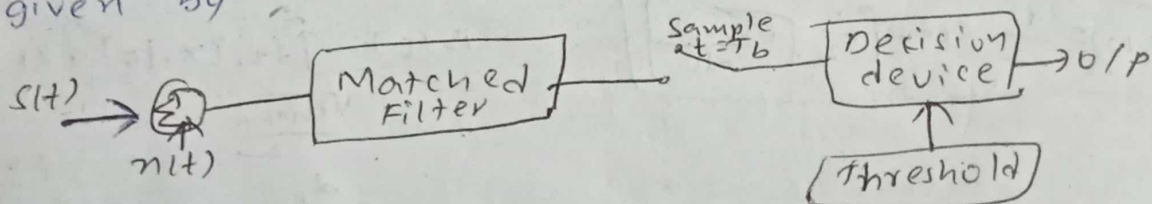
$$* s_2(t) = 0 \quad ; \quad 0 \leq t \leq T_b$$

Let channel be AWGN with zero mean and PSD by  $\frac{N_0}{2}$

→ Then, received signal is represented by,

$$x(t) = s_1(t) + n(t) \quad ; \quad 0 \leq t \leq T_b \quad \text{--- (2)}$$

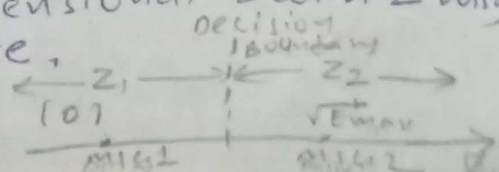
The optimum receiver uses a matched filter. The receiver structure for a binary encoded PCM is given by:



$$* s_1(t) = \sqrt{E_{max}} \cdot \phi_1(t) \quad \text{--- (3)}$$

$$\phi_1(t) = \text{Basis function} = \sqrt{\frac{1}{T_b}} \quad \text{--- (4)}$$

An ON/OFF PCM is characterized by having a signal space i.e., one dimensional & with 2 msg points as shown in figure.



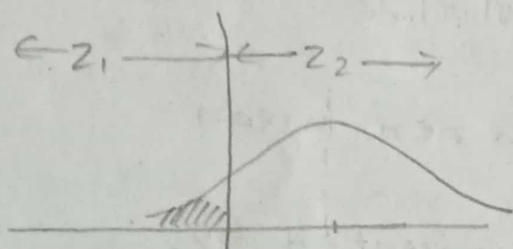
Two message points are represented as;

$$s_{11} = \int_0^{T_b} s_1(t) \cdot \phi_1(t) dt = \sqrt{E_{\max}}$$

$$s_{21} = \int_0^{T_b} s_2(t) \phi_1(t) dt = 0 \quad \text{--- (5)}$$

We assume that binary symbols '0' & '1' occur with equal probabilities and hence decision boundary is kept at Half-way point i.e.,  $\frac{\sqrt{E_{\max}}}{2}$ .

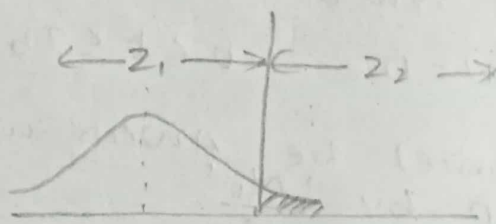
The 2 decision regions are denoted by  $z_1$  &  $z_2$ . The decision rule is to guess symbol-1 / signal  $s_{11}(t)$  was sent if received signal falls in region  $z_1$ . Similarly to guess symbol-0 / signal  $s_{21}(t)$  was sent if received signal point was in  $z_2$ .



$$\rightarrow f_{x_1}(x_1|1) \quad \text{--- (2)}$$

$$\rightarrow P_e(1) = \int_{-\infty}^{\frac{\sqrt{E_{\max}}}{2}} f_{x_1}(x_1|1) dx_1$$

$$P_e(1) = \frac{1}{2} \operatorname{erfc}\left[\frac{1}{2} \sqrt{\frac{E_{\max}}{N_0}}\right]$$



$$f_{x_1}(x_1|0) \quad \text{--- (1)}$$

$$\rightarrow f_{x_1}(x_1|0) = \frac{1}{\sqrt{\pi N_0}} \exp\left(-\frac{x_1^2}{N_0}\right) \quad \text{--- (6)}$$

Probability error,

$$P_e(0) = \int_{\frac{\sqrt{E_{\max}}}{2}}^{\infty} f_{x_1}(x_1|0) dx_1 \quad \text{--- (7)}$$

$$= \int_{\frac{\sqrt{E_{\max}}}{2}}^{\infty} \frac{1}{\sqrt{\pi N_0}} \exp\left(-\frac{x_1^2}{N_0}\right) dx_1$$

$$\text{Let, } z = \frac{x_1}{\sqrt{N_0}} \Rightarrow dx_1 = \sqrt{N_0} dz$$

$$P_e(0) = \int_{\frac{\sqrt{E_{\max}}}{2\sqrt{N_0}}}^{\infty} \frac{1}{\sqrt{\pi}} e^{-z^2} dz \quad \text{--- (8)}$$



$$= \frac{1}{\sqrt{\pi}} \left[ \frac{\bar{e}^2}{-2\bar{e}} \right]_{\frac{\sqrt{E_{\max}}}{2\sqrt{N_0}}}^{\alpha}$$

$$= \frac{1}{\sqrt{\pi}} \cdot \frac{e^{-\frac{E_{\max}}{4N_0}}}{2\sqrt{E_{\max}}} \cdot \frac{\sqrt{E_{\max}}}{2\sqrt{N_0}}$$

from (8),

$$P_{e(0)} = \frac{\text{erfc}\left(\frac{1}{2}\sqrt{\frac{E_{\max}}{N_0}}\right)}{2}$$

$$P_{e(0)} = \frac{1}{2} \text{erfc}\left(\frac{1}{2}\sqrt{\frac{E_{\max}}{N_0}}\right)$$

$$P_{e(0)} = \frac{\sqrt{N_0}}{\sqrt{\pi E_{\max}}} \cdot e^{-\frac{E_{\max}}{4N_0}}$$

$$\therefore \left[ \text{erfc}(u) = \frac{2}{\sqrt{\pi}} \int_u^{\infty} e^{-x^2/2} dx \right]$$

→ Similarly, to calculate making an error of second kind, assume symbol-1 is sent, then portion of likelihood function  $P_{e(1)}$  where limits are  $\frac{\sqrt{E_{\max}}}{2}$ .

→ The fact that " $P_{e(1)} = P_{e(0)}$ " is confirmation of symmetric nature of channel.  $P_{e(1)}$ ,  $P_{e(0)}$  are conditional probabilities. We assume that a priori probability of sending '0' is " $P_0$ " and a priori probability of sending '1' is " $P_1$ ". Hence, average probability of error is given by:

$$P_e = P_0 P_{e(0)} + P_1 P_{e(1)}$$

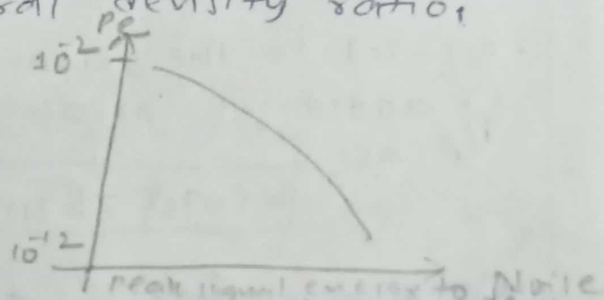
As  $P_{e(0)} = P_{e(1)}$ ,  $P_0 + P_1 = 1$  ∴ Hence, the average probability becomes half complementary error function.

$$\text{i.e., } P_e = \frac{1}{2} \text{erfc}\left(\frac{1}{2}\sqrt{\frac{E_{\max}}{N_0}}\right)$$

Here,  $\frac{E_{\max}}{N_0}$  = ratio of peak signal energy to Noise spectral density ratio.

$$\rightarrow E_{\max} = P_{\max} \cdot T_b$$

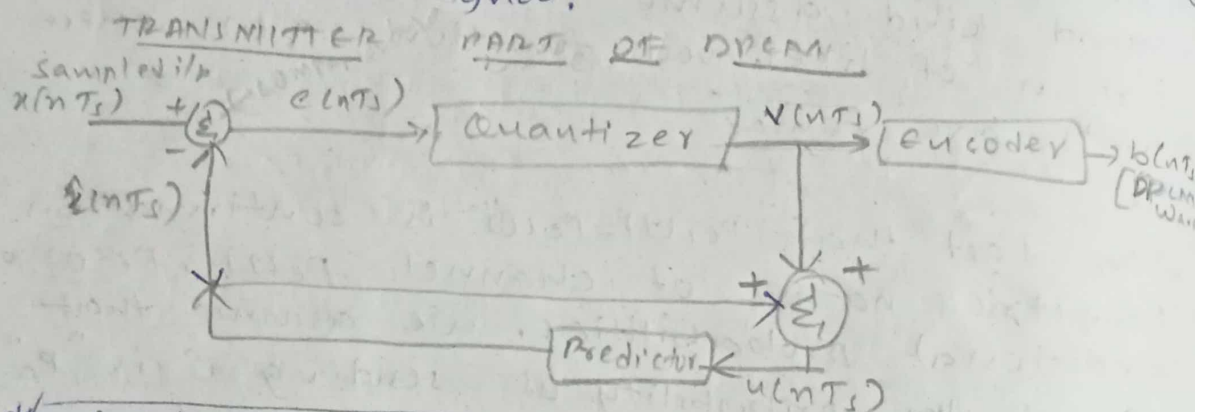
$$\Rightarrow \boxed{\frac{E_{\max}}{N_0} = \frac{P_{\max}}{(N_0/T_b)}}$$



" $\frac{N_0}{T_b}$ " may be taken as average noise power contained in transmission BW =  $1/T_b$ .

## \* DIFFERENTIAL PULSE CODE MODULATION:

In PCM, digitization of Noise/Video signal the signal is sampled at a rate slightly higher than Nyquist rate. The resulting signal doesn't change rapidly from one sample to next with result that the difference b/w adjacent samples is a variance i.e., smaller than the variance of signal itself. If these samples are encoded, it contains redundant information. By removing the redundancy before encoding we obtain a more efficient coded signal.



$$e(nTs) = x(nTs) - \hat{x}(nTs)$$

i.e.,  $e(nTs)$  is difference b/w unquantized i/p sample and prediction of it. The predicted value is produced by using a predictor whose i/p consists of a quantized version of i/p signal  $\hat{x}(nTs)$ .

→  $e(nTs)$  is called prediction error.

→ By encoding quantizer o/p we obtain "DPCM". The quantizer o/p is given by  $[b(nTs)]$ .

$$v(nTs) = Q[e(nTs)]$$

$$v(nTs) = e(nTs) + q(nTs)$$

•  $q(nTs)$  is the quantization error. The quantizer o/p is added to predicted value to produce predicted i/p as:

$$u(nTs) = \hat{x}(nTs) + v(nTs)$$

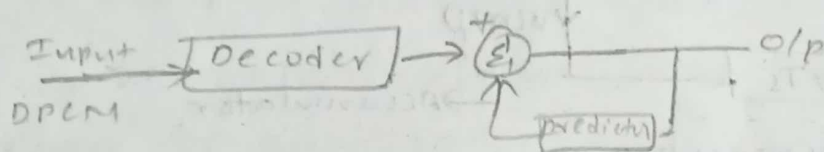
• We observe that  $\hat{x}(nTs) + v(nTs)$  is equal to the signal  $x(nTs)$ . Therefore, irrespective of properties



of predictor, the quantized signal  $x(nT_s)$  at predictor i/p differs from original signal by quantization error. Accordingly, your prediction is good.

→ The variance of prediction error will be smaller than the variance  $x(nT_s)$ , this leads to quantization error which is lesser than standard PCM. The receiver for constructing DPCM is given by:

### RECEIVER PART OF DPCM:



The receiver for constructing an analyzer version of i/p is decoder followed by a predictor. From this analysis, we can conclude that the predictor is a transmitter-receiver operate on same sequence of samples  $x(nT_s)$ . With this purpose in mind the f/b path is added to analyzer in transmitter.

~~SNR~~ The o/p signal to quantization ratio of a signal coder is defined by:

$$(SNR)_0 = \frac{\sigma_x^2}{\sigma_e^2}$$

$\sigma_x^2$  = variance of i/p  $x(nT_s)$   
 $\sigma_e^2$  = variance of quantization error  $e(nT_s)$

$$(SNR)_0 = \frac{\sigma_x^2}{\sigma_e^2} \times \frac{\sigma_e^2}{\sigma_{\epsilon}^2} ; \sigma_{\epsilon}^2 = \text{variance of prediction error } \epsilon(nT_s)$$

$$(SNR)_0 = (G)_p \times (SNR)_p$$

$G_p$  = Predictor gain  
 $(SNR)_p$  = Predictor error to quantization error ratio

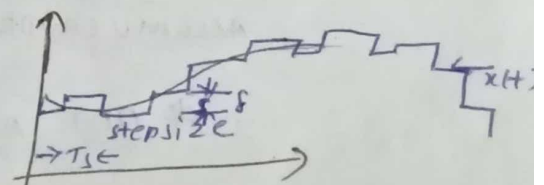
→ To design the predictor so as to minimize  $\sigma_{\epsilon}^2$

### \*DELTA MODULATION:

It is one-bit version of DPCM (Two level version)

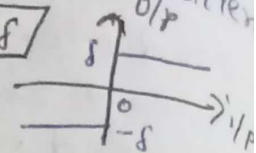
'd' can be +ve/-ve.

's' denotes absolute value of the 2 representation levels of the one bit quantizer, used in delta modulator.

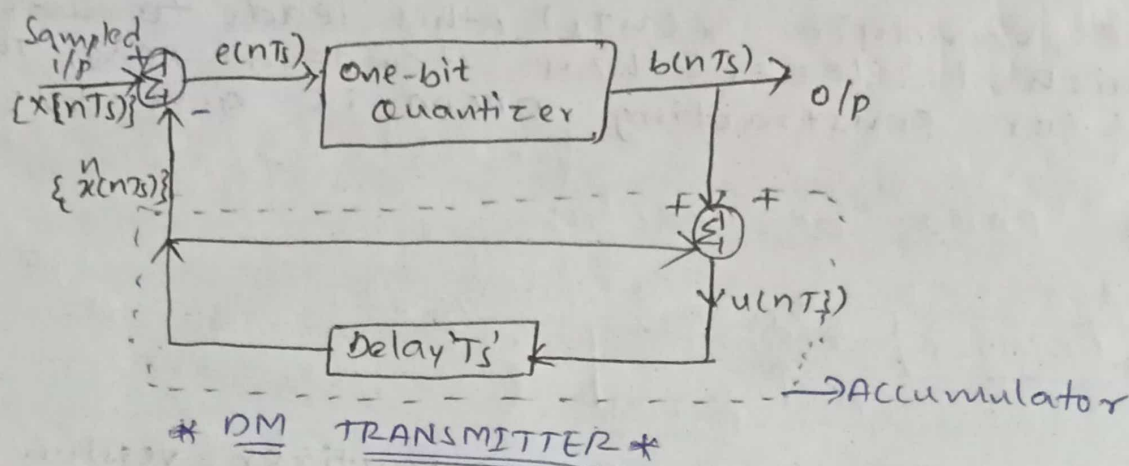




These 2 levels are indicated in the Transfer character as shown in figure and the step size,  $\Delta = 2f$



Block diagram of delta modulation is shown below:



$$* e(nTs) = x(nTs) - \hat{x}(nTs) \quad \text{--- (1)}$$

$$* e(nTs) = x(nTs) - u(nTs - Ts) \quad \text{--- (2)}$$

$$* u(nTs) = b(nTs) + u(nTs - Ts) \quad \text{--- (3)}$$

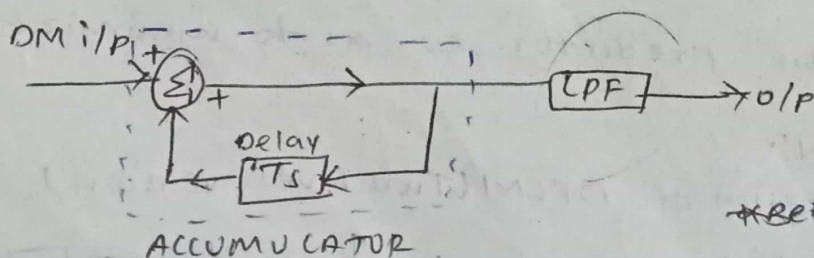
$$* b(nTs) = \text{sgn}\{e(nTs)\} \quad \text{--- (4)}$$

$T_s$  = Sampling period  
 $e(nTs)$  = Prediction error

- The binary quantity  $b(nTs)$  is the algebraic sign of the error  $e(nTs)$  except for scaling factor 'f'. In fact  $b(nTs)$  is the one bit word transmitted by the DM system.

- In comparing DPCM and DM networks except for an o/p LPF, Delta modulation is a special case of DPCM.

- DM offers 2 unique features: i) A one-bit code word for the o/p which eliminates the need for word framing  
 ii) Simplicity of design for both the transmitter & receiver



\* Better than DPCM \*

\* DM RECEIVER \*

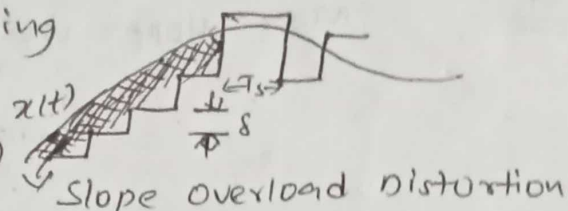
1. DM suffers from 2 types of Noise:

- i) Slope overload distortion
- ii) Granular Noise

i) SLOPE OVERLOAD DISTORTION:

We denote  $q(nT_s)$  as quantizing error. This  $q(nT_s)$  is

$$\rightarrow q(nT_s) = u(nT_s) - x(nT_s) \quad \text{--- (1)}$$



To eliminate  $u(nT_s - T_s)$  from (1) in DM, we write

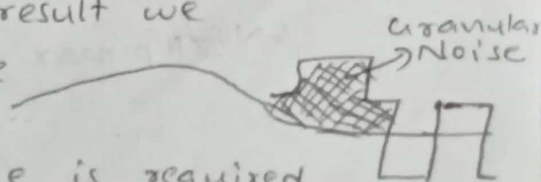
$$\rightarrow e(nT_s) = x(nT_s) - x(nT_s - T_s) - q(nT_s - T_s) \quad \text{--- (2)}$$

Except for quantization error,  $q(nT_s - T_s)$ , the quantizer i/p is a first backward difference of i/p signals, which may be viewed as digital approximation to the derivative of i/p signal (or) equivalently as a inverse of a digital integration process.

If we consider the maximum slope of the original i/p waveform of  $x(t)$ , it is clear that in order for the sequence of samples  $\{u(nT_s)\}$  to increase as fast as the i/p sequence of samples  $\{x(nT_s)\}$  is a region of maximum slope of, we require that the condition  $\frac{\delta}{T_s} \geq \max \left| \frac{dx(t)}{dt} \right|$ , otherwise slope overload will occur.

ii) GRANULAR NOISE:

Granular Noise occurs when stepsize  $\delta$  is too large relative to the slope characteristics of i/p waveform  $x(t)$ , thereby causing the staircase approximation  $u(t)$  to hunt around a relatively flat segment of i/p waveform. As a result we need to have a large step size to accommodate a wide dynamic range whereas a small step size is required for accurate representation of relatively low level signals.



It is therefore, clear that the choice of optimum step size that minimizes <sup>mean</sup> square value

of quantizer in linear DM will be a result of a compromise b/w slope overload distortion and granular noise.

Q: Find maximum slope of  $x(t) = a_0 \cos(2\pi f_0 t)$ .

Sol:

$$\text{Max slope} = \max \left| \frac{dx(t)}{dt} \right| = \max \left| \frac{d}{dt} (a_0 \cos(2\pi f_0 t)) \right|$$

$$= \max (-a_0 \sin(2\pi f_0 t) \cdot 2\pi f_0) = a_0 \cdot 2\pi f_0 \cdot 1$$

$$= a_0 \cdot 2\pi f_0$$

$$= \max | -2\pi f_0 a_0 \sin(2\pi f_0 t) |$$

$$\therefore \boxed{\text{max slope} = 2\pi f_0 a_0}$$

$$\therefore \frac{\delta}{T_s} \geq \max \left| \frac{dx(t)}{dt} \right| \Rightarrow \frac{\delta}{T_s} \geq 2\pi f_0 a_0 \quad \text{--- (3)}$$

$$\Rightarrow a_0 \leq \frac{\delta}{2\pi f_0 T_s} \quad \text{--- (4)}$$

$\Rightarrow a_0 \geq 2\pi f_0 T_s / \delta$   
Maximum permissible value of the o/p signal power

$$P_{\max} = \frac{a_0^2}{2} \quad \text{--- (5)} \Rightarrow P_{\max} = \frac{4\pi^2 f_0^2 T_s^2}{2\delta^2} = \frac{2\pi^2 f_0^2 T_s^2}{\delta^2}$$

$$\therefore \sigma_e^2 = \frac{\sigma^2}{12} \quad \text{--- (6)} \Rightarrow \sigma_e^2 = \frac{4\delta^2}{12} = \frac{\delta^2}{3}$$

Average power of quantization error is uniformly distributed over a frequency interval extending from  $-\frac{1}{T_s}$  to  $+\frac{1}{T_s}$ .

$$\therefore \text{Average o/p Noise power, } P_{\text{avg}} = \underbrace{W}_{\text{cut-off f. of LPF}} \cdot T_s \cdot \frac{\delta^2}{3} \quad \text{--- (7)}$$

$$\therefore (\text{SNR})_{\text{o/p max}} = \frac{P_{\max}}{\text{avg. noise power}}$$

$$= \frac{2\pi^2 f_0^2 T_s^2}{\delta^2} \times \frac{3}{W T_s \delta^2}$$

$$= \frac{6\pi^2 f_0^2 T_s}{\delta^4}$$