COL215 DIGITAL LOGIC AND SYSTEM DESIGN Divider Design

13 September 2017

Division: example

Unsigned Division

$$A = Q \times B + R$$

```
R = A; Q = 0; D = B
for i in 0 to n - 1 loop
if (D \times 2^{n-i-1} \le R) then
R = R - D \times 2^{n-i-1}
Q_{n-i-1} = 1
else Q_{n-i-1} = 0
end if
end loop
```

Unsigned Division

$$A = Q \times B + R$$

R = A; Q = 0; D = B
for i in 0 to n - 1 loop
if (D x
$$2^{n-i-1}

R = R - D x 2^{n-i-1}
 $Q_{n-i-1} = 1$
else $Q_{n-i-1} = 0$
end if
end loop$$

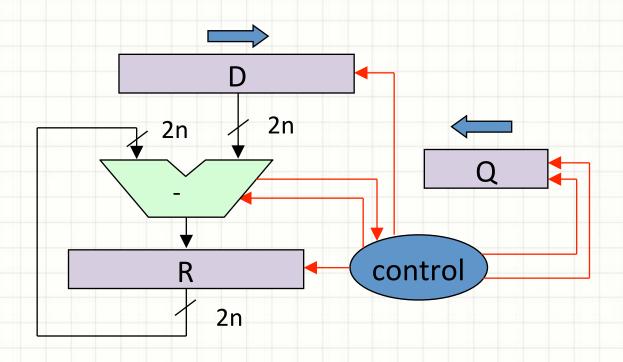
can this be avoided?

Introducing shift registers

$$A = Q \times B + R$$

```
R = A; Q = 0; D = B \times 2^{n-1}
for i in 0 to n – 1 loop
  if (D \le R) then
     R = R - D
     Q = 2 \times Q + 1
  else Q = 2 \times Q
  end if
  D = D / 2
end loop
```

Divider design - 1

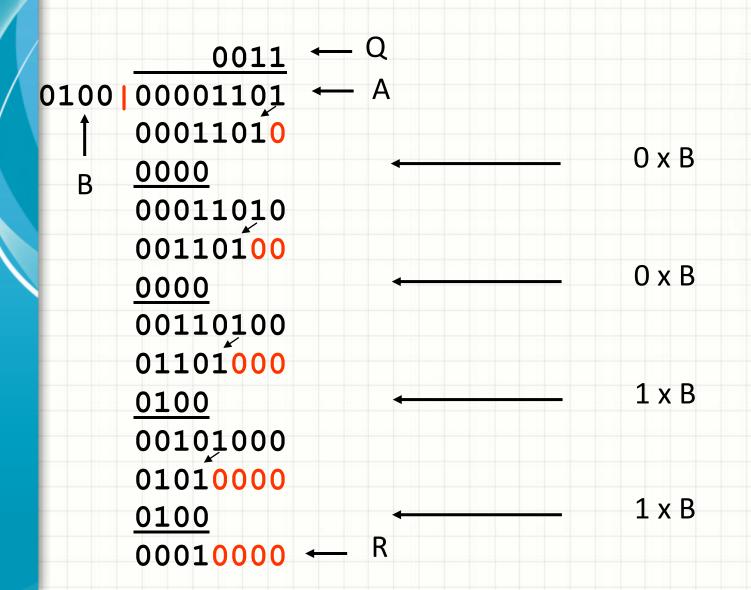


Reducing subtractor size

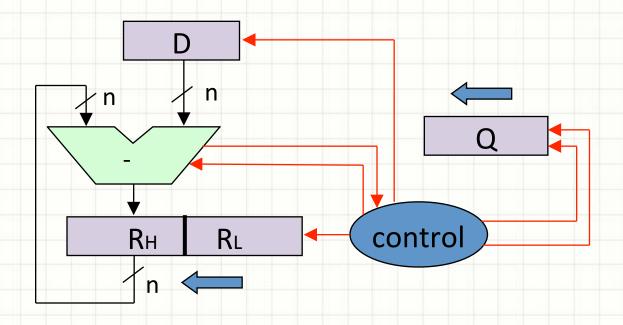
$$A = Q \times B + R$$

```
R = A; Q = 0; D = B
for i in 0 to n – 1 loop
  R = 2 \times R
  if (D \le R_H) then
     R_H = R_H - D
    Q = 2 \times Q + 1
  else Q = 2 \times Q
  end if
end loop
```

Division: example



Divider design - 2

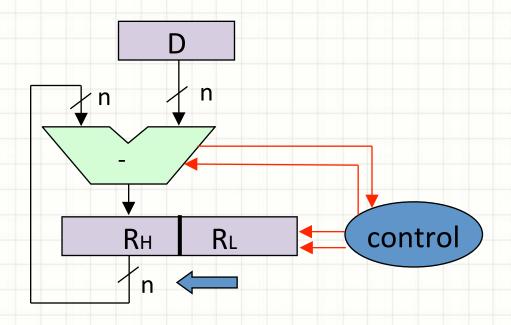


Reducing registers

$$A = Q \times B + R$$

```
R = A; D = B
for i in 0 to n – 1 loop
  R = 2 \times R
  if (D \le R_H) then
    R_H = R_H - D
    R = R + 1
  end if
end loop
\# R_H = remainder, R_I = quotient
```

Divider design - 3



Restoring division

- Non-restoring division
 - First compare the remaining dividend with divisor
 - Subtract only if dividend is large enough
- Restoring division
 - Subtract without comparing
 - Add (restore) the divisor if subtraction result negative

Signed multiplication/division

- Handle sign and magnitude separately
- Directly multiply/divide signed integers

Dividend	Divisor	Quotient	Remainder
+	+	+	+
-	+	-	-
+	-		+
-	-	+	-

Direct signed multiplication

 Use a common expression representing the values of positive as well as negative integers

$$B = -B_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} B_i \cdot 2^i$$

Common expr for +/- integers

for
$$B \ge 0$$
, $B = \sum_{i=0}^{n-1} B_i \cdot 2^i = -B_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} B_i \cdot 2^i$ (:: $B_{n-1} = 0$)

for B < 0, B =
$$-|B|$$

now
$$|B| = 2^n - \sum_{i=0}^{n-1} B_i \cdot 2^i$$

$$\therefore B = -2^{n} + B_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} B_{i} \cdot 2^{i}$$

$$= -B_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} B_i \cdot 2^i \quad (:: B_{n-1} = 1)$$

Direct signed multiplication

$$\begin{split} B &= -B_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} B_i \cdot 2^i \\ &= -B_{n-1} \cdot 2^{n-1} + B_{n-2} \cdot 2^{n-2} + \dots + B_0 \cdot 2^0 \\ &= -B_{n-1} \cdot 2^{n-1} + 2B_{n-2} \cdot 2^{n-2} - B_{n-2} \cdot 2^{n-2} + \dots + 2B_0 \cdot 2^0 - B_0 \cdot 2^0 \\ &= -B_{n-1} \cdot 2^{n-1} + B_{n-2} \cdot 2^{n-1} - B_{n-2} \cdot 2^{n-2} + \dots + B_0 \cdot 2^1 - B_0 \cdot 2^0 \\ &= \sum_{i=0}^{n-1} (B_{i-1} - B_i) \cdot 2^i \qquad \text{where } B_{-1} = 0 \\ &\therefore A \cdot B = \sum_{i=0}^{n-1} A \cdot (B_{i-1} - B_i) \cdot 2^i \end{split}$$

Booth's algorithm

Comparing with unsigned case

unsigned multiplication

signed multiplication

B_i operation

B_i, B_{i-1} operation

0 no addition

1 add A

0 0 no addtion

0 1 add A

1 0 subtract A

1 1 no addition

Original motivation for Booth's algorithm

† † add subtract

- +1000000
- <u>0000001</u> 0111111

What have we learnt?

- Logic design (combinational circuits) [2, 4]
 - truth tables, expressions, circuits, VHDL
- Logic design (sequential circuits) [7, 8]
 - state transition tables, diagrams, circuits, VHDL
- Combinational & sequential modules [6, 7]
 - mux, demux, decoders, encoders, VHDL
 - flip-flops, registers, counters, VHDL
- From logic to arithmetic

[5]

- representations, conversions
- operations and operators (add, subtract, compare, multiply, divide)

What lies ahead?

- Technology [3]
 - transistor to FPGA and things in between
- A little more theory [4, 8, 9]
 - minimizing logic, minimizing states
- System design [10]
 - from algorithmic description to circuits
 - control-data partition
- Testing [11]
 - testing tools, design for testability

