

CJFS: Concurrent Journaling for Better Scalability

Joontaek Oh*, Seung Won Yoo*, Hojin Nam*, Changwoo Min†, Youjip Won*

*KAIST

†Virginea Tech





Outline

- ✓ Background and Motivation
- ✓ Design
 - Dual Thread Journaling
 - Multi-Version Shadow Paging
 - Opportunistic Coalescing
 - Compound Flush
- Evaluation
- ✓ Conclusion

Background and Motivation

Hardware: Software:



2 cores
Intel Core 2 Duo
@2006



877 IOPS
Western Digital Caviar SE16
@2006

Hardware:



2 cores
Intel Core 2 Duo
@2006

128 cores AMD EPYC 7763 @2021







877 IOPS
Western Digital Caviar SE16
@2006

700K IOPS Seagate FireCuda 530 @2021

Hardware:







2 cores Intel Core 2 Duo @2006

128 cores AMD EPYC 7763 @2021



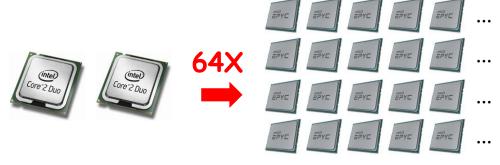




877 IOPS
Western Digital Caviar SE16
@2006

700K IOPS Seagate FireCuda 530 @2021

Hardware:



2 cores
Intel Core 2 Duo
@2006

128 cores AMD EPYC 7763 @2021

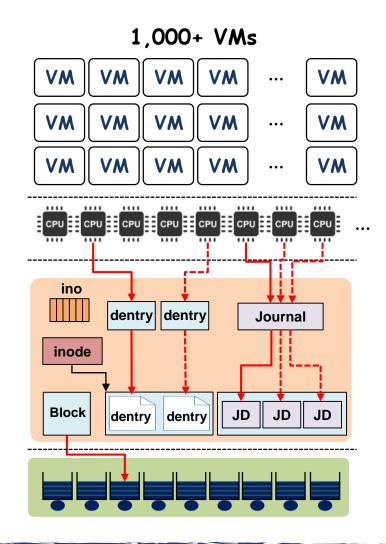


798X →

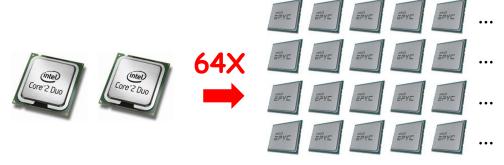


877 IOPS
Western Digital Caviar SE16
@2006

700K IOPS
Seagate FireCuda 530
@2021



Hardware:



2 cores
Intel Core 2 Duo
@2006

128 cores AMD EPYC 7763 @2021

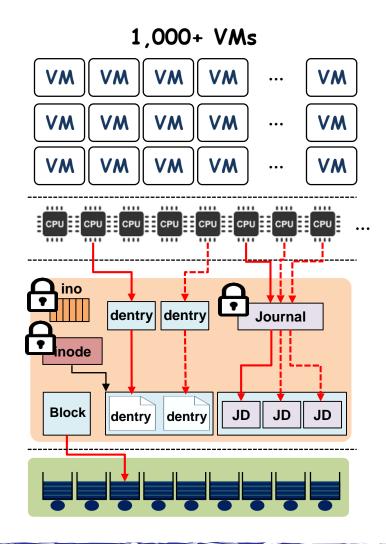


798X →

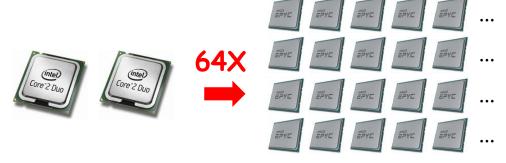


877 IOPS
Western Digital Caviar SE16
@2006

700K IOPS Seagate FireCuda 530 @2021



Hardware:



2 cores
Intel Core 2 Duo
@2006

128 cores AMD EPYC 7763 @2021

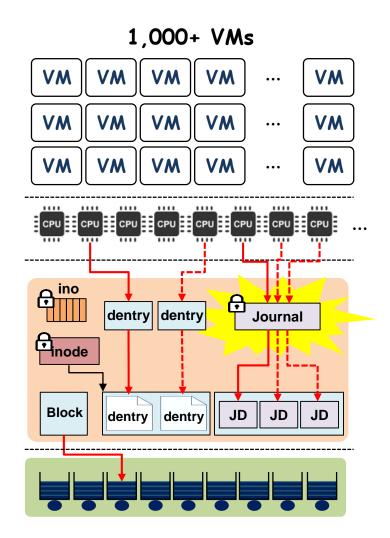


798X →

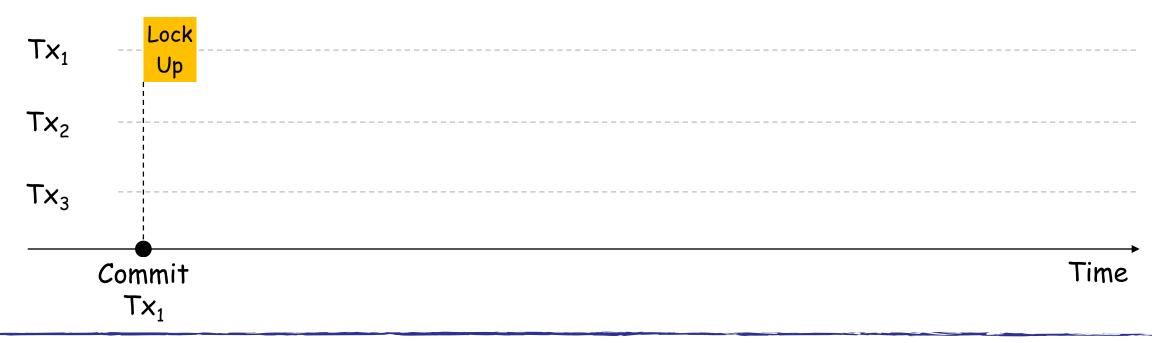


877 IOPS
Western Digital Caviar SE16
@2006

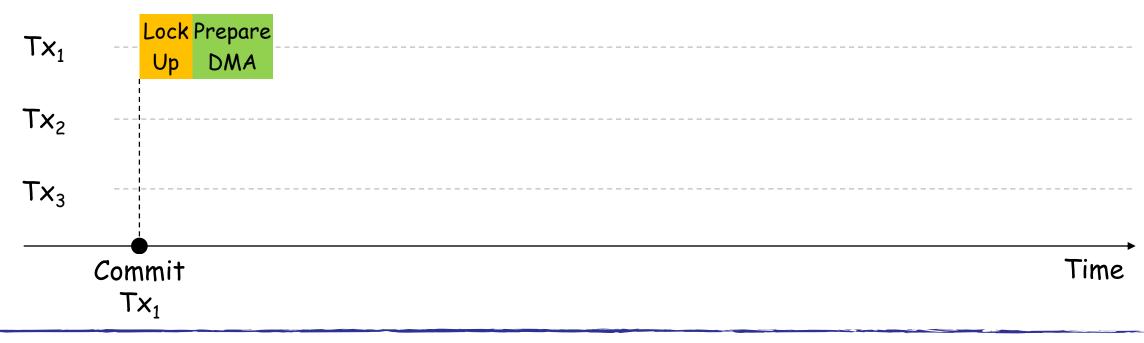
700K IOPS
Seagate FireCuda 530
@2021



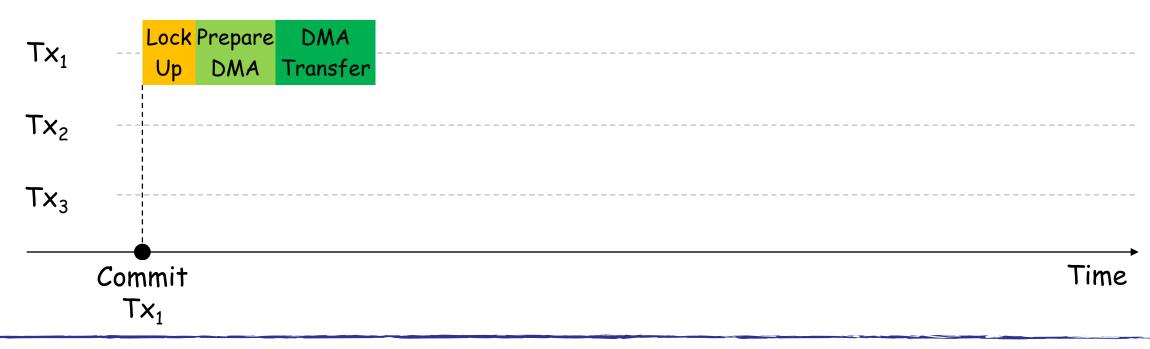
- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation



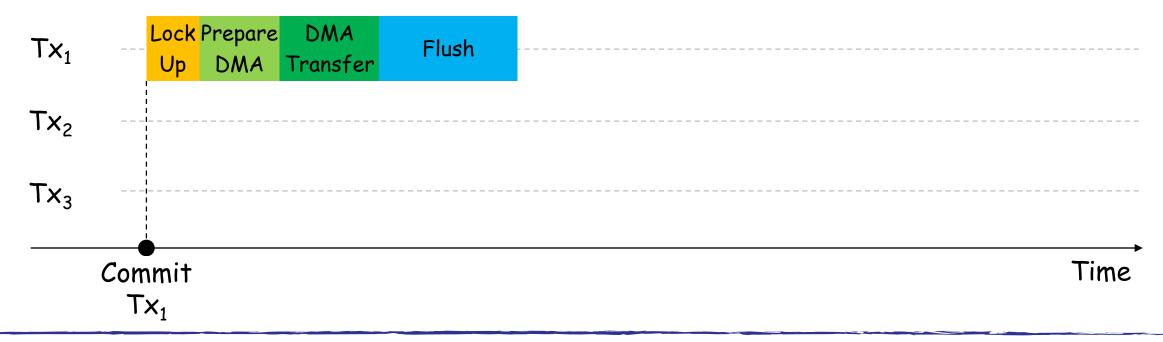
- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction



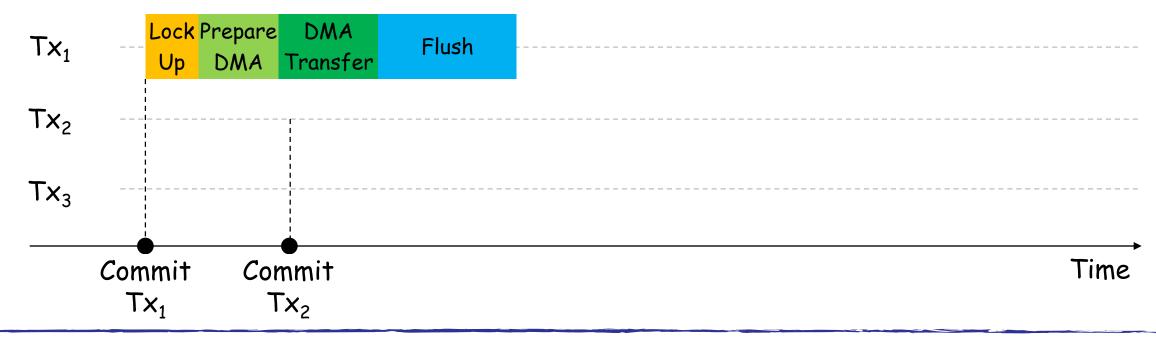
- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction
 - DMA Transfer: Waiting for the completion of DMA Transfer of the transaction



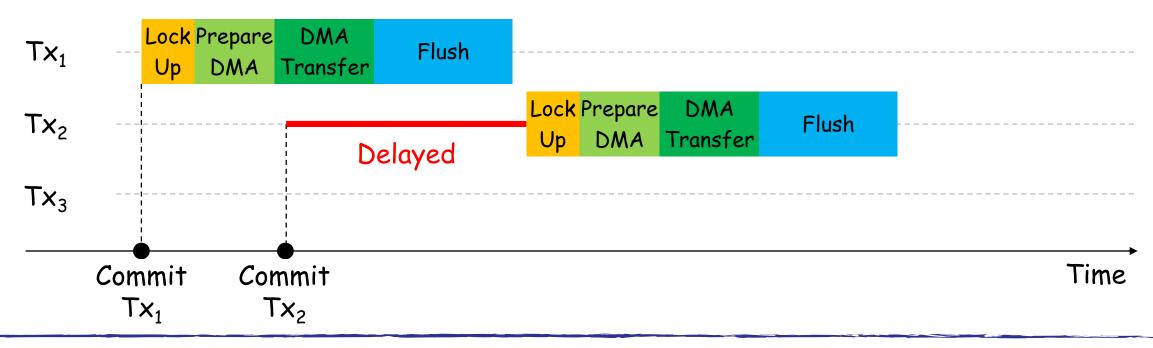
- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction
 - DMA Transfer: Waiting for the completion of DMA Transfer of the transaction
 - Flush: Flush transferred data



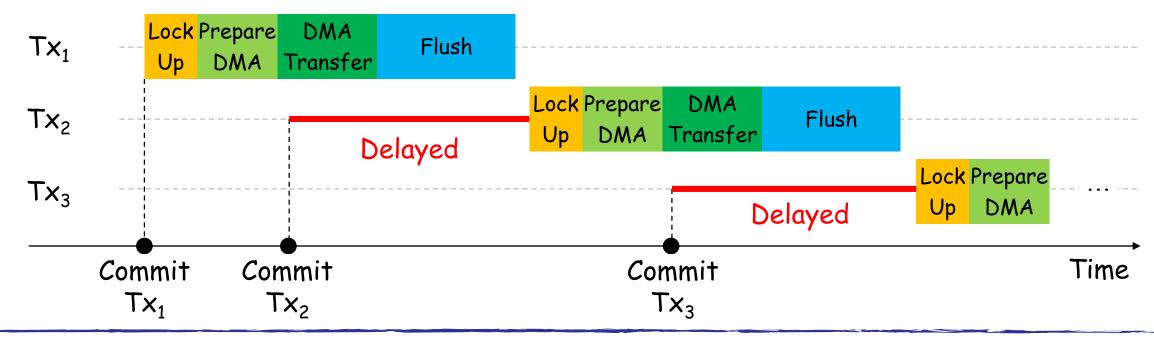
- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction
 - DMA Transfer: Waiting for the completion of DMA Transfer of the transaction
 - Flush: Flush transferred data



- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction
 - DMA Transfer: Waiting for the completion of DMA Transfer of the transaction
 - Flush: Flush transferred data

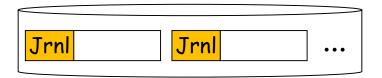


- All steps of journal commit are serialized
 - Lock-Up: Lock the running transaction and waiting for remained file operation
 - Prepare DMA: Create and dispatch write command for the transaction
 - DMA Transfer: Waiting for the completion of DMA Transfer of the transaction
 - Flush: Flush transferred data



Multiple journal regions:

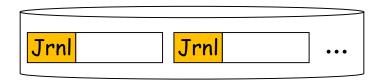
IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)



Multiple journal regions:

IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)

Still serial transaction commit in each journal region



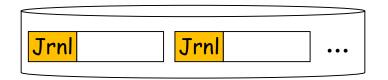
Multiple journal regions:

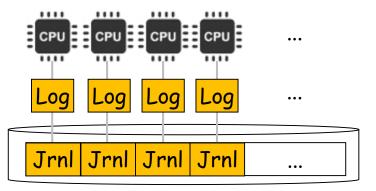
IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)

Still serial transaction commit in each journal region

Per-core running transaction:

ScaleFS (SOSP '17), MQFS (SOSP '21)





Multiple journal regions:

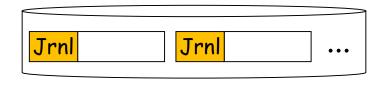
IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)

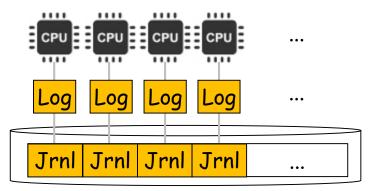
Still serial transaction commit in each journal region

Per-core running transaction:

ScaleFS (SOSP '17), MQFS (SOSP '21)

Conflict between multiple transactions and Still serial commit





Multiple journal regions:

IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)

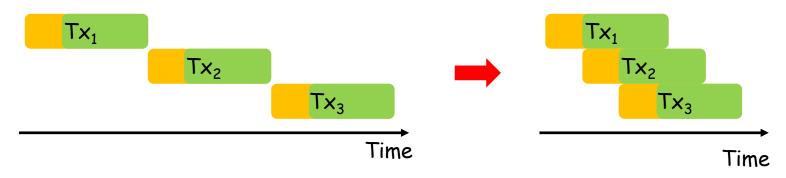
Still serial transaction commit in each journal region

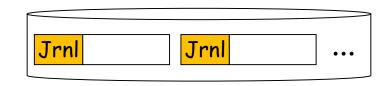


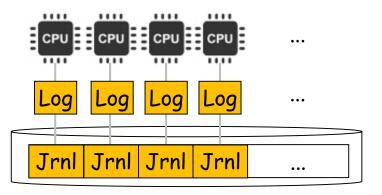
ScaleFS (SOSP '17), MQFS (SOSP '21)

Conflict between multiple transactions and Still serial commit

Parallel journal commit: BarrierFS (FAST '18)







Multiple journal regions:

IceFS (OSDI '14), SpanFS (ATC '15), Z-journal (ATC'21)

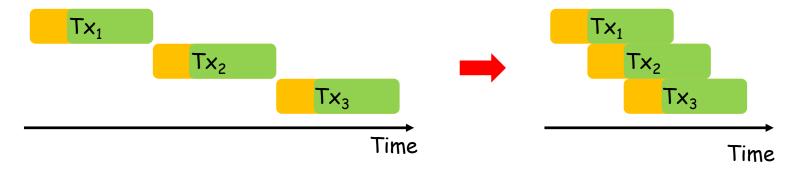
Still serial transaction commit in each journal region

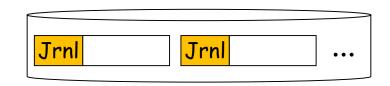


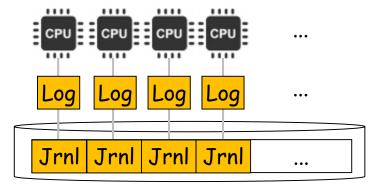
ScaleFS (SOSP '17), MQFS (SOSP '21)

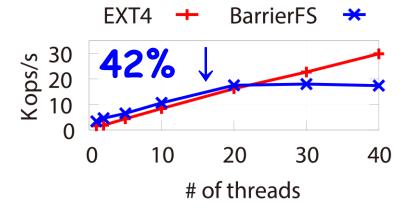
Conflict between multiple transactions and Still serial commit

Parallel journal commit: BarrierFS (FAST '18)









Main reasons

- Transaction conflict
- Transaction Lock-Up

The situation that a file operation modifies a page which is being committed

S User	
∃ JBD	
Tx_1	

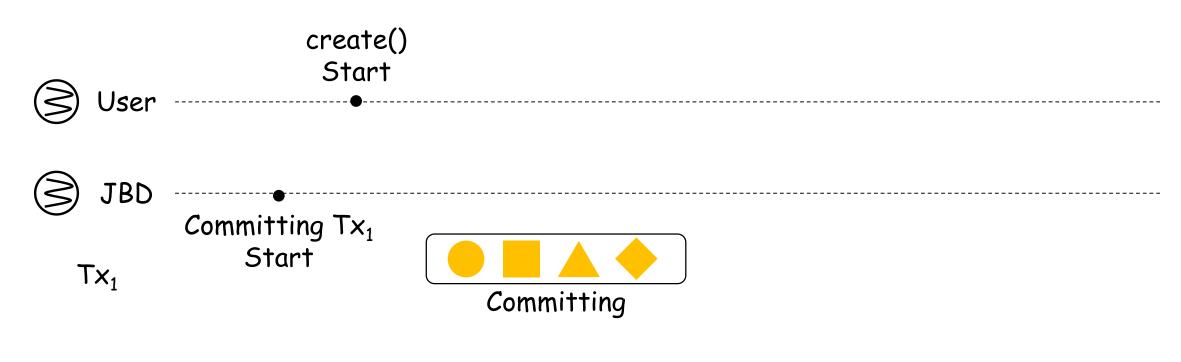
 Tx_2

The situation that a file operation modifies a page which is being committed



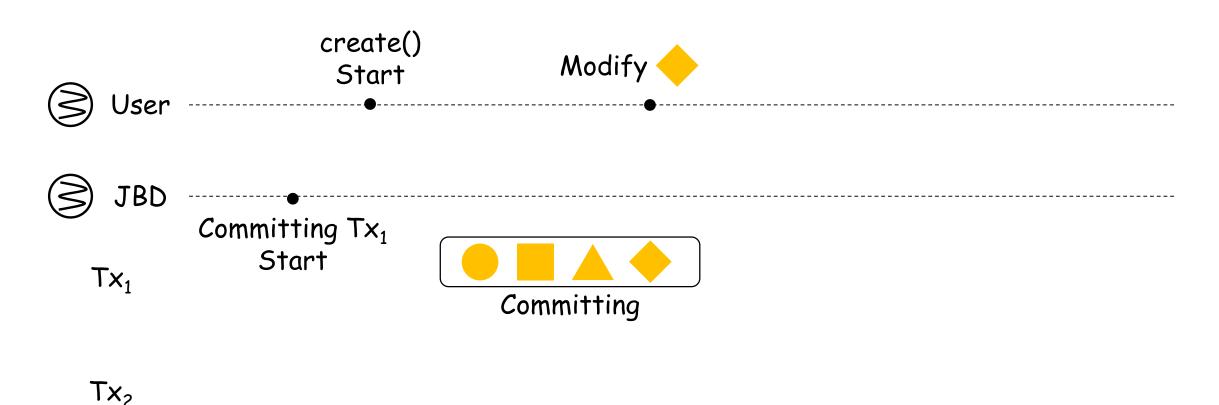
Time

The situation that a file operation modifies a page which is being committed



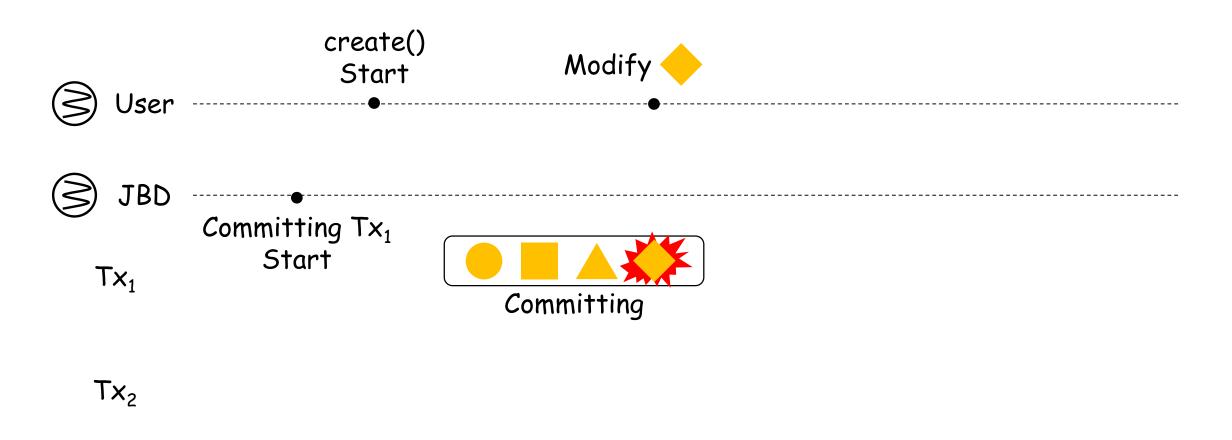
 Tx_2

The situation that a file operation modifies a page which is being committed



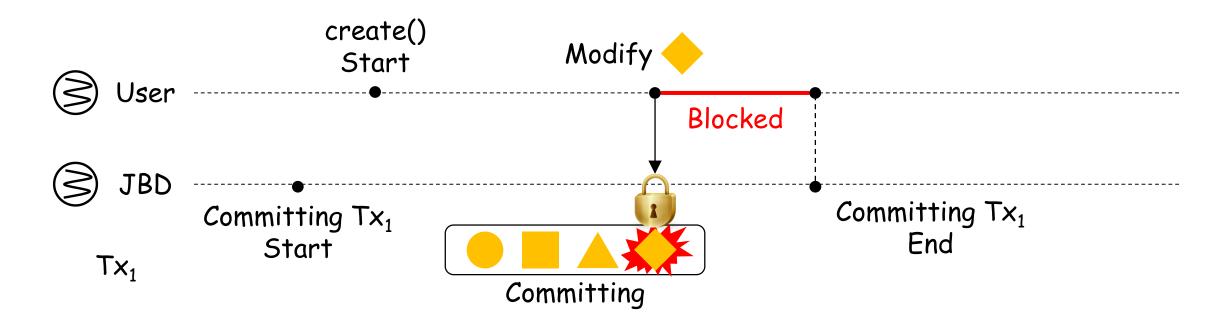
Time

The situation that a file operation modifies a page which is being committed



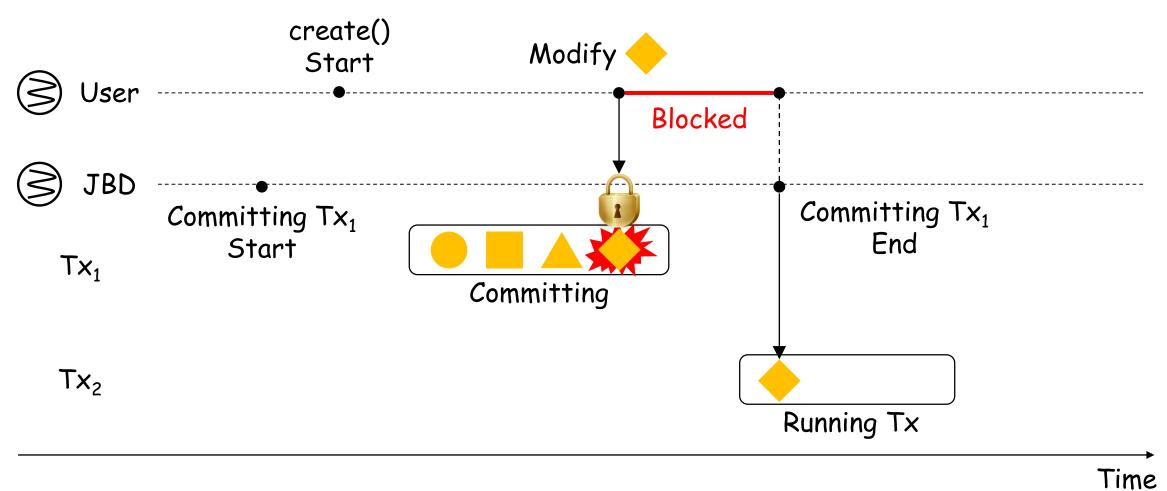
Time

The situation that a file operation modifies a page which is being committed

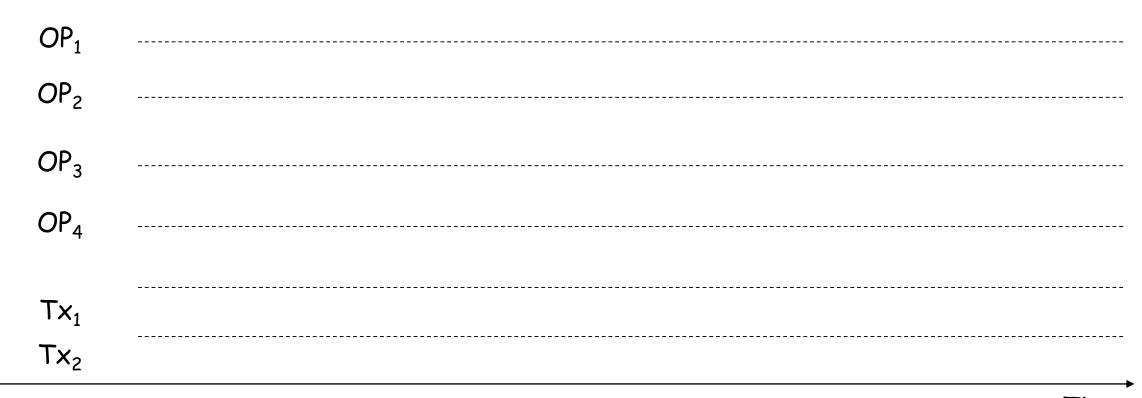


 Tx_2

The situation that a file operation modifies a page which is being committed

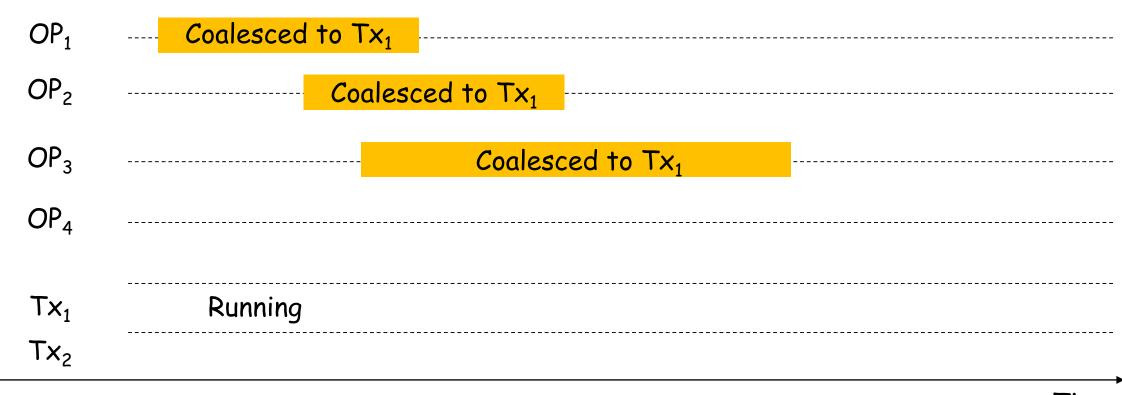


the locked period for isolating the running transaction from file operations



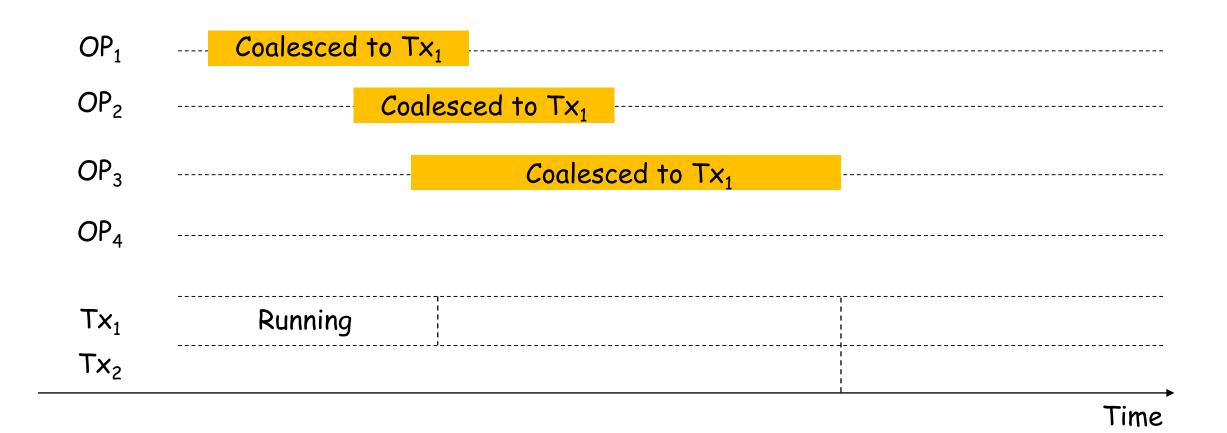
Time

the locked period for isolating the running transaction from file operations

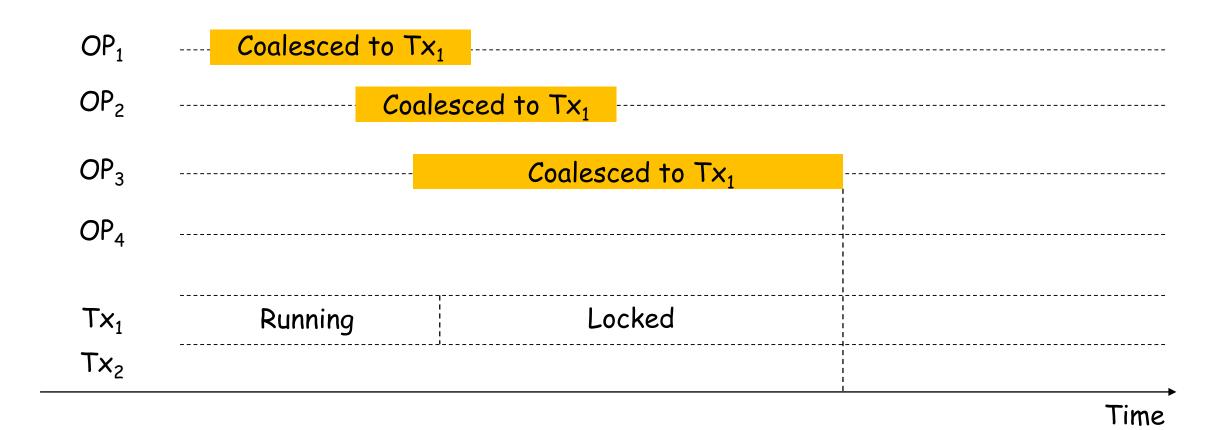


Time

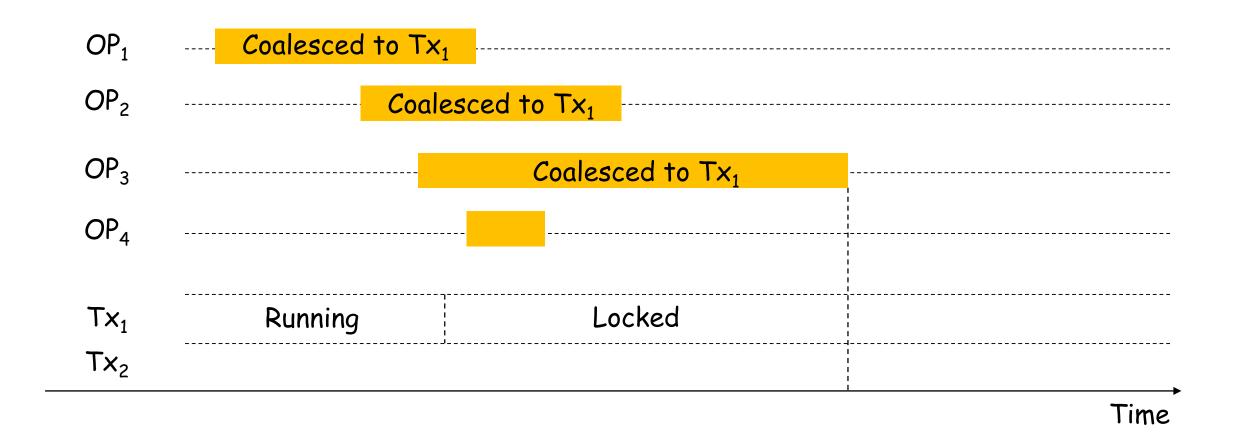
the locked period for isolating the running transaction from file operations



the locked period for isolating the running transaction from file operations

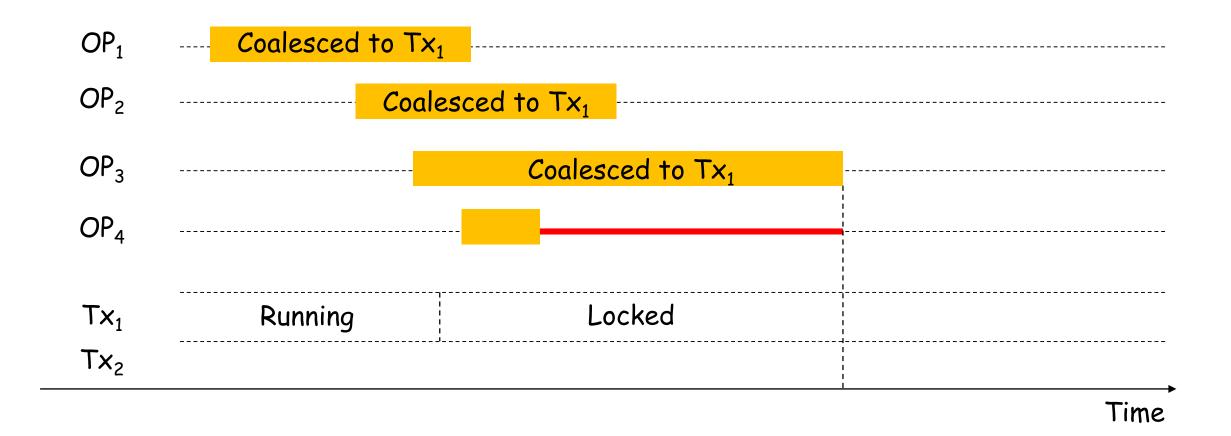


the locked period for isolating the running transaction from file operations



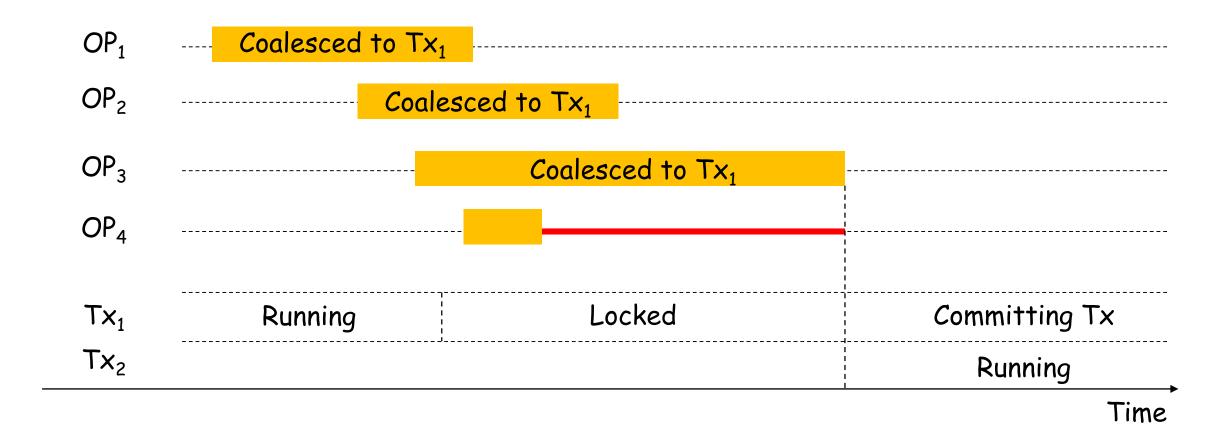
Transaction Lock-Up

the locked period for isolating the running transaction from file operations



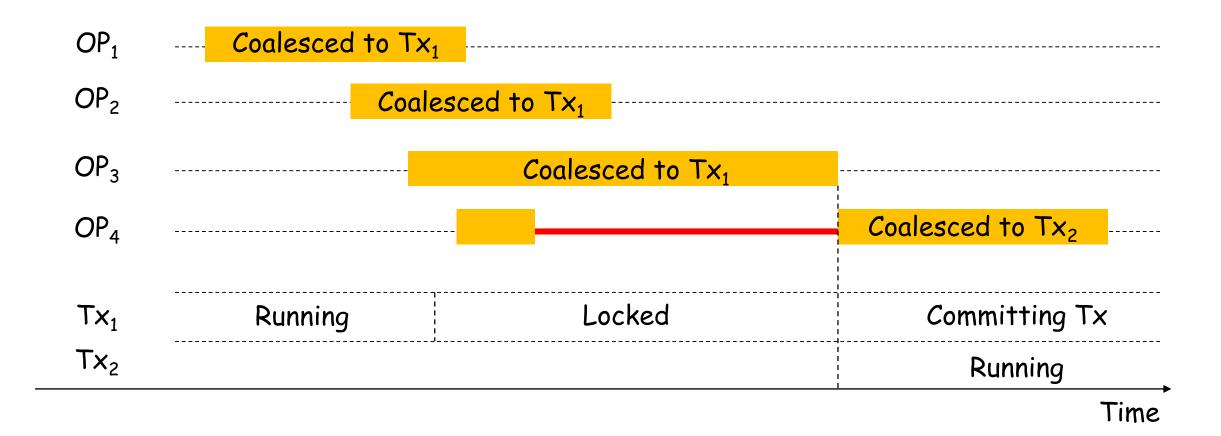
Transaction Lock-Up

the locked period for isolating the running transaction from file operations



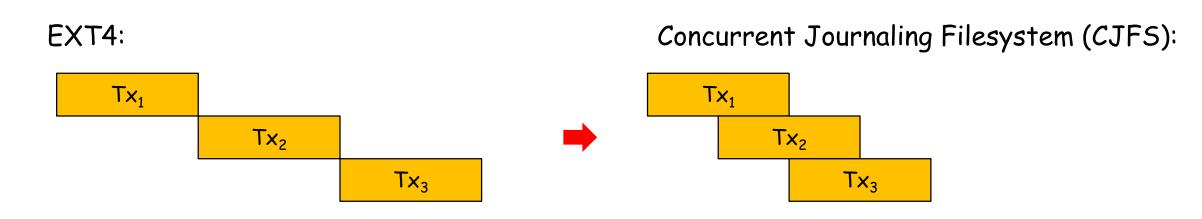
Transaction Lock-Up

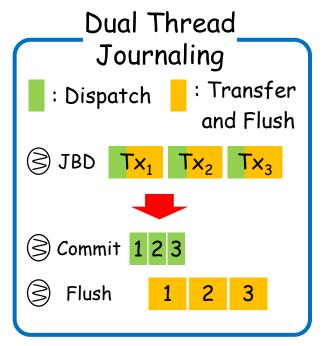
the locked period for isolating the running transaction from file operations

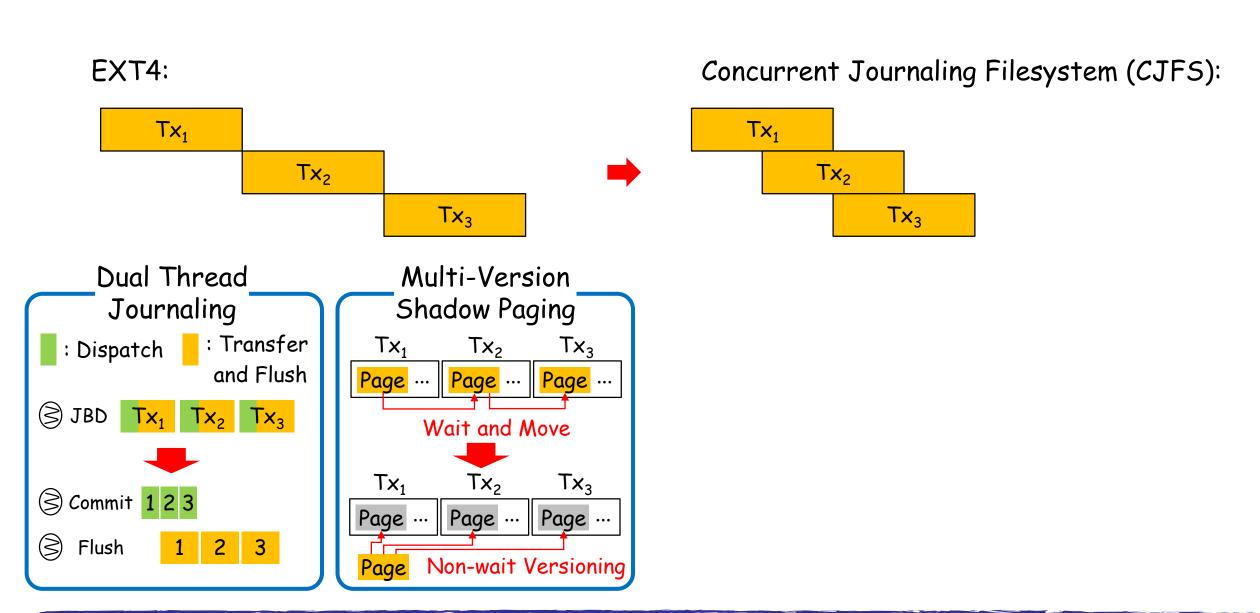


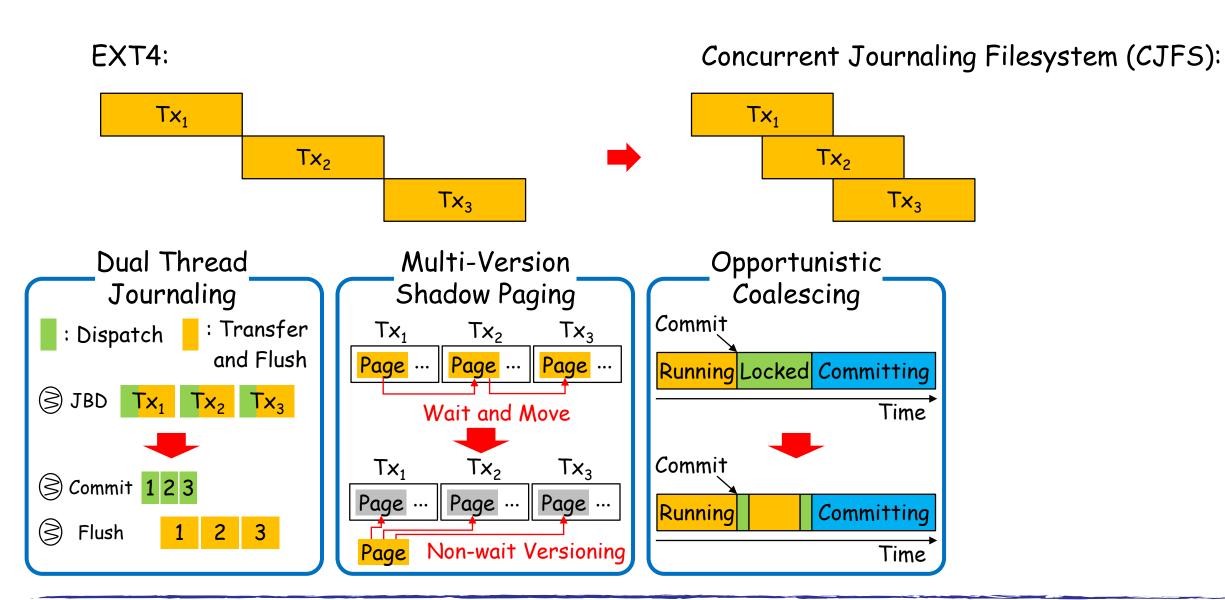
39

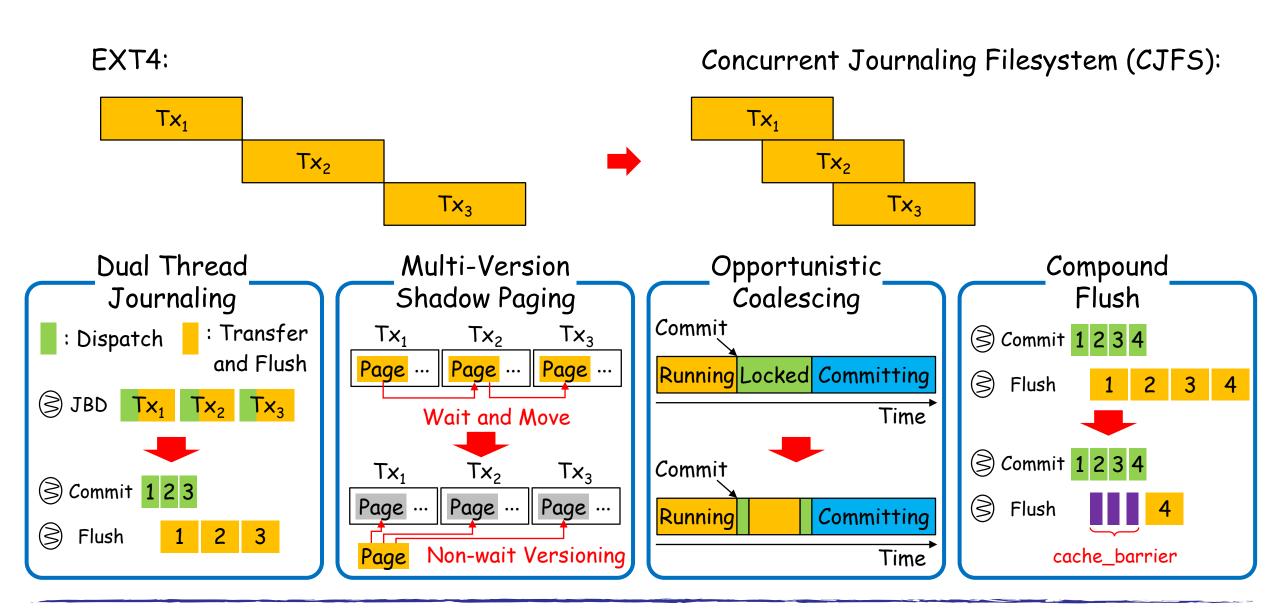
Design: Concurrent Journaling Filesystem (CJFS)



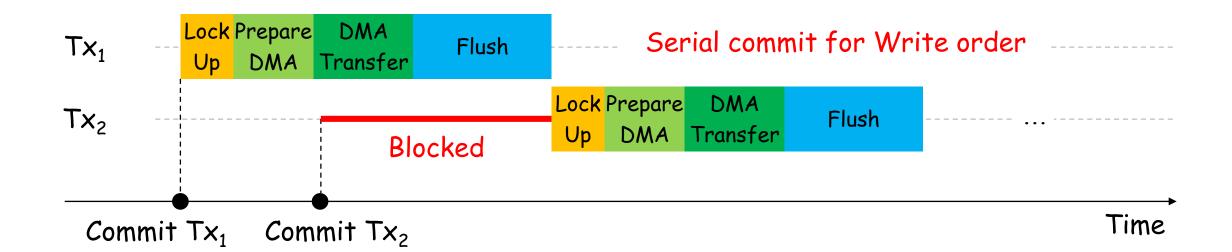




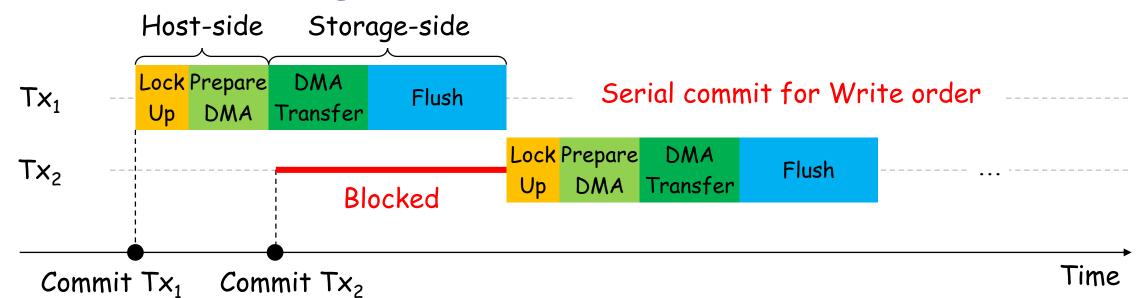




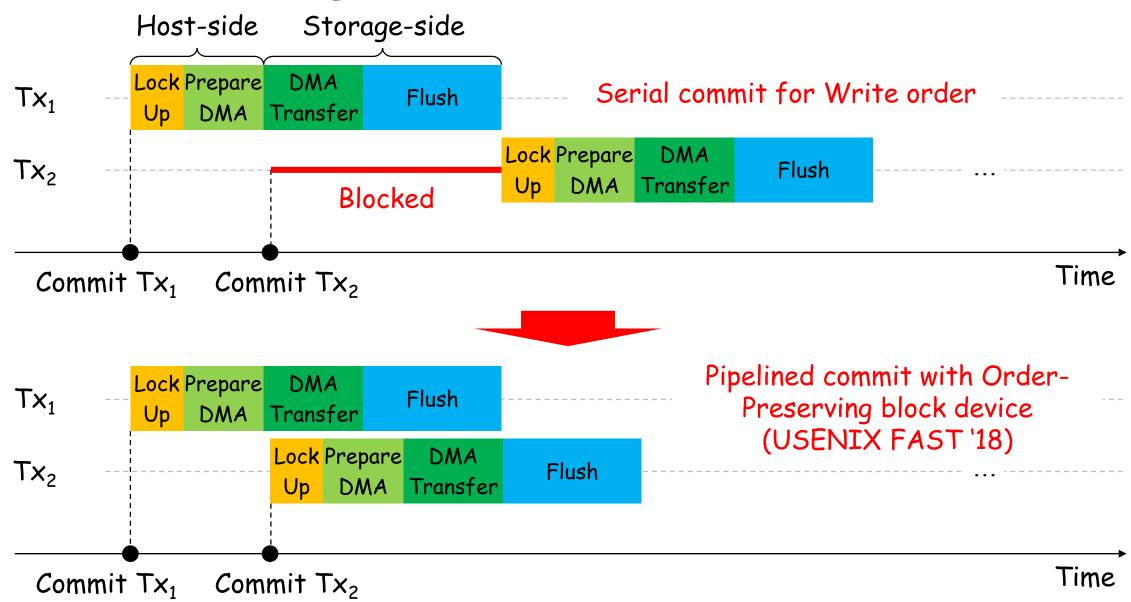
Dual Thread Journaling

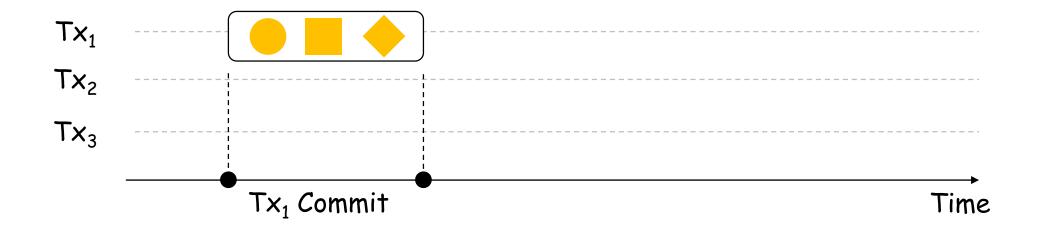


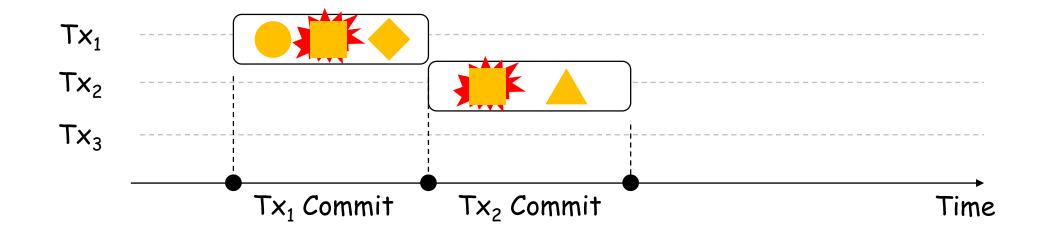
Dual Thread Journaling

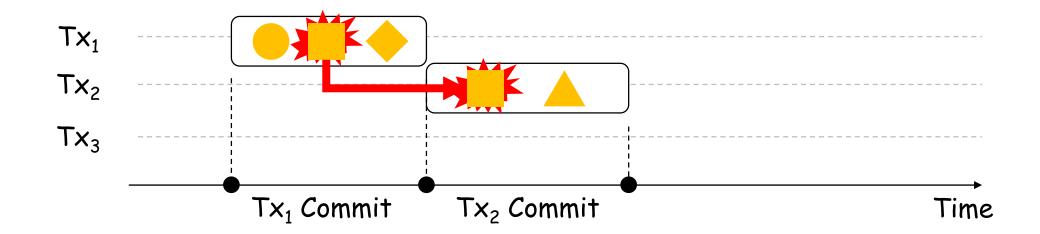


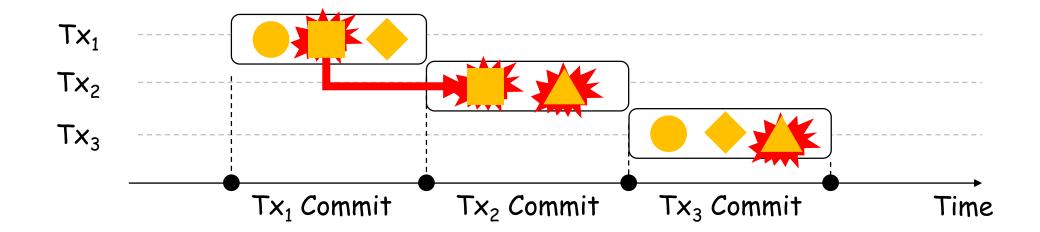
Dual Thread Journaling

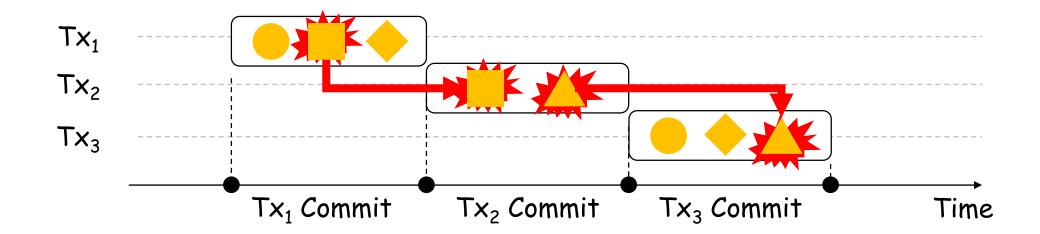


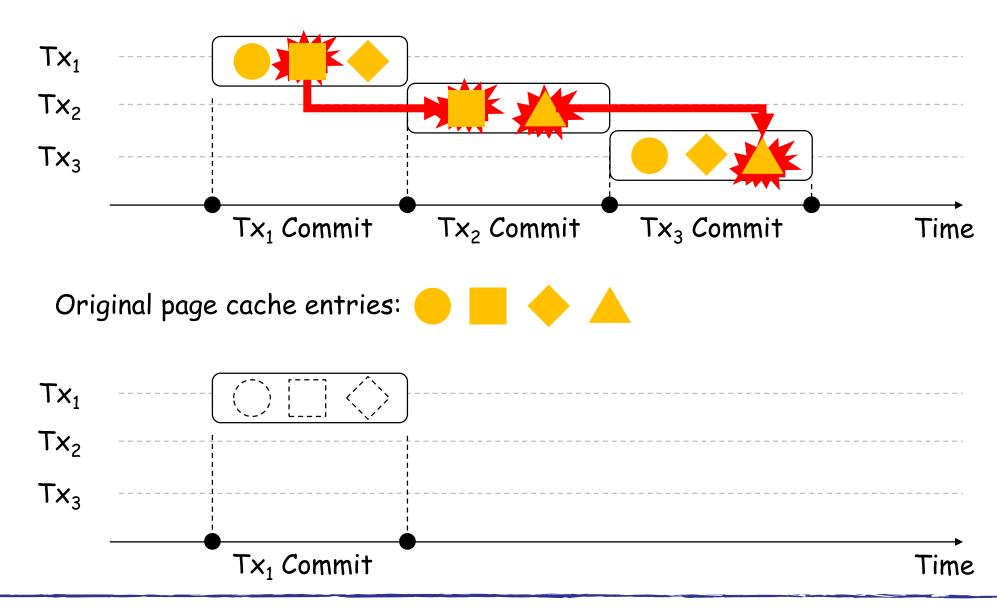


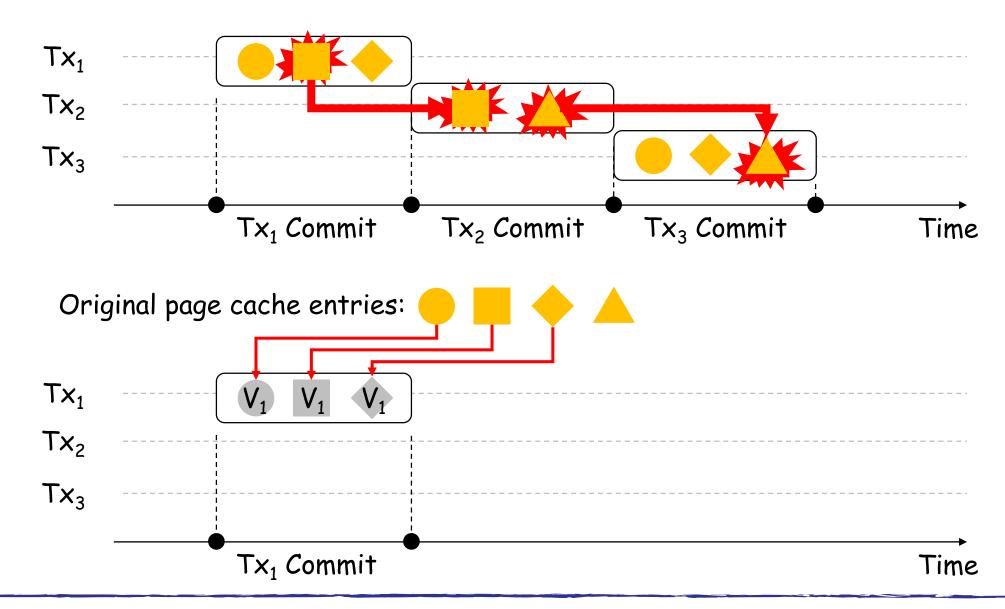


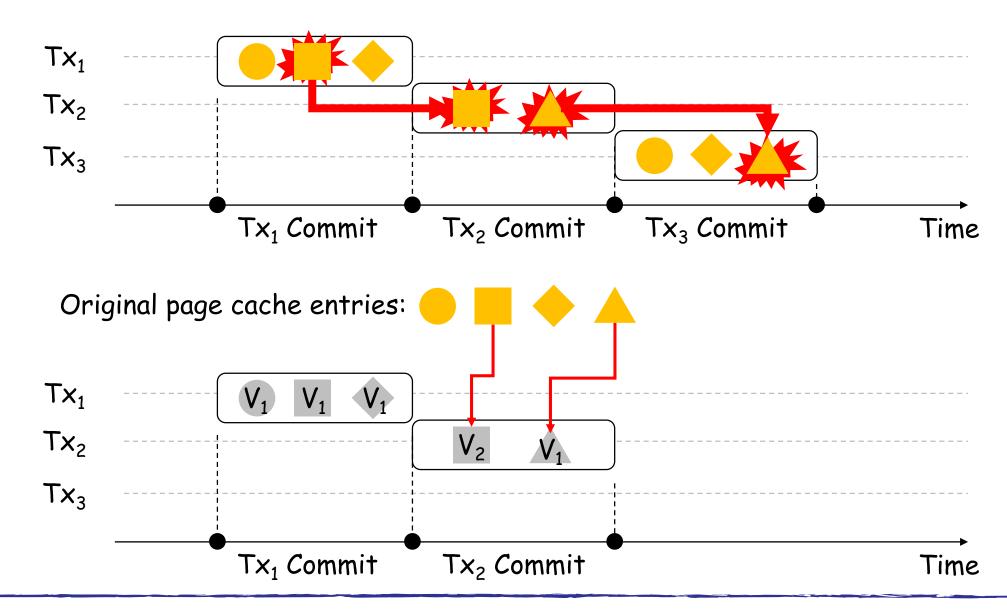


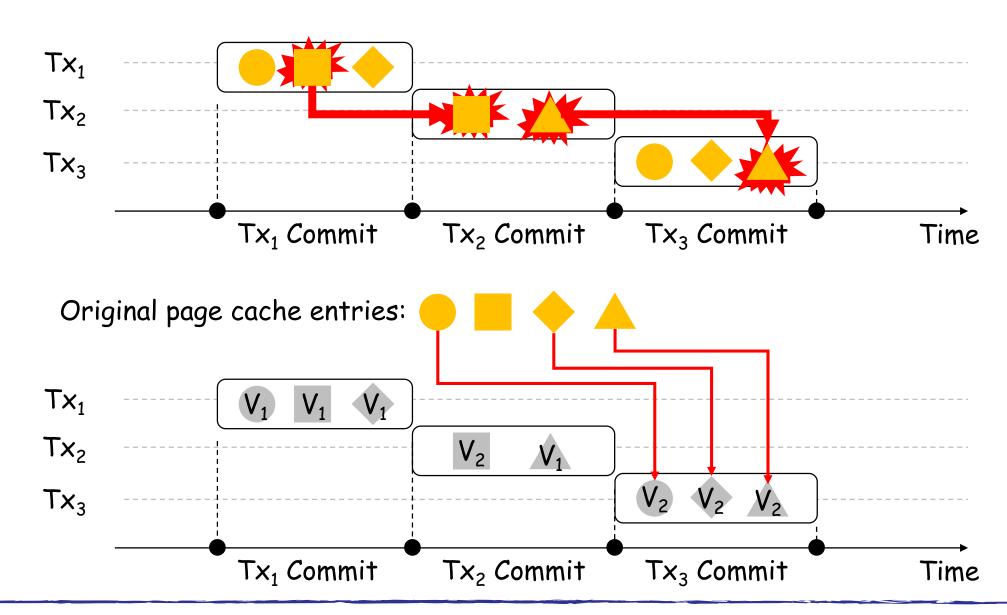


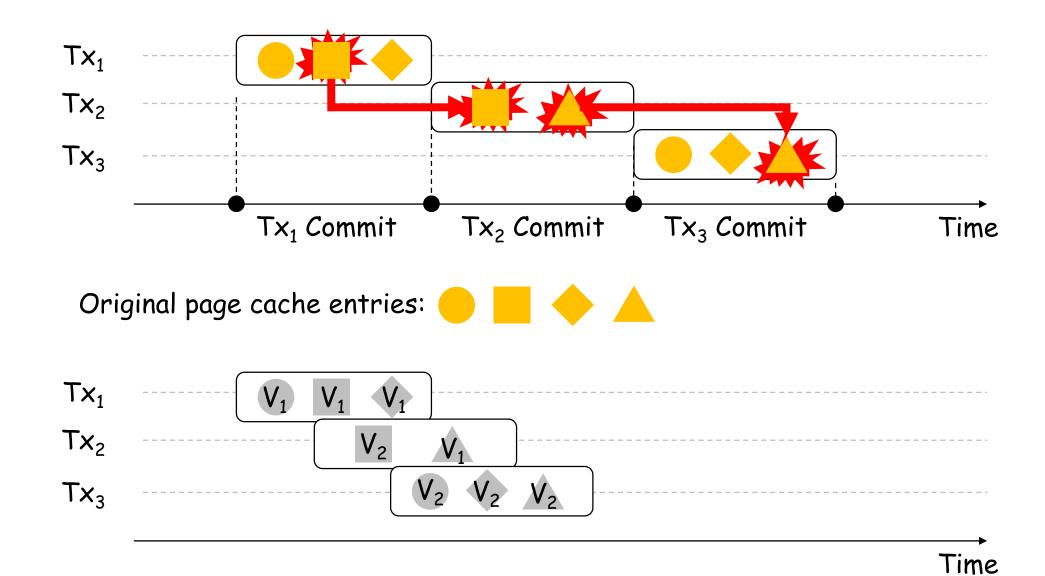


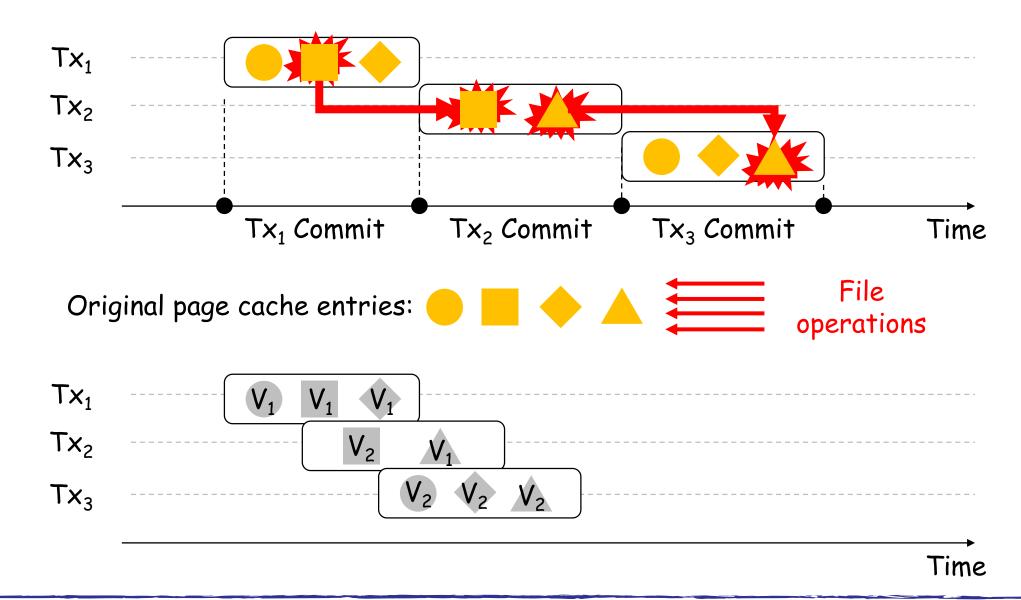




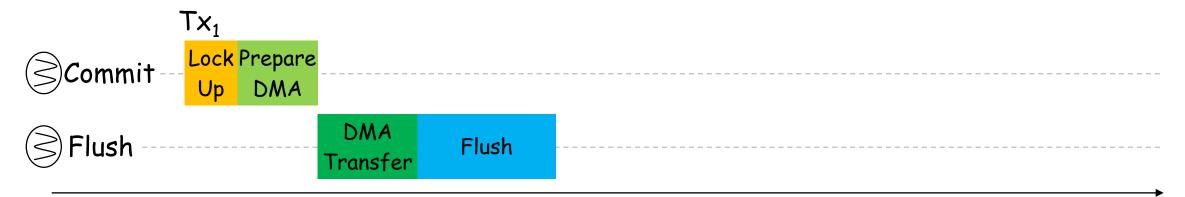






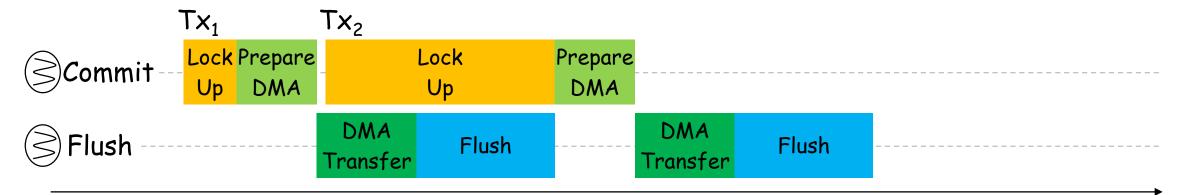


- When versions are exhausted, transaction commits are serialized
- The running transaction is locked and waits for preceding transaction commits



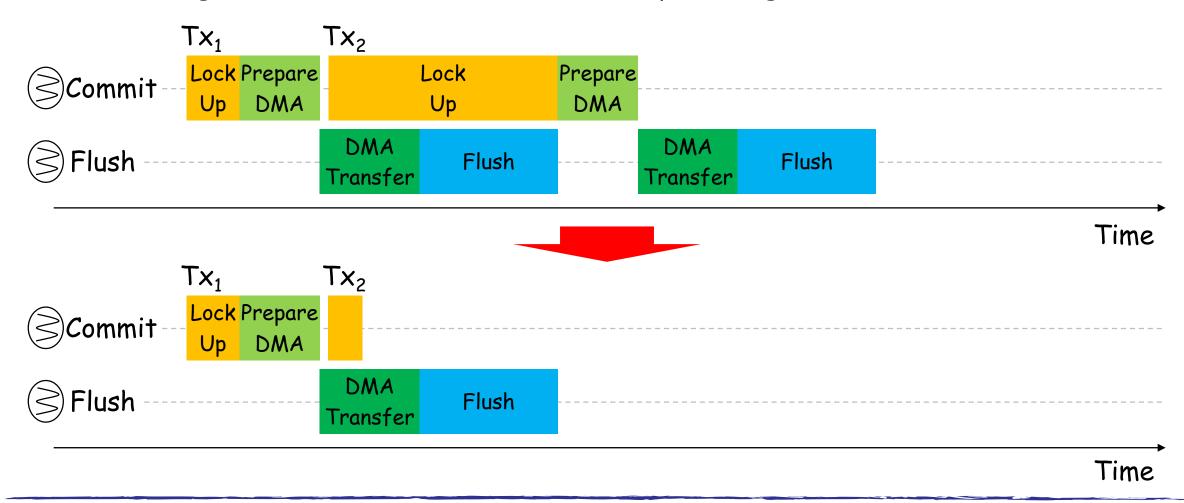
Time

- When versions are exhausted, transaction commits are serialized
- The running transaction is locked and waits for preceding transaction commits

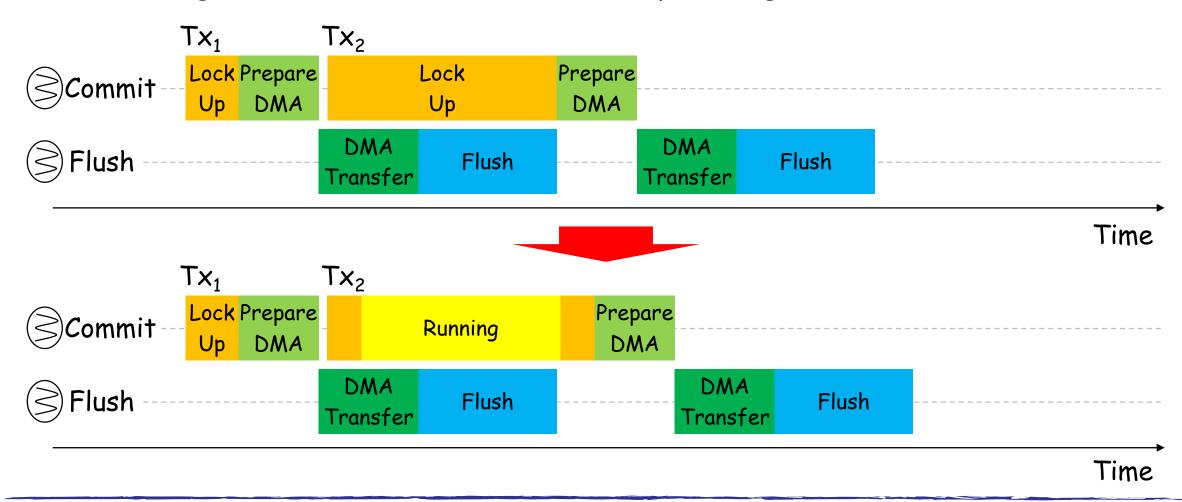


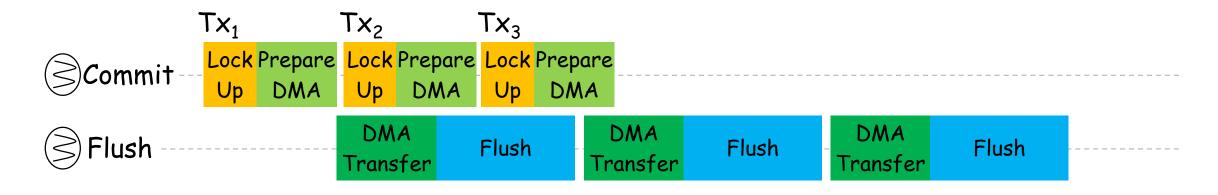
Time

- When versions are exhausted, transaction commits are serialized
- The running transaction is locked and waits for preceding transaction commits

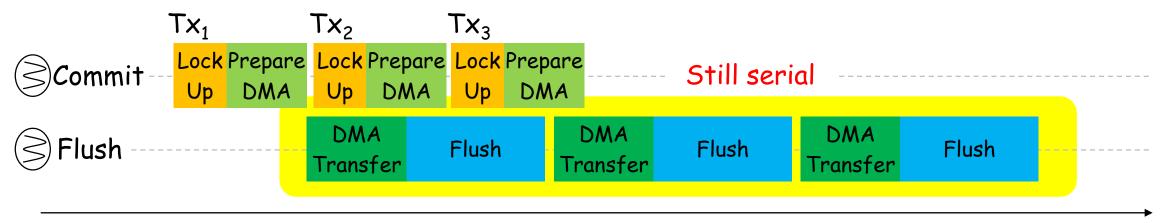


- When versions are exhausted, transaction commits are serialized
- The running transaction is locked and waits for preceding transaction commits

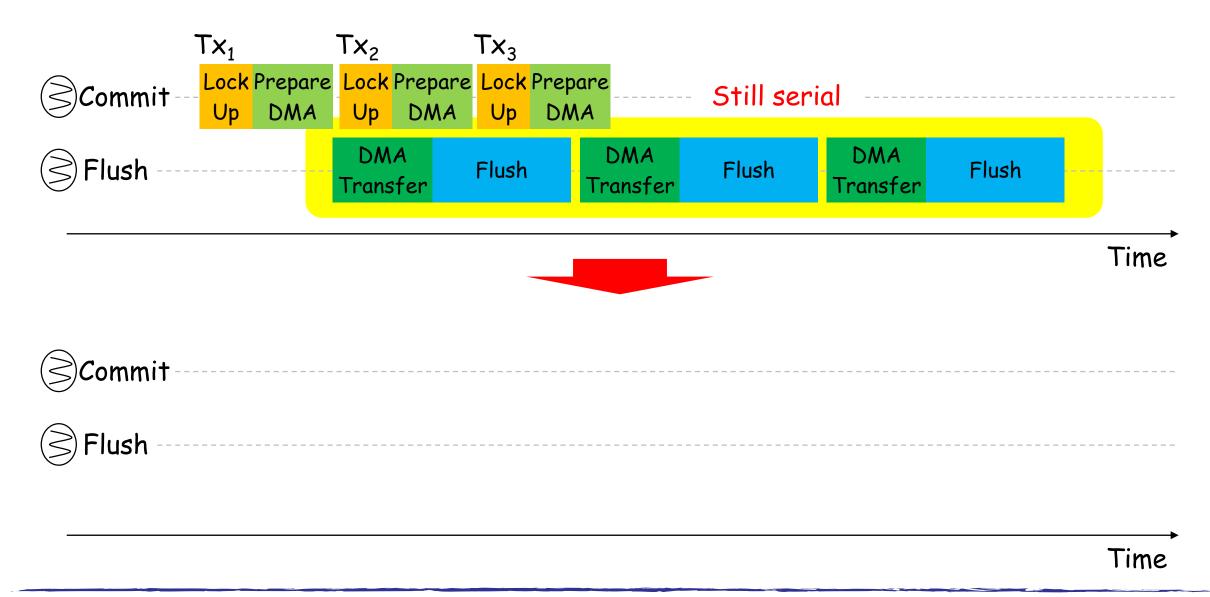


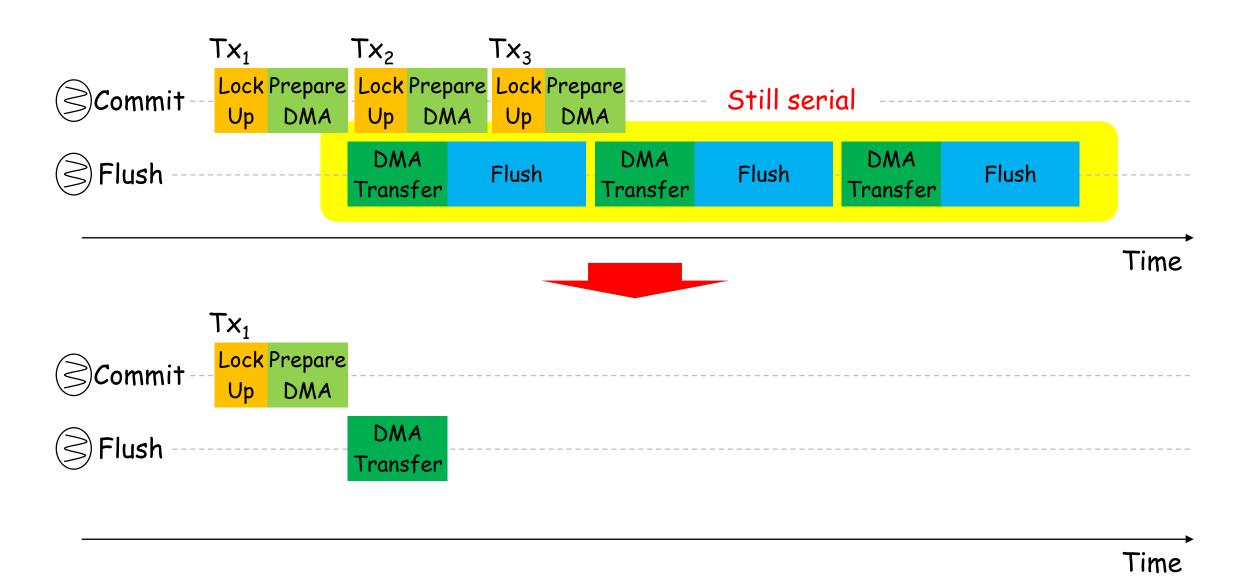


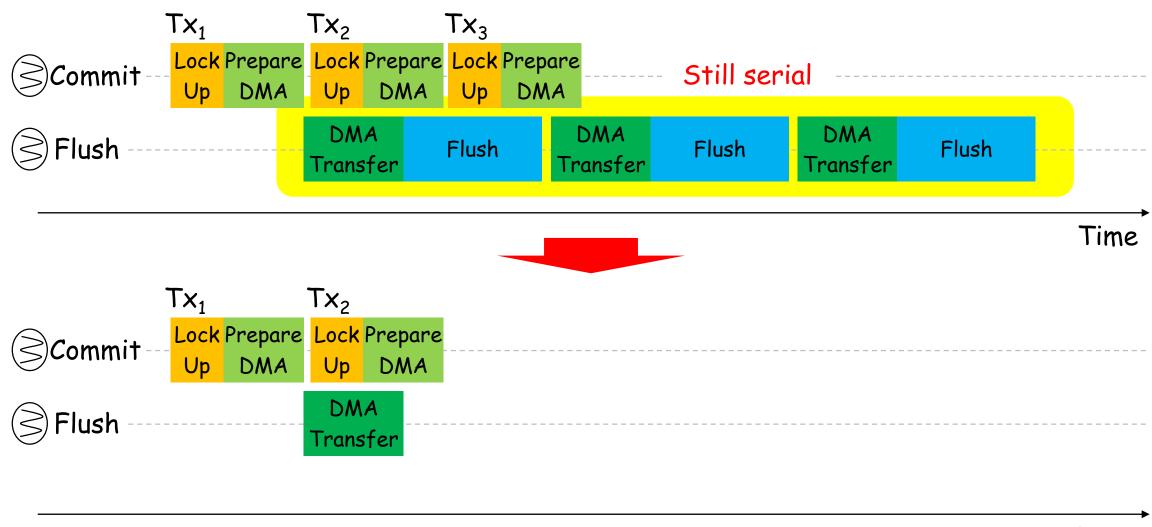
Time



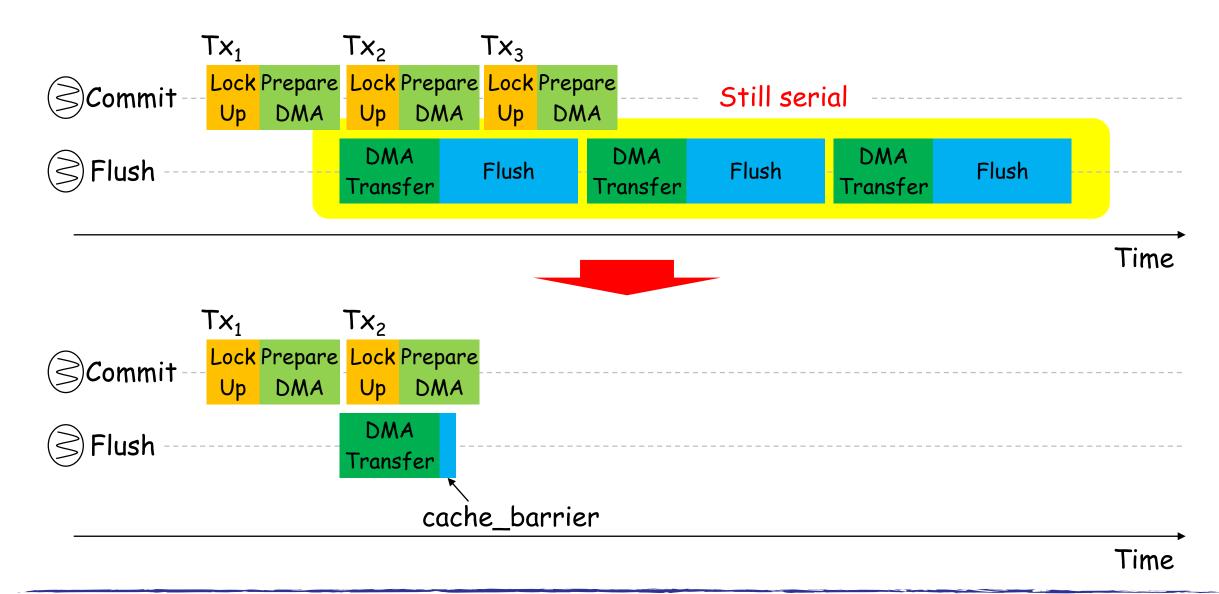
Time

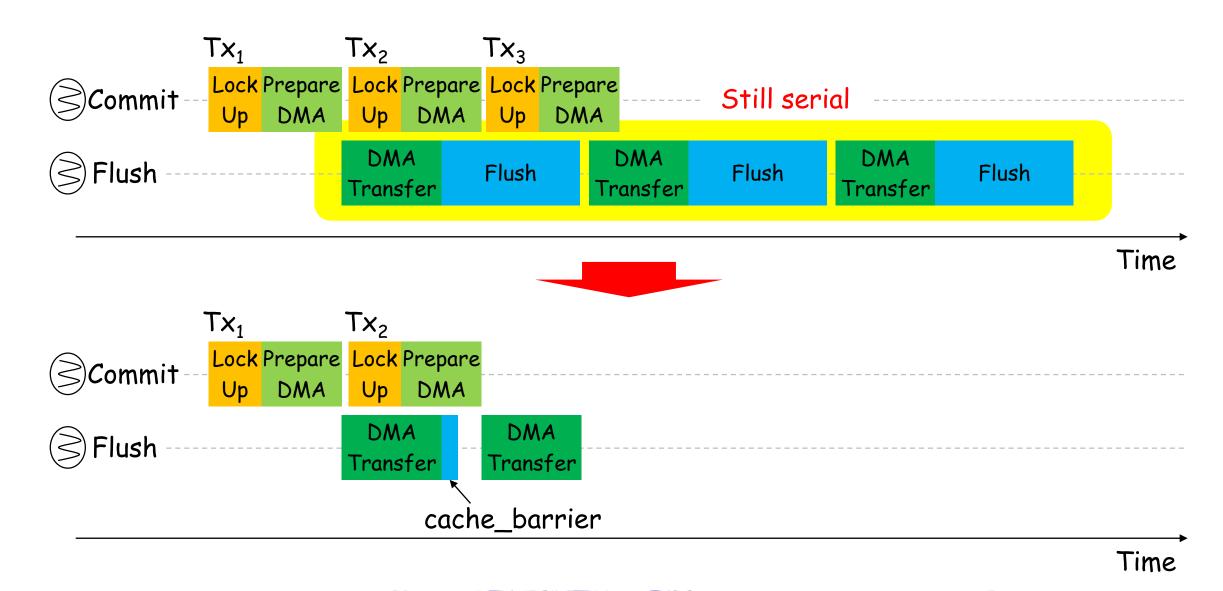


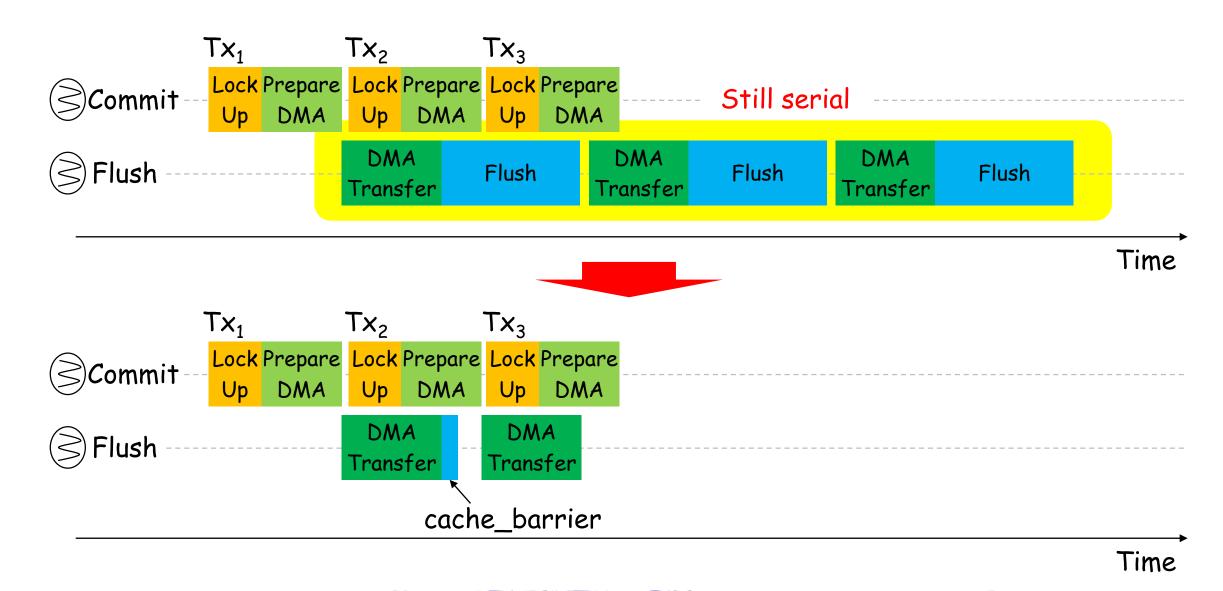


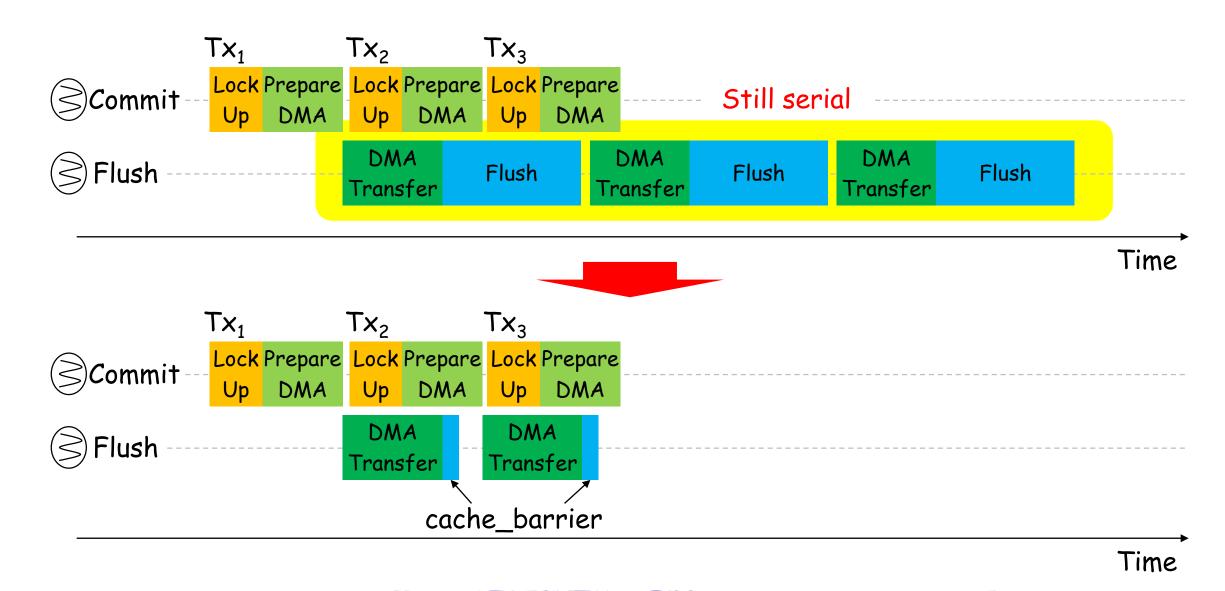


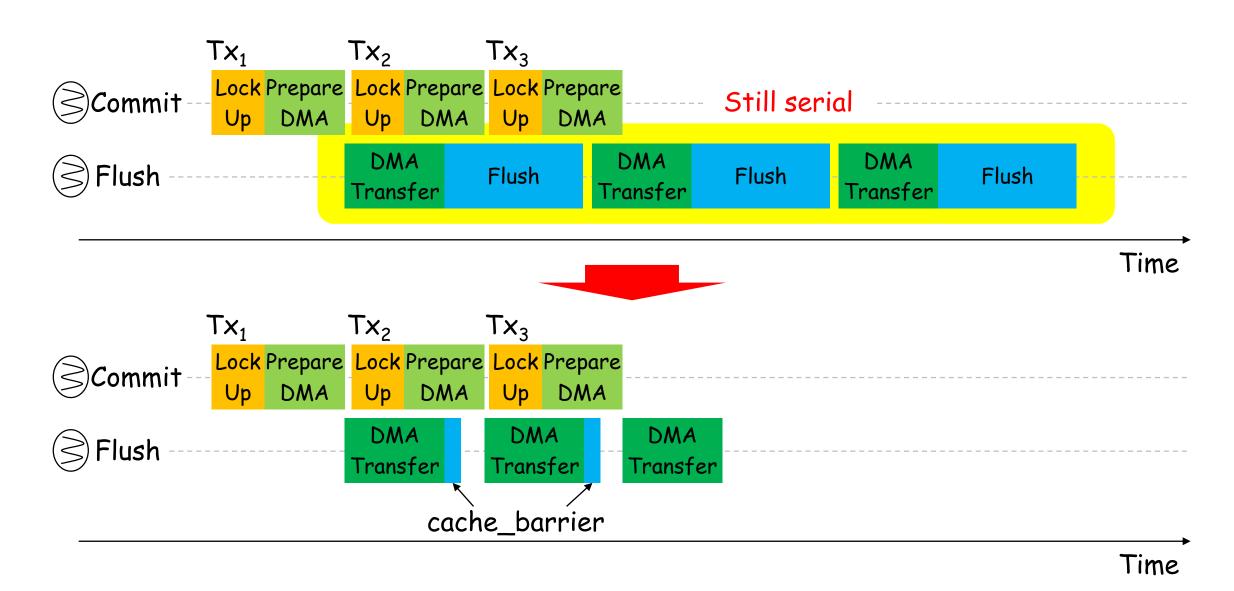
Time

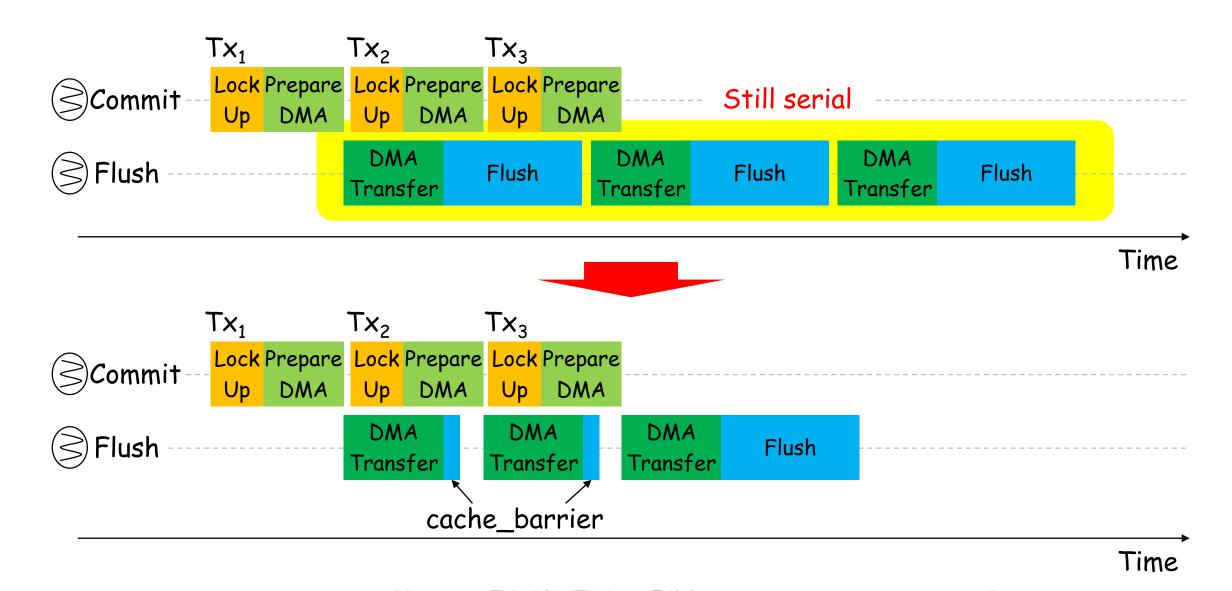










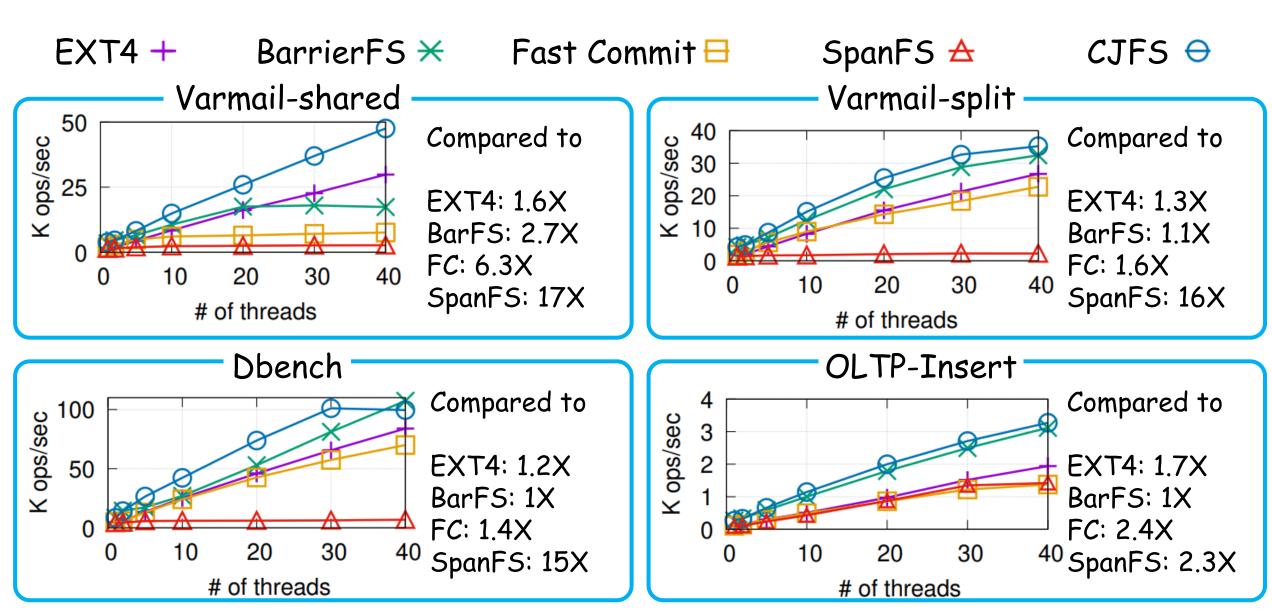


Evaluation

Evaluation Setup

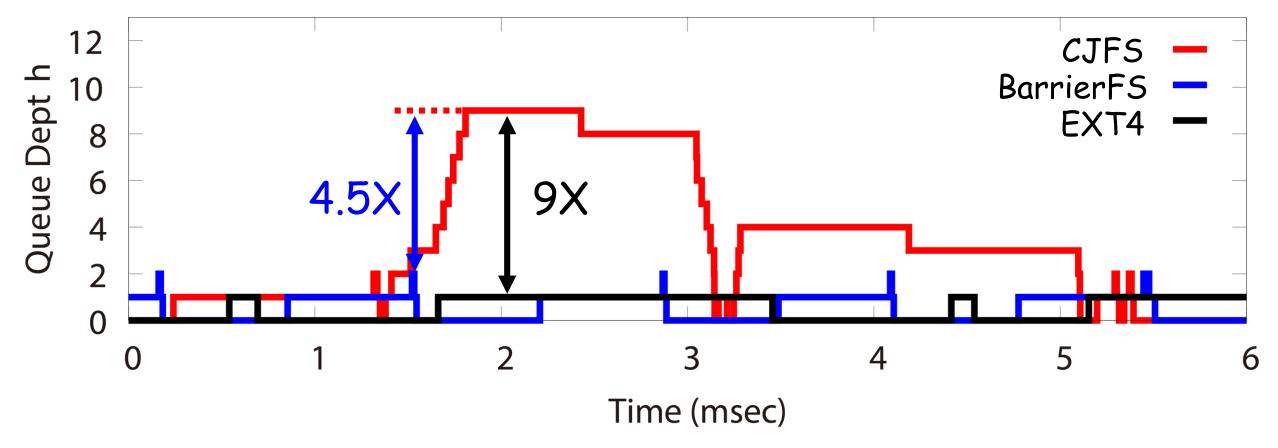
- CPU: Intel Xeon Gold 6320 (2.1 GHz, 2 Socket X 20 core = 40 core)
- Memory: 512GB DRAM
- Storage: Samsung 970 Pro (MLC, NVMe)
- OS (Kernel)
 - CentOS 7.4 (Linux Kernel 5.18.18)
- Filesystem: EXT4, BarrierFS, EXT4 with fast commit, SpanFS, CJFS
- Workloads: Varmail-shared, Varmail-split, Dbench, OLTP-Insert
 - Varmail-shared: Varmail with a shared directory
 - Varmail-split: Varmail with a per-thread directory

Macro Benchmarks



Command Queue Depth

Workload: Varmail with 40 threads



Transactions are transferred and flushed concurrently

Conclusion

- We propose CJFS, Concurrent Journaling Filesystem
- CJFS achieves concurrent transaction commit with four techniques
 - Dual Thread Journaling
 - Multi-Version Shadow Paging
 - Opportunistic Coalescing
 - Compound Flush
- CJFS improves the throughput in macro benchmarks
 - Varmail-shared: 1.6X, Varmail-split: 1.3X, Dbench: 1.2X, OLTP-Insert: 1.7X

Question & Answer

https://github.com/ESOS-Lab/cjfs

