Backtracking and Forward-Checking | N-Queens.

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BACKTRACKING PSEUDO-CODE

```
FUNCTION backtrack(self, col): # col is current col being checked
           OUTPUT "NO SOLUTIONS EXIST FOR THE GIVEN N"
            RETURN False
       ENDIF
        IF N=col:
           RETURN True #done
       ENDIF
       FOR r in range(0, N): #loop over every row trying to get a valid solution
            queen_pos[col] <- r</pre>
            IF checkPositions(r, col) = True: # current board setup is valid
                IF backtrack(col+1) = True:
                   RETURN True
                ENDIF
           ELSE:
                queen_pos[col] <- -2* N; #failed, go back to default value</pre>
           ENDIF
    ENDFUNCTION
```

FORWARD-CHECKING PSEUDO-CODE

```
FUNCTION forwardtrack(self,col, domains):
        # invalid game size 2 OR 3.
        IF N = 2 OR N = 3:
            RETURN False
        # IF N = col, N queens finished
        ENDIF
            RETURN True
        IF len(domains[col])==0:
            RETURN False
        WHILE length of (domains[col]) > 0 :
            r \leftarrow (domains[col]).pop() \# top of set, and pops it off
            queen_pos[col] <- r</pre>
            IF current board setup is ok:
                # domain wipeout!
                new_domains <- domain_wipeout(board setup)</pre>
                IF a queen domain size == 0 :
                    break
                IF forwardtrack(col+1, new_domains) = True:
                        ENDFOR
                    RETURN True
                ENDIF
            ELSE:
                queen_pos[col] <- -2* N</pre>
            ENDIF
    ENDFUNCTION
```

CODE:

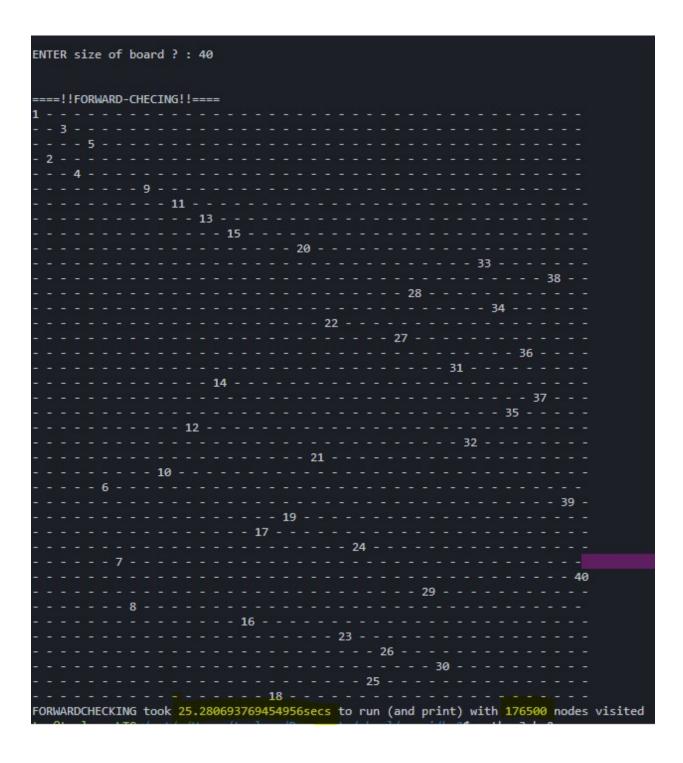
```
import time
import random
class n_queens:
   def __init__(self, N, queen_pos):
    self.N = N # SIZE OF BOARD
        self.queen pos = queen pos #ARR[i] RETURNS ROW POS OF COL I QUEEN
        self.nodes_visited = 0
        for i in range(0,N):
            queen_pos.append(-2*N) #fill array with junk for now
    def checkPositions(self, row, col):
        # given a queen in column 'col' and in row 'row', is it valid vs the rest of the board
        row_to_check = row
        diag1_check = row - col
        diag2_check = row + col
        for c in range(0, len(self.queen_pos) ):
           if c == col:
               continue #dont want to validate against self
            cur_row = self.queen_pos[c]
            cur_diag1 = self.queen_pos[c] - c
            cur diag2 = self.queen pos[c] + c
            if row_to_check == cur_row or diag1_check == cur_diag1 or diag2_check==cur_diag2:
               return False
    def print_sol(self):
        res = [ [0] * self.N for _ in range(self.N) ]
        for i in range(0, self.N):
           res[i][self.queen_pos[i]] = self.queen_pos[i]+1
        for j in range(0, self.N):
            for k in range(0, self.N):
                if res[j][k] == 0:
                   s +=
                            " #print("- ")
                    s+= str(res[j][k]) + " " #print(str(res[j][k]) + " ")
           print(s)
class backtracking(n_queens):
    def backtrack(self, col): # col is current col being checked
        self.nodes_visited += 1
        if self.N==2 or self.N==3:
            print("NO SOLUTIONS EXIST FOR THE GIVEN N")
        if self.N==col:
            self.print_sol()
        for r in range(0, self.N,): #loop over every row trying to get a valid solution
            self.queen_pos[col] = r
            if self.checkPositions(r, col) == True:
                if self.backtrack(col+1) == True:
                self.queen_pos[col] = -2*self.N; #failed, go back to def value
class forwardtracking(n_queens):
   def __init__(self, N, queen_pos):
```

```
n_queens.__init__(self, N, queen_pos)
    self.domains = dict()
    temp = [x for x in range(N)]
    for i in range(N):
        random.shuffle(temp)
        self.domains[i]=set(temp[:]) # domains is a dict int->set,
def domain_wipeout(self, old_domains, row, col):
    for i in range(self.N):
        if i==col:
        old_domains[i] = (old_domains[i]).difference({row})
    dif = row - col
    for i in range(self.N):
        if i==col:
        old_domains[i] = (old_domains[i]).difference({i+dif})
    #3. diag2, row+col
add = row + col
    for i in range(self.N):
        if i==col:
        old_domains[i] = (old_domains[i]).difference({add-i})
    return old_domains
def forwardcheck(self,col, domains):
    self.nodes_visited += 1
    # invalid game size 2 or 3.
    if self.N == 2 or self.N == 3:
    if self.N == col:
        self.print_sol()
    if len(domains[col])==0:
    while (domains[col]).__len__() >0 :
    r = (domains[col]).pop()
        self.queen_pos[col] = r
        if self.checkPositions(r, col) == True:
            temp = domains.copy()
            new_domains = self.domain_wipeout(temp, r, col )
            valid = True
            for i in range(col+1,self.N):
                if new_domains[i].__len__() == 0 :
                    valid = False
            if not valid:
```

```
break
                   if self.forwardcheck(col+1, new_domains) == True:
                   self.queen_pos[col] = -2*self.N
def main():
    print("n_queens mulvey\n")
    n = int(input("ENTER size of board ? : "))
print("\n====!!BACKTRACKING!!====")
    start_time = time.time()
    a = backtracking(n, [])
    a.backtrack(0)
    end_time = time.time() - start_time
  print("BACKTRACKING took " + str(end_time) + "secs to run (and print) with " + str(a.nodes_visited) +
nodes visited")
    print("\n\n====!!FORWARD-CHECING!!====")
    start_time = time.time()
    b = forwardtracking(n, [])
    b.forwardcheck(0, b.domains)
end_time = time.time() - start_time
print("FORWARDCHECKING took " + str(end_time) + "secs to run (and print) with " + str(b.nodes_visited)
  " nodes visited")
main()
```

SAMPLE OUTPUTS

```
ENTER size of board ? : 30
====!!BACKTRACKING!!====
- - - - - - - 26 - - - -
- - - - - - - - - 28 - -
 - - - - - - - - - - - - - 30
BACKTRACKING took 624.1868057250977secs to run (and print) with 4164548 nodes visited
====!!FORWARD-CHECING!!====
------25 ----
- - - - - - 26 - - -
- - - - - - - - - - - - - - 30
- - - - - 10 - - - - - -
FORWARDCHECKING took 1.6253600120544434secs to run (and print) with 20797 nodes visited
```



COMPARISONS:

For Backtracking vs Forward-Checking, I only ran tests for sizes 10, 20, and 30. Since backtracking took over 10 minutes for N=30, I decided that was enough of that. I ran My Forward-Checking for N=40 and N=50. Here is a table of N, Time Elapsed, and Nodes Visited (with a 1 CPU core and 0.5GB RAM VM).

ALGORITHM	SIZE, N	TIME RAN (S)	NODES VISITED
BackTracking	10	0.011760235	148
Forward-Checking	10	0.005118132	72
BackTracking	20	2.160827875	24814
Forward-Checking	20	0.134341002	1020
BackTracking	30	624.1868057	4164548
Forward-Checking	30	1.625360012	20797
Forward-Checking	40	25.28069377	176500
Forward-Checking	50	2878.868532	15426815

For smaller Ns, N<10, Backtracking seems to visit about 2x as many nodes as the forward-checking will. When N got a little larger, 20, backtracking visited 24x more nodes than forward checking. At the largest tested N, 30, backtracking visited over **200x** more nodes than its counterpart! With my forward-checking algorithm, 40x40 board seemed to be the limit because N=50 took 48 minutes to run and visited 15,426,815 nodes. The total # of possible nodes for a NxN board is $1 + n + n^2 + n^3 + ... + n^n = \frac{n^{n+1}-1}{n-1}$. Considering for N=50, there are 9.0630451*10^84 nodes, and only 1.54x10^7 were visited, that is pretty good, but can be improved upon.

CONCLUSION:

Forward checking was always faster than backtracking due to the fewer nodes it had to visit. Much more calculations had to be done for forward-checking, but they obviously paid off as you can see in the chart above. Backtracking basically brute forced its way through the NxN chessboard, while forward-checking did the same but was able to tell when a future Queen would have nowhere to go, and it would go back before going through many other possibilities where that same queen would also fail. The forward checking can still be improved upon by implementing a Priority Queue of the domain lengths of each queen. The priority queue would try the smallest domain (most constrained variable). Within that domain, it would choose the value/row that would eliminate the least amount of values from other domains (Least constraining value heuristic). The most constrained variable and least constraining value heuristic would make the forward-checking even faster.

Both backtracking and forward-checking are exponential in terms of time, and the forward-checking with the two heuristics is also exponential time complexity. Space complexity wise, backtracking will be better off than its forward-checking counterparts because they will require extra space for domains and advanced data structures. The extra space used by forward-checking does result in many fewer nodes visited than backtracking though, and in CSP like N-queens, every solution is a valid one, so why not go get the fastest one? When implementing N-Queens recursively, it does not explicitly show the backtracking. Following the stack traces though, does show giving up on a certain queen spot and trying a different one. Like mentioned earlier, the Forward-checking can be improved upon with a PriorityQueue().

There is also another approach called "Hill Climbing" which is a very rudimentary heuristic. It goes as follows:

- Select an initial (random or pre-determined) assignment of Queens
- WHILE there are threatened queens
 - Select Random Threatened Queen
 - Move that gueen to another row that minimizes conflicts

The heuristic is # of queens threatening that piece, and obviously the lower the better. This is much faster than the forward checking or backtracking but is a heuristic-based approach. Russel and Norvig claim in their book "Artificial Intelligence: A Modern Approach" that they had an average of 50 moves for 1 million queens with this method. I would assume they used more efficient methods (for example, break a recursive stack after x iterations of trying to move the queen), but this method would be the fastest to solve many queens.

My forward-checking algorithm was not too great, and it would probably benefit greatly with a priority queue implemented with the two heuristics. It would still be slower than hill climbing. Hill climbing will "pick best random" and forward-checking will brute-force, but take some of the branching depth and width down. A simple Hill-Climbing algorithm for me was able to perform N=50 in 27 seconds and N=100 in 521 seconds, so not as good as Russel and Norvig's implementation.

Total number of Steps Climbed: 5854 Number of random restarts: 130 Steps Climbed after last restart: 51 Took 521 Secs to run