

Left-leaning Red Black Trees

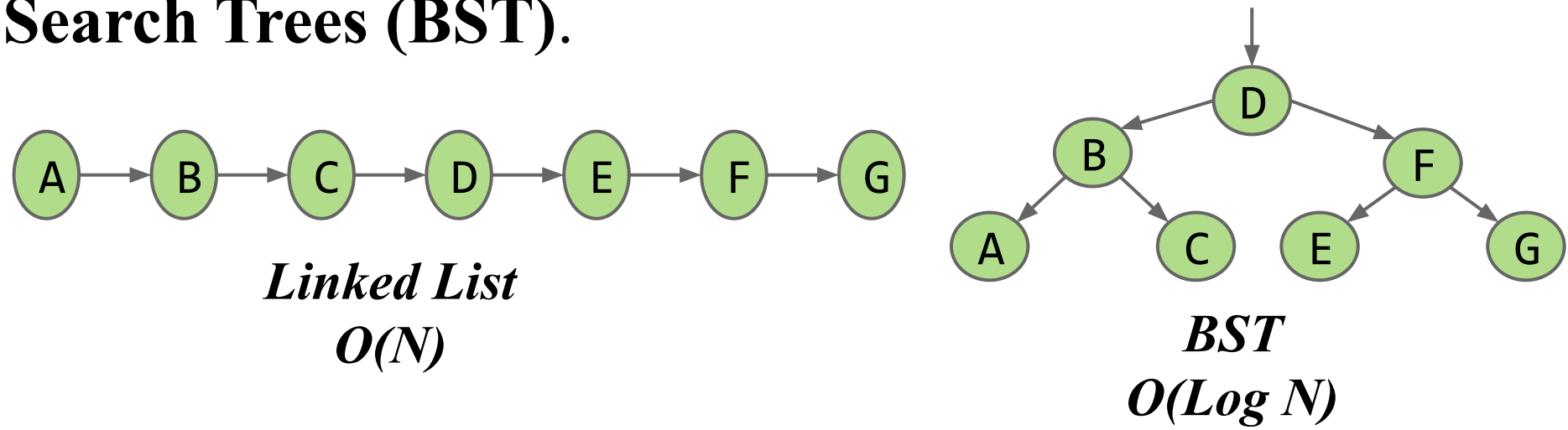
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Background

In order to overcome slow search problems in ordered data structures such as Linked Lists, we use **Binary Search Trees (BST)**.

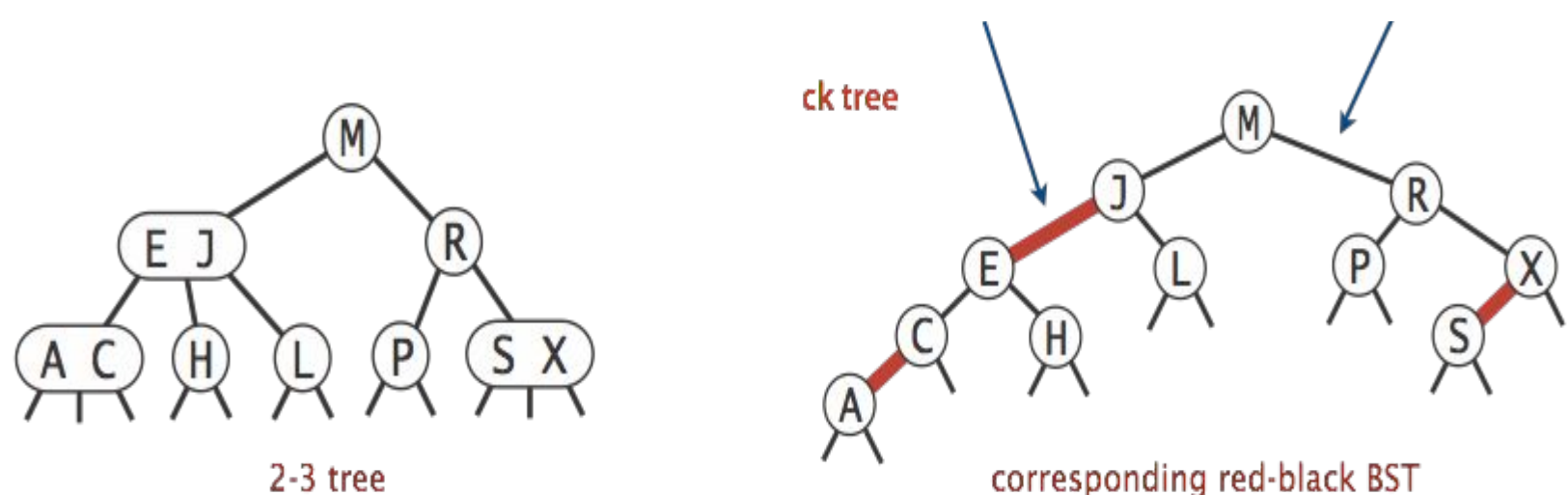


However, BSTs can become **spindly** and the performance of a spindly tree can be just as bad as a Linked List. This can be prevented by using **Self-Balancing Trees (SBT)**. A SBT is any BST that automatically keeps its height small in the face of arbitrary item insertions and deletions.

Red Black Trees and **2-3 Trees** are two of the many types of SBTs and while they're efficient at what they do, their implementations are byzantine. This brings us to the topic of Left-Leaning Red Black Trees.

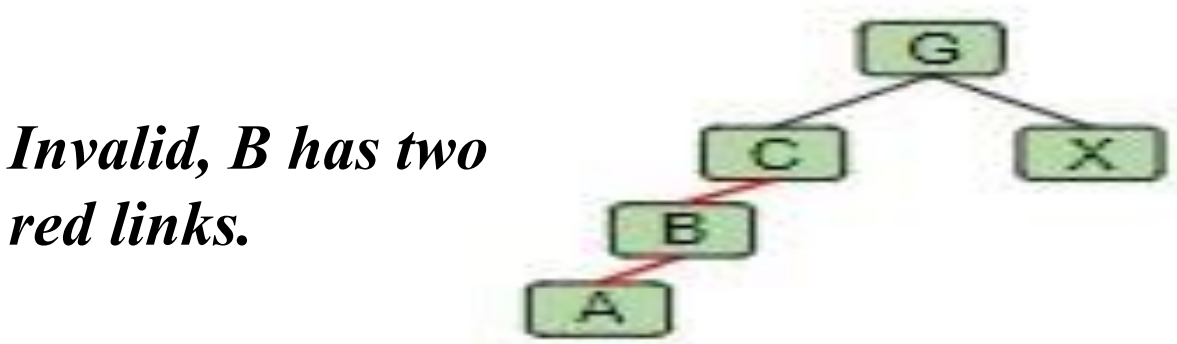
Introduction

LLRBs are structurally identical to 2-3 Trees yet meet all the original design goals set by RBTs. A BST with **left** red links that represents a 2-3 tree called a LLRB. There is a 1-1 correspondence between an LLRB and an equivalent 2-3 tree.



Properties

- No node has **two red links** otherwise it'd become a node with 3 items and 4 children i.e. a 4 node which are not allowed in 2-3 Trees i.e. a tree with nodes that have either 2 or 3 children.



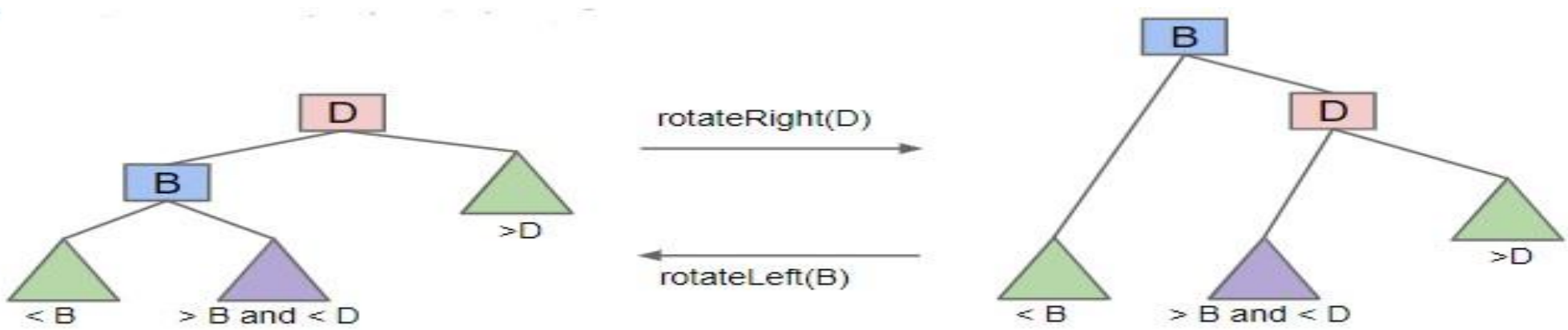
- Every path from the root to a leaf has the **same number of black links**. This is borrowed from 2-3 Trees as this is the mechanism they use to keep trees balanced.



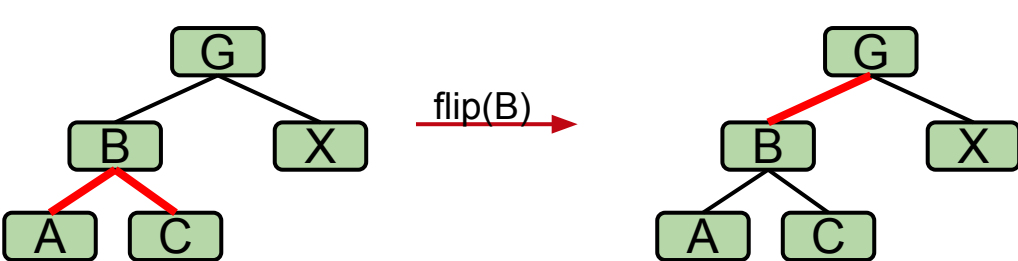
Operations

LLRBs use the following 3 operations to restore the above stated invariants after every insertion and deletion.

Rotate Right & Rotate Left:

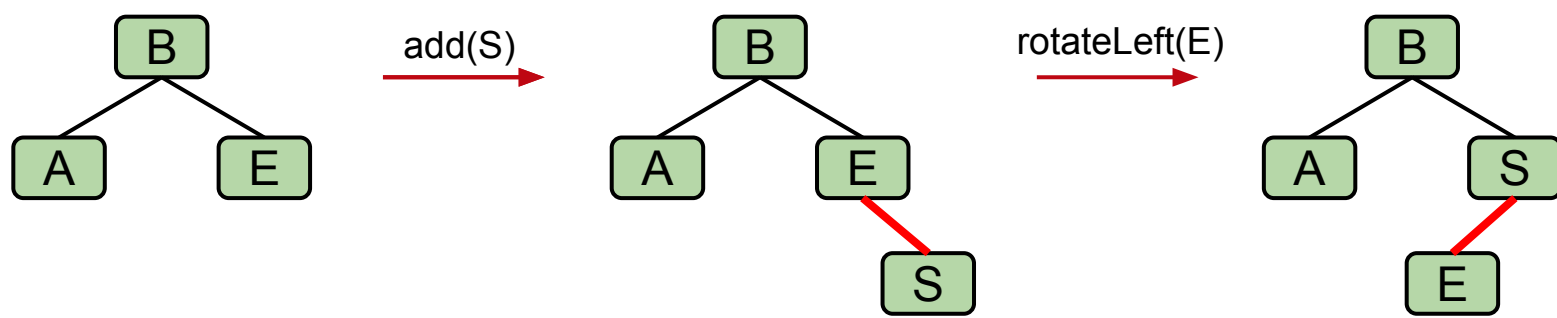


Color Flip:



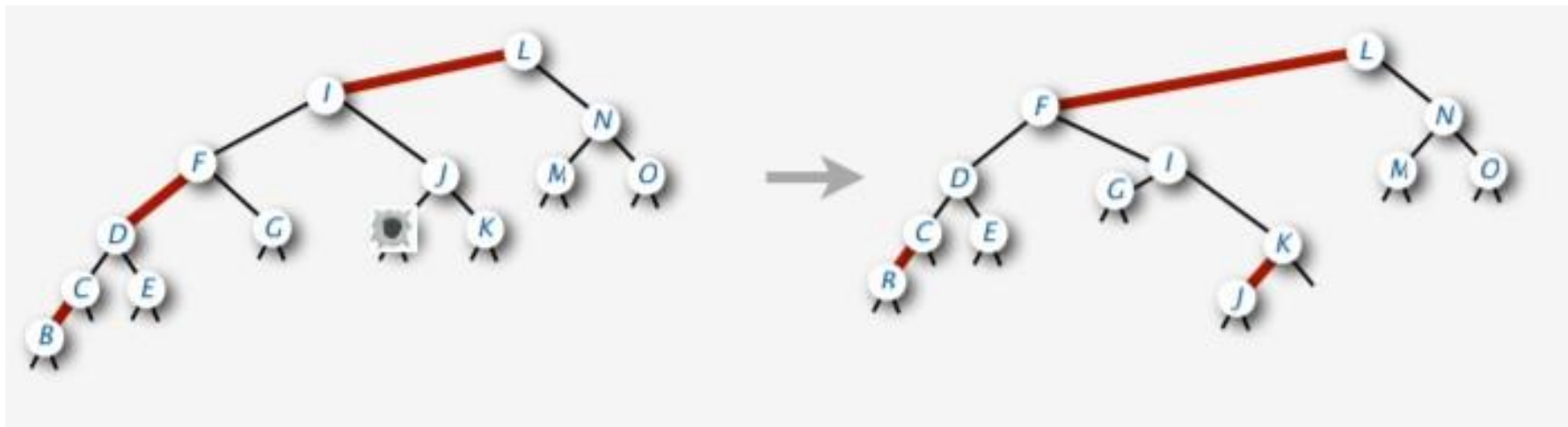
Insertion

Insertion operation is the same as that of a BST. When inserting use a red link. Then travel upwards the length of the tree and apply any of the aforementioned operations repeatedly to restore the inherent properties of LLRBs.



Deletion

When deleting a node, first travel down its left or right spine while shifting the red links downwards. Deleting a red node helps maintain 1-1 correspondence with 2-3 trees and consequently maintain the balance. Afterwards travel up the the length of the tree to correct any violations of the invariants. The searching of a node and the deletion of a leaf are the same as that in a BST.



Runtime

Since 2-3 Trees have $O(\log N)$ runtime, searching, insertion and deletion in an LLRB are all $O(\log N)$.

Reference: Sedgewick, Robert. *Left-Leaning Red Black Trees*. (2008)