

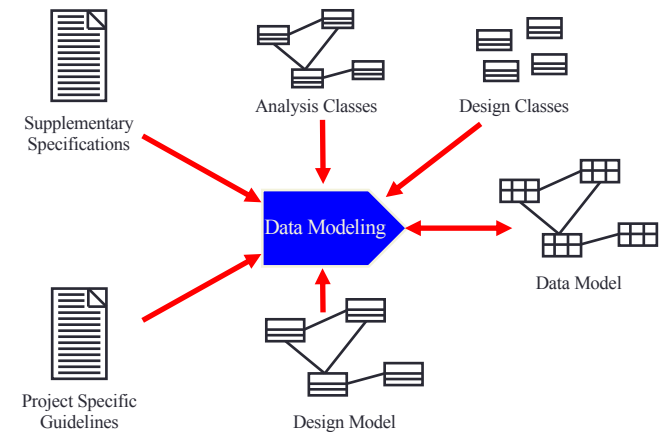
07. DATA MODELING

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Some slides extracted from IBM coursewares

Data Modeling Overview



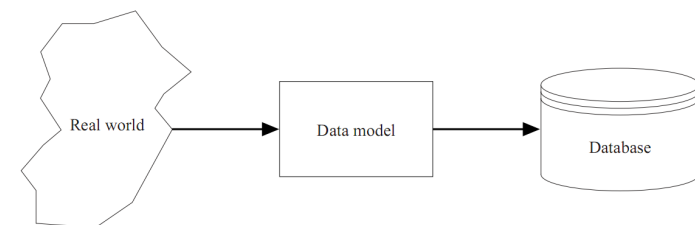
Content

- ➡ 1. Data models
- 2. Object model and Rational Data Model
- 3. Mapping class diagram to E-R diagram
- 4. Normalization

1. Data Models

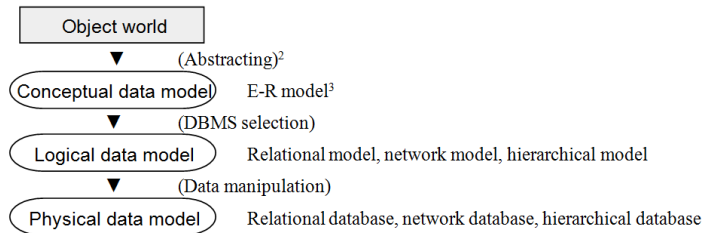
◆ Data modeling:

- Abstracting and organizing the structure of real-world information, which is the object to be made into a database, and then expressing it



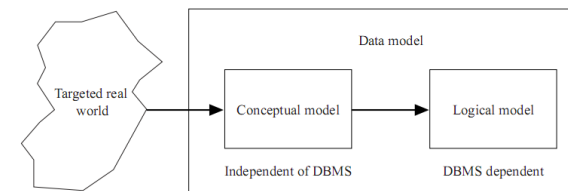
1. Data Models (2)

- 3 types of data models



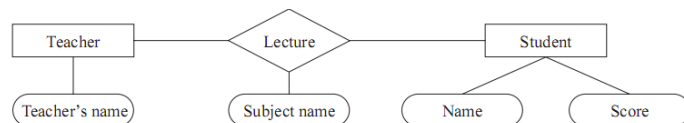
1.1. Conceptual data model

- Natural expressions without constraints imposed by DBMS
- E-R model
 - Expressed by E-R diagram



E-R Diagram

- Three elements
 - Entities
 - Relationships
 - Attributes



1.2. Logical Data Model

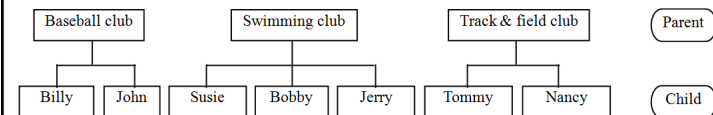
- 3 types
 - relational model,
 - network model,
 - and hierarchical model

1.3. Physical Data Model

- Logical data models, when they are implemented, become physical data models:
 - relational databases,
 - network databases,
 - or hierarchical databases

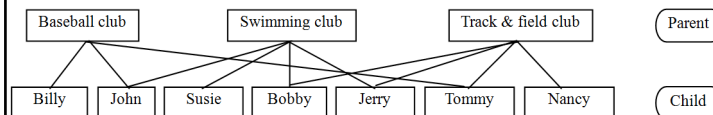
1.3.1. Hierarchical Database (Tree-Structure Database)

- Divides records into parents and children and shows the relationship with a hierarchical structure
- 1-to-many (1:n) correspondences between parent records and child records



1.3.2. Network Database

- Parent records and child records do not have 1-to-n (1:n) correspondences; rather, they are in many-to-many (m:n) correspondence
- Sometimes called CODASYL database



1.3.3. Rational database

- Data is expressed in a two-dimensional table.
 - Each row of the table corresponds to a record, and each column is an item of the records.
 - The underlined columns indicate the primary key

Name of the table: Employee_tbl

Columns (items, attributes,)

<u>Employee_number</u>	Name	<u>Tel_number</u>
00100	Paul Smith	03-3456-0001
00200	Rick Martin	03-3456-0011
00300	Billy Graham	03-3456-0010
00400	John Wilson	03-3456-0200

← Row (pair, tuple, record)



NOSQL DATABASES

Overview, Models, Concepts, Examples

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What is NoSQL Database?

- NoSQL (cloud) databases
 - Use document-based model (non-relational)
 - Schema-free document storage
 - Still support indexing and querying
 - Still support CRUD operations (create, read, update, delete)
 - Still supports concurrency and transactions
 - Highly optimized for append / retrieve
 - Great performance and scalability
 - NoSQL == “No SQL” or “Not Only SQL”?

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Relational vs. NoSQL Databases

- Relational databases
 - Data stored as table rows
 - Relationships between related rows
 - Single entity spans multiple tables
 - RDBMS systems are very mature, rock solid
- NoSQL databases
 - Data stored as documents
 - Single entity (document) is a single record
 - Documents do not have a fixed structure

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Relational vs. NoSQL Models

Relational Model

Name	Svetlin Nakov
Gender	male
Phone	+359333777555
Email	nakov@abv.bg
Site	www.nakov.com

*
1

Street	Al. Malinov 31
Post Code	1729

*
1

Town	Sofia
------	-------

*
1

Country	Bulgaria
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Document Model

Name: [Svetlin Nakov](#)

Gender: [male](#)

Phone: [359333777555](#)

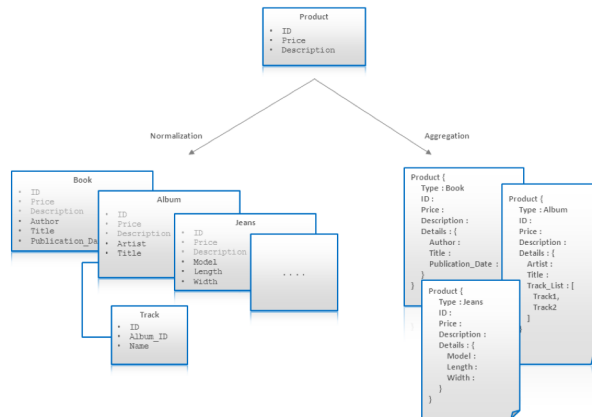
Address:

- Street: [Al. Malinov 31](#)
- Post Code: [1729](#)
- Town: [Sofia](#)
- Country: [Bulgaria](#)

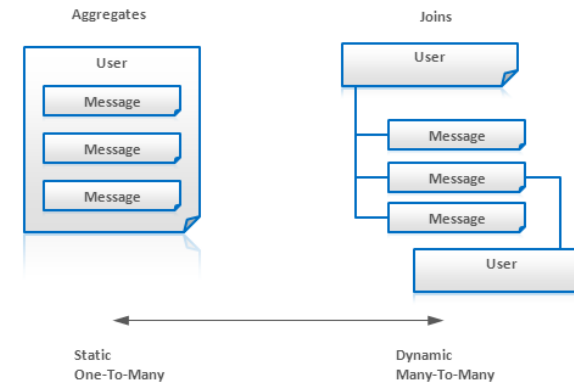
Email: [nakov@abv.bg](#)

Site: [www.nakov.com](#)

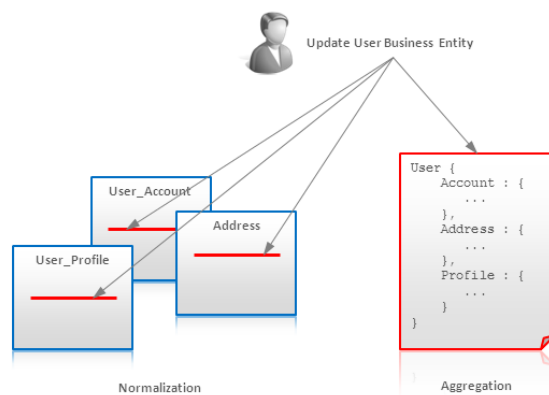
Normalization vs. Aggregation



Aggregates vs. Joins

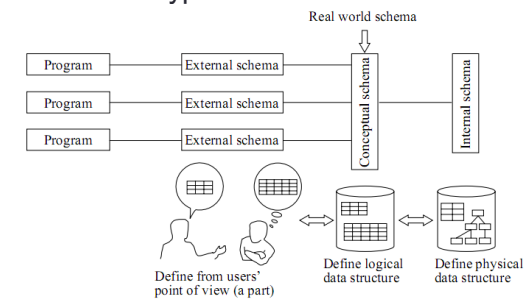


Atomic Aggregates



1.3.4. Three-layer schema

- A schema is a description of the framework of a database
- Classified into 3 types:



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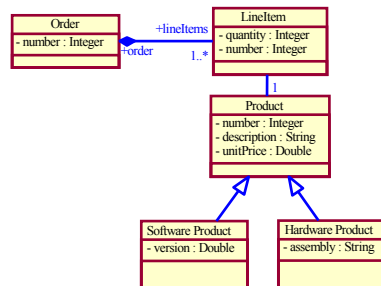
1. Data models
- ➡ 2. Object model and Rational Data Model
3. Mapping class diagram to E-R diagram
4. Normalization

2.1. Relational Databases and OO

- RDBMS and Object Orientation are not entirely compatible
 - RDBMS
 - Focus is on data
 - Better suited for ad-hoc relationships and reporting application
 - Expose data (column values)
 - Object Oriented system
 - Focus is on behavior
 - Better suited to handle state-specific behavior where data is secondary
 - Hide data (encapsulation)

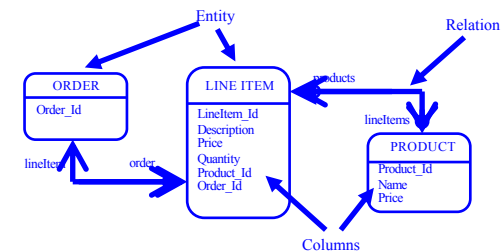
2.2. The Object Model

- The Object Model is composed of
 - Classes (attributes)
 - Relationships
 - Associations
 - Generalization



2.3. The Relational Data Model

- Relational data model is composed of
 - Entities - Table
 - Relations - Relationship
- ➔ Also called E-R model



2.3.1. Entities/Tables

- Entities is mapped to table when design physical database
- Including
 - Columns: Attributes
 - Rows: Concrete values of attributes

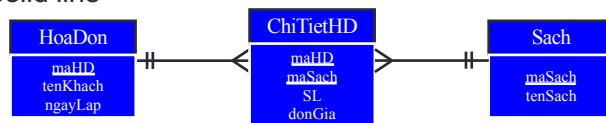
courseID	description	startDate	endDate	location
2008.11.001	This course...	12 Nov 2008	30 Nov 2008	D3-405
2008.11.002	This course...	22 Nov 2008	10 Dec 2008	T-403

2.3.2. Relations/Relationships

- Relations between entities or relationship between tables
 - Multiplicity/Cardinality
 - One-to-one (1:1)
 - One-to-many (1:m)
 - Many-to-one (m:1)
 - Many-to-many (m:n)
- (Normally, many-to-many relation is divided to one-to-many and many-to-one relations)

Dependency relationships

- The child entity can exist only when the parent entity exists
- The child entity has a foreign key referencing to the primary key of the parent entity
- This foreign key is included in the primary key of the child
- Solid line



Independency relationships

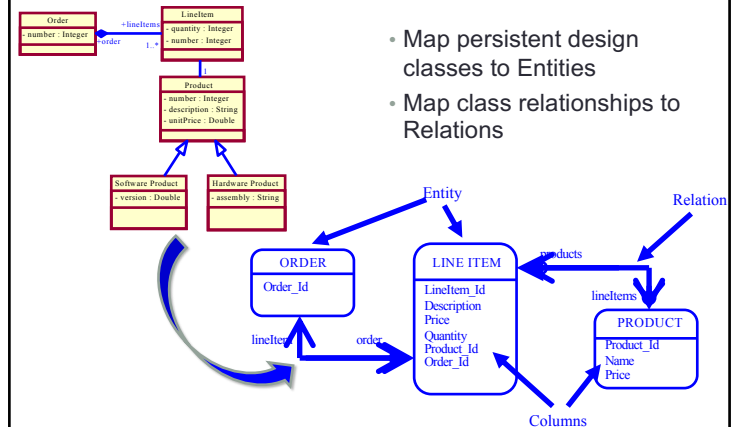
- The child entity can exist even if the parent entity does not exist
- The child entity has a foreign key referencing to the primary key of the parent entity
- This foreign key is not included in the primary key of the child
- Dash line



Content

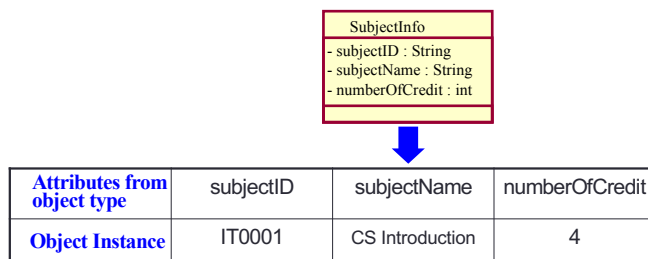
1. Data models
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3. Mapping class diagram to E-R diagram



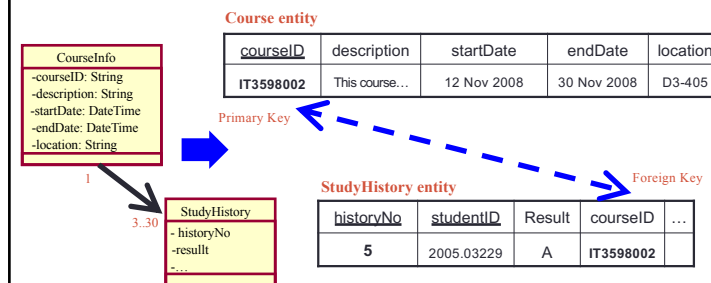
3.1. Mapping Persistent Design Classes to Entities

- In a relational database
 - Every row is regarded as an object
 - A column in a table is equivalent to a persistent attribute of a class



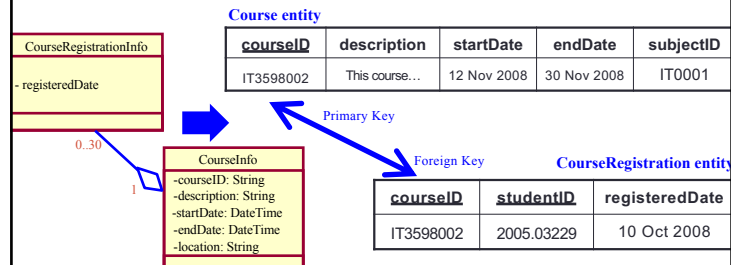
3.2. Mapping Associations Between Persistent Objects

- Associations between two persistent objects are realized as foreign keys to the associated objects.
 - A foreign key (not in primary key) is a column in one table that contains the primary key value of associated object
 - → Independence relationship



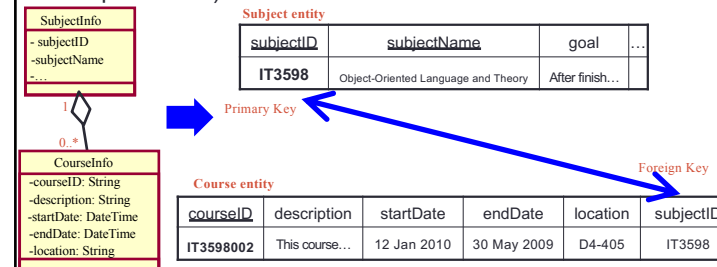
3.3. Mapping Aggregation to the Data Model

- Aggregation is also modeled to dependency relationship using foreign key relationships
 - The use of composition implements a cascading delete constraint

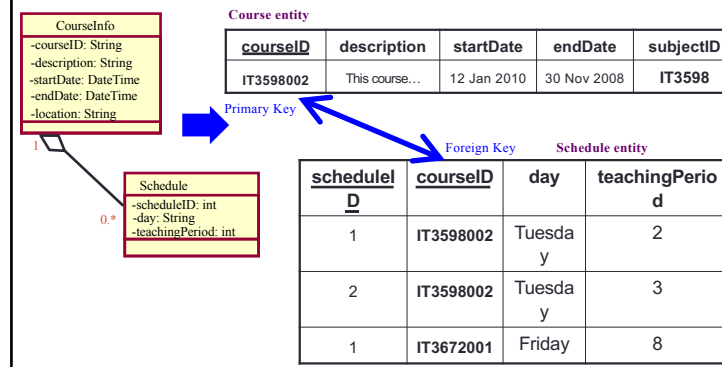


3.3. Mapping Aggregation to the Data Model (2)

- In some case, we can map to independency relationship to simplify the primary key.
- Example: `CourseID` is the primary key (according the requirements)

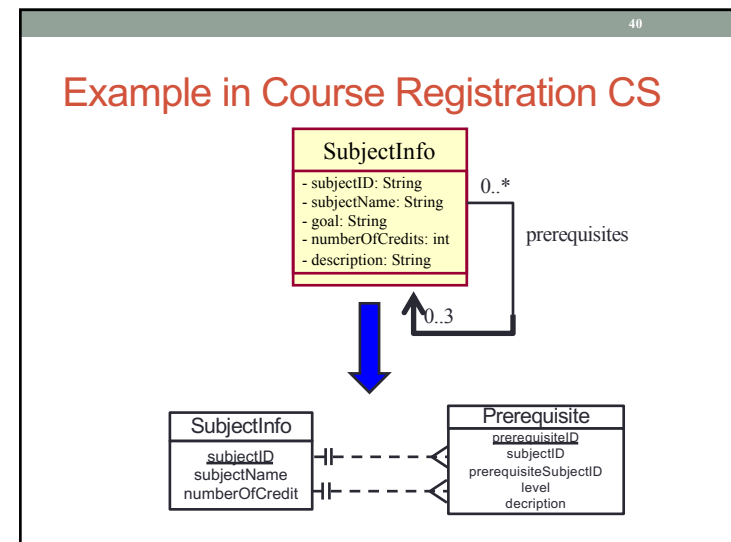
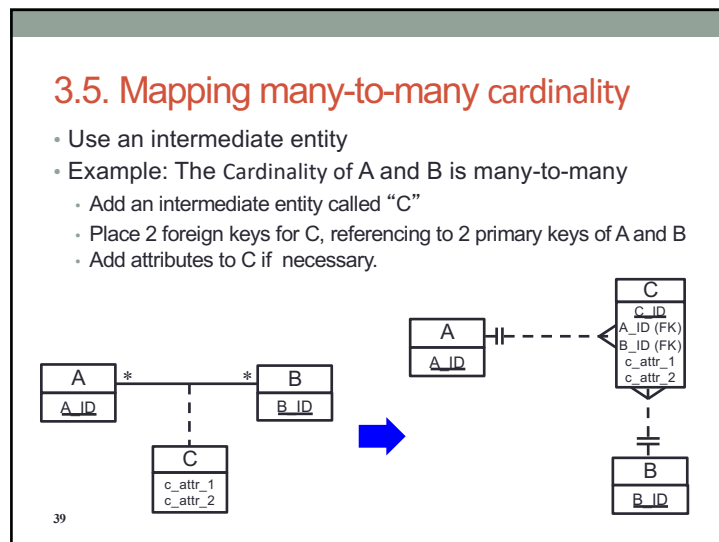
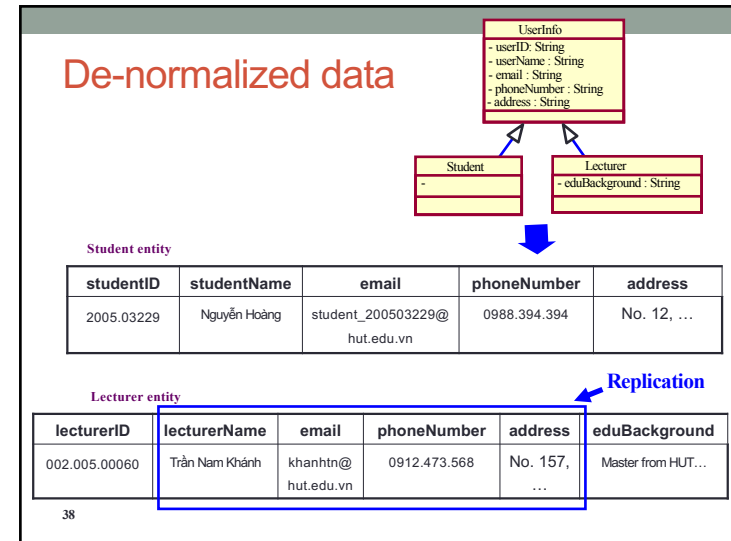
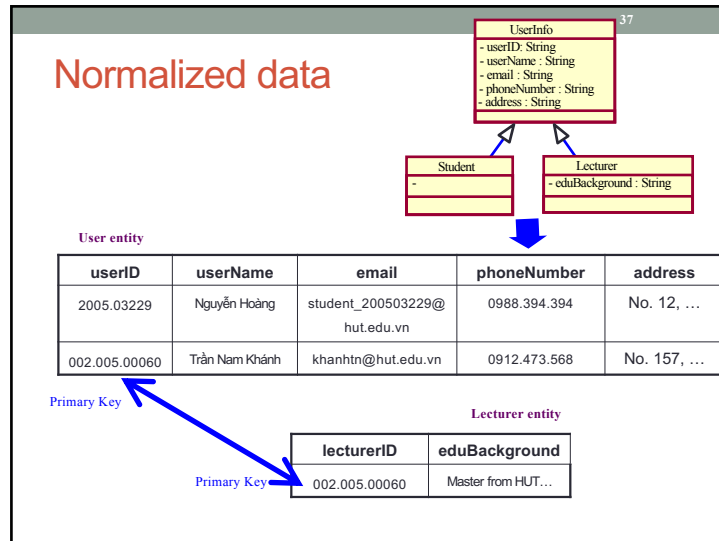


More example in Course Registration

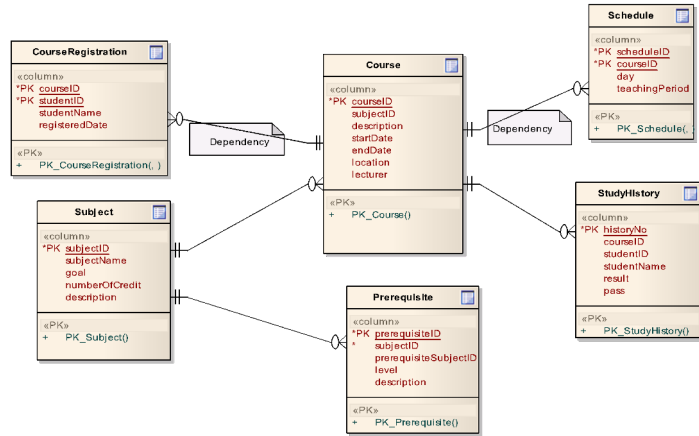


3.4. Modeling Inheritance in the Data Model

- A Data Model does not support modeling inheritance in a direct way
- Two options:
 - Use separate tables (normalized data)
 - Duplicate all inherited associations and attributes (de-normalized data)



E-R diagram



Content

1. Data models
2. Object model and Rational Data Model
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- ➔ 4. Normalization

4.1. Overview of Normalization

- Normalization: the process of steps that will identify, for elimination, redundancies in a database design.
- Purpose of Normalization: to improve
 - storage efficiency
 - data integrity
 - and scalability

4.1. Overview of Normalization (2)

- In relational model, methods exist for quantifying how efficient a database is.
- These classifications are called **normal forms** (or **NF**), and there are algorithms for converting a given database between them.
- Normalization generally involves splitting existing tables into multiple ones, which must be re-joined or linked each time a query is issued

4.2. History



- Edgar F. Codd first proposed the process of normalization and what came to be known as the 1st normal form in his paper *A Relational Model of Data for Large Shared Data Banks* Codd stated:

“There is, in fact, a very simple elimination procedure which we shall call normalization. Through decomposition nonsimple domains are replaced by ‘domains whose elements are atomic (nondecomposable) values’.”

4.3. Normal Forms

- Edgar F. Codd originally established three normal forms: 1NF, 2NF and 3NF.
- There are now others that are generally accepted, but 3NF is widely considered to be sufficient for most applications.
- Most tables when reaching 3NF are also in BCNF (Boyce-Codd Normal Form).



Functionally determines

- In a table, a set of columns X, **functionally determines** another column Y...

$X \rightarrow Y$

... if and only if each X value is associated with at most one Y value in a table.

- i.e. if you know X then there is only **one** possibility for Y.

Normal forms so Far...

◆ First normal form

- All data values are atomic, and so everything fits into a mathematical relation.

◆ Second normal form

- As 1NF plus no *non-primary-key attribute* is partially dependant on the primary key

◆ Third normal form

- As 2NF plus no non-primary-key attribute depends transitively on the primary key

Normalization Example

◆ Consider a table representing orders in an online store

◆ Each entry in the table represents an item on a particular order. (thinking in terms of records. Yuk.)

◆ Columns

- Order
- Product
- Customer
- Address
- Quantity
- UnitPrice

◆ Primary key is {Order, Product}

Functional Dependencies

- Each order is for a **single** customer $\{\text{Order}\} \rightarrow \{\text{Customer}\}$
- Each customer has a **single** address $\{\text{Customer}\} \rightarrow \{\text{Address}\}$
- Each product has a **single** price $\{\text{Product}\} \rightarrow \{\text{UnitPrice}\}$
- FD's 1 and 2 are transitive $\{\text{Order}\} \rightarrow \{\text{Address}\}$

Example – FD Diagram

INF



Normalization to 2NF

◆ Remember 2nd normal form means no partial dependencies on the key. But we have:

$\{\text{Order}\} \rightarrow \{\text{Customer, Address}\}$

$\{\text{Product}\} \rightarrow \{\text{UnitPrice}\}$

And a primary key of: {Order, Product}

- So to get rid of the first FD we *project* over:

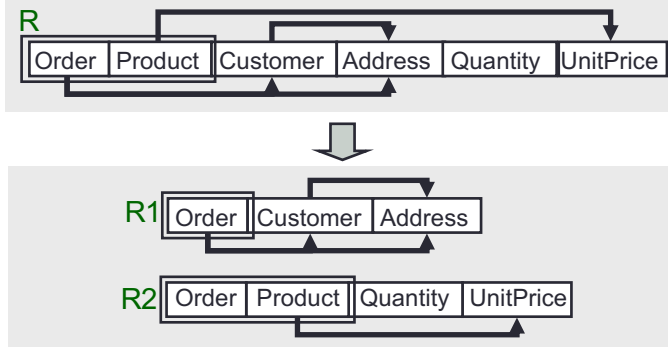
$\{\text{Order, Customer, Address}\}$

and

$\{\text{Order, Product, Quantity and UnitPrice}\}$

Normalization to 2NF

1NF



Normalization to 2NF

- ◆ R1 is now in 2NF, but there is still a partial FD in R2:

$\{Product\} \rightarrow \{UnitPrice\}$

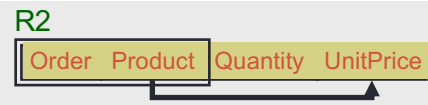


- To remove this we project over:

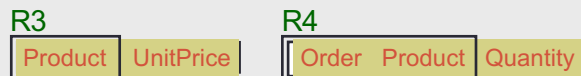
$\{Product, UnitPrice\}$ and $\{Order, Product, Quantity\}$

Normalization to 2NF

1NF



2NF



Now let's go 3NF...

- R has now been split into 3 relations - R1, R3, and R4... but R1 has a transitive FD on its key...

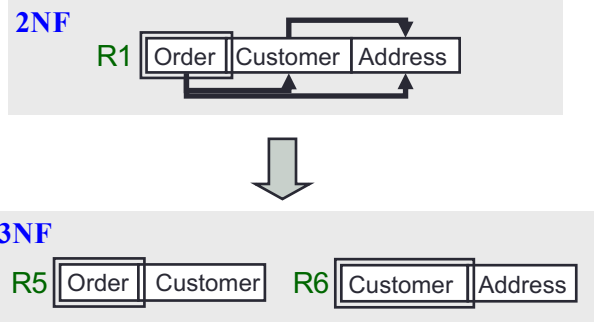


$\{Order\} \rightarrow \{Customer\} \rightarrow \{Address\}$

- To remove this problem we project R1 over:

$\{Order, Customer\}$ and $\{Customer, Address\}$

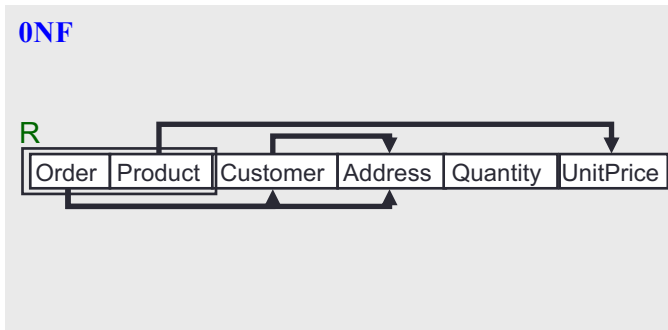
So more chopping...



Let's summarize that:

- **1NF:**
{Order, Product, Customer, Address, Quantity, UnitPrice}
- **2NF:**
{Order, Customer, Address}
{Product, UnitPrice}
{Order, Product, Quantity}
- **3NF:**
{Product, UnitPrice}
{Order, Product, Quantity}
{Order, Customer}
{Customer, Address}

So this...



has become this...



“Register for course” use case

- Make the E-R diagram from the previous step for “Register for course” use case to become:
 - The first normal form
 - The second normal form
 - The third normal form

