

Add two numbers represented as Linked Lists

Problem Statement: Given the **heads** of two non-empty linked lists representing two non-negative integers. The digits are stored in **reverse order**, and each of their nodes contains a single digit. Add the two numbers and return the **sum** as a linked list.

Examples:

Input Format:

(Pointer/Access to the head of the two linked lists)

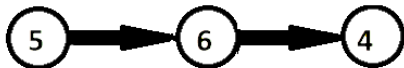
num1 = 342, num2 = 564

l1 = [2,4,3]

l2 = [5,6,4]

Result: sum = 807; L = [7,0,8]

Explanation: Since the digits are stored in reverse order, reverse the numbers first to get the original number and then add them as
the or
→ 342 + 465 = 807. *Refer to the image below.*



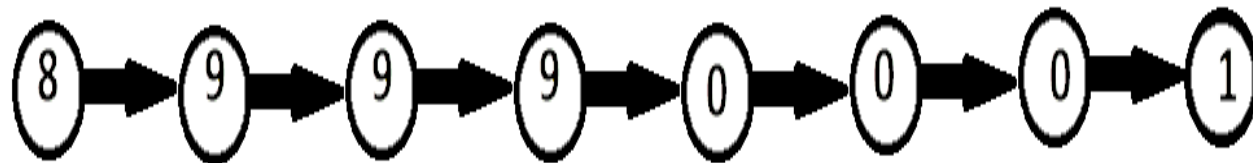
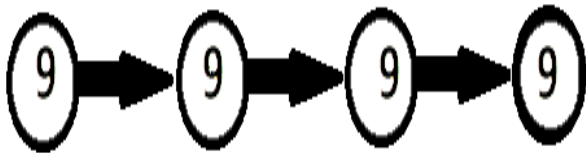
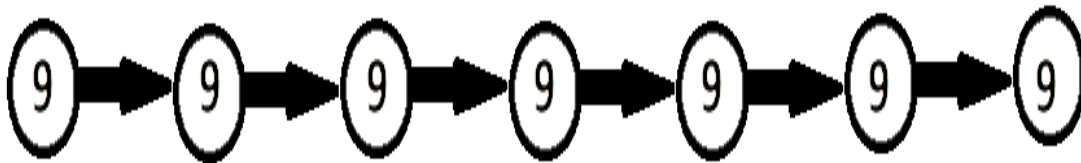
Input Format:

(Pointer/Access to the head of the two linked lists)

l1 = [9,9,9,9,9,9,9], l2 = [9,9,9,9]

Result: [8,9,9,9,0,0,0,1]

Explanation: Since the digits are stored in reverse order, reverse the numbers first to get the original number and then add them as $\rightarrow 9999999 + 9999 = 8999001$. *Refer to the image below.*



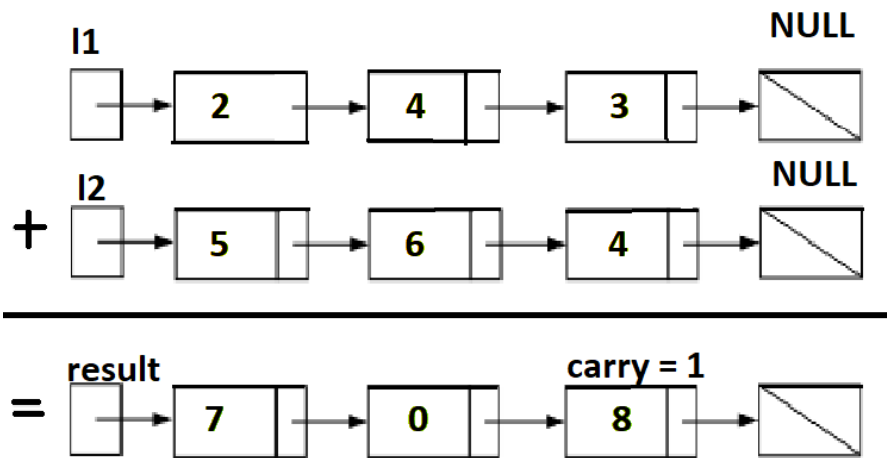
Solution

Disclaimer: Don't jump directly to the solution, try it out yourself first.

Solution 1: Elementary Math

Intuition: Keep track of the **carry** using a variable and simulate digits-by-digits sum starting from the head of the list, which contains the least significant digit.

Approach:



Visualization of the addition of two numbers:

$$342 + 465 = 807$$

$$342 + 465 = 807.$$

Each node contains a single digit and the digits are stored in reverse order.

Just like how you would sum two numbers on a piece of paper, we begin by summing the least significant digits, which is the head of **l1** and **l2**. Since each digit is in the range of 0...9, summing two digits may "overflow". For example

$5 + 7 = 12$. In this case, we set the current digit to 2 and bring over the $\text{carry} = 1$ to the next iteration.

carry must be either 0 or 1 because the largest possible sum of two digits (including the carry) is $9 + 9 + 1 = 19$.

Pseudocode:

- Create a dummy node which is the head of new linked list.
- Create a node temp, initialise it with dummy.
- Initialize carry to 0.
- Loop through lists **l1** and **l2** until you reach both ends, and until carry is present.
 - Set $\text{sum} = \text{l1.val} + \text{l2.val} + \text{carry}$.
 - Update $\text{carry} = \text{sum} / 10$.
 - Create a new node with the digit value of $(\text{sum} \% 10)$ and set it to temp node's next, then advance temp node to next.
 - Advance both **l1** and **l2**.

- Return dummy's next node.

Note that we use a dummy head to simplify the code. Without a dummy head, you would have to write extra conditional statements to initialize the head's value.

Take extra caution in the following cases:

Test case	Explanation
l1=[0,1], l2=[0,1,2]	When one list is longer than the other.
l1=[], l2=[0,1]	When one list is null, which means an empty list.
l1=[9,9], l2=[1]	The sum could have an extra carry of one at the end, which is easy to forget.

Source Code:

```
public ListNode addTwoNumbers(ListNode l1, ListNode l2) {  
    int carry=0;  
    ListNode start=new ListNode();  
    ListNode tail=start;  
    while(l1!=null||l2!=null||carry>0)  
    {  
        int sum=0;  
        if(l1!=null)  
        {  
            sum+=l1.val;  
            l1=l1.next;  
        }  
        if(l2!=null)  
        {  
            sum+=l2.val;  
            l2=l2.next;  
        }  
    }  
}
```

```
        sum+=carry;
        carry=sum/10;
        ListNode temp=new ListNode(sum%10);
        tail.next=temp;
        tail=tail.next;
    }

    return start.next;
}
```

Time Complexity: $O(\max(m,n))$. Assume that m and n represent the length of $l1$ and $l2$ respectively, the algorithm above iterates at most $\max(m,n)$ times.

Space Complexity: $O(\max(m,n))$. The length of the new list is at most $\max(m,n)+1$.