

Longest Consecutive Sequence in an Array

Problem Statement: You are given an array of 'N' integers. You need to find the length of the longest sequence which contains the consecutive elements.

Examples:

Example 1:

Input: [100, 200, 1, 3, 2, 4]

Output: 4

Explanation: The longest consecutive subsequence is 1, 2, 3, and 4.

Input: [3, 8, 5, 7, 6]

Output: 4

Explanation: The longest consecutive subsequence is 5, 6, 7, and 8.

Solution

Disclaimer: Don't jump directly to the solution, try it out yourself first.

Solution 1: (Brute force)

Approach: We can simply sort the array and run a for loop to find the longest consecutive sequence

```
public static int longestConsecutive(int[] nums) {  
    if(nums.length == 0 || nums == null){  
        return 0;  
    }  
    Arrays.sort(nums);  
    int ans = 1;  
    int prev = nums[0];  
    int cur = 1;  
  
    for(int i = 1; i < nums.length; i++){  
        if(nums[i] == prev+1){  
            cur++;  
        }  
        else if(nums[i] != prev){  
            cur = 1;  
        }  
        prev = nums[i];  
        ans = Math.max(ans, cur);  
    }  
    return ans;  
}
```

Time Complexity: We are first sorting the array which will take $O(N * \log(N))$ time and then we are running a for loop which will take $O(N)$ time. Hence, the overall time complexity will be $O(N * \log(N))$.

Space Complexity: The space complexity for the above approach is $O(1)$ because we are not using any auxiliary space

Solution 2: (Optimal Approach)

Approach: We will first push all are elements in the HashSet. Then we will run a for loop and check for any number(x) if it is the starting number of the consecutive sequence by checking if the HashSet contains (x-1) or not. If 'x' is the starting number of the consecutive sequence we will keep searching for the numbers $y = x+1, x+2, x+3, \dots$ And stop at the first 'y' which is not present in the HashSet. Using this we can calculate the length of the longest consecutive subsequence.

```
public static int longestConsecutive(int[] nums) {  
    Set < Integer > hashSet = new HashSet < Integer > ();  
    for (int num: nums) {  
        hashSet.add(num);  
    }  
    int longestStreak = 0;  
    for (int num: nums) {  
        if (!hashSet.contains(num - 1)) {  
            int currentNum = num;  
            int currentStreak = 1;  
            while (hashSet.contains(currentNum + 1)) {  
                currentNum += 1;  
                currentStreak += 1;  
            }  
            longestStreak = Math.max(longestStreak, currentStreak);  
        }  
    }  
    return longestStreak;  
}
```

Time Complexity: The time complexity of the above approach is $O(N)$ because we traverse each consecutive subsequence only once.

Space Complexity: The space complexity of the above approach is $O(N)$ because we are maintaining a HashSet.