

Weekly Assignment 2

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Algorithm Design and Analysis

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Solution 2:

Given in the textbook, the BinSearch computes

$$mid = \frac{lo + hi}{2}$$

at every iteration and checks if the item i is in the array A . We have to prove that the algorithm terminates. If we can show that in every iteration the search interval decreases until it no longer satisfies the loop condition $lo \leq hi$, that is, at some points we will have $lo > hi$, then we will have the proof.

Let's the values of lo and hi in two successive iterations be $[lo_1, hi_1]$ and $[lo_2, hi_2]$. We need to show, $hi_2 - lo_2 < hi_1 - lo_1$.

- **Case-1:** $A[mid] = i$

The element i is found and the algorithm returns **True**, so the **while** loop terminates.

- **Case 2:** $A[mid] < i$

i is in the right half of the array, so $lo_2 = mid + 1$ and $hi_2 = hi_1$. The size of the interval becomes, $hi_2 - lo_2 = hi_1 - (mid + 1)$. We can write, $mid \geq \frac{lo_1 + hi_1}{2}$ and substituting that in the new interval size

$$hi_2 - lo_2 = \frac{hi_1 - lo_1}{2}$$

we get exactly half of the previous interval size, so we have $hi_2 - lo_2 < hi_1 - lo_1$

- **Case 3:** $A[mid] > i$

In this case, i is in the left half of the array, the new size becomes, $hi_2 = mid - 1$ and $lo_2 = lo_1$. Using $mid \leq \frac{lo_1 + hi_1}{2}$ we get

$$hi_2 - lo_2 = \frac{hi_1 - lo_1}{2} - 1,$$

which shows that the new interval size is smaller than the half of the interval. So, we again can say that $hi_2 - lo_2 < hi_1 - lo_1$

In all the cases, the size of the search interval $hi - lo$ decreases by at least half in each iteration. Thus, the algorithm will eventually terminate.