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4. Filtering Algorithms

Discrete linear Consolution

•
$$X * h = y$$

• $F((x_j)_{j \in \mathbb{Z}})_k = y_k = \sum_{j=0}^{n-1} x_j \cdot h_{k-j}$

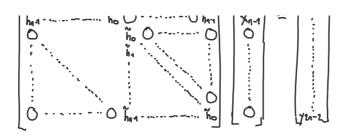
Case m=n:

$$\begin{array}{c|cccc}
h_0 & O & \cdots & O \\
h_1 & h_0 & O & \cdots & O \\
h_1 & h_0 & O & \cdots & O \\
h_{n-1} & \cdots & h_1 & h_0 & \vdots \\
O & \cdots & O & h_{n-1}
\end{array}$$

Maximal doration of output: n+m-1

Discrete circular/perodic convolution

In order that the dimensions work out we need to pad x and h to the length of y here to a length L:= n+m-1



Note: Linear convolution and perodic convolution give the same results. Just the way to get it is different The reason we do periodic consolution is the composition. We and of it vay fast with FFT.

Competion of Circular Convolution (Normal)

1.) Zero Pad

2.) Compute their DFT's

3.) Pointwise multiplication

4.) Take Invase DFT

To = \frac{1}{n} \cdot \tau_1^H

O(n^2)

Complexity: Not fister than linear convolution. But now do FFT:

Competion of Circular Convolution (FFT)

Complexity: $O\left(\frac{3n}{2} \cdot \log_2(n)\right) = O(n \cdot \log(n))$

Eigen: 103, 110

Frequency Filtering with DFT

Denoising: Cutoff No. of Forer coefficients

Trick: Only 1 high value in the Fourierspettrom =) regular Graph

2D-DFT

$$U *_{n,n} X = \frac{1}{m \cdot n} \overline{f_n} \left[(\overline{f_n} U \overline{f_n}) \cdot (\overline{f_n} X \overline{f_n}) \right] \overline{f_n}$$
 $\overline{f_n} = \overline{f_n}^H$
Eigen: 5.113