

1. Computing with Matrices/Vectors

If-Statement: $\text{if}(|\text{abs}(x)| < \text{numeric_limits}(\text{double})::\text{epsilon})$

Absolute Error: $\epsilon_{\text{abs}} = |x - \tilde{x}|$

Relative Error: $\epsilon_{\text{rel}} = \frac{|x - \tilde{x}|}{|x|}$ \tilde{x} has l correct digits if $\epsilon_{\text{rel}} \leq 10^{-l}$

EPS: $\text{EPS} = \max_{x \in \mathbb{R}} \frac{|\text{rd}(x) - x|}{|x|}$

Rounding: $\text{rd}(x) : \begin{cases} \mathbb{R} \rightarrow \mathbb{M} \\ x \rightarrow \max(\arg\min_{\tilde{x} \in \mathbb{M}} |x - \tilde{x}|) \end{cases}$

Trick: EPS is smallest number such that $1 \mp \text{EPS} \neq 1$ ($10 \mp \text{EPS} = 10$)

Assume you solve $Ax = b$ with x_{ex} being the exact value:

Forward error: $|x_{\text{app}} - x_{\text{ex}}|$ not computable

Backward error: $|b_{\text{app}} - b|$ computable

Small backward error \nRightarrow Small forward error

S.p.d. Matrix: Matrix a is s.p.d. if

$$\begin{cases} A^H = A \\ \forall x \in \mathbb{R}^n : x^H A x \in \mathbb{R} \quad x^H A x > 0 \Leftrightarrow x \neq 0 \\ \text{All eigenvalues positive} \end{cases}$$

Complexity:

| what | what | # of \times operations | # of $+$ operations | complexity |
|----------------|---------|--------------------------|-------------------------|------------|
| Dot product | $x^H y$ | n | $n-1$ | $O(n)$ |
| Tensor product | $x y^H$ | $n \cdot m$ | 0 | $O(nm)$ |
| Matrix product | $A B$ | $n \cdot m \cdot k$ | $m \cdot k \cdot (n-1)$ | $O(nmk)$ |

Stability: An algorithm \tilde{F} for finding y for a given b is numerically stable if:

$$\forall b \in \mathbb{M} : \frac{\|b - b_{\text{app}}\|}{\|b\|} = O(\text{cost}(\tilde{F}) \cdot \text{EPS})$$

How close is $\tilde{F}(x) (= F(\tilde{x}))$ to $F(x)$? \Rightarrow Condition Number

$$\text{cond}(A) := \|A\| \cdot \|A^{-1}\| = \frac{\sigma_{\max}}{\sigma_{\min}} \quad (\text{singular values})$$

Trick: Impact of roundoff errors within \tilde{F} is not worse than impact of rounding the input for \tilde{F}