

3. Direct methods for linear least Squares

Introduction: Want to solve $Ax=b$ but the system is over determined. So you won't find an exact solution

$x \in \text{lsq}(A,b) \iff Ax$ is closest element in $R(A)$ to b

Note: If $\text{rank}(A) = n$ you still have the problem that b just isn't in $R(A)$. Therefore Least squares and not Gauss / LU

Normal Equation: $A^T A x = A^T b$

To be a unique solution we would need: $N(A^T A) = N(A) \stackrel{!}{=} \{0\}$

1.) Compute $C := A^T A \quad O(m \cdot n^2)$

2.) Compute $c := A^T b \quad O(n \cdot m)$

3.) Solve s.p.d. LSE $Cx=c \quad O(n^3) \implies O(n^2 m + n^3)$

Sparse A:

Problem: $A^T A$ not sparse anymore, Squaring cond(A)

Fix: $\begin{bmatrix} -\alpha I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$ Use α to do good conditioning.
 $[O(n^{3/2}), O(n^{5/2})]$

Generalized Solution:

If $N(A^T A) \neq \{0\}$ we get a unique solution by taking that x , such that:

$x^\dagger = \arg\min \|x\|_2 : x \in \text{lsq}(A,b)$

$x^\dagger = V(V^T A^T A V)^{-1} V^T A^T$ Moore-Penrose-Pseudoinverse A^\dagger of A .
 \hookrightarrow Does not depend on V

Note: V is a Basis for the $\text{Kern}(A)$.

Note: Alternative way to compute the Moore-Penrose-inv. \rightarrow SVD

Orthogonal Transformation (\Rightarrow Householder QR) ($\text{Rank}(A) = n$)

$A = QR$ with Q unitary, R upper Triangular.

$Rx = Q^T b$ Backwards substitution: $O(n^2)$

To compute the QR factorization:

Gramm Schmidt: unstable, $O(n^3)$

Householder: stable, $O(m \cdot n^2)$

Matrix $X \in \mathbb{R}^{n \times n}$ is an householder Q : // applies all $\|x\|_2 = \|x\|_2$

Singular Value Decomposition

- + If $\text{cond}(A) \gg 0$
- + Construction of Moore-Penrose

$$A = U \Sigma V^H \quad \begin{array}{l} \bullet \text{ Exists for all } A \\ \bullet U, V \text{ are unitary matrices} \end{array}$$

Get Factorization: $O(m \cdot n^2)$ S. 84


To solve least squares:

$$A = \begin{bmatrix} U_1 & U_2 \end{bmatrix} \begin{bmatrix} \Sigma_r & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} V_1^H \\ V_2^H \end{bmatrix} \quad r = \text{Rank}(A)$$

Not generalized: $\sum_r V_1^H x = U_1^H b$ (not unique)

generalized: $x = \underbrace{V_1 \cdot \sum_r^{-1} U_1^H b}_{\text{Moore-Penrose-Pseudoinverse}}$ (unique)

SVD-Based Optimization / Approximation

- $\|Ax\|_2 \rightarrow \min$ with $\|x\| = 1$
 $\hookrightarrow x = (V)_{1:n} = V \cdot e_1$
- Find \tilde{A} that's close to A and has Rank k .
 $\hookrightarrow \tilde{A} = \sum_{i=1}^k \sigma_i \cdot (U)_{:,i} \cdot (V)_{:,i}^H$
- Principle Component Analysis ($\mathcal{P}(A)$)
 $\hookrightarrow \sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_p \gg \sigma_{p+1} \approx \dots \approx \sigma_{\min(m,n)} \approx 0$
 \hookrightarrow z.B.  \Rightarrow 2 diodes (electricity)