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1. If $t(n) \in O(g_1(n))$ and $t_2(n) \in O(g_2(n))$, then $t_1(n) + t_2(n) \in O(\max\{g_1(n), g_2(n)\})$ prove that assertions

sol We need to show that $t_1(n)+t_2(n) \in O\max\{g_1(n), g_2(n)\}$.

This means there exists a positive constant c and c and n_0 such that $t_1(n) + t_2(n) \leq c$

$t_1(n) \leq c_1 g_1(n)$ for all $n \geq n_1$

$t_2(n) \leq c_2\, g_2$ for all $n \geq n_2$

let $n_0 = \max\{n_1, n_2\}$ for all $n \geq n_0$.

$t_1(n) + t_2(n) \leq c_1 g_1(n) + c_2 g_2(n)$

we need to Rotate $g_1(n)$ and $g_2(n)$ to $\max\{g_1(n), g_2(n)\}$

Thus

$c_1 g_1(n) \leq c_1 \max\{g_1(n), g_2(n)\}$

$c_2 g_2(n) \leq c_2 \max\{g_1(n), g_2 g_2(n)$

$c_1 g_1(n) + c_2 g_2(n) \leq c_1 \max\{g_1(n), g_2(n)\} + c_2 \max\{g_1(n), g_2(n)\}$

$c_1 g_1(n) + c_2 g_2(n) \leq (c_1 + c_2) \max\{g_1(n), g_2(n)\} \max\{g_1(n), g_2(n)\}$

$t_1(n) + t_2(n) \leq (c_1 + c_2) \max\{g_1(n), g_2(n)\}$ for all $n \geq n_0$.

By the defination of Big 'O' Notation

$t_1(n) + t_2(n) \in O\max\{g_1(n), g_2(n)\}$

$t_1(n) + t_2(n) \in O\max\{g_1(n), g_2(n)\}$.

Thus, the assertion is proved

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(2) Find the time complexity of the Recurrence equation

Sol Let us consider such that Recurrence for merge sort

$T(n) = 2T(n/2) + n$

By using master theorem

$T(n) = aT(n/b) + f(n)$

where $a \geq 1, b \geq 1$ and $f(n)$ is Positive function

Ex: $T(n) = 2T(n/2) + n$

$a = 2, b = 2, f(n) = n$

By comparing of $f(n)$ with $n \log_b a$

$\log_b a = \log_2 2 = 1$

compare $f(n)$ with $n^{\log_b a}$

$f(n) = n$

$n^{\log_b a} = n^1 = n$

* $f(n) = O(n^{\log_b a})$ then $T(n) = O(n^{\log_b a} \log n)$

In our case;

$\log_b a = 1$

$T(n) = O(n^1 \log n) = O(n \log n)$

The time complexity of recurrence Relation is $T(n) = 2T(n/2) + n$ is $O(n \log n)$.

③ $T(n) = \begin{cases} 2T(n/2)+1 & \text{if } n>1 \\ 1 & \text{otherwise} \end{cases}$

अ By applying of master theorem

$T(n) = aT(n/b) + f(n)$ where $a \geq 1$, $b > 1$

$T(n) = 2T(n/2)+1$

Here $a=2, b=2, f(n)=1$

By comparision of $f(n)$ and $n^{\log_b a}$

If $f(n) = O(n^c)$ where $c < \log_b a$, then $T(n) = O(n^{\log_b a})$

If $f(n) = O(n^{\log_b a})$, then $T(n) = O(n^{\log_b a} \log n)$.

If $f(n) = \Omega(n^c)$ where $c > \log_b a$ then $T(n) = O(f(n))$

lets calculate $\log_b a$:

$\log_b a = \log_2 2 = 1$

$f(n) = 1$

$n^{\log_b a} = n^1 = n$

$f(n) = O(n^c)$ with $c < \log_b a$ (case 1)

In this case $c=0$ and $\log_b a = 1$

$c < 1$, so $T(n) = O(n^{\log_b a}) = O(n^1) = O(n)$

Time complexity of Recorrence Relation

$T(n) = 2T(n/2)+1$ is (On).

(4) $T(n) = \begin{cases} 2T(n-1) & \text{if } n>0 \\ 1 & \text{Otherwise} \end{cases}$

Sol Here, where $n=0$

$T(0)=1$

Recurrence Relation Analysis :

for $n>0$

$T(n) = 2T(n-1)$

$T(n) = 2T(n-1)$

$T(n-1) = 2T(n-2)$

$T(n-2) = 2T(n-3)$

$T(1) = 2T(0)$

from this pattern

$T(n) = 2\cdot 2\cdot 2 \cdots 2T(0) = 2^n T(0)$

Since $T(0)=1$, we have.

$T(n) = 2^n$

The recurrence relation is

$T(n) = 2T(n-1)$ for $n>0$ and $T(0)=1$ is $T(n) = 2^n$.

(5) Big O notation show that $f(n) = n^2 + 3n + 5$ is $O(n^2)$

$f(n) = O(g(n))$ means $c>0$ and $n_0 \geq 0$

$f(n) \leq c \cdot g(n)$ for all $n \geq n_0$

Given is $f(n) = n^2 + 3n + 5$

$c>0, n_0 \geq 0$ such that $f(n) \leq cn^2$

$f(n) = n^2 + 3n + 5$

Lets choose $c=2$

$f(n) \leq 2 \cdot n^2$

$f(n) = n^2 + 3n + 5 \leq n^2 + 3n^2 + 5n^2 = 9n^2$

So, $c=9$, $n_0=1$ $f(n) \leq 9n^2$ for all $n \geq 1$

$f(n) = n^2 + 3n + 5$ is $O(n^2)$.