Procedure Msort
$$(A, l, h)$$
 {

if (h) l l
 $m = \lfloor \frac{l+h}{2} \rfloor \longrightarrow D$
 $m = \lfloor \frac{l+h}{2} \rfloor \longrightarrow T(\frac{n}{2})$
 $T(h)$
 $Msort(A, M, h) \longrightarrow T(\frac{n}{2})$
 $Msorf(A, m+l, h) \longrightarrow T(\frac{n}{2})$
 $Merge(A, l, m, h) \longrightarrow Cn$

End Procedure

Procedure Merge
$$(A, l, m, h)$$
?

while $loop \rightarrow if \cdot else \rightarrow c_1 \frac{n}{2} \leftarrow c_2 \frac{n}{2} \leftarrow c_1 \frac{n}{2} \leftarrow c_2 \frac{n}{2} + c_2 \frac{n}{2} + c_3 n$

if $\cdot \cdot else \rightarrow for loop \rightarrow c_2 \frac{n}{2} \leftarrow c_1 \frac{n}{2} + c_2 \frac{n}{2} + c_3 n$

$$= \left(\frac{c_1}{L} + \frac{c_2}{L} + c_3\right) n$$

$$= cn.$$

Total =
$$\frac{(2 + 2 + 3)}{(2 + 3)}$$
 = $\frac{(2 + 2 + 3)}{(2 + 3)}$ = $\frac{(2 + 2 + 3)}{(2 + 3)}$

Recurrence Relation for Time

$$T(n) = \begin{cases} A & n=1 \\ aT(\frac{n}{L}) + cn + D & n > 1 \text{ and } n = 2^{K} \end{cases}$$

$$\begin{array}{l} = 2 T \left(\frac{n}{L} \right) + cn + D \\ = 2 \left[2 T \left(\frac{n}{L} \right) + c \frac{n}{L} + D \right] \\ + cn + D \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + 2D + D \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + D \left[2 + 1 \right] \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + D \left[2 + 1 \right] \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + D \left[2 + 1 \right] \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + D \left[2 + 1 \right] \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 2cn + D \left[2 + 1 \right] \\ = 2^{n} T \left(\frac{n}{L^{2}} \right) + 3cn + D \left[2^{n} + 2^{n$$