

## Fibonacci Series

1	1	2	3	5	8	13
0	1	2	<u>3</u>	4	5	6

$$f_{ib}(n) = \frac{f_{ib}(n-1)}{\text{last}} + \frac{f_{ib}(n-2)}{\text{second last}}$$

# PMI - Extended Form

1. Proof for base case  $f(0) \quad f(1)$

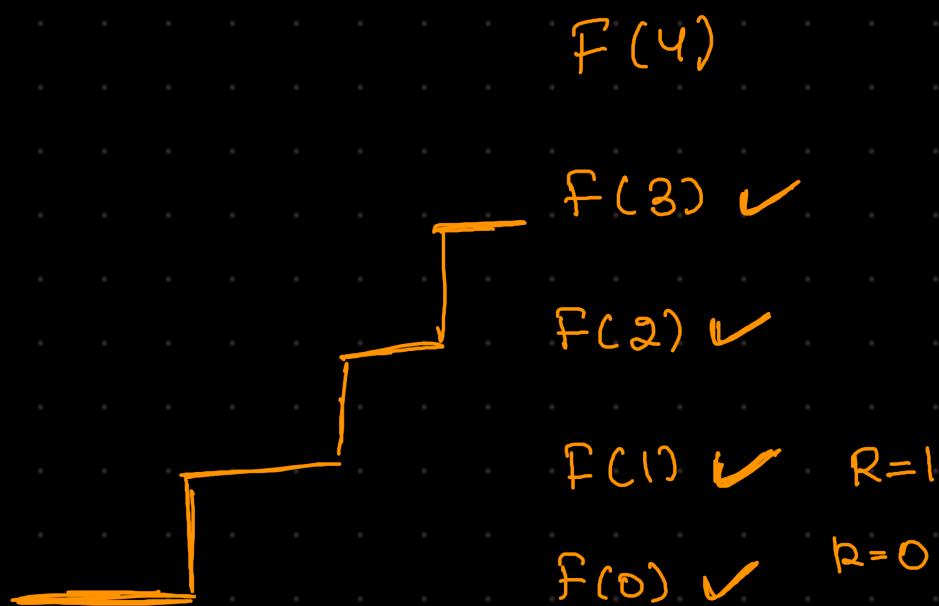
2. Assume for  $f(i)$  equal to true

where you  $i$  can be from

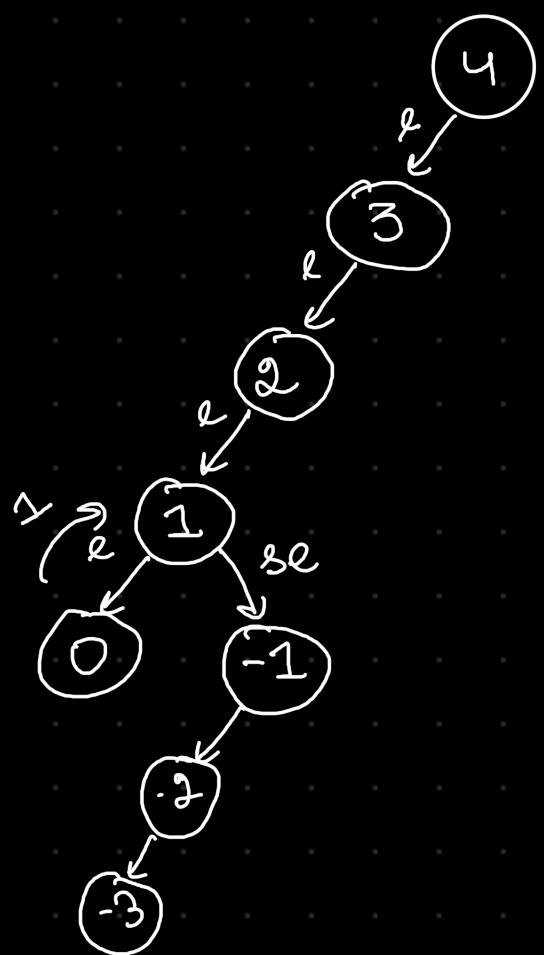
$0 \leq i \leq R \quad f(R)$

so that means we can assume  
for

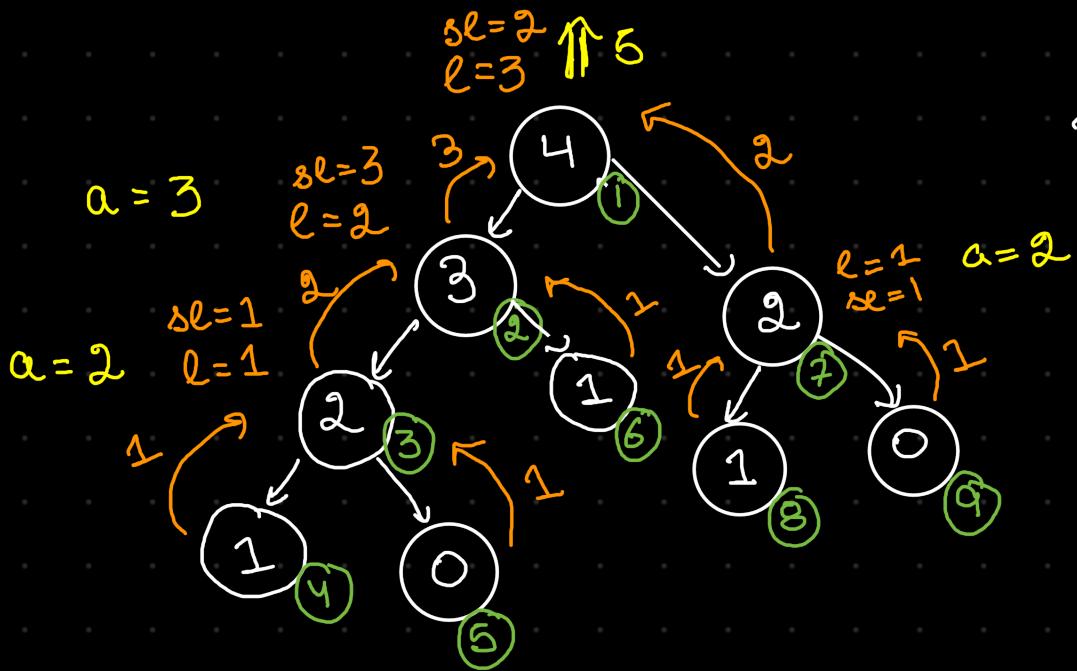
$f(R)$  true  
 $f(R-1)$  true }  $\Rightarrow$  all assume true  
for less than  $n$



last  $f(n-1)$  to be true  
second last  $f(n-2)$  to be true



# Fibonacci Series : Recursion Tree



```
def fibonacci(n):
    # print(n)
    if(n==0):
        return 1
    if(n==1):
        return 1

    last = fibonacci(n-1)
    secondLast = fibonacci(n-2)

    ans = last + secondLast

    return ans
```

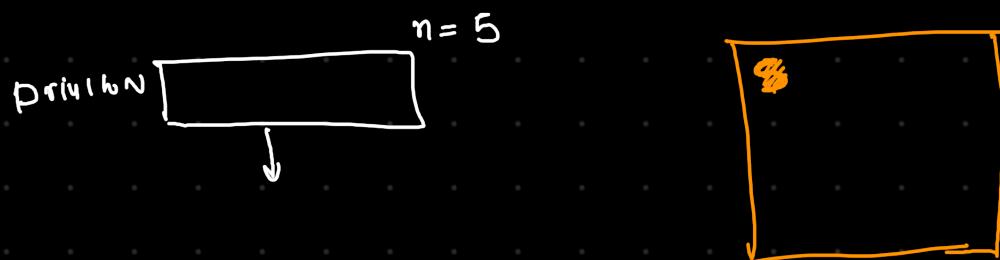
```
The function call is 1 for 4
The function call is 2 for 3
The function call is 3 for 2
The function call is 4 for 1
The function call is 5 for 0
The function call is 6 for 1
The function call is 7 for 2
The function call is 8 for 1
The function call is 9 for 0
```



Q: Write a program for a given number n.

1. Print 1 to N.
2. Print N to 1

Print 1 to N :-



Print N to 1



1. Base Case
2. Recursive Call ↗
3. Our work ↘

Head

Tail

## Assignment

1. Sum of digit of a number
2. Power of a number (base, exp)  
3      2

# Head vs Tail Recursion

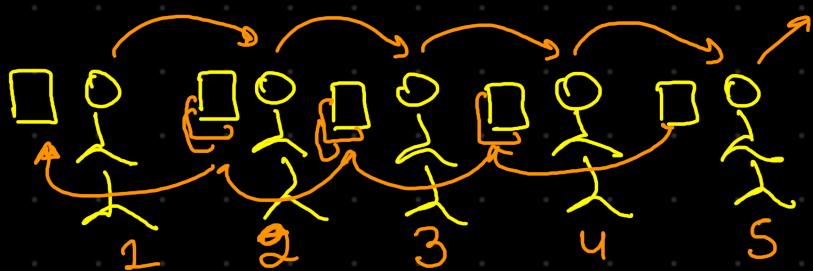
## Head

We make recursive call  
at the beginning of  
our fn implementation.

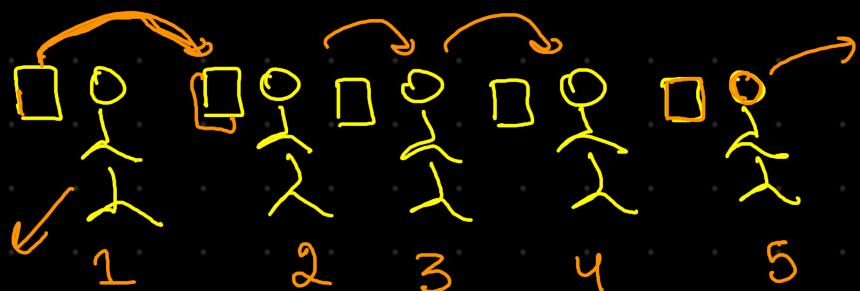
## Tail

When we make recursive  
call at the end of our  
implementation

head



tail



## Recursion with Arrays/List

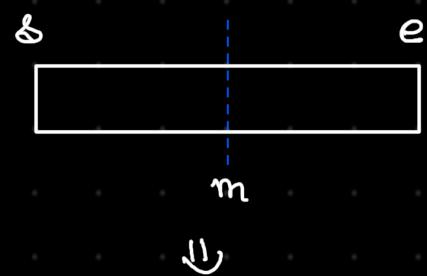
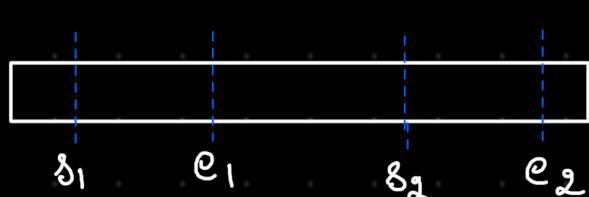
Whenever in arrays, we are dealing with a problem that can be broken & individually be applied to left over part, we can use recursion there.

Below are major ways we break our array/list.

1. First or last element removed



2. Using start and end index (e.g. binary search)



divide from middle

3. Copy part into a new array and send to further calls.

Q: Write a Program to check if array is sorted?

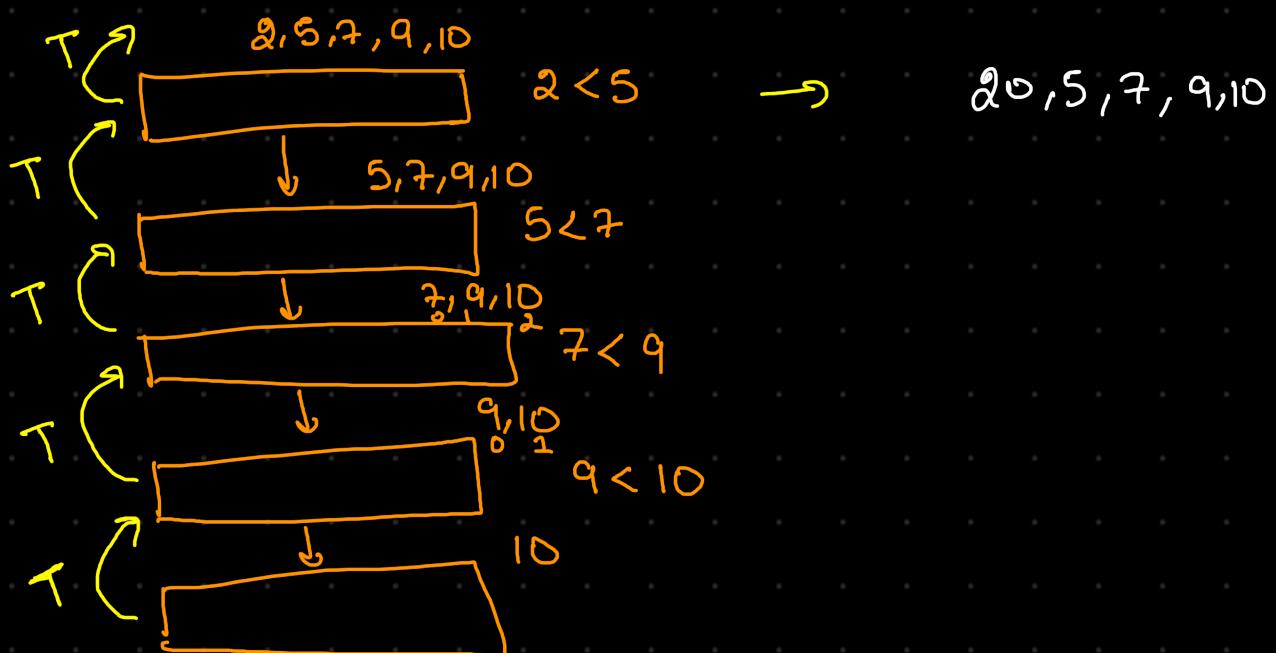
recurrence relation

recursion tree

head / tail recursion.

1. Base case : size == 1 or size == 0

2. Recursion call



check sorted ( $l1$ ) = ( $l1[0] < l1[1]$ ) and check sorted ( $l1[1:]$ )

Q: Sum of an array using Recursion.

Please specify recursion relation & draw recursion tree as well.

Give head and tail recursion solution.

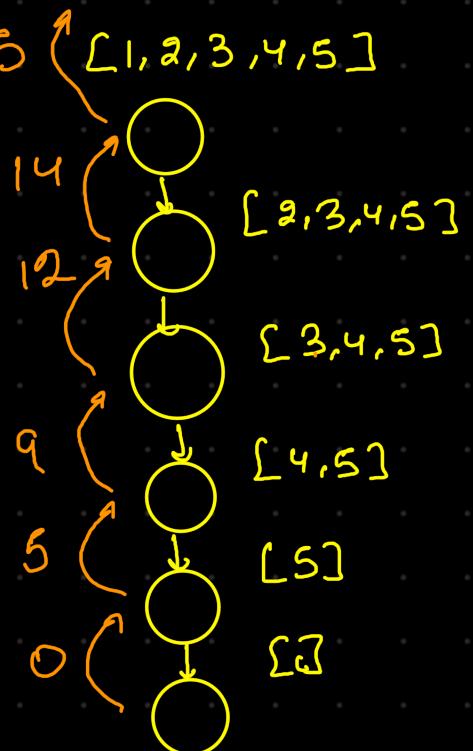


1. Base Case      size=1      l[0]

size = 0      0

2. sum (l1[1:])

3. l1[0] + sum of l1



So, above question showed us that either we can do our work first or after recursive call.

- 1] Base Case
- 2] Recursive call / Small answer
- 3] Our calculation



Q: Given an array and an element , find

a) first index of element.

b) last index of element.

[ 3, 2, 5, 2, 8, 2, 1 ] , 2

-1 if not found

first : 1  
last : 5

first index :-

- 1] Base case
- 2] if  $l1[0] == x$   
return 0
- 3] recursive call

Q: Write a program to find all index of a given element.

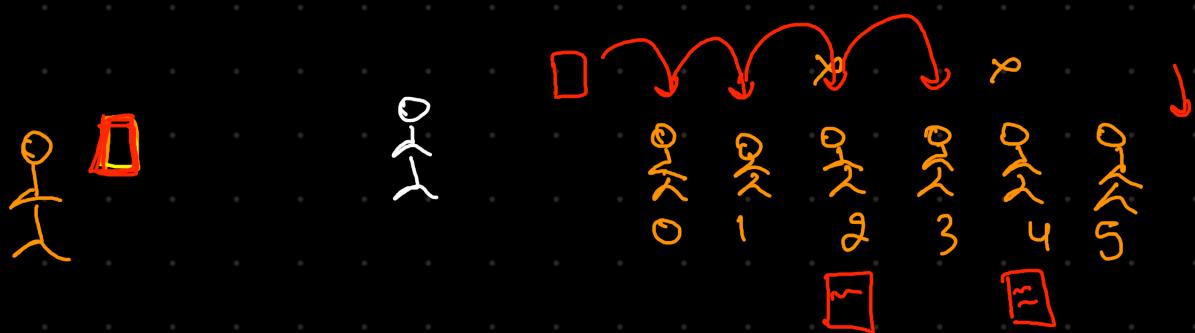
Result have to be as below :-

- 1] Print the indexes
- 2] Update the answer in list provided
- 3] Update indices in the global list
- 4] Return answer as a new list.

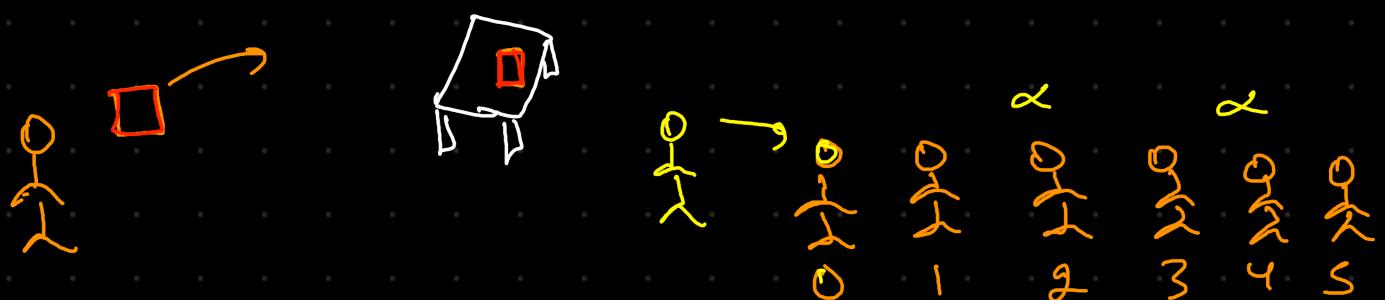
1]  $\{3, 2, 5, 2, 8, 2, 1\}$ , 2

1  
3  
5

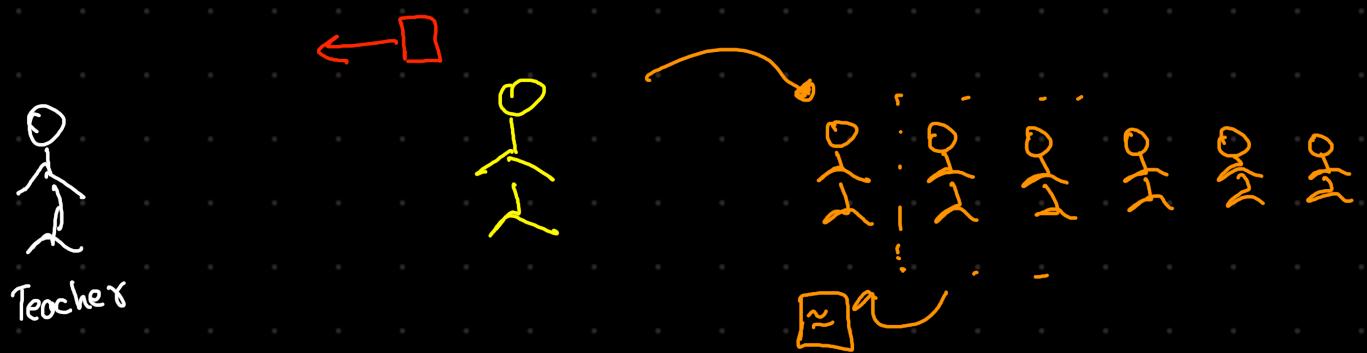
Update all indices in a provided list



Update indices in a global list



Return a list as an answer



1. Base case