Isomorphism & Hamiltonian Graphs

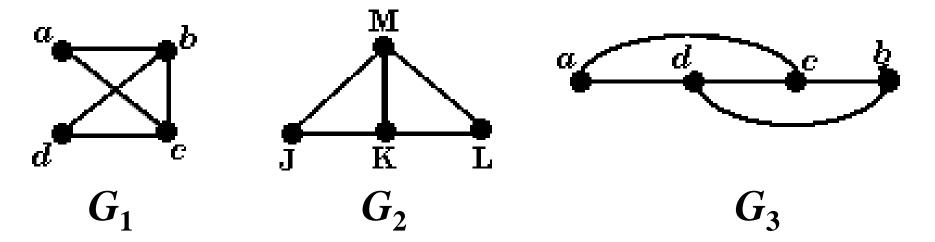
Isomorphism

- Two graphs are isomorphic when the vertices of one can be re-labeled to match the vertices of the other in a way that preserves adjacency.
- Formally, Graphs G_1 and G_2 are isomorphic if there exists a one-to-one function, called an isomorphism,

$$f: V(G_1) \rightarrow V(G_2)$$

such that uv is an element of $E(G_1)$ if and only if f(u)f(v) is an element of $E(G_2)$.

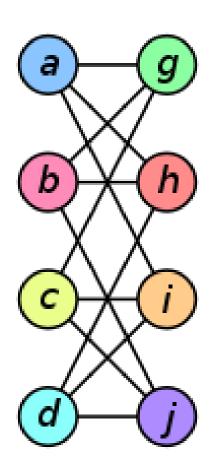
Example

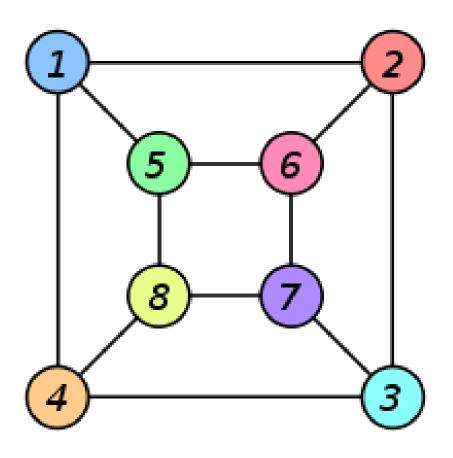


- G_1 and G_2 f(a) = J, f(b) = K, f(c) = M, and f(d) = L.
- G_1 and G_3 $f(a)=a, f(b)=d, f(c)=c, \, \mathrm{and} \, f(d)=b.$

Example

- f(a) = 1
- f(b) = 6
- f(c) = 8
- f(d) = 3
- f(g) = 5
- f(h) = 2
- f(i) = 4
- f(j) = 7





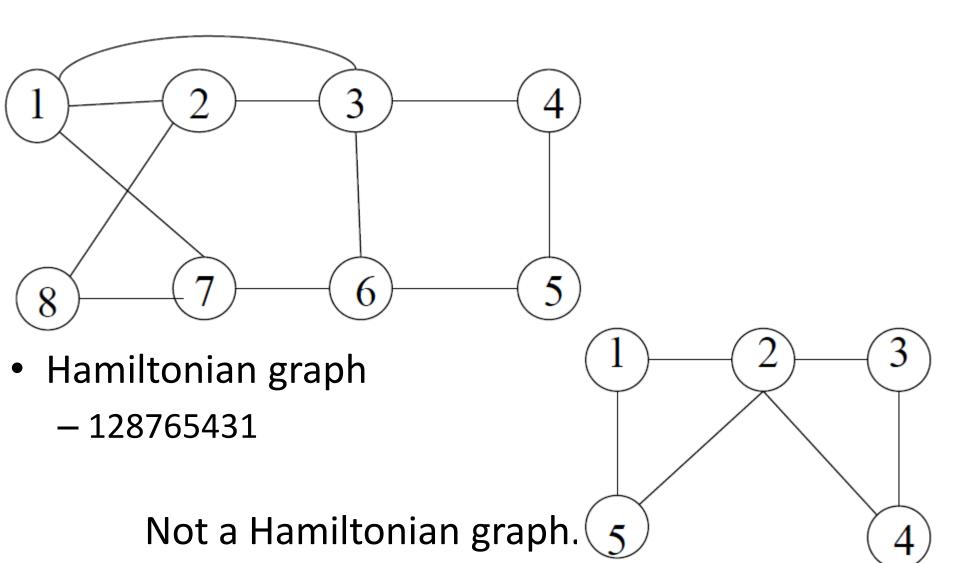
Graph Isomorphism Problem

- Determining whether two finite graphs are isomorphic or not?
- Applications:
 - Cheminformatics,
 - Mathematical chemistry (identification of chemical compounds), and
 - Electronic design automation (verification of equivalence of various representations of the design of an electronic circuit).
 - And Similar other problems

Hamiltonian Graphs

- A path passing through all the vertices of a graph is called a Hamiltonian path.
 - A graph containing a Hamiltonian path is said to be traceable.
- A cycle passing through all the vertices of a graph is called a Hamiltonian cycle (or Hamiltonian circuit).
- A graph containing a Hamiltonian cycle is called a **Hamiltonian graph**.
- Hamiltonian path problem: Determine the existence of Hamiltonian paths and cycles in graphs (NP-complete).

Example



Solution 1 – Brute Force Search Algorithm

 A Hamiltonian Path in a graph having N vertices is nothing but a permutation of the vertices of the graph

$$[v_1, v_2, v_3,v_{N-1}, v_N]$$

such that there is an edge between v_i and v_{i+1} where $1 \le i \le N-1$.

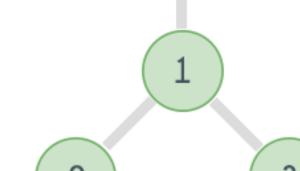
 So it can be checked for all permutations of the vertices whether any of them represents a Hamiltonian Path or not.

Contd...

```
function check_all_permutations(adj[][], n)
  for i = 0 to n
       p[i]=i
  while next permutation is possible
       valid = true
       for i = 0 to n-1
              if adj[p[i]][p[i+1]] == false
                      valid = false
                      break
       if valid == true
              print_permutation(p)
              return true
       p = get_next_permutation(p)
  return false
```

Example

Not a Hamiltonian graph as path exists and not a circuit.



Solution 2 – Backtracking

- 1. Create an empty path array and add vertex 0 to it.
- 2. Add other vertices, starting from the vertex 1.
- 3. Before adding a vertex, check for whether it is adjacent to the previously added vertex and not already added.
- 4. If such a vertex is found, add that vertex as part of the solution.
- 5. Otherwise, return false.

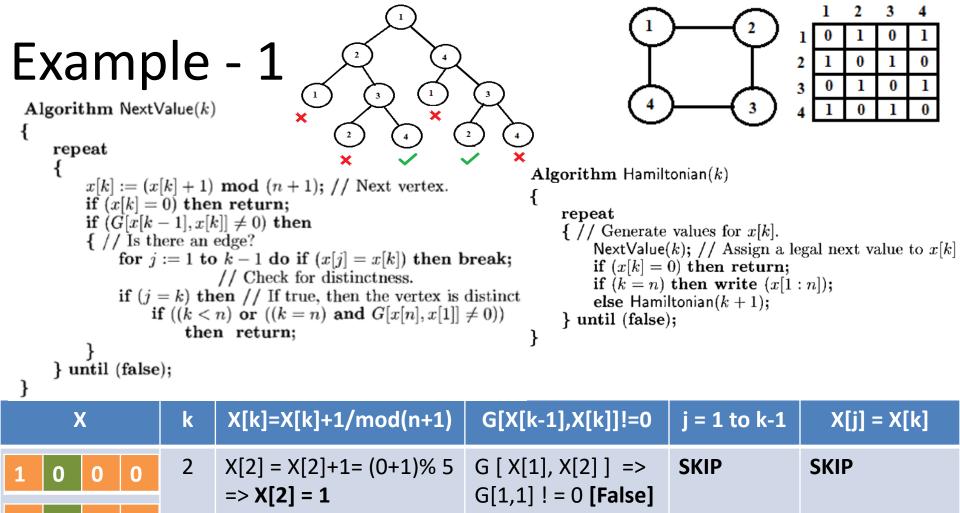
```
Contd...1
                  Algorithm Hamiltonian(k)
                  // This algorithm uses the recursive formulation of
                  // backtracking to find all the Hamiltonian cycles
                  // of a graph. The graph is stored as an adjacency
                   // matrix G[1:n,1:n]. All cycles begin at node 1.
                       repeat
                       \{ // \text{ Generate values for } x[k]. \}
                           NextValue(k); // Assign a legal next value to x[k].
                           if (x[k] = 0) then return;
               10
                           if (k = n) then write (x[1:n]);
                           else Hamiltonian(k+1);
              12
                       } until (false);
              13
```

- Array x is a solution vector. x[i] represents the ith visited vertex of the proposed cycle.
- -G[1:n,1:n] is the adjacency matrix of the given graph.
- Initialize x[1] to 1 (as vertex 1 is the starting vertex) and x[2:n] to zero, then call Hamiltonian(2).

Let,

14 }

```
Algorithm NextValue(k)
   //x[1:k-1] is a path of k-1 distinct vertices. If x[k]=0, then
   // no vertex has as yet been assigned to x[k]. After execution,
   //x[k] is assigned to the next highest numbered vertex which
   // does not already appear in x[1:k-1] and is connected by
   // an edge to x[k-1]. Otherwise x[k] = 0. If k = n, then
6
7
8
    // in addition x[k] is connected to x[1].
9
         repeat
10
             x[k] := (x[k] + 1) \mod (n + 1); // \text{ Next vertex.}
11
             if (x[k] = 0) then return;
12
             if (G[x[k-1], x[k]] \neq 0) then
13
14
             { // Is there an edge?
                  for j := 1 to k - 1 do if (x[j] = x[k]) then break;
15
                               // Check for distinctness.
16
                  if (j = k) then // If true, then the vertex is distinct
17
                      if ((k < n) \text{ or } ((k = n) \text{ and } G[x[n], x[1]] \neq 0))
18
19
                           then return;
20
         } until (false);
21
22
```



Now, j == k i.e., 2 == 2 (distinct) and 2<4, so NextValue(2) returns. In Hamiltonian(2) => X[2] !=0 and 2 != 4, so Hamiltonian(k+1) i.e., Hamiltonian(3) is recursively called.

X[2] = X[2]+1=(1+1)%5

=> X[2] = 2

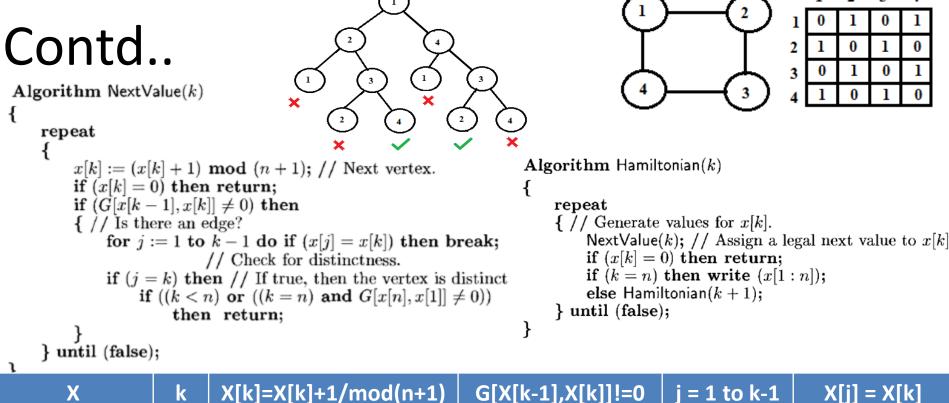
G [X[1], X[2]] =>

G[1,2]! = 0 [True]

i = 1 to 1;

X[1] = X[2]

=> 1 == 2



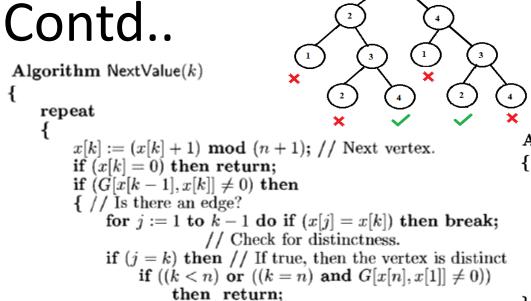
ı	if $(j = k)$ then // If true, then the vertex is distinct if $((k < n)$ or $((k = n)$ and $G[x[n], x[1]] \neq 0)$) then return; } until (false); if $(k = n)$ then write $(x[1 : n])$; else Hamiltonian $(k + 1)$; until (false);									
	X			k	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]		
1	2	0	0	3	X[3] = X[3]+1= (0+1)% 5 => $X[3] = 1$	G [X[2], X[3]] => G[2,1]!=0 [True]	j = 1	X[1] = X[3] [T] Break		
1	2	1	0							

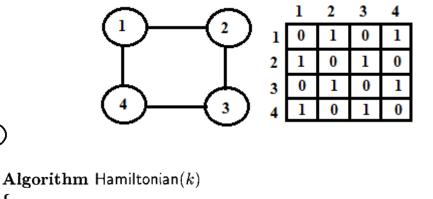
X[3] = X[3]+1= (1+1)%5 G[X[2], X[3]] => **SKIP SKIP** => X[3] = 2G[2,2]! = 0 [False] X[3] = X[3]+1= (2+1)%5 G[X[2], X[3]] => j = 1 X[1] = X[3] [F]

=> X[3] = 3G[2,3]! = 0 [True] i = 2X[2] = X[3] [F]

Now, j == k i.e., 3 == 3 (distinct) and 3 < 4, so NextValue(3) returns. In Hamiltonian(3) => X[3] !=0 and 3!= 4, so Hamiltonian(k+1) i.e., Hamiltonain(4) is recursively called.

} until (false);





NextValue(k); // Assign a legal next value to x[k]

X				k	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]
1	2	3	0	4	X[4] = X[4]+1= (0+1)% 5 => $X[4] = 1$	G [X[3], X[4]] => G[3,1] ! = 0 [False]	SKIP	SKIP
1	2	3	2		X[4] = X[4]+1= (1+1)%5 => X[4] = 2	G [X[3], X[4]] => G[3,2] ! = 0 [True]	j = 1 j = 2	X[1] = X[4] [F] X[2] = X[4] [T] BREAK
1	2	3	3		X[4] = X[4]+1= (2+1)%5 => $X[4] = 3$	G [X[3], X[4]] => G[3,3]!=0 [False]	SKIP	SKIP

repeat

} until (false);

 $\{ // \text{ Generate values for } x[k].$

if (x[k] = 0) then return;

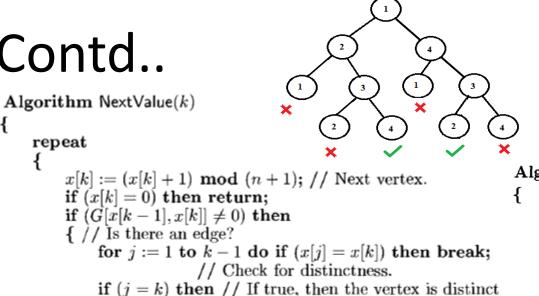
else Hamiltonian(k+1);

if (k = n) then write (x[1:n]);

Contd..

repeat

} until (false);



if $((k < n) \text{ or } ((k = n) \text{ and } G[x[n], x[1]] \neq 0))$

then return:

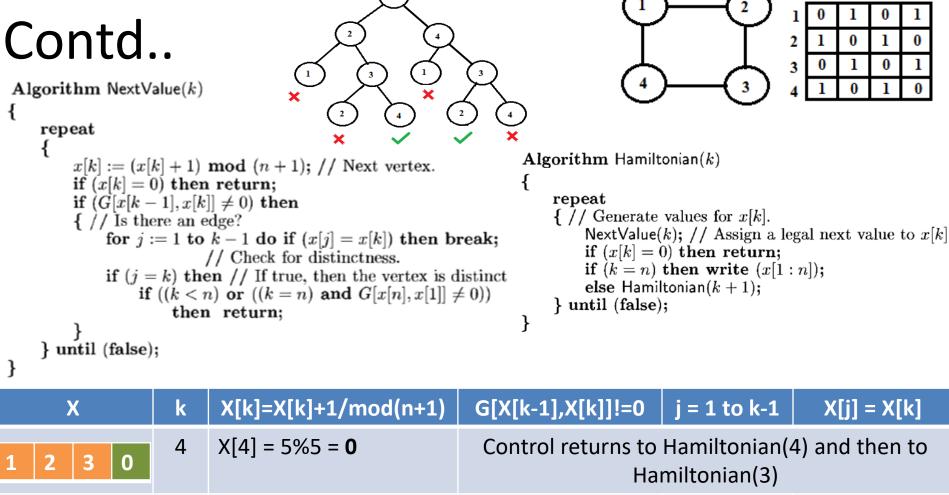
```
2
Algorithm Hamiltonian(k)
    repeat
    \{ // \text{ Generate values for } x[k]. \}
         NextValue(k); // Assign a legal next value to x[k]
         if (x[k] = 0) then return;
        if (k = n) then write (x[1:n]);
        else Hamiltonian(k+1);
    } until (false);
```

X				k	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]
1	2	3	4	4	X[4] = X[4]+1= (3+1)%5			X[1] = X[4] [F]
					=> X[4] = 4	G[3,4] ! = 0 [True]	j = 2 j = 3	X[2] = X[4] [F] X[3] = X[4] [F]

Now, j == k i.e., 4 == 4 (distinct) and k == n (4 == 4) and G[X[n], X[1]]i.e., G[4,1] != 0 (there is an edge between last and first vertex),

So, NextValue(4) returns. In Hamiltonian(4) => as X[4] != 0 and k == 4, So, First solution is printed as: 1, 2, 3, 4

And then cycle continues for finding out other solutions in the state space search tree.



} until (false) }	} until (false);								
X	k	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]				
1 2 3 0	4	X[4] = 5%5 = 0	Control returns to Ha	Hamiltonian(4 miltonian(3)	4) and then to				
1 2 4 0	3	X[3]=4%5 = 4	G[X[2],X[3]] i.e., G[2,4] != 0 [False]	SKIP	SKIP				

Hamiltonian(2)

G[X[1],X[2]] i.e.,

G[1,3] != 0 [False]

Control returns to Hamiltonian(3) and then to

SKIP

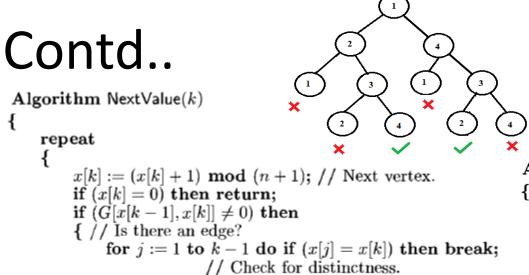
SKIP

X[3]=5%5=0

X[2] = 3%5 = 3

2

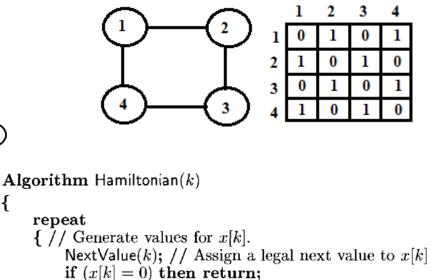
} until (false);



then return:

if (j = k) then // If true, then the vertex is distinct

if $((k < n) \text{ or } ((k = n) \text{ and } G[x[n], x[1]] \neq 0))$



X[2] = X[3] [F]

if (k = n) then write (x[1:n]);

else Hamiltonian(k+1);

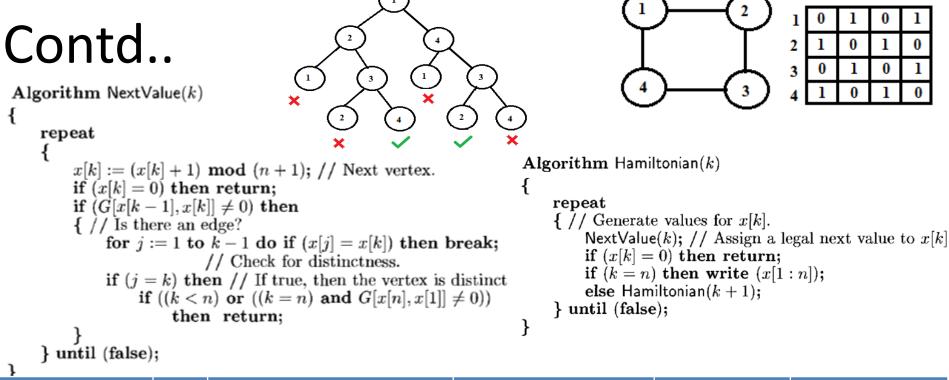
i = 2

} until (false);

X				K	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]
1	4	0	0	2	X[2] = 4%5 = 4	G[X[1],X[2]] i.e., G[1,4] != 0 [True]	j = 1	X[1] = X[2] [F]

Now, j == k i.e., 2 == 2 (distinct) and 2 < 4, so NextValue(2) returns. In Hamiltonian(2) => X[2] != 0and 2!= 4, so Hamiltonian(k+1) i.e., Hamiltonain(3) is recursively called.

1	4	1	0	3	X[3] = 1%5 = 1	G[4,1] != 0 [True]	j = 1 to 2	X[1] = X[3] [T]
1	4	2	0		X[3] = 2%5 = 2	G[4,2] != 0 [False]	SKIP	SKIP
			0		X[3] = 3%5 = 3	G[4,3] != 0 [True]	j = 1	X[1] = X[3] [F]

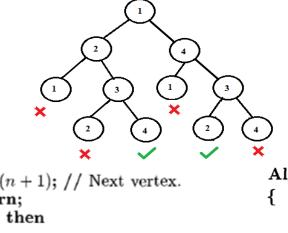


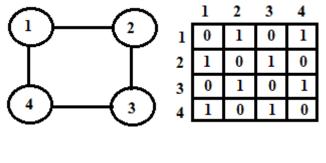
X[4] = 1%5 = 1G[3,1] != 0 [False]**SKIP SKIP** i = 1X[1] = X[4] [F] X[4] = 2%5 = 2G[3,2] != 0 **[True]** i = 2X[2] = X[4] [F] X[3] = X[4] [F] j = 3

i.e., G[2,1] != 0, So, NextValue(4) returns. In Hamiltonian(4) => as X[4] != 0 and K== 4,

Now, j == k i.e., 4 == 4 (distinct) and k == n (4 == 4) and G[X[n],X[1]]So, Second solution is printed as: 1, 4, 3, 2

Contd..



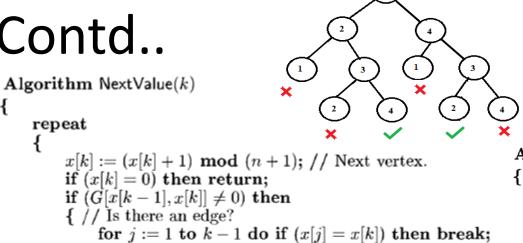


	3 0 1 0 1
Algorithm NextValue(k) { repeat * 2 4 2 4 2 4 7 7 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	4 3 4 1 0 1 0
$x[k] := (x[k] + 1) \mod (n + 1); // \text{ Next vertex.}$ if $(x[k] = 0)$ then return:	f Algorithm Hamiltonian (k)
if $(x[k] = 0)$ then return; if $(G[x[k-1], x[k]] \neq 0)$ then { // Is there an edge? for $j := 1$ to $k-1$ do if $(x[j] = x[k])$ then break; // Check for distinctness. if $(j = k)$ then // If true, then the vertex is distinct if $((k < n))$ or $((k = n))$ and $G[x[n], x[1]] \neq 0)$) then return; } until (false);	repeat { // Generate values for $x[k]$. NextValue(k); // Assign a legal next value to $x[k]$ if $(x[k] = 0)$ then return; if $(k = n)$ then write $(x[1:n])$; else Hamiltonian($k + 1$); } until (false); }

}	then return; } until (false); } until (false);										
	X k			k	X[k]=X[k]+1/mod(n+1)	G[X[k-1],X[k]]!=0	j = 1 to k-1	X[j] = X[k]			
1 1	4	3 3	4 0	4	X[4] = 3%5 = 3 X[4] = 4%5 = 4 X[4] = 5%5 = 0	G[3,3] != 0 [False] G[3,4] != 0 [True] RETURN	SKIP j = 1 j = 2	SKIP X[1] = X[4] [F] X[2] = X[4] [T] BREAK			
1	4	0	0	3	X[4] = 4%5 = 4 X[4] = 5%5 = 0	G[4,4] != 0 [False] RETURN	SKIP	SKIP			



} until (false);



// Check for distinctness.

then return:

if (j = k) then // If true, then the vertex is distinct

if $((k < n) \text{ or } ((k = n) \text{ and } G[x[n], x[1]] \neq 0))$

```
2
Algorithm Hamiltonian(k)
    repeat
    \{ // \text{ Generate values for } x[k].
```

NextValue(k); // Assign a legal next value to x[k]

if (x[k] = 0) then return;

else Hamiltonian(k+1);

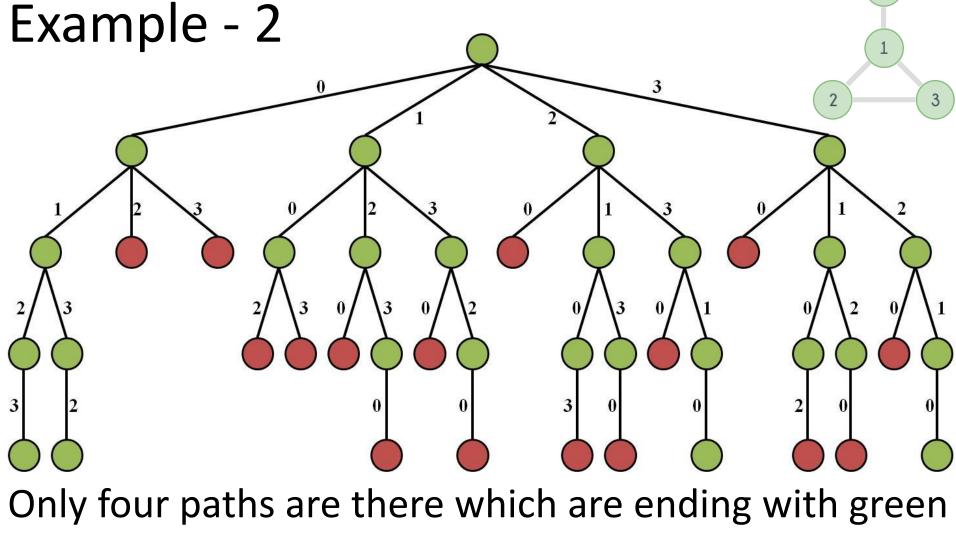
} until (false);

if (k = n) then write (x[1:n]);

```
X[k]=X[k]+1/mod(n+1)
                                            G[X[k-1],X[k]]!=0
                                                                   i = 1 \text{ to } k-1
                                                                                    X[j] = X[k]
X
              X[4] = 5\%5 = 0
                                            RETURN
      0
```

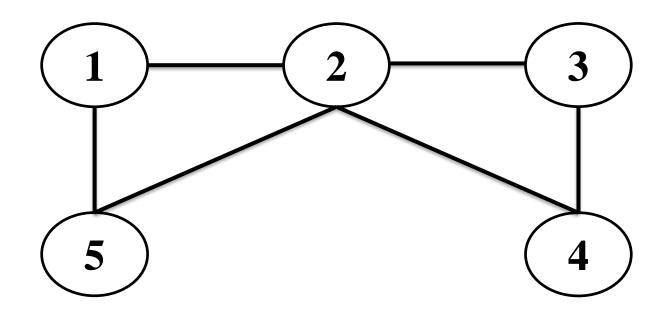
Control finally returns from Hamiltonian(2) [The original initial function call] and

Algorithm Stops !!



Only four paths are there which are ending with green leaf nodes. None of them is a circuit as starting and ending vertices are not adjacent in any of the obtained path, thus this graph is non-hamiltonian.

Home Work



Applications

- Painting road lines,
- Plowing roads after a snowstorm,
- Checking meters along roads,
- Garbage pickup routes, etc.

Thank you