

PRACTICAL - 01.

Aim : Implement linear search to find an item in the list.

Theory :- Linear Search

Linear search is one of the simplest searching algorithms in which targeted item is sequentially matched with each item in list.

It is worst searching algorithm with worst case time complexity. It is a forced approach.

On the other hand in case of an ordered list instead of searching the list in sequence binary search is used which will start by examining middle terms.

Linear search is a technique to compare each & every element with the key element, to be found, if both matches, the algorithm returns that elements found & its position is also found.

if Unsolved :-

Algorithm :-

- (1) Create an empty list & assign it to a variable
- (2) Accept the total no. of elements to be inserted

into the list from the user say 'x'.

- M ③ Use for loop for adding the elements in the list.
- ④ Print the new list.
- ⑤ Accept an element from the user that to be searched in the list.
- ⑥ Use for loop in a range from '0' to the total no. of elements to search the elements from the list.
- ⑦ Use if loop that the elements in the list is equal to the element from user.
- ⑧ If the element is found then print the statement with the element position.
- ⑨ Draw the output of given algorithm.

```

input ("linear Search")
a = []
n = int(input("Enter the range"))
for s in range(0, n):
    s = int(input("Enter the no. = "))
    a.append(s)
print(a)
c = int(input("Enter a search no. = "))
for i in range(0, n):
    if (a[i] == c):
        print("found at position = ", i)
else:
    print("not found")

```

Output :-

```

>>> Linear Search
Enter the range = 3
Enter the number = 2
[2]
Enter the number = 1
[1, 2]

```

Code :-

```

print("Linear Search")
a = []
n = int(input("Enter a range = "))
for s in range(0, n):
    s = int(input("Enter a number = "))
    a.append(s)
a.sort()
print(a)

c = int(input("Enter a search number = "))
for i in range(0, n):
    if (a[i] == c):
        print("found at position = ", i)
        break
    else:
        print("not found")

```

Q) Sorted :-

Ans :- Sorting means to arrange the elements in increasing / decreasing order.

* Algorithm :-

Step 1 : Create an empty list and assign it a variable.

Step 2 : Accept total number of elements to be inserted into the list from user.

Step 3 :- Use for loop to add element in the list using append method().

Step 4 :- Use the sort method to sort the accepted elements and assign in increasing order.

Step 5 :- Use if statement to give a range in which elements needs to be found.

Step 6 :- The use else statement, if statement is not found in range then satisfy the condition.

880

Step 7 : Use for loop in range from 0 to the total no. of elements to be searched. will print minimum

Step 8 : Attach the input and output above Algorithm.

to find the minimum value:

id is already for minimum value is 1000
and result will be 1000

1000 because prime number and all
will be divisible by 1000

between 1000 and below than all
prime numbers are divisible by 1000

Suppose if we take 1000 for all
keep id at their places which is

1000 for 1000 and so on
After suppose if we keep the same
operations at

Output :-

>>> Linear Search.

Enter a range = 8

Enter a number = 2

[2] found in position = 1

Enter a number = 1

[1,2]

Enter a search number = 2

found in position = 1

```
a = []
n = int(input("Enter a range = "))

for b in range(0, n):
    b = int(input("Enter a number = "))
    a.append(b)
print(a)
```

```
s = int(input("Enter a number = "))

if (s < a[0] or s > a[n - 1]):
    print("Element not found")
```

else :

$$f = 0$$

$$l = n - 1$$

for i in range(0, n):

~~$$m = \text{int}((f+l)/2)$$~~

~~print(m)~~

if (s == a[m]):

print("Element found at = ", m)

break.

else :

$$f = m + 1$$

Practical - 02.

~~in small int~~ Aim :- Implement ~~using~~ Binary Search to find a ~~given~~ ~~number~~ in ~~list~~.

Theory :-

Binary Search.

Binary Search is also known as half-interval search, logarithmic search or binary chop is a search algo that finds the position of a target value within a sorted array.

Algorithm :-

Step 1 :- Create an empty list and assign it a variable.

Step 2 :- Using input method, accept the range of given list.

Step 3 :- Use for loops, add elements in list using the append method.

Step 4 :- Use sort() method to sort accept element & assign it. Print the list.

Algorithm

Step 5 :- Use if loop to give the range in which the elements is found.

Step 6 :- Use else statement , if statement is found .

Step 7 :- Accept an argument & key of element that element has to be found with respect to a list .

Step 8 :- Initialize first to 0 and last to element of list .

Step 9 :- Use for loop and assign the range .

Step 10 :- Repeat until you found the element .

Output

>>> Enter a range =4.
Enter a number =2.

[2]

Enter a number =1

[1, 2]

Enter a number =4.

[1, 2, 4]

Enter a search number =2

Element found at =1

`print("Bubble Sort")`

`a = []`

`b = int(input("Enter no. of elements"))`

`for s in range(0, b) :`

`s = int(input("Enter element = "))`

`a.append(s)`

`print(a)`

`n = len(a)`

`for i in range(0, b) :`

`for j in range(n-1) :`

`if a[i] < a[j] :`

`temp = a[j]`

`a[j] = a[i]`

`a[i] = temp`

~~`print("Elements after sorting")`~~

Practical - 03

Aim : Implementation of Bubble Sort programme of given list.

Theory :

Bubble sort

In this sort, we sort the given element in ascending or descending order by comparing two adjacent elements at a time.

Algorithm :-

Step 1 : Bubble sort algorithm starts by comparing the first two elements of an array and swapping if necessary.

Step 2 : If we want to sort the elements of array in ascending order then first element is greater than second, then we need to swap element.

so - Inserted

Step 3 :- If first element is smaller than second then we do not swap the elements.

Step 4 :- Again second and third elements are compared and swapped if it necessary and this process go on until last and second last element is compared and swapped so problem of insertion sort is to break insertion sort whenever

Step 5 :- If there are n elements to be sorted then used $(n-1)$ time to get required output / result

Step 6 : Stick out the input and output

Output:-

038

Bubble sort

Enter no. of elements = 4
Enter elements : 6

[6]

Et Enter elements : 5.

[6, 5]

Enter elements : 8.

[6, 5, 8]

Elements after sorting : [2, 5, 6, 8]

Practical - 04.

Aim :- Implement quick sort to sort the given list.

Theory :- The given sort is a recursion algo based on the divide and conquer technique.

Algorithm :-

Step 1 :- Quick sort first selects a value, which is called value. first element serves our first point value since we know that first will eventually end up as last in that list.

Step 2 :- The partition process will happen next. It will find the split point and at some time move other times to appropriate side of list.

Step 3 :- Partitioning begins by locating two position numbers call them left marks and right marks at the end and beginning of remaining items in list.

Practical - 04.

Aim :- Implement quick sort to sort the given list.

Theory :- The given sort is a recursion algo based on the divide and conquer technique.

Algorithm :-

Step 1 :- Quick sort first selects a value, which is called value. first element serves our first point value since we know that first will eventually end up as last in that list.

Step 2 :- The partition process will happen next. It will find the split point and at some time move other times to appropriate side of list.

Step 3 :- Partitioning begins by locating two position numbers call them left marks and right marks at the end and beginning of remaining items in list.

Step 4 :-

Step 4 :- We began by incrementing left mark until we located a value that is greater than previous. At this point we have discovered two items that are out of place with respect to eventual split point.

Step 5 :- At the point when right mark becomes less than left mark, we stop.

Step 6 :- The pivot value will be exchanged with the content of split point.

Step 7 :- In addition of all items to left of split point are less than p.v and all items to right are greater than p.v.

~~and pivoted p.v swapped. Since left and right areas are now below median, it is required to swap both sets to change address of each item as most items~~

def quick_sort (olist):
 help (olist, 0, len (a list)-1)

def help (a list, first, last):

if first < last:

split = part (list, first, last)

help (0, list, first+1, last)

def part (a list, first, last):

pivot = a list [first]

l = first + 1

r = last

done = False

while not done:

while l <= r and list[l] <= pivot:

l = l + 1

while a list[r] >= pivot and r >= l:

r = r - 1

if r < l:

done = True

x = int (input ("Enter a number = "))

a list []

for b in range (0, x):

b = int (input ("Enter a number = "))

a list.append (b)

o = len a (list)

040

```

180
Point ("mustafa")
class stack :
    global tos
    def __init__(self) :
        self.tos = -1
        self.l = [0, 0, 0, 0, 0]
    def push(self, data) :
        n = len(self.l)
        if self.tos == n + 1 :
            print("Stack is full")
        else :
            self.tos = self.tos + 1
            self.l[self.tos] = data
    def pop(self) :
        if self.tos < 0 :
            print("Stack empty")
        else :
            k = self.l[self.tos]
            del self.l[self.tos]
    def peek(self) :
        if self.tos < 0 :
            print("Stack empty")
        else :
            p = self.l[self.tos]
            print("Top element = ", p)
S = stack()

```

would be in the form of a physical stack or a pile! The element in the stack can be added or removed at

Algorithm :-

Create a class stack with instance variable
into.

Define the init method with self argument
and initialize the initial value and then
initialize empty list.

Define method push and pop under the
class stack.

Use if statement to give the condition
that if length of given list is greater
given range.

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- 180

5) The condition to give list is given if statement of given length of list that range of more than than range of list.

6) Push methods used to insert the element but pop methods used to delete the element. In stack, value is less than in pop-method, value is less than return the stack is empty or else it will be the element.

8) Attach input and output by above algorithm.

Output:

mustafa.

s.push(10)
s.push(20)

s.pop()

[10, 20, 0, 0, 10]

s.pop()

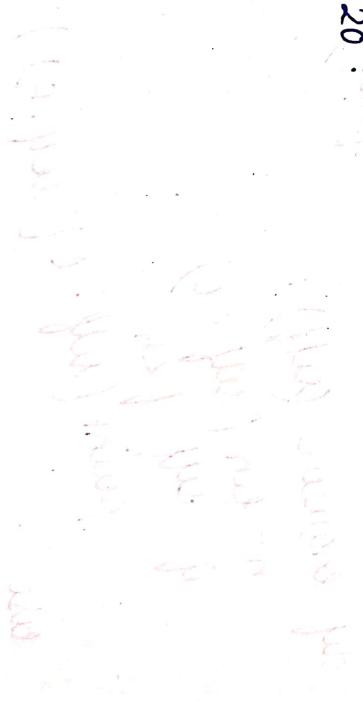
data = 20

s.pop()

[10, 0, 0, 0, 10]

s.pop()

Top element = 20.



The program is used by mustafa.
The program is used by mustafa.

```

class queue:
    global r
    global f
    def __init__(self):
        self.r = 0
        self.f = 0
        self.l = [0, 0, 0]
    def enqueue(self, data):
        n = len(self.l)
        if self.r < n:
            self.l[self.r] = data
            self.r = self.r + 1
            print("Element inserted", data)
        else:
            print("Queue is full")
    def enqueue(self):
        n = len(self.l)
        if self.r < n:
            print(self.l[self.r])
        else:
            print("Queue is empty")
q = queue()

```

Aim : Implementing a queue python with

Theory :- Queue is a data structure which has 2 front and rear implementation a queue using python list in implemented as a python list's inbuilt function to perform specified operation of a queue.

Queue() : Create a new empty queue.

Enqueue(): Insert an element at the rear of

queue and similar to that of insertion of limit using stack using

Dequeue(): Returns the element which was at

the front . The front is moved to successive element of dequeue operation always as cannot remove elements of the queue if it is empty

Algorithm :-

- 1) Define a class queue and assign global variable than define int method() with self argument.
- 2) Define an empty list and define enqueue() with 2 arguments.
- 3) we if statement that length is equal to more than queue is full insert the element in empty list or display that queue element added successfully by 1.
- 4) Define enqueue() with self argument under this we if statement that front this equal to length of list.
- 5) Now call the queue() and give the element that has to be added in empty list by using enqueue() and print the list.

Output

044

```
>>> q.add(10)
element eniut = 10
>>> q.add(20)
element eniut = 20
>>> q.add(30)
element eniut = 30
>>> q.add(40)
element eniut = 40
queue is full
>>> q.remove()
10 elements deleted.
empty queue
```

i: 10

def evaluate () :

k = s.split (' ')

n = len(k)

stack = []

for i in range (n) :

if k[i] is digit () :

stack.append (int (k[i]))

else :

if k[i] == "+" :

a = stack.pop()

b = stack.pop()

stack.append (int(b) + int(a))

elif k[i] == "-" :

a = stack.pop()

b = stack.pop()

stack.append (int(b) - int(a))

elif k[i] == "*" :

a = stack.pop()

b = stack.pop()

stack.append (int(b) * int(a))

else :

a = stack.pop()

b = stack.pop()

print (a)

Practical no. 07.

Aim :- Program of stack in python environment given string by

Theory :- The parenthesis postfix expression is free of any priorities of further use. took care of the operation in the program.

Algorithm :-

Define evaluate a function, then create an empty stack in python then call an empty

convert the string to list by using the string () 'split'.

Calculate the length of string and print.

use for loop to assign the range of string then give condition we if statement

Scanning the token list from left to right
if token is an operand convert it from string to an integer and push the value.

三

10

- g) If the token is at operation $+$, $-$, $*$, $/$, $!$ will need two operands. Pop the first operand. The second pop is first operand.

and

7) Perform the arithmetic operation. Push the result back on 'm'.

- g) Paint the string ' result of the string

- ~~g) Attach the input and output of above algorithm.~~

$s = "869 * + "$

046

$x = \text{evaluator}^{(s)}$
 $\text{print}("The evaluator value is ", r)$
 $\text{print}("AP")$

Output :-

The evaluated value is 62.

AP

```

class node:
    global data
    def __init__(self, item):
        self.data = item
        self.next = None
class linkedlist:
    global s
    def __init__(self):
        self.s = None
    def add(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
    def add(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            newnode.next = self.s
            self.s = newnode

```

Ques: Implementation of single linked list by adding the node from last position.

Algorithm:

1) Traversing of the node in linked list means visiting all operations on them in order to perform some

2) The entire linked list means can be accessed by the first node of linked list i.e. first node of the linked list in turn of referred by the head pointer of linked list.

Thus the entire linked list can be traversed by using the node which is referred by head pointer.

4) Now that we know that we can traverse the entire linked list using head pointer we should only use it to refer the first node of list only.

5) We should not use the head pointer to traverse the entire linked list because the head pointer is our reference.

6) We may lose the reference to the 1st node in our linked list and hence most of the one list so in order, we will use temporary function.

if we will use those temporary node or a copy of current node we are making temporary node a copy of current node should also be node type of temporary node.

8) But the 1st node is referred by current no we can traverse to 2nd nodes as $h = h \cdot next$.

9) Similarly we can traverse rest of node in linked list using same method by while loop.

10) Our concern now is to find terminating condition for while loop.

11) we can refer to last of linked list
say $s = none$

12) ~~Attack the coding on input & output of above algorithm~~

def

display (self):

head = self.s

while head != None:

print ("head.next! = None :")

head = head.next

print ("head.data")

delte (self):

if self.s == None:

print ("list is empty")

else:

head = self.s

while True:

if head.next == None:

d = head

head = head.next

else:

d.next = None

break.

break.

s = linked list()

s.add (10)

s.add (20)

s.add (30)

s.add (40)

s.add (50)

s.add (60)

s.add (70)

s.add (80)

s.add (90)

s.display ()

Output :

10

20

30

40

50

60

70

80

Q1

```
print ("mustafa")
set1 = set()
set2 = set()
for i in range(8, 15):
    set1.add(i)
    print("set1:", set1)
```

```
print("set2:", set2)
for i in range(11, 12):
    set2.add(i)
    print("set2:", set2)
```

```
set3 = set1 / set2
print("Intersection of set1 and set2: set4 & set4")
```

set4

1)

```
print("set1 & set2")
print("set3 & set4")
```

if set3 > set4:

print("set3 is superset of set4")

if set3 < set4:

print("set4 is subset of set3")

set5 = set3 - set4

print("set5")

set5

clear()

print("After applying clear, set5 is empty set")

Output :

Aim :- Practical - 10.

Algorithm :- Implementation of set using Python

049.

i) Define two empty set now
of above statement as set 1 and set 2

Now add () method used for addition the

element according to given range than print —

find union and intersection of above of 2 set —
using (and) , (or) () print the set .

Use if statement to find out the subset
and superset of 3 and u .

~~Display the element in set 3 is not in 4
using mathematical operation~~

Use disjoint () to check that anything is
common in present or not . It is mutually
exclusive .

Q. 10

- Q. 10. ~~Explain the difference between insertion and deletion in a set.~~
- Ans. In a set we can insert or delete the element present in set.
- 2) We clear to remove after clearing the point present in set.

Q. 11. Explain the difference between insertion and deletion in a queue.

Ans. In a queue we can insert at front and delete at rear.

In a stack we can delete at top and insert at bottom.

Q. 12. Explain the difference between insertion and deletion in a linked list.

Ans. In a linked list we can insert at any position and delete at any position.

Q. 13. Explain the difference between insertion and deletion in a binary tree.

Ans. In a binary tree we can insert at any position and delete at any position.

Q. 14. Explain the difference between insertion and deletion in a binary search tree.

Ans. In a binary search tree we can insert at any position and delete at any position.

Q. 15. Explain the difference between insertion and deletion in a binary search tree.

Ans. In a binary search tree we can insert at any position and delete at any position.

Practical - 11.

Algo :- Program implementing traversal.
by Postorder

Ex:- Postorder

Theory :-

Binary tree is a tree which supports maximum of 2 children for any node with the tree. Thus any particular node can have 0 or 1 or 2 children. There is another identity of binary tree that is another child is identified a ordered such that 2 child is identified a left child & other as right child.

1. Inorder :

Traverse the left subtree. The left subtree in turn may have left and right subtrees.

Visit the root node.

Traverse the right subtree and repeat.

Inorder :- Visit the root node.

Traverse the left subtree. The left subtree in turn may have left and right subtrees.

Postorder :- Traverse the left subtree. The left subtree in turn may have left

def: newnode .data, " added on right of ",
new = newnode .data, " added on right of ",
print (newnode)

3.0

book :

def postorder (self , start) :

if start ! = None :

 start .data

 print (start .data)

 self .postorder (start .r)

 self .postorder (start .l)

def inorder (self , start) :

if start ! = None :

 inorder (start .l)

 self .inorder (start .r)

 self .inorder (start .l)

 self .inorder (start .r)

 self .inorder (start .l)

 self .inorder (start .r)

 self .inorder (start .l)

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 self .inorder (start .l)

 self .inorder (start .r)

 self .inorder (start .l)

 self .inorder (start .r)

and μ
True
False

Algorithm

1) Define
with
meth

Ag
with

2) If
with

3)

4)

5)

6)

7)

8)

9)

10)

11)

12)

13)

14)

15)

16)

17)

18)

19)

20)

print ("postorder")

T . postorder (T . root)

and right subtree

Traverse subtree
Visit the root right subtree.

Algorithm :-

1) Define class ~~bst~~ node with 2 arguments and define init() method

2) Again, define a class BST with init method() with 2 arguments and assign root is none.

3) Define add() method for adding the node
Define a variable p that p = node

4) Use if statement for checking the condition
that root is none then use else statement
for if node then main node than arrange

5) Use if statement within else statement for
checking that node is greater than main
node then put on average

6) After this left subtree and right subtree

method to arrange the node
repeat this process to
arranging to

class

g) In 'inorder' and 'postorder' with 'if statement' that is all 'inorder', 'postorder' and 'inversion' that is all arrangement and and root is more and

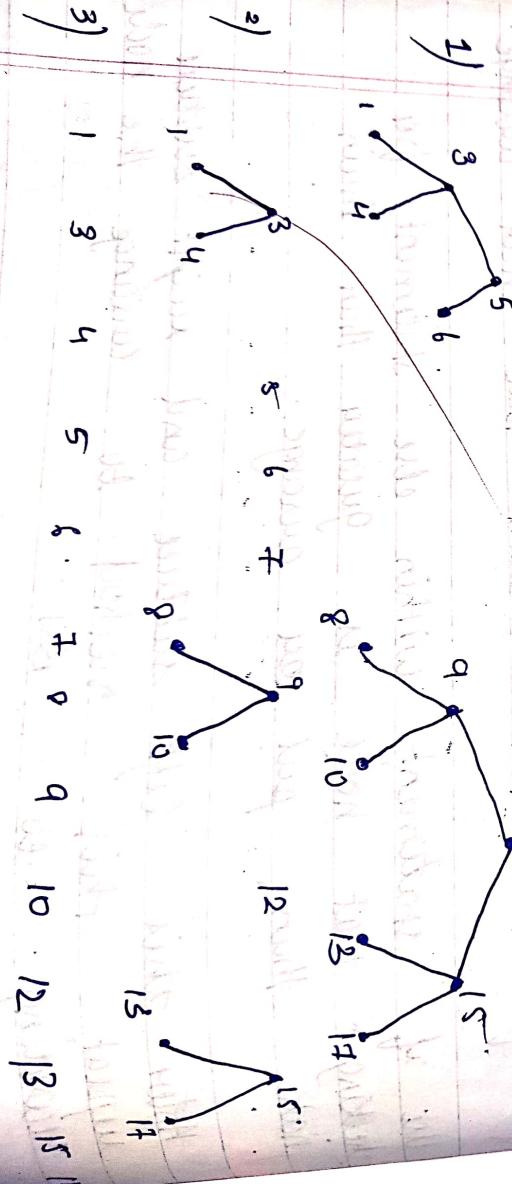
g) In 'inorder' condition 'else' statement used 'for giving that right'.

10) For 'postorder', we have to give condition in 'else' than 'first root'.

11) For 'postorder', in 'else' I assign 'left node' right and than 'go for root'.

14) Display the Q/O of above algorithm.

Procedure: (LVR)



Pseudocode: (VRV)

#

1

2



1

2

5



3

7

5

13

14

6

12

9

8

10

15

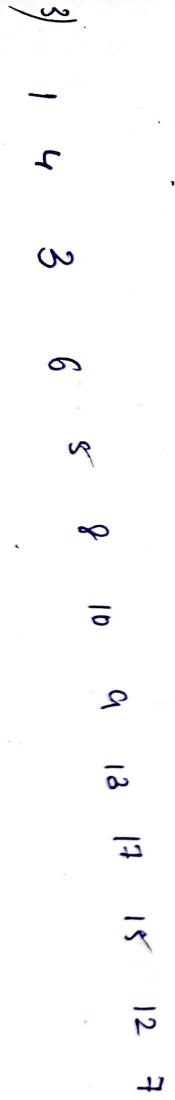
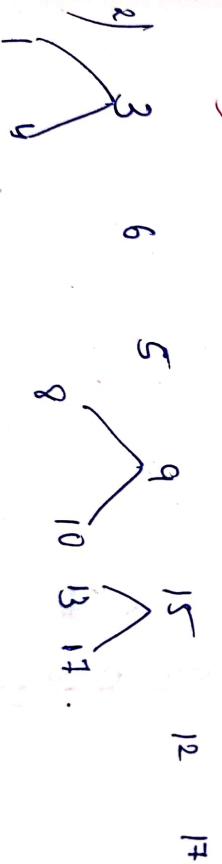
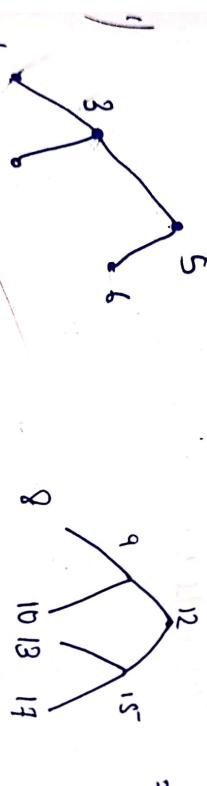
13

17

12

17

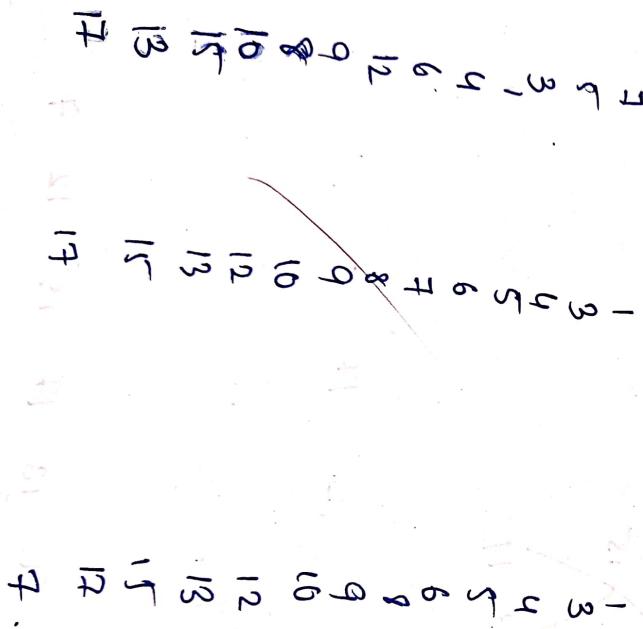
Postorder: (LRV)



Test Output:-

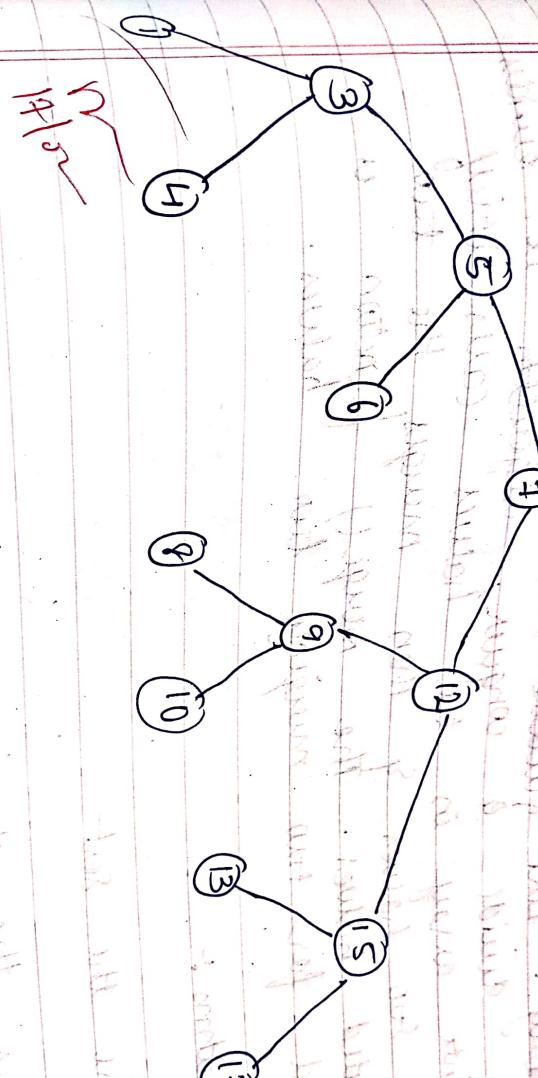
- 5 added on left of 7.
- 12 added on right of 7.
- 3 added on left of 5.
- 6 added on right of 5.
- 9 added on left of 12.
- 15 added on right of 3.
- 1 added on left of 3.
- 4 added on right of 3.
- 8 added on left of 9.
- 10 added on right of 9.
- 13 added on left of 15.
- 17 added on right of 15.

Pswonder, swonder, Postwonder.



Binary Search Tree

• Implementation :



• Time complexity

• Search :
 Need to compare current node value with target value.

• Insert :
 Need to insert new node in appropriate position.

• Delete :
 Need to handle three cases.

• Traversal :

• Height :

• Implementation :

Practical.09

20

Algo : Implementation of merge sort

Theory : Like Quicksort , merge sort is a divide & conquer algorithm . It divides array into two halves . Then merges the two sorted halves . The merge() function is used for two merging two halves .

Algorithm :-

- 1) Define the sort
- 2) Shows the starting position of both parts in temporary variables
- 3) Checks if first part comes to an end or not .
- 4) Checks if second part comes to an end or not .
- 5) Check which part has smaller elements

- 1) Now the main array has element in manner including both parts .

$$\begin{aligned}
 n_1 &= \lfloor \alpha_1, \gamma_1 m_1 \rfloor \\
 n_2 &= m - \ell_{1,2} \\
 \ell &= \lfloor 0 \rfloor * \lfloor n_2 \rfloor
 \end{aligned}$$

for i in $\lfloor n_2 \rfloor$

$$U[i] = \text{range}(0, n_2)$$

for i in $\text{range}[\ell, \ell + 1]$

$$R[j] = \text{arr}[\lfloor m_{1,1} + i \rfloor]$$

$$\begin{cases} \ell = 0 \\ k = \ell \end{cases}$$

while

$$i < n_1 \quad \& \quad i < n_2$$

$$if \ U[i] < R[i]$$

$$\text{arr}[k] = U[i]$$

$$i + = 1$$

else :

$$\text{arr}[k] = R[i]$$

$$j + = 1$$

$$k + = 1$$

while $k < n_1$:

$$\text{arr}[k] = U[i]$$

$$i + = 1$$

$$k + = 1$$

while $j < n_2$:

$$\text{arr}[k] = R[j]$$

$$j + = 1$$

$$k + = 1$$

and $\text{arr}[\lfloor m_{1,1} \rfloor]$!

10

```
if l <= m < r:  
    m = int((l + r) / 2)  
    mergsort(ans, l, m)  
    mergsort(ans, m, r)
```

```
merge sort (ans, m+1, r)  
sort (ans, l, m, r)  
print (ans)  
n = len (ans)  
mergesort (ans, 0, n-1)  
print (ans)
```

Output :-

~~[12 11 13 5 6 7]~~
[5 6 7 11 12 13]

Defined the context

057

Refined the
cured away in 2 parts.

Sent the end Paul I was of away.

卷之三

part of array.

Merge both parts by comparing elements of both parts.

03/03/2001