Rounding

→ we need to do rounding, because computers can read upto certain bits/significant figures

- To round a number, we need to draw a number line.

It The mapping from R (real no.) to F is called rounding, and it is denoted by fl(x).

Example: A computer has $\beta = 2$, m = 3. Now, suppose a number is given, $x = (0.1000100)_2$ We need to round this number.

1 Draw number line.

- We need to see where this number lies.

· find the middle point

$$(0.1001)_{2}$$

$$(0.100)_{2}$$
How did we find?
$$(0.100)_{2} = \frac{1}{2}$$

$$(0.100)_{2} = \frac{5}{8}$$

$$(0.101)_{2} = \frac{8}{16} + \frac{1}{16}$$

$$= \frac{1}{2} + \frac{1}{24}$$

$$= 2^{-1} + 2^{-4}$$

$$= (0.1001)_{1}$$

· if number lies on left side of the number line, then round the number to left number.

$$E_{X}$$
. $(0.1000100)_{2} = (0.100)_{2}$

middle point is (0.1001)2, so all other no. (0.1000....)

if number lies on right side of the number line,
then round the number to right number.

$$E_{X}$$
. $(0.100101001)_{2} = (0.101)_{2}$

. if number is exactly at middle, then it will get rounded to the nearest number.

Now, how to know the even numbers in binary? if it ends in 0, it is even

→ if it ends in 1, it is odd.

Now, in our example, x = (0.1000 100)2

X (real number) is converted / rounded to fl(x) [rounded no.]

Rounding Error

actual
$$x = 2.0 \text{cm}$$

measured $x = 1.8 \text{cm}$

Error =
$$(2-1.8)$$

= 0.2 cm.
: Error = $|fL(x) - x|$

fl(x) -> rounded value

$$x = \begin{bmatrix} 999.8 \\ cm \end{bmatrix}$$
 a cheal = 1000 cm
 $measured = 999.8 cm$ $fl(x) - x > modulus$.
 $error = 0.2 cm. = | 999.8 - 1000|$
 $= 0.2 cm$

It is difficult to understand the impact using Error only, hence, we find scale invariant rounding error, denoted by:

$$\begin{cases} 2 & \frac{|f\lambda(x) - x|}{|x|} \\ \delta \cdot x = f(x) - x & \text{[multiplication]} \end{cases}$$

$$f(x) = \delta x + x$$

$$f(x) = x(1+\delta)$$

We deal with maximum scaled invariant error called Hachine Epsilon, \in $\delta_{max} == \epsilon$

Maximum scale invariant error will be:

$$\delta = \frac{|f(x) - x|}{|x|} + \delta \propto |f(x) - x|$$

-> & will change according to conventions.

Derivation

$$f(x)$$
, $d/2$ $d = (0.101)_2 \times 2^e$
 $d = (0.001)_2 \times 2^e$
 $|f(x) - x|_{max} = \frac{1}{2}(0.001)_2 \times 2^e$
 $= \frac{1}{2} \times (1 \times 2^{-3}) \times 2^e$
 $= \frac{1}{2} \times \beta^{-m} \times \beta^e$

$$|x|_{min} = (0.100)_{2} \times 2^{e}$$

$$= (1 \times 2^{-1}) \times 2^{e}$$

$$= \beta^{-1} \times \beta^{e}$$

.. Machine Epsilon
$$(\epsilon) = \frac{|f(x) - x|}{|x|} = \frac{\frac{1}{2} \times \beta^{-m} \times \beta^{e}}{\beta^{-1} \times \beta^{e}}$$

$$= \frac{\frac{1}{2} \beta^{1-m} [Axed]}{\beta^{-1} \times \beta^{e}}$$

Questions

-> calc.
$$\in$$
 for conv. $1 = \frac{1}{2}\beta^{1-m}$

-> min. value of x for conv. $1 = \beta^{-1}\beta^{e}$

Similar derivation concept:

Convention 3 (Denormalised form) (0.1did2..dm) px pe

Similar derivation concept:

For denormalised form, we need 1 extra bit Hence,

$$|f|(x) - x|_{max} = \frac{1}{2} \beta^{m-1} \beta^{e}$$

Hence,

$$|f|(x) - x| = 1/2 B^{n} - (m+1) B^{n} = 1/2 B^{n} - m - 1 B^{n} = 1/$$

Point to note: • There is no exponent (e) in E.

- so value of machine epsilon (E)
 won't be affected due to exponent.
- . 8 ≤ En

$$\beta = 2$$
, $m = 3$, $\lambda = \frac{5}{8}$, $y = \frac{7}{8}$

1 Convert to birary

$$x = \frac{5}{8}$$

$$= \frac{4}{8} + \frac{1}{8}$$

$$= \frac{1}{2} + \frac{1}{8}$$

$$= \frac{1}{2} + \frac{1}{4}$$

$$= \frac{1}{2} + \frac{1}{2^{3}}$$

$$= 2^{-1} + 2^{-3}$$
 (there will be 1 in -1 and -3 patition)
$$= (0.101)_{2}$$

(1) Actual
$$(x * y) = (0.101)_2 * (0.111)_2$$

$$= \frac{\pi}{8} \times \frac{\pi}{8} = \frac{35}{64}$$

$$= \frac{35}{64}$$

$$= \frac{32}{64} + \frac{3}{64}$$

$$= \frac{1}{2} + \frac{1}{64} + \frac{1}{64}$$

$$= \frac{1}{2} + \frac{1}{2^5} + \frac{1}{2^6}$$

$$= 2^4 + 2^5 + 2^{-6}$$

We need to round to the value since m=3

= (0·100011)₂

$$\begin{array}{c|c}
(0 \cdot 100011)_{2} \\
\hline
(2 * y) & (0 \cdot 1001)_{2} \\
\hline
(0 \cdot 100)_{2} & (0 \cdot 101)_{2}
\end{array}$$

0.
$$100011 =$$
 round to left side
So, $(0.100011)_2 = (0.100)_2$

rounded-value:
$$f(x) * f(y) = (0.100)_2$$

= $\frac{1}{4}$

:.
$$\frac{35}{64} \neq \frac{1}{2}$$
 [rounding error]

A citual rounded.

Loss of significance

if,
$$x \neq f(x)$$
 $y \neq f(y)$
then $f(x) = x(1+\delta)$ $f(y) = y(1+\delta_2)$
now, if we want to calc. $x \pm y$
 $x \pm y = f(x) \pm f(y)$
 $= x(1+\delta_1) \pm y(1+\delta_2)$
 $= (x \pm y) \pm x\delta_1 \pm y\delta_2$
 $= (x \pm y) \left(1 + \frac{x\delta_1 \pm y\delta_2}{x \pm y}\right)$
Scale invariant error

if we want to coulc. 2-y For scale invariant error, we have

Now, if x and y are closer values (x≈y), scale invariant error becomes significantly high. If x = = y, then silve becomes infinite (∞)

> This phenomenon is called loss of significance

- -> subtract two close numbers
- -> denominator will be very small approximately zero.
- -> scale invariant error will be very high.

Example
$$x^{2} - 56x + 1 = 0$$

$$x = -b \pm \sqrt{b^{2} - 4ac}$$

$$2a$$

 $\rightarrow x_1 = 28 + \sqrt{783} = 55.98$

 $x_1 = 28 + 27.98 = 55.98$ $x_2 = 28 - 27.98 = 0.02000$ The rotequal $x_1 = 28 + 27.98 = 0.02000$

Closs of significance) [28 and 27.98 -> very close values]

0.01786 \$ 0.02000 > loss of significance

- -> denominator very small
- -> rounding error/scale invariant error very high

Solution

$$x^{2} - 56x + 1$$

 $x^{2} - (d+\beta)x + d\beta$

$$d\beta = 1$$

$$\beta = \frac{1}{\kappa}$$

$$= \frac{1}{55.98}$$

= 0.01786 (same as original 2)

1 target => avoid subtraction.

Example: Average of 5.01 and 5.02

let's say, computer can read to 3 sf.

so,
$$f_{\lambda}\left(\frac{5\cdot01+5\cdot02}{2}\right)$$

$$=\int \int \left(\frac{10.03}{2}\right)$$

rounding error