

# Module #3: **The Theory of Sets**

Rosen 5<sup>th</sup> ed., §§1.6-1.7  
~43 slides, ~2 lectures

## Introduction to Set Theory (§1.6)

- A *set* is a new type of structure, representing an *unordered* collection (group, plurality) of zero or more *distinct* (different) objects.
- Set theory deals with operations between, relations among, and statements about sets.
- Sets are ubiquitous in computer software systems.
- *All* of mathematics can be defined in terms of some form of set theory (using predicate logic).

# Naïve set theory

- Basic premise: Any collection or class of objects (*elements*) that we can describe (by any means whatsoever) constitutes a set.
- But, the resulting theory turns out to be *logically inconsistent*!
  - This means, there exist naïve set theory propositions  $p$  such that you can prove that both  $p$  and  $\neg p$  follow logically from the postulates of the theory!
  - $\therefore$  The conjunction of the postulates is a contradiction!
  - This theory is fundamentally uninteresting, because any possible statement in it can be (very trivially) “proved” by contradiction!
- More sophisticated set theories fix this problem.



## Basic notations for sets

- For sets, we'll use variables  $S, T, U, \dots$
- We can denote a set  $S$  in writing by listing all of its elements in curly braces:
  - $\{a, b, c\}$  is the set of whatever 3 objects are denoted by  $a, b, c$ .
- *Set builder notation*: For any proposition  $P(x)$  over any universe of discourse,  $\{x|P(x)\}$  is *the set of all  $x$  such that  $P(x)$* .

# Basic properties of sets

- Sets are inherently *unordered*:
  - No matter what objects  $a$ ,  $b$ , and  $c$  denote,  
 $\{a, b, c\} = \{a, c, b\} = \{b, a, c\} =$   
 $\{b, c, a\} = \{c, a, b\} = \{c, b, a\}.$
- All elements are *distinct* (unequal); multiple listings make no difference!
  - If  $a=b$ , then  $\{a, b, c\} = \{a, c\} = \{b, c\} =$   
 $\{a, a, b, a, b, c, c, c, c\}.$
  - This set contains at most 2 elements!

## Definition of Set Equality

- Two sets are declared to be equal *if and only if* they contain exactly the same elements.
- In particular, it does not matter *how the set is defined or denoted*.
- For example: The set  $\{1, 2, 3, 4\} =$   
 $\{x \mid x \text{ is an integer where } x > 0 \text{ and } x < 5\} =$   
 $\{x \mid x \text{ is a positive integer whose square}$   
 $\text{is } > 0 \text{ and } < 25\}$

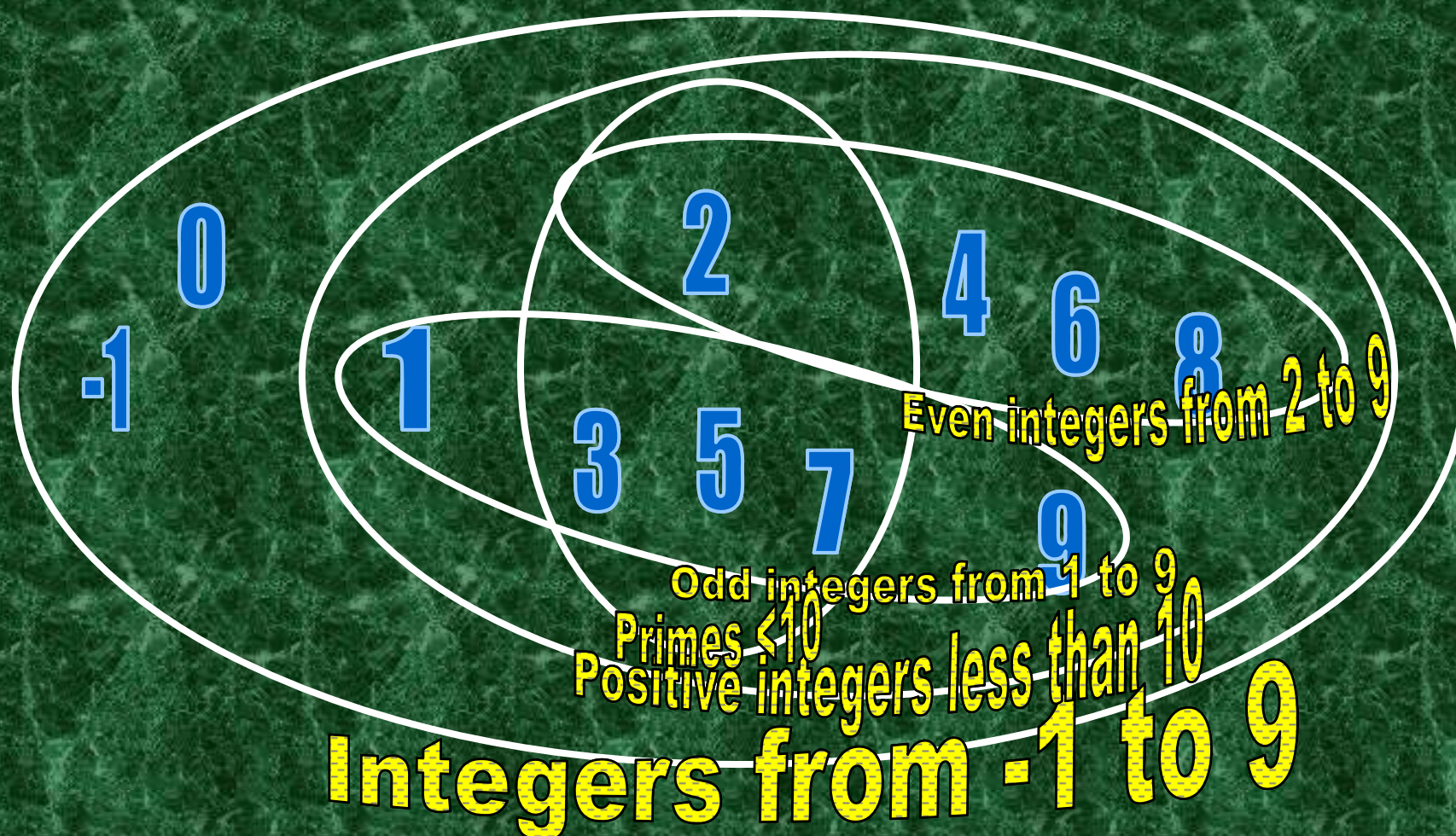
# Infinite Sets

- Conceptually, sets may be *infinite* (i.e., not *finite*, without end, unending).
- Symbols for some special infinite sets:  
 $\mathbf{N} = \{0, 1, 2, \dots\}$  The **N**atural numbers.  
 $\mathbf{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$  The **Z**ntegers.  
 $\mathbf{R}$  = The “**R**real” numbers, such as  
374.1828471929498181917281943125...
- Infinite sets come in different sizes!

More on this after module #4 (functions).



# Venn Diagrams





## Basic Set Relations: Member of

- $x \in S$  (“ $x$  is in  $S$ ”) is the proposition that object  $x$  is an *element* or *member* of set  $S$ .
  - e.g.  $3 \in \mathbf{N}$ , “a”  $\in \{x \mid x \text{ is a letter of the alphabet}\}$
  - Can define set equality in terms of  $\in$  relation:  
$$\forall S, T: S = T \leftrightarrow (\forall x: x \in S \leftrightarrow x \in T)$$
“Two sets are equal iff they have all the same members.”
- $x \notin S \equiv \neg(x \in S)$       “ $x$  is not in  $S$ ”

## The Empty Set

- $\emptyset$  (“null”, “the empty set”) is the unique set that contains no elements whatsoever.
- $\emptyset = \{ \} = \{x/\mathbf{False}\}$
- No matter the domain of discourse, we have the axiom  $\neg \exists x: x \in \emptyset$ .

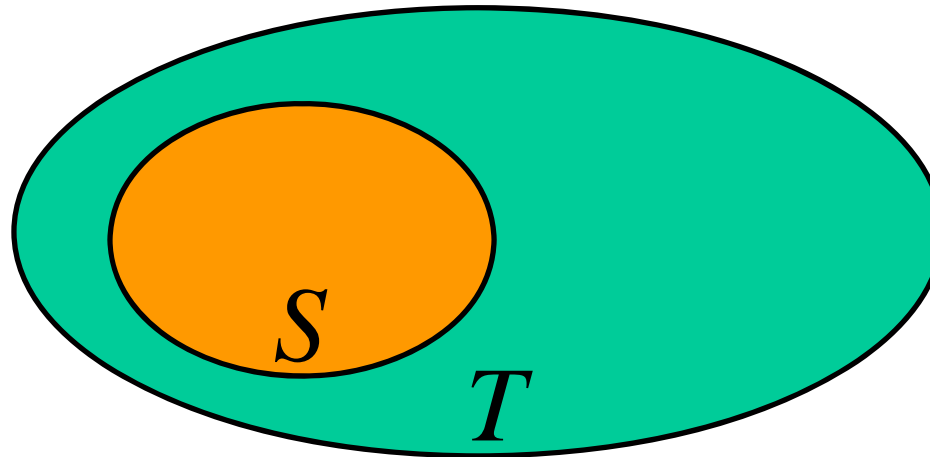
# Subset and Superset Relations

- $S \subseteq T$  (“ $S$  is a subset of  $T$ ”) means that every element of  $S$  is also an element of  $T$ .
- $S \subseteq T \Leftrightarrow \forall x (x \in S \rightarrow x \in T)$
- $\emptyset \subseteq S, S \subseteq S$ .
- $S \supseteq T$  (“ $S$  is a superset of  $T$ ”) means  $T \subseteq S$ .
- Note  $S = T \Leftrightarrow S \subseteq T \wedge S \supseteq T$ .
- $S \not\subseteq T$  means  $\neg(S \subseteq T)$ , *i.e.*  $\exists x(x \in S \wedge x \notin T)$



## Proper (Strict) Subsets & Supersets

- $S \subset T$  (“ $S$  is a proper subset of  $T$ ”) means that  $S \subseteq T$  but  $T \not\subseteq S$ . Similar for  $S \supset T$ .



Example:

$$\{1,2\} \subset \{1,2,3\}$$

Venn Diagram equivalent of  $S \subset T$

## Sets Are Objects, Too!

- The objects that are elements of a set may *themselves* be sets.
- *E.g.* let  $S = \{x \mid x \subseteq \{1,2,3\}\}$   
then  $S = \{\emptyset,$   
           $\{1\}, \{2\}, \{3\},$   
           $\{1,2\}, \{1,3\}, \{2,3\},$   
           $\{1,2,3\}\}$
- Note that  $1 \neq \{1\} \neq \{\{1\}\}$  !!!!

 **Very  
Important!**

# Cardinality and Finiteness

- $|S|$  (read “the *cardinality* of  $S$ ”) is a measure of how many different elements  $S$  has.
- *E.g.*,  $|\emptyset| = 0$ ,  $|\{1,2,3\}| = 3$ ,  $|\{a,b\}| = 2$ ,  
 $|\{\{1,2,3\},\{4,5\}\}| = \underline{2}$
- If  $|S| \in \mathbf{N}$ , then we say  $S$  is *finite*.  
Otherwise, we say  $S$  is *infinite*.
- What are some infinite sets we’ve seen?

**NZR**



## The *Power Set* Operation

- The *power set*  $P(S)$  of a set  $S$  is the set of all subsets of  $S$ .  $P(S) = \{x \mid x \subseteq S\}$ .
- *E.g.*  $P(\{a,b\}) = \{\emptyset, \{a\}, \{b\}, \{a,b\}\}$ .
- Sometimes  $P(S)$  is written  $2^S$ .  
Note that for finite  $S$ ,  $|P(S)| = 2^{|S|}$ .
- It turns out that  $|P(\mathbf{N})| > |\mathbf{N}|$ .  
*There are different sizes of infinite sets!*

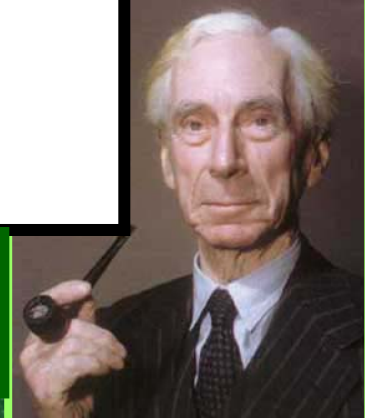
## Review: Set Notations So Far

- Variable objects  $x, y, z$ ; sets  $S, T, U$ .
- Literal set  $\{a, b, c\}$  and set-builder  $\{x|P(x)\}$ .
- $\in$  relational operator, and the empty set  $\emptyset$ .
- Set relations  $=, \subseteq, \supseteq, \subset, \supset, \not\subset$ , etc.
- Venn diagrams.
- Cardinality  $|S|$  and infinite sets  $\mathbf{N}, \mathbf{Z}, \mathbf{R}$ .
- Power sets  $P(S)$ .

# Naïve Set Theory is Inconsistent

- There are some naïve set *descriptions* that lead pathologically to structures that are not *well-defined*. (That do not have consistent properties.)
- These “sets” mathematically *cannot* exist.
- *E.g.* let  $S = \{x \mid x \notin x\}$ . Is  $S \in S$ ?
- Therefore, consistent set theories must restrict the language that can be used to describe sets.
- For purposes of this class, don't worry about it!

Bertrand Russell  
1872-1970





## Ordered $n$ -tuples

- These are like sets, except that duplicates matter, and the order makes a difference.
- For  $n \in \mathbf{N}$ , an *ordered  $n$ -tuple* or a *sequence of length  $n$*  is written  $(a_1, a_2, \dots, a_n)$ . The *first* element is  $a_1$ , *etc.*
- Note  $(1, 2) \neq (2, 1) \neq (2, 1, 1)$ .
- Empty sequence, singlets, pairs, triples, quadruples, quintuples, ...,  $n$ -tuples.

# Cartesian Products of Sets

- For sets  $A, B$ , their *Cartesian product*  $A \times B \equiv \{ (a, b) \mid a \in A \wedge b \in B \}$ .
- *E.g.*  $\{a, b\} \times \{1, 2\} = \{(a, 1), (a, 2), (b, 1), (b, 2)\}$
- Note that for finite  $A, B$ ,  $|A \times B| = |A| |B|$ .
- Note that the Cartesian product is *not* commutative:  $\neg \forall A, B: A \times B = B \times A$ .
- Extends to  $A_1 \times A_2 \times \dots \times A_n \dots$



René Descartes  
(1596-1650)

## Review of §1.6

- Sets  $S, T, U...$  Special sets  $\mathbf{N}, \mathbf{Z}, \mathbf{R}$ .
- Set notations  $\{a,b,...\}, \{x|P(x)\}...$
- Set relation operators  $x \in S, S \subseteq T, S \supseteq T, S = T, S \subset T, S \supset T$ . (These form propositions.)
- Finite vs. infinite sets.
- Set operations  $|S|, P(S), S \times T$ .
- Next up: §1.5: More set ops:  $\cup, \cap, -$ .



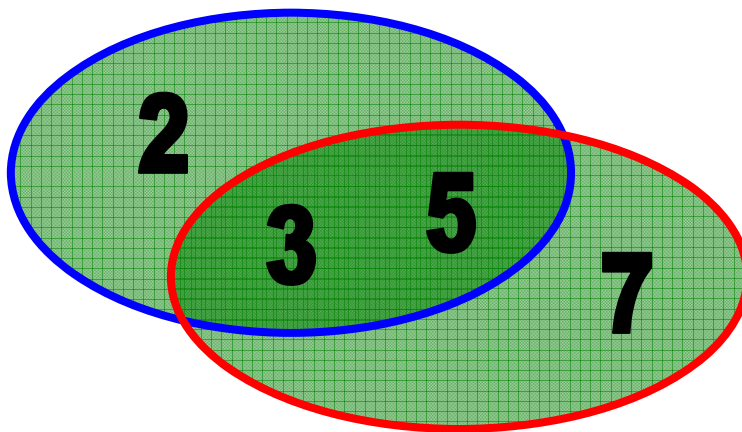
## Start §1.7: The Union Operator

- For sets  $A, B$ , their *union*  $A \cup B$  is the set containing all elements that are either in  $A$ , **or** (“ $\vee$ ”) in  $B$  (or, of course, in both).
- Formally,  $\forall A, B: A \cup B = \{x \mid x \in A \vee x \in B\}$ .
- Note that  $A \cup B$  contains all the elements of  $A$  **and** it contains all the elements of  $B$ :  
$$\forall A, B: (A \cup B \supseteq A) \wedge (A \cup B \supseteq B)$$

## Union Examples

- $\{a,b,c\} \cup \{2,3\} = \{a,b,c,2,3\}$
- $\{2,3,5\} \cup \{3,5,7\} = \{2,3,5,3,5,7\} = \{2,3,5,7\}$

Required Form



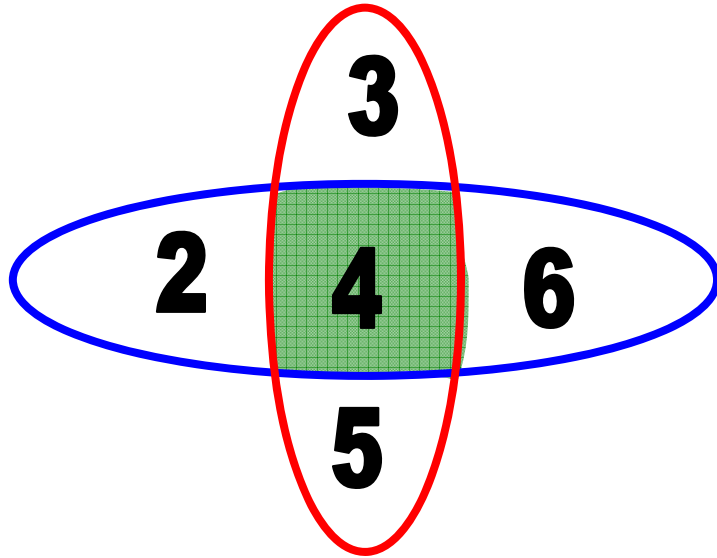
Think “The United States of America includes every person who worked in any U.S. state last year.” (This is how the IRS sees it...)

# The Intersection Operator

- For sets  $A, B$ , their *intersection*  $A \cap B$  is the set containing all elements that are simultaneously in  $A$  **and** (“ $\wedge$ ”) in  $B$ .
- Formally,  $\forall A, B: A \cap B \equiv \{x \mid x \in A \wedge x \in B\}$ .
- Note that  $A \cap B$  is a subset of  $A$  **and** it is a subset of  $B$ :  
$$\forall A, B: (A \cap B \subseteq A) \wedge (A \cap B \subseteq B)$$

# Intersection Examples

- $\{a,b,c\} \cap \{2,3\} = \underline{\emptyset}$
- $\{2,4,6\} \cap \{3,4,5\} = \underline{\{4\}}$



Think “The intersection of University Ave. and W 13th St. is just that part of the road surface that lies on *both* streets.”



# Disjointedness

- Two sets  $A, B$  are called *disjoint* (i.e., unjoined) iff their intersection is empty. ( $A \cap B = \emptyset$ )
- Example: the set of even integers is disjoint with the set of odd integers.



# Inclusion-Exclusion Principle

- How many elements are in  $A_1 \cup \dots \cup A_n$ ?
- Example: How many students in a class? Consider set  $E = I \cup M$ ,  
 $I = \{s \mid s \text{ turned in an assignment}\}$   
 $M = \{s \mid s \text{ has their email address}\}$
- Some students are in both sets.  
 $|E| = |I| + |M| - |I \cap M|$

## Set Difference

- For sets  $A, B$ , the *difference of  $A$  and  $B$* , written  $A - B$ , is the set of all elements that are in  $A$  but not  $B$ .
- $A - B \equiv \{x \mid x \in A \wedge x \notin B\}$   
 $= \{x \mid \neg(x \in A \rightarrow x \in B) \}$
- Also called:  
*The complement of  $B$  with respect to  $A$ .*

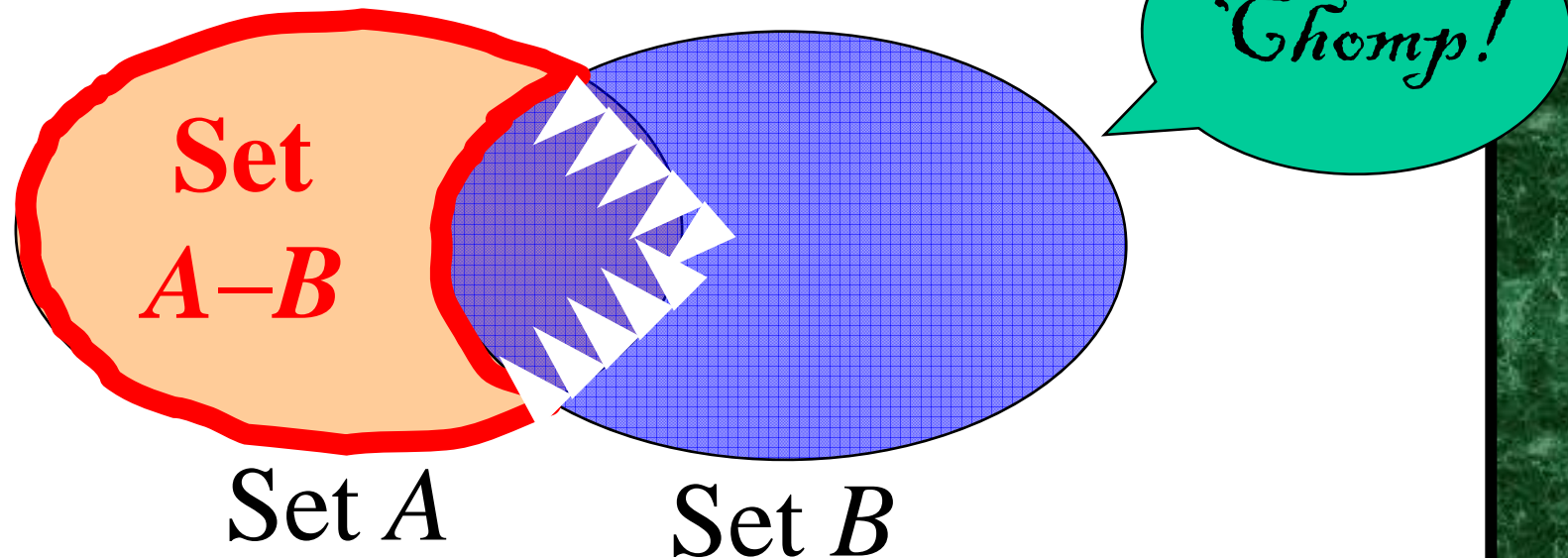
# Set Difference Examples

- $\{ \textcircled{1}, \textcircled{2}, \textcircled{3}, \textcircled{4}, \textcircled{5}, \textcircled{6} \} - \{ 2, 3, 5, 7, 9, 11 \} =$   
 $\underline{\textcircled{1}, \textcircled{4}, \textcircled{6}}$
- $\mathbf{Z} - \mathbf{N} = \{ \dots, -1, 0, 1, 2, \dots \} - \{ 0, 1, \dots \}$   
 $= \{ x \mid x \text{ is an integer but not a nat. \#} \}$   
 $= \{ x \mid x \text{ is a negative integer} \}$   
 $= \{ \dots, -3, -2, -1 \}$



# Set Difference - Venn Diagram

- $A-B$  is what's left after  $B$   
“takes a bite out of  $A$ ”



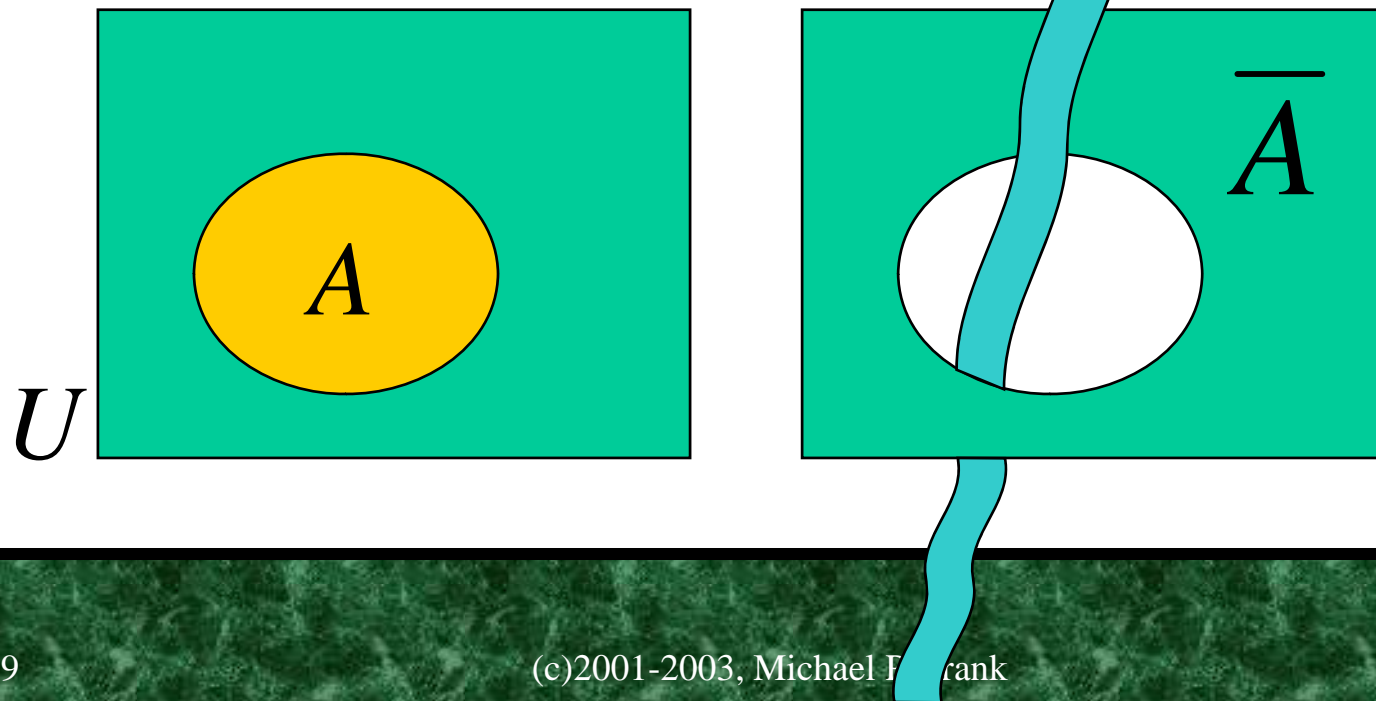
## Set Complements

- The *universe of discourse* can itself be considered a set, call it  $U$ .
- When the context clearly defines  $U$ , we say that for any set  $A \subseteq U$ , the *complement* of  $A$ , written  $\overline{A}$ , is the complement of  $A$  w.r.t.  $U$ , *i.e.*, it is  $U - A$ .
- *E.g.*, If  $U = \mathbb{N}$ ,  $\overline{\{3,5\}} = \{0,1,2,4,6,7,\dots\}$

## More on Set Complements

- An equivalent definition, when  $U$  is clear:

$$\overline{A} = \{x \mid x \notin A\}$$



## Set Identities

- Identity:  $A \cup \emptyset = A$   $A \cap U = A$
- Domination:  $A \cup U = U$   $A \cap \emptyset = \emptyset$
- Idempotent:  $A \cup A = A = A \cap A$
- Double complement:  $\overline{\overline{A}} = A$
- Commutative:  $A \cup B = B \cup A$   $A \cap B = B \cap A$
- Associative:  $A \cup (B \cup C) = (A \cup B) \cup C$   
 $A \cap (B \cap C) = (A \cap B) \cap C$



## DeMorgan's Law for Sets

- Exactly analogous to (and derivable from) DeMorgan's Law for propositions.

$$\overline{A \cup B} = \bar{A} \cap \bar{B}$$

$$\overline{A \cap B} = \bar{A} \cup \bar{B}$$

## Proving Set Identities

To prove statements about sets, of the form  $E_1 = E_2$  (where  $E$ s are set expressions), here are three useful techniques:

- Prove  $E_1 \subseteq E_2$  and  $E_2 \subseteq E_1$  separately.
- Use set builder notation & logical equivalences.
- Use a *membership table*.

## Method 1: Mutual subsets

Example: Show  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

- Show  $A \cap (B \cup C) \subseteq (A \cap B) \cup (A \cap C)$ .
  - Assume  $x \in A \cap (B \cup C)$ , & show  $x \in (A \cap B) \cup (A \cap C)$ .
  - We know that  $x \in A$ , and either  $x \in B$  or  $x \in C$ .
    - Case 1:  $x \in B$ . Then  $x \in A \cap B$ , so  $x \in (A \cap B) \cup (A \cap C)$ .
    - Case 2:  $x \in C$ . Then  $x \in A \cap C$ , so  $x \in (A \cap B) \cup (A \cap C)$ .
  - Therefore,  $x \in (A \cap B) \cup (A \cap C)$ .
  - Therefore,  $A \cap (B \cup C) \subseteq (A \cap B) \cup (A \cap C)$ .
- Show  $(A \cap B) \cup (A \cap C) \subseteq A \cap (B \cup C)$ . ...

## Method 3: Membership Tables

- Just like truth tables for propositional logic.
- Columns for different set expressions.
- Rows for all combinations of memberships in constituent sets.
- Use “1” to indicate membership in the derived set, “0” for non-membership.
- Prove equivalence with identical columns.



# Membership Table Example

Prove  $(A \cup B) - B = A - B$ .

$A$	$B$	$A \cup B$	$(A \cup B) - B$	$A - B$
0	0	0	0	0
0	1	1	0	0
1	0	1	1	1
1	1	1	0	0

# Membership Table Exercise

Prove  $(A \cup B) - C = (A - C) \cup (B - C)$ .

$A$	$B$	$C$	$A \cup B$	$(A \cup B) - C$	$A - C$	$B - C$	$(A - C) \cup (B - C)$
0	0	0					
0	0	1					
0	1	0					
0	1	1					
1	0	0					
1	0	1					
1	1	0					
1	1	1					

## Review of §1.6-1.7

- Sets  $S, T, U...$  Special sets  $\mathbf{N}, \mathbf{Z}, \mathbf{R}$ .
- Set notations  $\{a,b,...\}, \{x|P(x)\}...$
- Relations  $x \in S, S \subseteq T, S \supseteq T, S = T, S \subset T, S \supset T$ .
- Operations  $|S|, P(S), \times, \cup, \cap, -, \bar{S}$
- Set equality proof techniques:
  - Mutual subsets.
  - Derivation using logical equivalences.

## Generalized Unions & Intersections

- Since union & intersection are commutative and associative, we can extend them from operating on *ordered pairs* of sets  $(A,B)$  to operating on sequences of sets  $(A_1, \dots, A_n)$ , or even unordered *sets* of sets,  $X = \{A / P(A)\}$ .



# Generalized Union

- Binary union operator:  $A \cup B$
- $n$ -ary union:  
 $A \cup A_2 \cup \dots \cup A_n \equiv ((\dots((A_1 \cup A_2) \cup \dots) \cup A_n)$   
(grouping & order is irrelevant)
- “Big U” notation:  $\bigcup_{i=1}^n A_i$
- Or for infinite sets of sets:  $\bigcup_{A \in X} A$

# Generalized Intersection

- Binary intersection operator:  $A \cap B$
- $n$ -ary intersection:  
 $A \cap A_2 \cap \dots \cap A_n \equiv ((\dots((A_1 \cap A_2) \cap \dots) \cap A_n)$   
(grouping & order is irrelevant)
- “Big Arch” notation:  $\bigcap_{i=1}^n A_i$
- Or for infinite sets of sets:  $\bigcap_{A \in X} A$

# Representations

- A frequent theme of this course will be methods of *representing* one discrete structure using another discrete structure of a different type.
- *E.g.*, one can represent natural numbers as
  - Sets:  $0 \equiv \emptyset$ ,  $1 \equiv \{0\}$ ,  $2 \equiv \{0, 1\}$ ,  $3 \equiv \{0, 1, 2\}$ , ...
  - Bit strings:  
 $0 \equiv 0$ ,  $1 \equiv 1$ ,  $2 \equiv 10$ ,  $3 \equiv 11$ ,  $4 \equiv 100$ , ...

## Representing Sets with Bit Strings

For an enumerable u.d.  $U$  with ordering  $\{x_1, x_2, \dots\}$ , represent a finite set  $S \subseteq U$  as the finite bit string  $B = b_1 b_2 \dots b_n$  where  $\forall i: x_i \in S \leftrightarrow (i < n \wedge b_i = 1)$ .

E.g.  $U = \mathbb{N}$ ,  $S = \{2, 3, 5, 7, 11\}$ ,  $B = 001101010001$ .

In this representation, the set operators “ $\cup$ ”, “ $\cap$ ”, “ $\bar{\phantom{x}}$ ” are implemented directly by bitwise OR, AND, NOT!