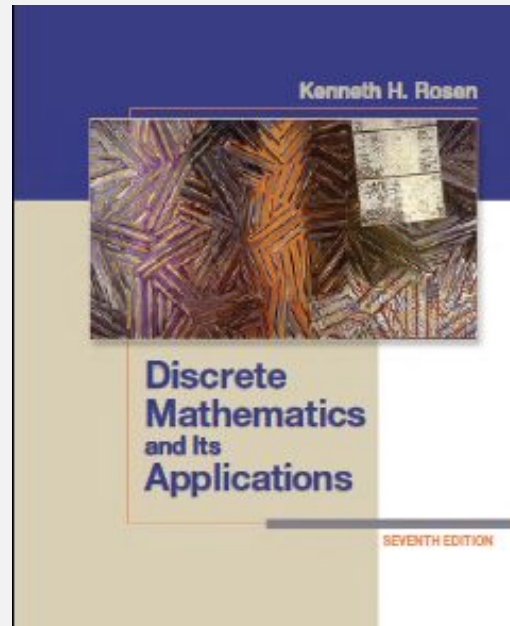


# CSE 105: Discrete Mathematics



# Course Evaluation

Topic	Marks
Attendance	05
Individual Presentations	05
Group Assignments	05
Quizzes (3/4)	15
Mid term	30
Final	40

# Course Materials

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## **Text Book:**

- “*Discrete Mathematics and Its Application*”, Kenneth H. Rosen, 7<sup>th</sup> Edition, McGraw-Hill.
- Lecture notes

## **Reference Materials:**

- “*Schaum’s Outlines Discrete Mathematics*”, Seymore Lipschutz & Marc Lipson, 3<sup>rd</sup> Edition, McGraw-Hill.
- [http://en.wikiversity.org/wiki/Introductory\\_Discrete\\_Mathematics\\_for\\_Computer\\_Science](http://en.wikiversity.org/wiki/Introductory_Discrete_Mathematics_for_Computer_Science)

# Course Contents

No	Topic	Exams/Quiz
Lectures 1-2	Introduction to Discrete Mathematics+ Propositional Logic	
Lectures 3-4	Propositional Logic and Introduction to Set	
Lectures 5-6	Quiz-1 + Set	Quiz-1
Lectures 7-8	Function	
Lectures 9-10	Quiz-2 + Algorithm	Quiz-2
Lectures 11-12	Midterm Exam	
Lectures 13-14	Induction + Discrete Probability	GA-1
Lectures 15-16	Counting	Presentations
Lectures 17-18	Quiz-3+Relation	Quiz-3
Lectures 19-20	Number Theory	GA-2
Lectures 21-22	Quiz-4+Graph-Tree	Quiz-4
Lectures 23-24	Graph-Tree	Presentations
	Final Exam	

# Discrete mathematics

## Discrete mathematics

- study of **mathematical structures and objects** that are fundamentally **discrete rather than continuous**.
- **Examples of objects with discrete values are**
  - **integers, graphs, or statements in logic.**
- **Discrete mathematics and computer science.**
  - Concepts from discrete mathematics are useful for describing **objects and problems in computer algorithms and programming languages**. **These** have applications in cryptography, automated theorem proving, and software development.

# Logic

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- Logic defines a **formal language** for representing **knowledge** and for making logical inference
- It helps us to understand **how to construct a valid argument**

## **Logic defines:**

- Syntax of statements
- The meaning of statements
- The rules of logical inference (manipulation)

# Propositional logic

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## The simplest logic

- **Definition:**

- A proposition is a statement that is either true or false.

- **Examples:**

- Airport is located in the North Part of Dhaka City.

- **(T)**

- $5 + 2 = 8$ .

- **(F)**

- It is raining today.

- **(either T or F)**

# Propositional logic

## Examples (cont.):

– How are you?

- a question is not a proposition

–  $x + 5 = 3$

- since  $x$  is not specified, neither true nor false

– 2 is a prime number.

- (T)

– She is very talented.

- since she is not specified, neither true nor false

– There are other life forms on other planets in the universe.

- either T or F



# Composite Statements

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- More complex propositional statements can be built from elementary statements using **logical connectives**.

## Example:

- Proposition A: **It rains outside**
- Proposition B: **We will see a movie**
- A new (combined) proposition:  
**If it rains outside then we will see a movie**

# Composite Statements

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- More complex propositional statements can be built from elementary statements using **logical connectives**.
- **Logical connectives:**
  - Negation
  - Conjunction
  - Disjunction
  - Exclusive or
  - Implication
  - Biconditional

# Logical connectives: Negation

**Definition:** Let **p** be a proposition. The statement "**It is not the case that p.**" is another proposition, called the **negation of p**. The **negation of p** is denoted by  $\neg p$  and read as "**not p.**"

## Example:

- Airport is located in the North Part of Dhaka City.



- **It is not the case that** Airport is located in the North Part of Dhaka City.

## Other examples:

- $5 + 2 \neq 8$ .
- 10 is **not** a prime number.
- It is **not the case that** buses stop running at 9:00pm.

# Logical connectives: Negation

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**Negate the following propositions:**

- It is raining today.
  - It is **not** raining today.
- 2 is a prime number.
  - 2 is **not** a prime number
- There are other life forms on other planets in the universe.
  - **It is not the case that** there are other life forms on other planets in the universe.

# Logical connectives: Negation

A **truth table** displays the **relationships between truth values** (T or F) of different propositions.

P	$\neg p$
T	F
F	T

Rows: all possible values  
of elementary proposition

# Logical connectives: Conjunction

**Definition:** Let  $p$  and  $q$  be propositions. The proposition “ $p$  and  $q$ ” denoted by  $p \wedge q$ , is true when **both  $p$  and  $q$  are true** and is false otherwise. The proposition  $p \wedge q$  is called the **conjunction of  $p$  and  $q$** .

- **Examples:**

- It is raining today **and** 2 is a prime number.
- 2 is a prime number **and**  $5 + 2 \neq 8$ .
- 13 is a perfect square **and** 9 is prime.

# Logical connectives: Disjunction

**Definition:** Let  $p$  and  $q$  be propositions. The proposition “ $p$  or  $q$ ” denoted by  $p \vee q$ , is false when both  $p$  and  $q$  are false and is true otherwise. The proposition  $p \vee q$  is called the **disjunction of  $p$  or  $q$** .

- **Examples:**

- It is raining today or 2 is a prime number.
- 2 is a prime number or  $5 + 2 \neq 8$ .
- 13 is a perfect square or 9 is a prime.

# Truth Tables

- **Conjunction and disjunction**
  - Four different combinations of values for p and q

p	q	$p \wedge q$	$p \vee q$
F	F		
F	T		
T	F		
T	T		

Rows: all possible combinations of values for elementary propositions:  $2^n$  values



# Truth Tables

- **Conjunction and disjunction**
  - Four different combinations of values for p and q

p	q	$p \wedge q$	$p \vee q$
F	F	F	
F	T	F	
T	F	F	
T	T	T	

# Truth Tables

- **Conjunction and disjunction**
  - Four different combinations of values for p and q

p	q	$p \wedge q$	$p \vee q$
F	F	F	F
F	T	F	T
T	F	F	T
T	T	T	T

NB:  $p \vee q$  (the **or** is used **inclusively**, i.e.,  $p \vee q$  is true when either p or q or both are true).

# Exclusive or

**Definition:** Let  $p$  and  $q$  be propositions. The proposition " $p$  exclusive or  $q$ " denoted by  $p \oplus q$ , is true when exactly one of  $p$  and  $q$  is true and it is false otherwise.

$p$	$q$	$p \oplus q$
F	F	F
F	T	T
T	F	T
T	T	F

# Implications

**Definition:** Let  $p$  and  $q$  be propositions. The proposition " $p$  implies  $q$ " denoted by  $p \rightarrow q$  is called implication. It is **false** when  $p$  is true and  $q$  is false and is true otherwise.

In  $p \rightarrow q$ ,  $p$  is called the **hypothesis** and  $q$  is called the **conclusion**.

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	<b>F</b>
T	T	T

# Implications

$p \rightarrow q$  is read in a variety of equivalent ways:

- if  $p$  then  $q$
- $p$  only if  $q$
- $p$  is sufficient for  $q$
- $q$  whenever  $p$

## Examples:

– if Germany won 2010 world cup then 2 is a prime.

- If F then T ? **T**

– if today is Monday then  $2 * 3 = 8$ .

- $T \rightarrow F$ ? **F**

# Implications

The **converse** of  $p \rightarrow q$  is  $q \rightarrow p$ .

- The **contrapositive** of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$
- The **inverse** of  $p \rightarrow q$  is  $\neg p \rightarrow \neg q$

## Examples:

- If it jams, the traffic moves slowly.
- $p$ : it jams
- $q$ : traffic moves slowly.
- $p \rightarrow q$

# Implications

## The converse:

If the traffic moves slowly then it jams.

- $q \rightarrow p$

## The contrapositive:

• If the traffic does not move slowly then it does not jams.

- $\neg q \rightarrow \neg p$

## The inverse:

• If it does not jams the traffic moves quickly.

- $\neg p \rightarrow \neg q$

# Biconditional

Definition: Let  $p$  and  $q$  be propositions. The biconditional  $p \leftrightarrow q$  (read  $p$  if and only if  $q$ ), is true when  $p$  and  $q$  have the same truth values and is false otherwise.

$p$	$q$	$p \leftrightarrow q$
F	F	T
F	T	F
T	F	F
T	T	T

Note: two truth values always agree.



# Constructing the truth table

**Example: Construct a truth table for**  $(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$

- Simpler if we decompose the sentence to elementary and intermediate propositions

p	q	$\neg p$	$p \rightarrow q$	$\neg p \leftrightarrow q$	$(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$
F	F				
F	T				
T	F				
T	T				

# Constructing the truth table

**Example: Construct a truth table for**  $(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$

- Simpler if we decompose the sentence to elementary and intermediate propositions

p	q	$\neg p$	$p \rightarrow q$	$\neg p \leftrightarrow q$	$(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$
F	F	T			
F	T	T			
T	F	F			
T	T	F			

# Constructing the truth table

**Example: Construct a truth table for**  $(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$

- Simpler if we decompose the sentence to elementary and intermediate propositions

p	q	$\neg p$	$p \rightarrow q$	$\neg p \leftrightarrow q$	$(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$
F	F	T	T		
F	T	T	T		
T	F	F	F		
T	T	F	T		

# Constructing the truth table

**Example: Construct a truth table for**  $(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$

- Simpler if we decompose the sentence to elementary and intermediate propositions

p	q	$\neg p$	$p \rightarrow q$	$\neg p \leftrightarrow q$	$(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$
F	F	T	T	F	
F	T	T	T	T	
T	F	F	F	T	
T	T	F	T	F	

# Constructing the truth table

**Example: Construct a truth table for**  $(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$

- Simpler if we decompose the sentence to elementary and intermediate propositions

p	q	$\neg p$	$p \rightarrow q$	$\neg p \leftrightarrow q$	$(p \rightarrow q) \wedge (\neg p \leftrightarrow q)$
F	F	T	T	F	F
F	T	T	T	T	T
T	F	F	F	T	F
T	T	F	T	F	F

# Computer Representation of True and False

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We need to encode two values True and False:

- Computers represents data and programs using 0s and 1s
- Logical truth values – True and False
- A bit is sufficient to represent two possible values:
  - 0 (False) or 1(True)
- A variable that takes on values 0 or 1 is called a Boolean variable.
- Definition:A bit string is a sequence of zero or more bits. The length of this string is the number of bits in the string.

# Bitwise operation

- T and F replaced with 1 and 0

p	q	$p \wedge q$	$p \vee q$
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1

p	$\neg p$
0	1
1	0

# Bitwise operation

- T and F replaced with 1 and 0

1 0 1 1 0 0 1 1	1 0 1 1 0 0 1 1	1 0 1 1 0 0 1 1
∨ 0 1 0 0 1 0 0 1	∧ 0 1 0 0 1 0 0 1	⊕ 0 1 0 0 1 0 0 1
1 1 1 1 1 0 1 1	0 0 0 0 0 0 0 1	1 1 1 1 1 0 1 0



# Applications of propositional logic

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- **Translation of English sentences**
- **Inference and reasoning:**
  - new true propositions are inferred from existing ones
  - Used in **Artificial Intelligence**:
    - Builds programs that act intelligently
    - Programs often rely on symbolic manipulations
- **Design of logic circuit**

# Translation

Assume a sentence:

If you are older than 13 or you are with your parents then you can attend a PG-13 movie.

Parse:

- If ( you are older than 13 or you are with your parents ) then ( you can attend a PG-13 movie )

Atomic (elementary) propositions:

- A= you are older than 13
- B= you are with your parents
- C=you can attend a PG-13 movie

- Translation:  $A \vee B \rightarrow C$

# Application of Inference

Assume we know the following sentences are true:

If you are older than 13 or you are with your parents then you can attend a PG-13 movie. You are older than 13.

Translation:

- If ( you are older than 13 or you are with your parents ) then (you can attend a PG-13 movie) . (You are older than 13).
- A= you are older than 13
- B= you are with your parents
- C=you can attend a PG-13 movie
- $(A \vee B \rightarrow C)$ , A
- $(A \vee B \rightarrow C) \wedge A$  is true
- With the help of the logic we can infer the following statement (proposition):
- You can attend a PG-13 movie or C is True

# Tautology and Contradiction

Some propositions are interesting since their values in the truth table are always the same

## Definitions:

- A **compound proposition that is always true** for all possible truth values of the propositions is called a **tautology**.
- A **compound proposition that is always false** is called a **contradiction**.
- A proposition that is **neither a tautology nor contradiction** is called a **contingency**.

**Example:**  $p \vee \neg p$  is a **tautology**.  
 $p \wedge \neg p$  is a **contradiction**.

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
1	0	1	0
0	1	1	0

# Equivalence

How do we determine that **two propositions** are equivalent?

**Their truth values in the truth table are the same.**

- Example:  $p \rightarrow q$  is equivalent to  $\neg q \rightarrow \neg p$  (contrapositive)

p	q	$p \rightarrow q$	$\neg q \rightarrow \neg p$
F	F	T	T
F	T	T	T
T	F	F	F
T	T	T	T

**Equivalent statements are important for logical reasoning** since they can be substituted and can help us to make a logical argument.

# Logical Equivalence

**Definition:** The propositions  $p$  and  $q$  are called logically equivalent if  $p \leftrightarrow q$  is a tautology (alternately, if they have the same truth table). The notation  $p \Leftrightarrow q$  denotes  $p$  and  $q$  are logically equivalent.

$p$	$q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$(p \rightarrow q) \Leftrightarrow (\neg q \rightarrow \neg p)$
F	F	T	T	T
F	T	T	T	T
T	F	F	F	T
T	T	T	T	T

# Important Logical Equivalence

- **Double negation**

- $\neg(\neg p) \iff p$

- **Commutative**

- $p \vee q \iff q \vee p$

- $p \wedge q \iff q \wedge p$

- **Associative**

- $(p \vee q) \vee r \iff p \vee (q \vee r)$

- $(p \wedge q) \wedge r \iff p \wedge (q \wedge r)$

# Important Logical Equivalence

- **Distributive**

- $p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$
- $p \wedge (q \vee r) \Leftrightarrow (p \wedge q) \vee (p \wedge r)$

- **De Morgan**

- $\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$
- $\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$

- **Other useful equivalences**

- $p \vee \neg p \Leftrightarrow T$
- $p \wedge \neg p \Leftrightarrow F$
- $p \rightarrow q \Leftrightarrow (\neg p \vee q)$



# Showing Logical Equivalence

Show  $(p \wedge q) \rightarrow p$  is a tautology

In other words  $((p \wedge q) \rightarrow p \Leftrightarrow T)$

$(p \wedge q) \rightarrow p$	$\Leftrightarrow \neg(p \wedge q) \vee p$	Useful
	$\Leftrightarrow [\neg p \vee \neg q] \vee p$	DeMorgan
	$\Leftrightarrow [\neg q \vee \neg p] \vee p$	Commutative
	$\Leftrightarrow \neg q \vee [\neg p \vee p]$	Associative
	$\Leftrightarrow \neg q \vee [T]$	Useful
	$\Leftrightarrow T$	Domination

You can also use truth table to show this

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Thank You