

Università degli Studi Roma Tre Dipartimento di Informatica e Automazione Computer Networks Research Group

Interdomain routing

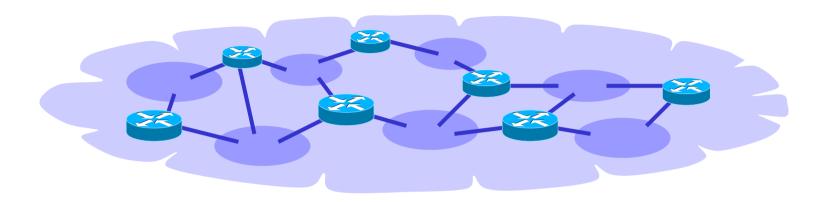
a brief introduction

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Description	the essentials of bgp

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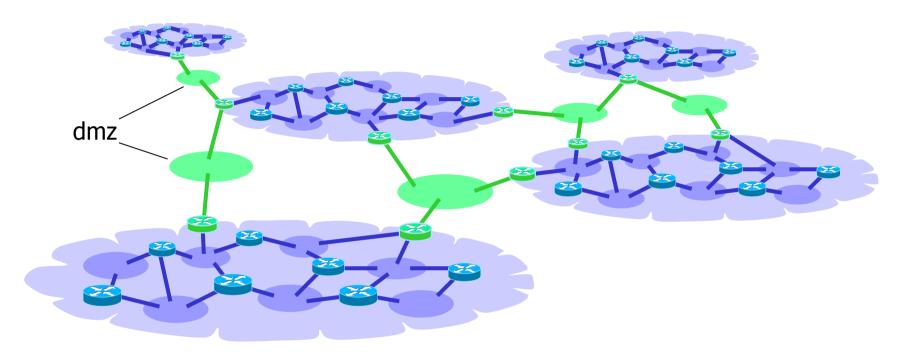
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why interdomain routing?



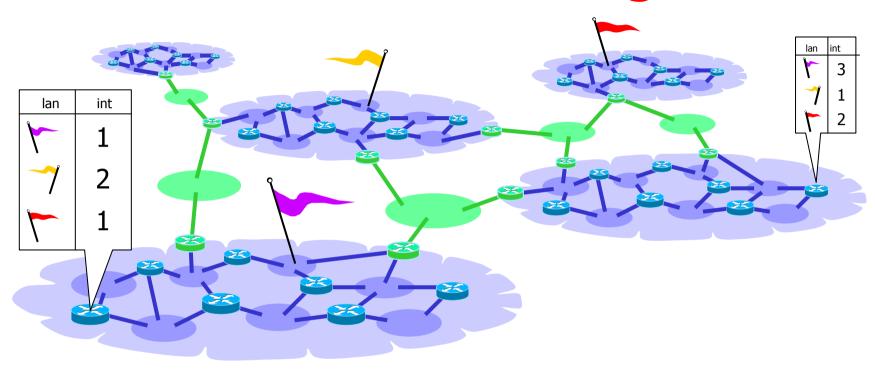
- each organization is a collection of routers and lans under a single administration
- a routing algorithm (rip, is-is, ospf,..) may be chosen to automatically update the routing tables of the routers of the administration

why interdomain routing?



- when several organizations join to form the internet they have to set up links between them
 - the added lans are called "demarcation zones"

what about the routing tables?

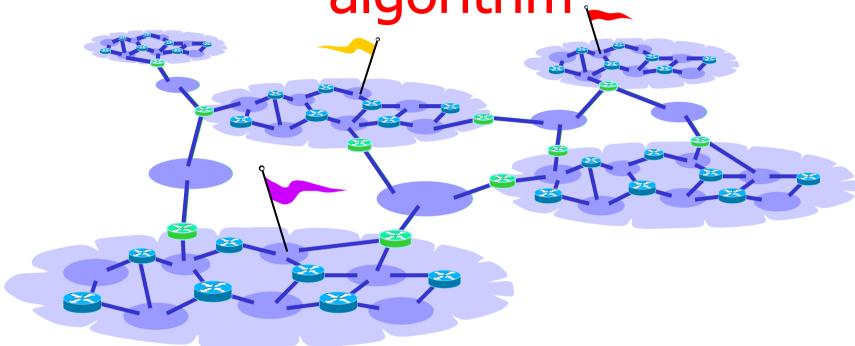


- in order to have global connectivity:
 - each router must have a routing entry (possibly the default one) that matches the destination address of the packet
 - this should be true for packets to be delivered locally as well as for packets to be delivered to remote lans

how to update the routing tables?

- in principle you have three options
 - 1 run a single routing algorithm along with adjacent organizations
 - 2 update the routing tables by hand, adding static routes to external lans
 - 3 combine an exterior gateway protocol with the interior gateway protocol of the networks

1 - run a single routing algorithm

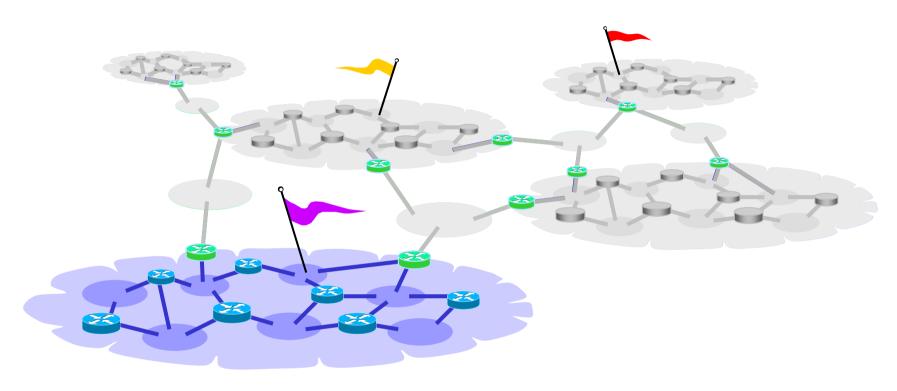


it was this way when egp (exterior gateway protocol) was introduced

single routing algorithm: problems

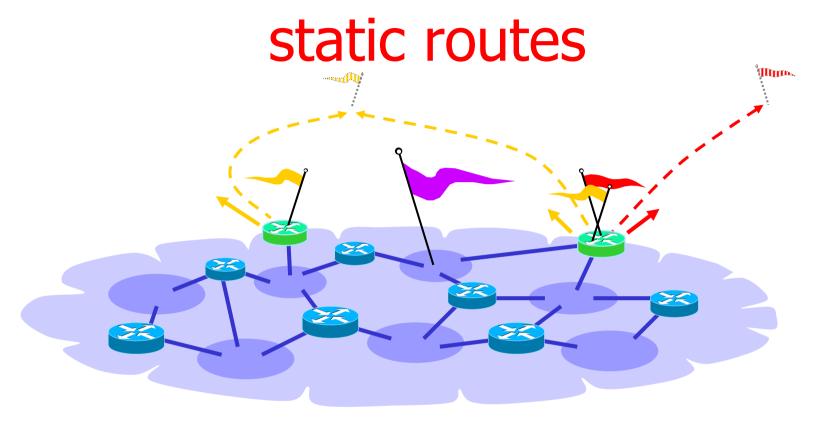
- technical
 - slow to converge
 - all the organizations are forced to use the same routing algorithm
 - difficult to deploy a new routing algorithm
 - difficult to debug and to configure
- political
 - routing is aimed at minimizing the global usage of network resources
 - does not take into account the ownership of the links

2 - use static routes



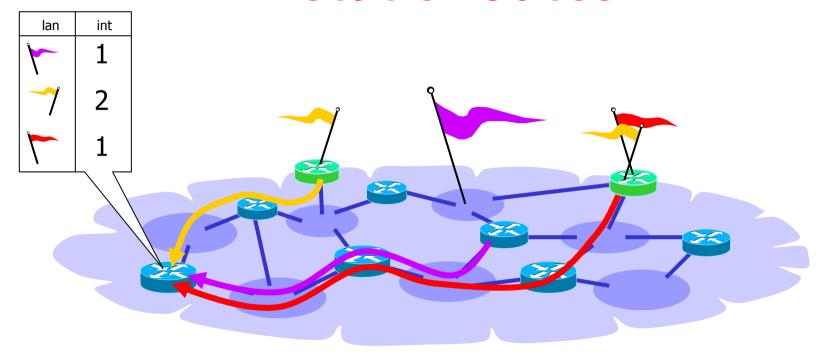
approach:

- hide the exterior part of your organization
 - hide demarcation zones
 - hide other organizations



- for each external target:
 - add a static route to some router on the border
 - you are stating that you can reach such target by sending packets to the specified next hop
 - you may use the default route

static routes

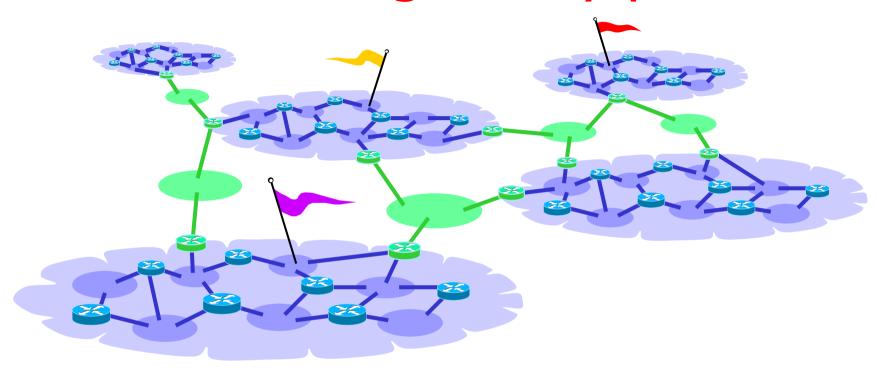


the igp routing algorithm will spread into the network the local targets as well as the statically added ones

static routes: problems

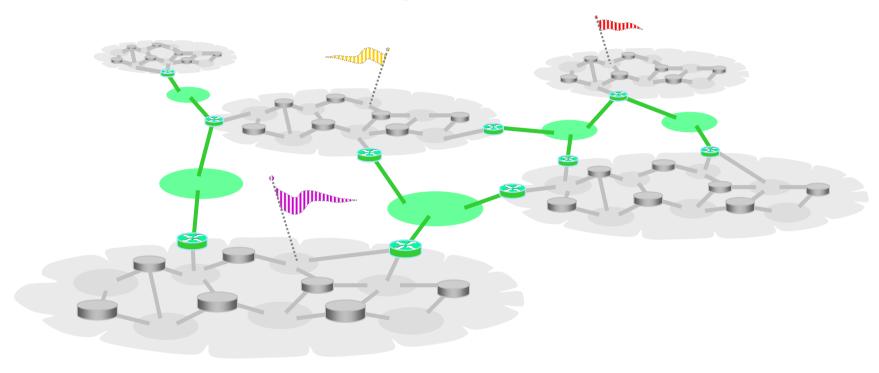
- technical
 - difficult to update and debug
 - faults are not handled
 - a static route is available even when the link is down
- political
 - no guarantee that next hop is willing to deliver the packets

3 - exterior gateway protocol



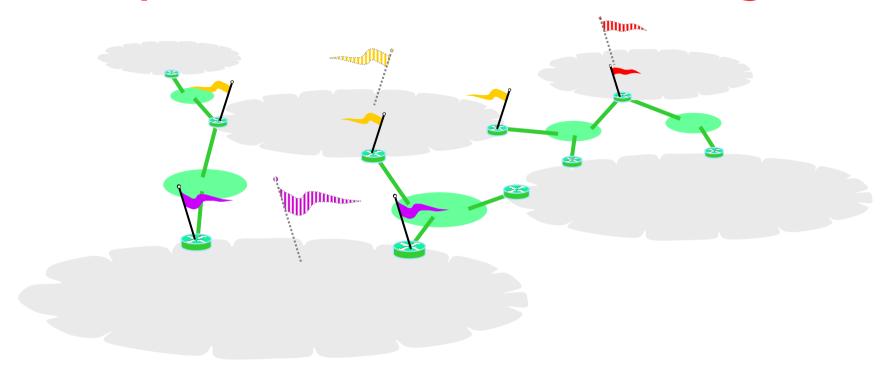
approach:

3 - exterior gateway protocol



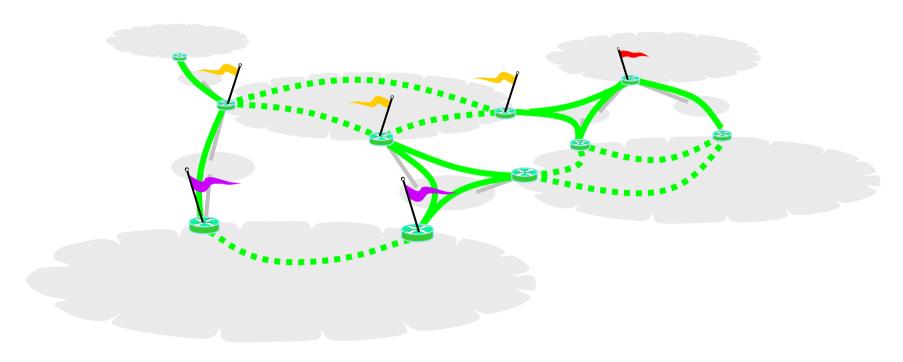
- approach:
 - hide the interior part of all organizations

represent the internal targets



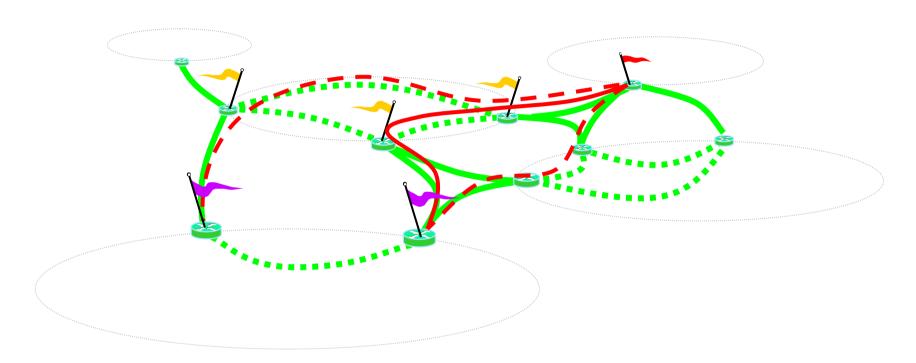
- each external router represents its internal targets as if they were local
 - you may use the default route

simplify the graph



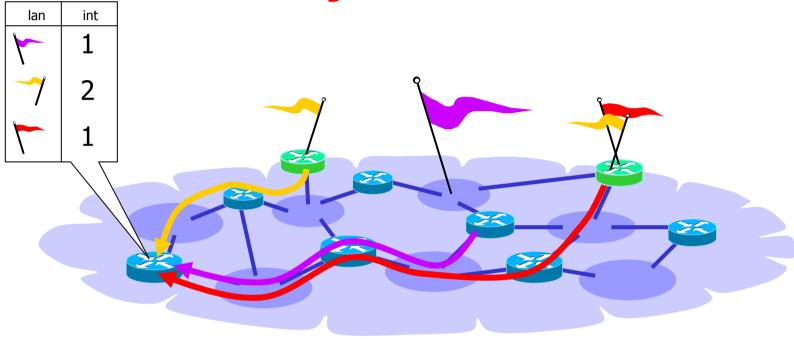
- simplify the graph
 - consider the external reachability of the routers
 - consider the internal reachability of the routers
 - the graph is actually managed through tcp connections called *peerings*

solve the routing problem



- solve the routing problem on the simplified graph
 - based on political considerations

inject routes



- analogous to the process used for static routes
 - this time no assumption is made on external routing!

border gateway protocol

- bgp is a routing protocol: it keeps the routing tables updated and propagates routing information
- takes into account the willingness of the organizations to cooperate in the routing process (commercial agreements, local preferences, priorities, legal issues, ...)

an exploration of the bgp world

- we will explore the bgp world step by step
- we will use practical examples to introduce new concepts
- we will begin from the edge of internet, supposing we are a new autonomous system joining the game
- we will consider the isp's perspective next

who uses bgp?

- bgp is used by:
 - customers connected to an Internet Service Provider (ISP)
 - customers connected to several ISPs
 - transit providers
 - ISPs that exchange traffic in an exchange point (NAP)
 - customers with very large networks

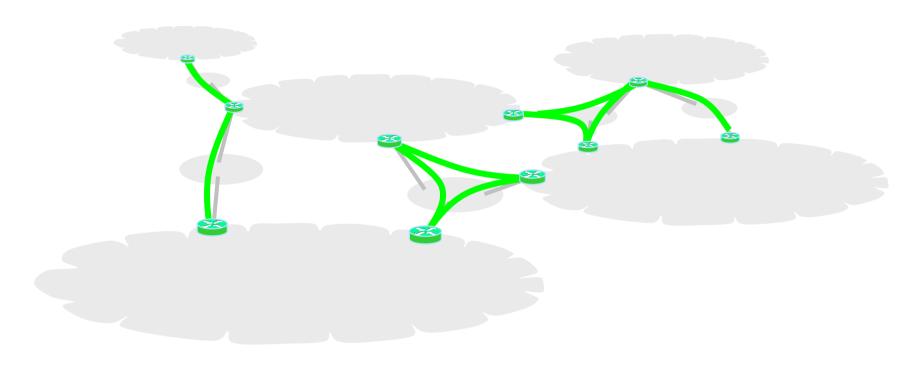
autonomous system number

- in order to use bgp you need an identification number called autonomous system number (asn)
 - between 1 and 65,535 (two bytes)
 - numbers greater than 64,511 are "private"
- as numbers may be requested:
 - global asn from your regional internet registry (rir): ripe, arin, apnic
 - private asn from your upstream isp

there are 2 bgp's

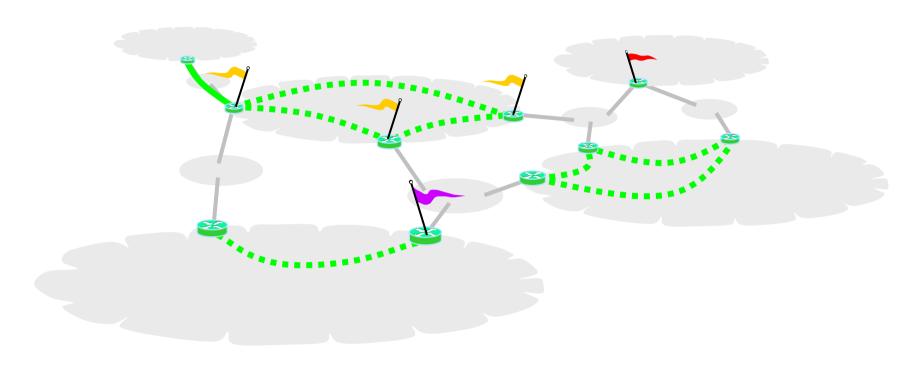
- e-bgp
 - external bgp
 - single-hop
- i-bgp
 - internal bgp
 - multi-hop
 - important: i-bgp is not an igp; rather it relies on an igp
- if it is self-evident we shall omit e- and i-

there are 2 bgp's



- e-bgp
 - used by pairs of routers in different ases for interdomain routing

there are 2 bgp's



i-bgp

- used by pairs of router inside the same as for letting the interdomain routing info traverse the as
- the routers of an as do i-bgp peering in full mesh