CS 350 Notes

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Oct 20, 2011

1 Assignment 2

First things to do:

- Write to console.
- Get _exit working.

1.1 _exit

- Clean up process
 - thread
 - Address space
 - more stuff as assignment progresses like process numbers, FDs, etc.
- Start by looking at thread_exit.

1.2 Implement Argument Passing

We need programs main to get the arguments:

```
int main (int argc, char **argv);
```

Note: look at the hints page. The command to run a program and give it arguments in os/161 is the p command arg1 arg2 command.

Issuing the p command calls the function in the kernel cmd_progthread() which calls the function run_program().

For cmd_progthread, it has two arguments. The first pointer is to argv. Make sure argv has a NULL at the end of it. argc is the count of the number of arguments (or the last index in argv.

We can place argv:

- On the stack
- In its own segment
- somewhere crazy (as long as it works)

Look at dumbvm and see how it creates a stack, especially if you're putting arguments on the stack.

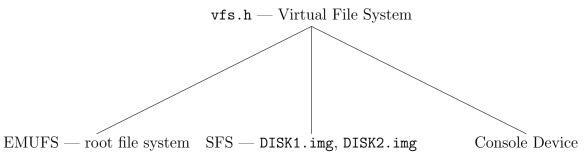
After copying argv and argc, we have to make the program aware of them by putting the arguments in a0,a1. To see how to do this, look at run_program(), and md_usermode(). The second function is the one that switches from kernel-mode to user-mode.

1.3 System Calls

- Files
 - open
 - close
 - write
- Processes
 - fork
 - waitpid
 - _exit
 - execv
- Exceptions

Do Files (without execv and processes first, in whatever order. Then do exceptions and execv.

1.3.1 File-Related System Calls



The File System Interface:

- vfs_open(path, flags, vnode)
 - path is the path to the file
 - flags are the read-write flags
 - vnode is returned, and is an abstraction of an open file.
- vfs_close(vnode) closes a file

Look at the file vnode.h for file operations. Kernel-level file operations are done for you, but the user level file operations must be implemented. Look at slides 1 to 7 of the file system notes.

1.3.2 Process System Calls

Need to write md_forkentry.

2 Address Space

2.1 Loading a Program into Address Space — ELF Files

- Format for executables is called ELF. It is used by the kernel to create the address space for a process.
- And ELF file contains
 - An ELF Header
 - Program headers
 - * Contains one header per segment.
 - * The header contains the following:
 - · Virtual address of segment
 - · Length of the segment
 - · Location of the start of the image in ELF file
 - · Length of the image in ELF file.
 - Segments $1 \dots n$.
 - Section Headers
- A segment is a chunk of data to be loaded into the address space.
- Kernel must fill rest of segment with 0's.
- OS/161 expects the ELF file to have two segments:
 - Program code
 - Data

2.2 ELF Sections and Segments

- The ELF file contains different sections
 - .text program code
 - .rodata read-only global data

- .data initialized global data
- $-\,$.bss uninitialized global data
- -.sbss small uninitialized global data
- Note all of these are present in every ELF file.
- $\bullet\,$.text and .rodata for the text segment. The rest for the data segment.

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