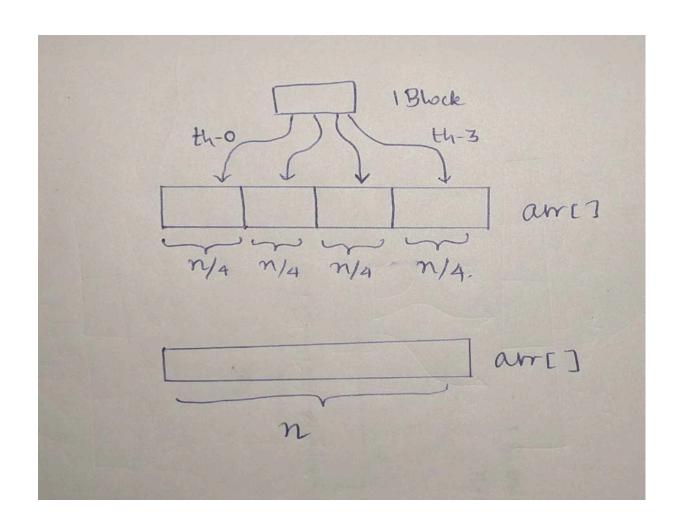
## CUDA Merge Sort

M.VENKATAKRISHNA

# MERGE SORT PARALLEL ALGORITHM (2-Stages)

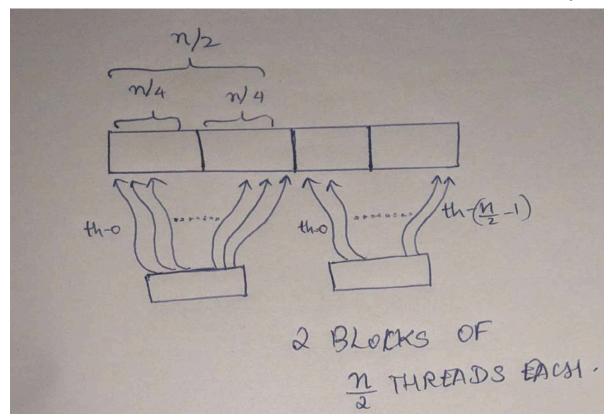
**EXPLAINED WITH 4 INITIAL SPLITS** 

#### Stage-1

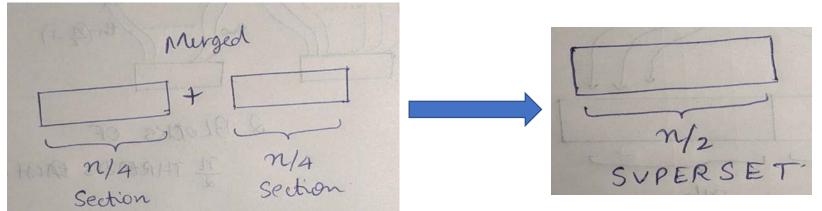


- The original array arr[] is of size n
- The array arr[] is initially split into 4 sections
- A block of 4 threads is launched: Each thread is responsible for sorting each of the 4 sections. (BUBBLE-SORT IS EMPLOYED FOR THIS)

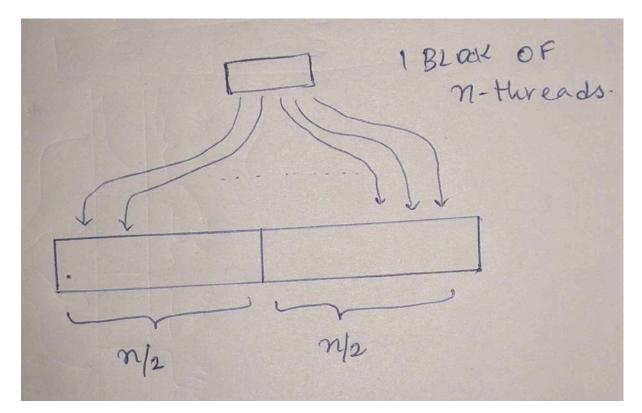
### STAGE-2 (Iteration-1)



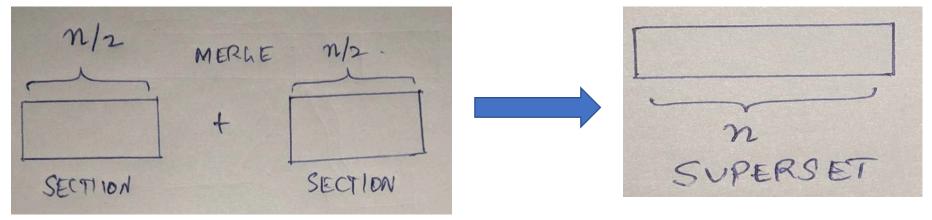
- Now, each section of n/4 elements has been sorted (STAGE-1)
- 2 blocks of n/2 threads are launched.
- Each block parallelly merges a pair of 2 sections of size: n/4 to produce a sorted SUPERSET of size: n/2
- Each Block of n/2 threads is responsible for producing 1 SUPERSET of size: n/2



### STAGE-2 (Iteration-2)



- Now, each section of n/2 elements has been sorted (STAGE-2 Iteration 1)
- 1 blocks of n threads are launched.
- This block merges a pair of 2 sections of size: n/2 to produce a sorted
   SUPERSET of size: n
- This Block of n threads is responsible for producing 1 SUPERSET of size-n

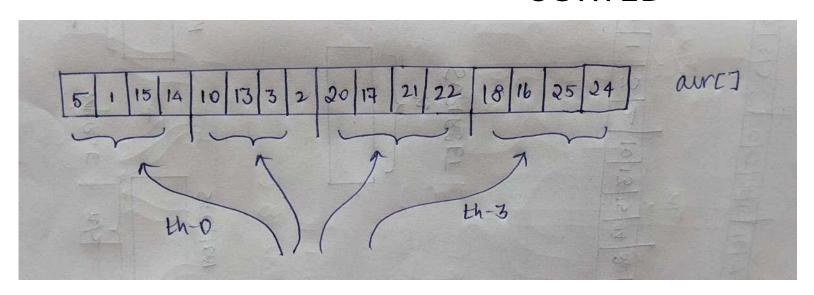


# MERGE SORT PARALLEL ALGORITHM

**EXAMPLE RUN WITH 4 INITIAL SPLITS** 

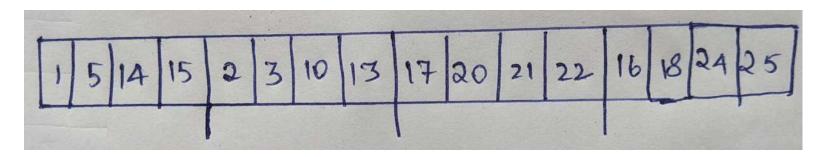
Example  $arr[] = \{5,1,15,14,10,13,3,2,20,17,21,22,18,16,25,24\}$ 

## STAGE-1: ARRAY IS DIVIDED INTO SECTIONS AND EACH SECTION IS SORTED

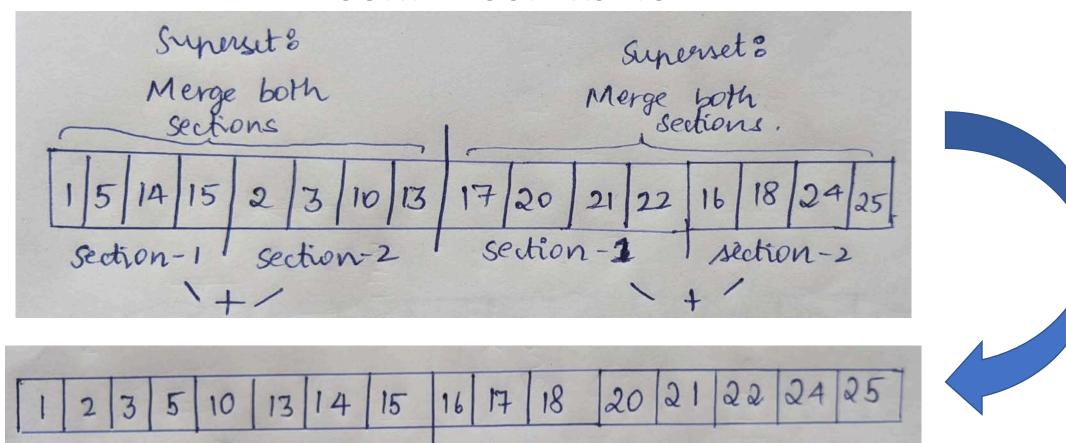


- 4 threads are launched
- Each thread
   individually sorts a
   section of 4 elements



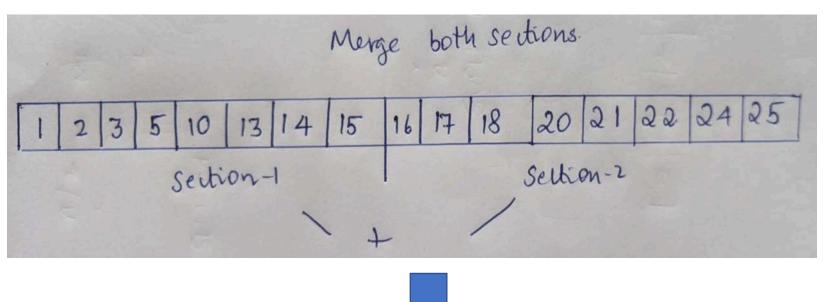


# STAGE 2: ITERATION-1 EACH SORTED SECTION IS MERGED IN PAIRS TO PRODUCE SORTED SUPERSETS



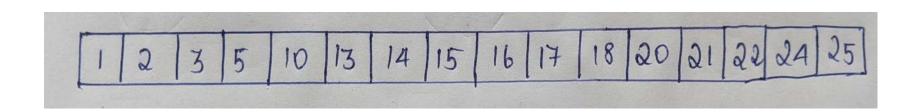
- 2 Sections on left are merged to form a sorted superset in left
- 2 sections in right are merged to form a sorted superset in right

## STAGE 2: ITERATION-2 EACH SORTED SECTION IS MERGED IN PAIRS TO PRODUCE SORTED SUPERSETS



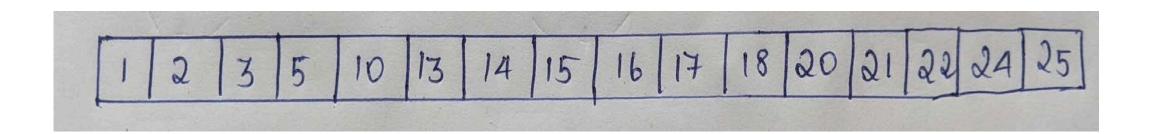


2 Sorted Sections of sizen/2 each are merged to form a sorted superset of size n



#### SORTED ARRAY IS FINALLY OBTAINED

arr[] = {1,2,3,5,10,13,14,15,16,17,18,20,21,22,24,25}



# MERGE SORT CUDA IMPLEMENTATION

### KERNEL : section\_sort()

```
//Entire block for the entire array. Each thread takes care of bubble-sorting an individual section of size : *section_length 
global_ void section_sort(int* nums, int section_size) //(, int n)

//The thread with thread index = idx will take care of nums[] from : [( section_size * idx ) to ( section_size * (idx + 1) - 1 )]

//For example: idx = 1 and section_size = 20, then, thread with idx = 1 will take care of nums[ 20: 39 ]

int idx = threadIdx.x;

//Bubble sort nums[] from index [ ( section_size * idx ) : ( section_size * (idx + 1) - 1 ) ]

bubblesort(nums, section_size * idx, ( section_size * (idx + 1) - 1));

}
```

#### Takes as INPUT:

- nums\* : The input array, which is to be divided into sections and sorted sectionally
- section\_size : The size (number of elements) of each section

#### **Produces OUTPUT:**

• nums[] Input array but with each section individually sorted

#### section sort(): KERNEL-CALL AND WORKING

```
section_sort <<<1, div_num>>> (d_arr, section_size);
```

- div\_num = Number of sections into which arr[] is split into.
- A single block of div\_num number of threads are launched.
- Each thread will be responsible for sorting one specific section of the arr[].
- Thread number idx will be responsible for the section:

```
[ idx * section_size : (idx+1) * section_size - 1 ]
```

 Bubble-Sort() is employed by each thread to sort the section its responsible for. Stage-2

Merge Sorted Sections

#### Merge Kernel Call

```
//Number of Threads = Size of Superset
//Number of blocks = Number of supersets
merge <<<ౖ superset_num, n/superset_num, n/superset_num >>> (d_arr, section_size, d_out_temp);
```

- **Number of Blocks** = Number of Supersets
- Number of threads = Length of each superset
- **Each Block**: Is responsible for a superset (merging the 2 sections of its superset)
- **Each thread**: Is responsible for 1 element of the superset

### KERNEL: merge()

```
global void merge(int* arr, int section length, int* d out temp)
  //int section length = *section size;
  int superset_length = section_length * 2;  //Block will be 2 * (size of 1 section). Because 2 sections are merged
  int idx = threadIdx.x;
  int b_idx = blockIdx.x;
  //Length of arr1[] and arr2[] are section size
  int len1 = section length;
  int len2 = section_length;
  //----Select *arr1 and *arr2 and *d out curr------
  int* arr1 = arr + (b idx * superset length);
  int* arr2 = arr1 + (section_length);
  int* d out curr = d out temp + (b idx * superset length); //Determine d out curr[], the output array for current merge
  //Dynamically allocated shared memory array.
  // scat ad[] from index [0 to n1-1] is for arr1[].
  //scat_ad[] from index [n1 to n2-1] is for arr2[]
   //Create a shared memory of size n1+n2 to accommodate the scatter-addresses corresonding to each element in arr1[] and arr2[]
  extern shared int scat ad[];
```

#### Takes as INPUT:

- arr\*: Input array
- section\_length: length of each section, which are to be merged in pairs
- **d\_out\_temp\*** : temporary array used to store intermediate output

#### **Produces As Output:**

arr[] but with pairs of sections
 of section\_length merged to
 form bigger sections of size
 section\_length\*2

#### Merge Kernel : Details

- section : Divisions in arr[] . They will be merged in pairs.
- **superset**: The result formed from merging to 2 sorted sections
- **superset\_length** : Length of each superset.
- arr1\*: Points to first section of superset.
- arr2\*: Points to second section of superset
- len1 = len2 = Length of arr1[] = Length of arr2[]
- scat\_ad[] : A shared memory array of length = superset-length.
  - It will hold the scatter addresses for each element of sorted superset

### Merge: Scatter Address Calculation

#### Each thread is responsible for producing scatter address of one element in the superset

- If (thread index < len1): Current thread is responsible for an element in arr1[]</li>
- Target = arr1[idx] (Current Number whose scatter address is to be calculated)
- Idx1 = Number of elements smaller than target in arr1[] = idx
- **Idx2** = Number of elements smaller than target in arr2[]
  - To calculate idx2, we perform a binary\_search for target in arr2[]
- <u>Scatter address of arr[idx] in sorted array will be (idx1 + idx2)</u>, i.e, the number of elements in superset smaller than target. Hence, store <u>scatt\_ad[idx] = idx1+idx2</u>

```
//-----These threads are responsible for arr2[]------
else if (idx \geq len1)
   //Number of elements in arr2[] that are lesser than arr2[idx].
   //idx1 = index of current element in arr2[]
   //(idx-len1) because threads with index n1 to n2-1 are responsible for arr2[] index [0: n2-1]
   int idx1 = idx - len1;
   int target = arr2[idx1]; //Target is current element in arr1[]
   //-----Find idx2------Binary Search Part------
   int idx2 = bin search(arr1, target, len1) + 1; //Number of elements in arr1[] that are lesser than arr2[idx]. +1 bcos we want appro
   //Calculate and store the scatter address in array
   //scat arr1[idx] = idx1 + idx2; //If there are 2 elements before a number in output array, its index will be 2
   scat_ad[idx] = idx1 + idx2; //Scatter address corresponding to arr2[idx - len1] = idx1 + idx2
__syncthreads(); //Barrier to ensure that all threads have finished writing scat_ad[].------Not necessary
```

If idx >= len1: Current thread is responsible for an element in arr2[]

Calculating scatter address is similar to previous case

\_\_syncthread(): Barrier to ensure all threads have finished calculating scatter addresses.

## Place the elements in Superset in appropriate sorted positions

- 1)Store the elements in current superset from arr[] to d\_out[] in appropriate sorted position using scatter address from scat\_ad[] array (d\_out[] is temporary array)
- 2)Barrier: Ensure all threads have finished copying to d\_out[]
- Copy back to arr[] from d\_out[]

## merge\_sort()

**CPU CODE** 

#### merge\_sort() Function

```
+ // ...
void merge sort()
    GpuTimer timer;
    int n = sizeof(h arr) / sizeof(int); //n = Total size of host array
    int div_num = 8;  //How many parts the array is initially split.
                                     // section size = Size of each section after splitting arr[] into div num parts (Stored in Hos
    int section size = n/div num;
    //-----Create input and output arrays in GPU-----
                                                                       // *d out2;
    int* d arr, * d out temp;
    cudaMalloc((void**)&d_arr, n * sizeof(int));
    cudaMemcpy((void*)d arr, (void*)h arr, n * sizeof(int), cudaMemcpyHostToDevice); //d arr[] is input array in device
    cudaMalloc((void**)&d out temp, n * sizeof(int));
                                                 //d out temp[] is temporarily used to store sorted block elements
```

- h\_arr[]: Input array to be sorted (Stored in HOST)
- d arr[]: Input array (Stored in Device)
- d\_out\_temp[] : temporary array (Device)

#### contd.

- div\_num: Number of sections into which the arr[] is initially SPLIT
- **section\_size:** Length of each section = Total array length / number of sections = n /div\_num

#### **STAGE – 1: SECTION-SORT KERNEL CALL**

```
//-----Stage-1: KERNEL CALL: Bubble Sort Each Section of section_size elements------section_size of section_size of section_si
```

- section-sort kernel call will bubble-sort all sections in arr[]. (The array is divided into div\_num number of sections each of section\_size)
- KERNEL calls 1 block of div\_num threads, 1 thread for each section in arr[].
- Each thread will sort its section

## STAGE-2: ITERATIVELY MERGE SECTIONS UNTIL ENTIRE ARRAY IS SORTED

- Iteratively call merge kernel
- In each iteration, merge sections into supersets.
- After each iteration, section size is doubled and number of superset is halved (Because after each iteration, each pair of section is merged, so that current superset is a section for the next iteration )
- In Final iteration, only a single pair of sections is to be merged and then, sorted array will be obtained

### **END**