CMPS 102 — Quarter Spring 2017 – Homework 1

VLADOI MARIAN

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I have read and agree to the collaboration policy. Vladoi Marian I want to choose homework heavy option. Name of students I worked with: Victor Shahbazian and Mitchell Etzel

Solution to Problem 2: Time Complexity

a) The following functions are ranked by increasing order of growth: gacrepqmobhdfsiltikn

$$\sum_{i=1}^{n} \frac{1}{2}^{i} < \ln(\ln n) < [\log(n^{2}), 14 \log_{3} n] < \log^{2}(n) < [\sqrt[2]{n}, 2^{\log_{4} n}] < [n, 2^{\log n}] < (n \log n, \log(n!)) < [n^{2}, 4^{\log_{2} n}, \sum_{x=5}^{n} \frac{x+1}{2}] < n^{\log 7} < 2^{\log^{2}(n)} < (\frac{5}{4})^{n} < [3^{n}, \sum_{i=1}^{n} 3^{i}] < n!$$

$$n!$$

$$\mathbf{g}) \sum_{i=1}^{n} (\frac{1}{2})^{i} = \int_{1}^{n} (\frac{1}{2})^{n} dn = \frac{\frac{1}{2} - 2^{-n}}{\ln 2} \approx \frac{1}{2^{n}} \rightarrow \sum_{i=1}^{n} (\frac{1}{2})^{i} < \ln \ln n \text{ because } \lim_{x \to \infty} \frac{\frac{1}{2} - 2^{-n}}{\ln 2} = 0$$

$$\mathbf{a}) \ln \ln n \rightarrow \lim_{x \to \infty} \frac{\ln \ln n}{14 \log_{3} n} = 0, \ln \ln n < 14 \log_{3} n$$

$$\mathbf{c}) 14 \log_{3} n \rightarrow \lim_{x \to \infty} \frac{14 \log_{3} n}{\log_{2}(n^{2})} = \frac{7 \ln 2}{\ln 3}, 14 \log_{3} n \approx \log_{2}(n^{2})$$

$$\mathbf{r}) \log_{2}(n^{2}) \rightarrow \lim_{x \to \infty} \frac{\log_{2}(n^{2})}{(\log_{2} n)^{2}} = 0, \log_{2}(n^{2}) < (\log_{2} n)^{2}$$

$$\mathbf{e}) (\log_{2} n)^{2} \rightarrow \lim_{x \to \infty} \frac{(\log_{2} n)^{2}}{2 \log_{4} n} = 0, (\log_{2} n)^{2} < 2^{\log_{4} n}$$

$$\mathbf{n}) 2^{\log_{4} n} - n^{\log_{4} 2} - n^{\frac{1}{2}}$$

a)
$$\ln \ln n \to \lim_{x \to \infty} \frac{\ln \ln n}{14 \log_2 n} = 0$$
, $\ln \ln n < 14 \log_3 n$

c)
$$14 \log_3 n \to \lim_{x \to \infty} \frac{14 \log_3 n}{\log_2(n^2)} = \frac{7 \ln 2}{\ln 3}, 14 \log_3 n \approx \log_2(n^2)$$

r)
$$\log_2(n^2) \to \lim_{x \to \infty} \frac{\log_2(n^2)}{(\log_2 n)^2} = 0$$
, $\log_2(n^2) < (\log_2 n)^2$

e)
$$(\log_2 n)^2 \to \lim_{x \to \infty} \frac{(\log_2 n)^2}{\log_4 n} = 0$$
, $(\log_2 n)^2 < 2^{\log_4 n}$

p)
$$2^{\log_4 n} = n^{\log_4 2} = n^{\frac{1}{2}}$$

q)
$$\sqrt[2]{n} = n^{\frac{1}{2}}$$

m)
$$2^{\log_2 n} = n^{\log_2 2} = n$$

- **o**) n
- **b)** nlog(n)
- **h)** $log(n!) \approx nlog(n) stirling formula$
- d) n^2
- **f)** $4^{\log_2 n} = n^{\log_2 4} = n^2$

s)
$$\sum_{x=5}^{n} \frac{x+1}{2} = \frac{1}{4}(n^2 + 3n - 28) \rightarrow \lim_{x \to \infty} \frac{\sum_{x=5}^{n} \frac{x+1}{2}}{n^2} = constant$$

j) $n^{log(7)}$, $log(7) > 2$

1)
$$2^{log^2(n)}$$
 exponential > polinominal
t) $\left(\frac{5}{4}\right)^n \to \lim_{x \to \infty} \frac{2^{log^2(n)}}{\left(\frac{5}{4}\right)^n} = 0$

i)
$$3^n$$
, because $3 > \frac{5}{4}$

k)
$$\sum_{i=1}^{n} 3^{i} = \frac{3}{2}(3^{n} - 1) \rightarrow \lim_{x \to \infty} \frac{\sum_{i=1}^{n} 3^{i}}{3^{n}} = constant$$
n) $n!$ $factorial > exponential$

n)
$$n!$$
 $factorial > exponential$

b)

Claim 1.
$$f(n) + g(n) = \Omega(max(f(n), g(n)))$$

This statement is ALWAYS TRUE

Proof. Let f (n) and g(n) be asymptotically non-negative functions.

Using the definition of Ω

 \exists a positive constant n_0 such that f(n) > 0 and $g(n) > 0 \ \forall n > n_0$

For such n we have:

$$0 \le \max(f(n), g(n)) \le \min(f(n), g(n)) + \max(f(n), g(n))$$

But
$$f(n) + g(n) = min(f(n), g(n)) + max(f(n), g(n))$$
, so $\forall n \ge 0$ we have :

$$0 \le 1 * max(f(n), g(n)) \le f(n) + g(n)$$

Thus
$$f(n) + g(n) = \Omega(\max(f(n), g(n)))$$
, as required.

Claim 2.
$$f(n) = \omega(g(n))$$
 and $f(n) = \mathcal{O}(g(n))$

This statement is ALWAYS FALSE

Proof. I will proof by contradiction.

If our statement is true, it means that $\omega((q(n)) \cap \mathcal{O}(q(n)) \neq \emptyset$.

Let g(n) be an asymptotically non negative function.

Assume that $f(n) \in \omega q(n) \cap \mathcal{O}q(n)$.

(1) Since
$$f(n) = \mathcal{O}g(n)$$
: $\exists c_1 > 0$, $\exists n_1 > 0$, $\forall n \ge n_1$: $0 \le f(n) \le c_1 g(n)$.

(2)Also since
$$f(n) = \omega g(n)$$
: $\exists c_2 > 0$, $\exists n_2 > 0$, $\forall n \ge n_2$: $0 \le c_2 g(n) < f(n)$.

Let $c_2 = c_1$ and $m = max(n_1, n_2)$. Then by (1) and (2) we have :

 $0 \le c_1 g(m) < f(m) \le c_1 g(m)$ which implies that $c_1 g(m) \le c_1 g(m)$, which is a contradiction.

Our assumption was false and $\omega((q(n)) \cap \mathcal{O}(q(n)) = \emptyset$.

$$f(n) = \omega(g(n))$$
 and $f(n) = \mathcal{O}(g(n))$ is false.

Claim 3. Either $f(n) = \mathcal{O}(n)$ or $f(n) = \Omega(g(n))$ or both.

This statement is ALWAYS TRUE

Proof. Let f (n) and g(n) be asymptotically non-negative functions.

Either is the case by definition of \mathcal{O}

$$\exists$$
 positive constants n_1, c_1 such that $f(n) \ge 0$ and $g(n) \ge 0 \ \forall n \ge n_1$

$$0 \le f(n) \le c_1 g(n)$$

Either is the case by definition of Ω

 \exists positive constants n_2, c_2 such that $f(n) \ge 0$ and $g(n) \ge 0 \ \forall n \ge n_2$

$$0 < c_2 q(n) < f(n)$$

Either both by the definition of Θ

 \exists positive constants $n_3 = max(n_1, n_2), c_1, c_2$ such that $f(n) \ge 0$ and $g(n) \ge 0 \ \forall n \ge n_3$

$$0 \le c_2 g(n) \le f(n) \le c_1 g(n)$$