

Context Dependant Environment Library Manual

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1 Introduction

This document is a manual for using the advanced parameter library within *Chisel*. For more general information regarding *Chisel* as a hardware construction language, please see the Getting Started documentation.

As hardware designs grow in complexity, modularity becomes necessary for maintainance and verification. The primary use case for *Chisel* is describing diverse and highly-parameterized hardware generators, and we quickly realized that the traditional parameterization method forces brittleness into a design's source code and limits component reuse.

The outline of this document is as follows: in Section 2, we describe the basic objects and methods for the advanced parameterization mechanism, as well as the required boilerplate to use it. In Section 3, a series of increasingly complex examples of design patterns are described. For each example, we propose the simplest parameterization scheme which solves the problem. As the examples build in complexity, so do the parameterization requirements, until we arrive at the described advanced parameterization mechanism. The next section, Section 4, introduces the concept of Knobs and their relationship to design constraints and the *Parameters* object. Finally, in Section 5, we explain multiple design heuristics which should be followed when using advanced parameterization.

2 Advanced Parameterization

In early versions of *Chisel*, every *Chisel* Module automatically had a member *params* of class *Parameters* that provides the mechanism for passing parameters between modules. *Chisel* 3 extracts this advanced parameterization into a library, the Context Dependent Environment (CDE) library.

This section describes the following features of

the CDE library: (1) the *Parameters* class and associated methods/members; (2) the basic usage model; (3) syntactic sugar; (4) boilerplate code for exposing parameters to external users/programs; (5) advanced functionality via *Views* (site, here, up).

2.1 Classes and Methods

The *Parameters* class has the following base methods:

```
class Parameters {  
  // returns a value of type T  
  def apply[T](key:Any):T  
  
  // returns new Parameters class  
  def alter(mask: (Any,View,View,View)=>Any):Parameters  
  
  % This is not true for Chisel 3, correct?  
  %// returns a Module's Parameters instance  
  %def params:Parameters  
}
```

View is a class containing a base method:

```
class View {  
  // returns a value of type T  
  def apply[T](key:Any):T  
}
```

Parameters has a factory object containing one basic method:

```
object Parameters {  
  // returns an empty Parameters instance  
  def empty:Parameters  
}
```

Chisel 3 no longer natively supports passing instances of *Parameter* via the *Module* factory. Authors of *Chisel* 3 code are responsible for passing the *Parameter* instances as they see fit. One suggested pattern is shown below:

```
object MyParameterizedModule (implicit val p :  
  Parameters) extends  
  Module {
```

```

    ...
}

object MyParameterizedTopModule(topParams :
    Parameters)
    extends Module {

    implicit val p = topParams
    val subModule = MyParameterizedModule()

}

```

2.2 Basic Usage Model

This example shows a simple usage of (1) querying params, (2) altering a Parameters object, and (3) passing a Parameters object to a Module:

```

class Tile (params : Parameters) extends Module {
    val width = params[Int]('width')
}

object Top {
    val parameters = Parameters.empty
    val tile_parameters = parameters.alter(
        (key,site,here,up) => {
            case 'width' => 64
        })
    def main(args:Array[String]) = {
        chiselMain(args,()=>Module(new
            Tile(tile_parameters)))
    }
}

```

Within the Module `Tile`, the `params` member is queried by calling `Parameters.apply` with the key and return value type.

In `Top`, an empty `parameters` is created by calling `Parameters.empty`; then it is altered with a function of type `(Any,View,View,View) => Any` to return a new `Parameters` instance, which is assigned to `tile_parameters`. The top-level `tile_parameters` are passed to the constructor to the `Tile` module, where they are renamed to the special name `p`.

2.3 Syntactic Sugar: Field[T]

The simple example requires the return type `Int` must be included as an argument to the `apply` method, otherwise the Scala compiler will throw an error:

```

class Tile (params : Parameters) extends Module {
    val width = params[Int]('width')
}

```

Alternatively, one can create a case object for each key which extends `Field[T]` and pass that directly into `params` `apply` method. Because `Field` contains

the return type information, the type does not need to be passed:

```

case object Width extends Field[Int]
class Tile (params : Parameters) extends Module {
    val width = params(Width)
}

```

For the rest of the document, assume the key to every query is a case class that extends `Field[T]` with the correct return type.

2.4 Syntactic Sugar: Passing and Altering

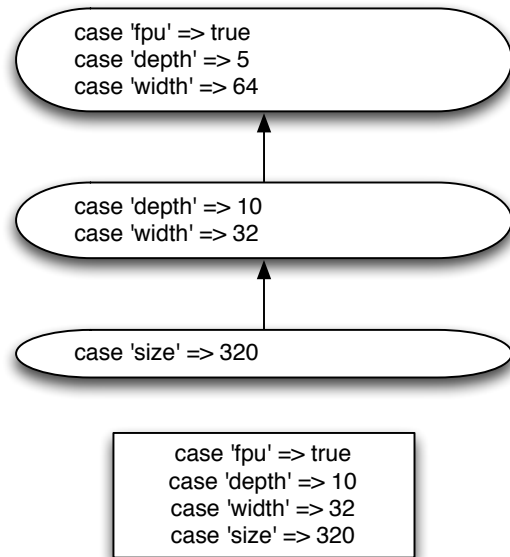


Figure 1: An example of Memory's key/value chain and flat map.

As a module hierarchy is formed, `Parameters` objects may be passed between a parent module and a child module. In Chisel 3, this is the responsibility of the Module author(s), as it is not natively supported by Chisel itself. If specified by the programmer, these objects can be copied and altered prior to instantiating the child.

Any time an alteration is performed, The *CDE* library internally copies the existing chain of key/-value mappings and attaches the provided key/-value mappings to the bottom of this chain. When a query is evaluated, it first queries the chain's bottom key/value mapping. If there is no match, the query is then evaluated on the next key/value mapping in the chain, and so forth. If a query reaches the top of

the chain with no matches, the *CDE* library triggers a `ParameterUndefinedException`.

When instantiating a child, the parent can pass its `Parameters` object as it sees fit. The suggested usage is one of two ways:

1. Explicitly pass its `Parameters` object to its child via an argument to its constructor:

```
class Core (p: Parameters) extends Module {
  val width = p(WIDTH)
}

class Tile (params : Parameters) extends Module {
  val width = params(WIDTH)
  val core = Module(new Core(params))
  // Explicit passing of Tile's params to Core
}
```

2. Implicitly pass its `Parameters` object to its child:

```
class Core (implicit val p : Parameters) extends
  Module {
  val width = p(WIDTH)
}

class Tile (params : Parameters) extends Module {
  implicit val p = params
  val width = p(WIDTH)
  val core = Module(new Core())
  // Implicit passing of Tile's params to Core
}
```

If a parent wants to copy/alter the child's dictionary, the parent has three methods to do so:

1. Call the `Parameter.alter` function, which returns a new `Parameters` object. This approach gives the programmer access to the new `Parameters` object, as well as the ability to use site, here, and up (see Sections 2.6, 2.7, 2.8):

```
class Tile (params : Parameters) extends Module {
  val width = params(WIDTH)
  val core_params = params.alter(
    (pname,site,here,up) => pname match {
      case WIDTH => 32
    })
  val core = Module(new Core(core_params))
  // Use the Parameter.alter method to return an
  // altered Parameter object. Only use when
  // site, here, or up mechanisms are needed.
}
```

2. Call the simpler `Parameter.alterPartial` function, which returns a new `Parameters` object. This approach still gives the programmer access to the new `Parameters` object.

```
class Tile (params : Parameters) extends Module {
  val width = params(WIDTH)
  val core_params = params.alterPartial ({
    case WIDTH => 32
  })
  val core = Module(new Core(core_params))
}
```

A more complicated example of a alteration chain is shown in Figure 1 and described below:

```
class Tile (params : Parameters) extends Module {
  ...
  implicit val p = params.alterPartial({case FPU =>
    true; case QDepth => 20; case WIDTH => 64})
  val core = Module(new Core())
}

class Core (implicit val p: Parameters) extends Module {
  val fpu = p(FPU)
  val width = p(WIDTH)
  val depth = p(DEPTH)
  val queue = Module(new Queue()(p.alterPartial({case
    DEPTH => depth*2;
    case WIDTH=>32})))
}

class Queue (implicit val p : Parameters) extends
  Module {
  val depth = p(DEPTH)
  val width = p(WIDTH)
  val mem = Module(new Memory()(p.alterPartial({case
    SIZE => depth * width})
  })
}

class Memory (implicit val p : Parameters) extends
  Module {
  val size = p(SIZE)
  val width = p(WIDTH)
}
```

2.5 ChiselConfig and Boilerplate

The *CDE* library's mechanism to seed the top-level parameters is through a `Config` object. `Config.topDefinitions` contains the highest parameter definitions and is of the following form:

```
case object Width extends Field[Int]
class DefaultConfig extends Config {
  val topDefinitions:World.TopDefs = {
    (pname,site,here) => pname match {
      case WIDTH => 32
    }
  }
}
```

Normally, a design calls `chiselMain.apply` to instantiate a design. To use the *CDE* library's parameterization mechanism and correctly seed a `Config`, one should instead call `chiselMain.run` with the design NOT surrounded by the `Module` factory. The reason for this change is to preserve backwards compatibil-

ity with existing designs, although we intend to fix this in future releases.

An example of calling `chiselMain.run` is as follows:

```
object Run {
  def main(args: Array[String]): Unit = {
    val configName = args(0)
    val config =
      Class.forName(configName).newInstance.asInstanceOf[Config]
    val world = config.toInstance
    val paramsFromConfig: Parameters =
      Parameters.root(world)
    chiselMain.run(args.drop(1), () => new
      Tile(paramsFromConfig))
  }
}
```

To instantiate a design with a specific `Config` in the example above, simply call the *Chisel* compiler with the `configClass_name` as the first argument.

2.6 Using site

To help the designer express dependencies between parameters, we added the *site* mechanism. To understand its function, remember that conceptually, a queried `Module`'s `params` member first looks at the bottom key/value mapping in its chain of key/value mappings. If there is no match, the query moves up the chain.

Suppose we have some modules which have following form:

```
class Core (implicit val p: Parameters) extends Module {
  val data_width = p(Width)
  ...
}
class Cache (implicit val p: Parameters) extends Module
{
  val line_width = p(Width)
  ...
}
```

Unfortunately, both have identical queries for `Width` but, for this example's sake, have different semantic meaning. Inside a `Core`, `Width` means the word size, while in the `Cache`, `Width` means the width of a cache line. We want to be able to easily tailor the parameter's response to either query.

The *site* mechanism allows a key/value mapping in the middle of the chain to make its own queries that start at the bottom of the chain.

Consider the following example:

```
class DefaultConfig extends ChiselConfig {
  val top:World.TopDefs = {
    (pname,site,here) => pname match {
      case Width => site(Location) match {
        case 'core' => 64 // data width
        case 'cache' => 128 // cache line width
      }
    }
  }
}
```

```
}
}
}
}
class Tile (params : Parameters) extends Module {
  val core = Module(new Core(params.alterPartial({case
    Location => 'core'})))
  val cache = Module(new
    Cache(params.alterPartial({case Location =>
      'cache'})))
}
```

The top-level key/value mapping is using *site* to query the bottom of the chain for `Location`. Depending on what value returns (either `'core'` or `'cache'`), the top-level key/value mapping produces a different value (Figure 2).

2.7 Using here

If a parameter is a deterministic function of other parameters expressed at the same group in the key/value mapping chain, one does not want to duplicate a value, as giving a new value would require multiple changes. Instead, one can use the *here* mechanism to query the same group of key/value mappings that *here* was called:

```
class Tile (params: Parameters) extends Module {
  val cache_params = params.alter(
    (pname, site, here, up) => pname match {
      case Sets => 128
      case Ways => 4
      case Size => here(Sets)*here(Ways)
    })
  val cache = Module(new Cache(cache_params))
}
```

2.8 Using up

The *up* mechanism enables the user to query the parent group of key/value mappings. It is equivalent to calling `Parameters.apply` directly, but can be done within calling `Parameters.alter`. For an example use, see Section 3.6.

3 Examples

The three goals of any parameterization scheme are: (1) all searchable parameters are exposed at the top level; (2) source code must never change when evaluating different points; (3) adding new parameters requires little source code change. After each example is described, we present the simplest parameterization scheme that supports the desired design space without violating any of the three goals. As

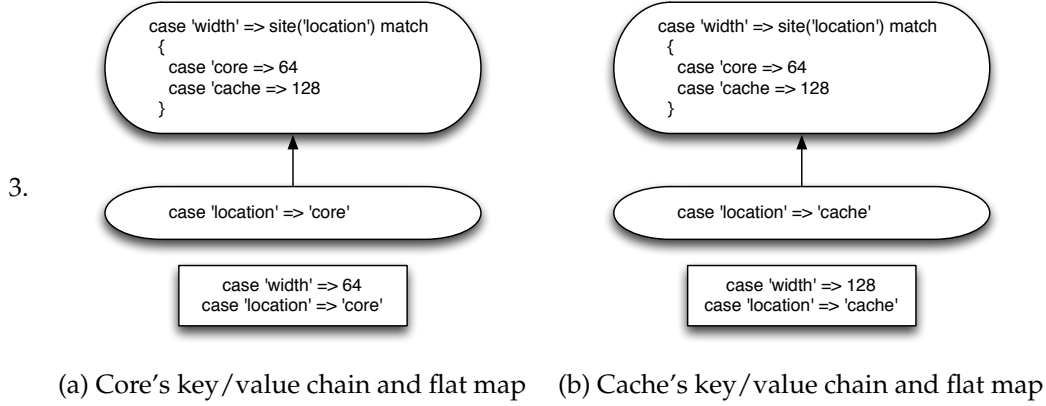


Figure 2: For (a), `site(Location)` will return Core, while in (b) `site(Location)` will return Cache.

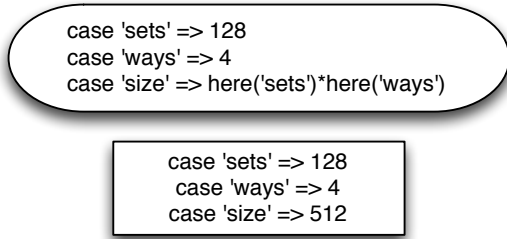


Figure 3: Instead of using 128 or 4 directly, we can access it via `here(Sets)` and `here(Ways)`, respectively.

examples grow in complexity, so too must the simplest parameterization scheme, until we arrive at the current advanced parameterization method.

3.1 Simple Parameters

In this simple design, we only vary core and cache-specific parameters. The most straightforward parameterization scheme is passing all parameters via arguments to `Tile`'s constructor. These values are then passed to `Core` and `Cache` via their respective constructors:

```
class Tile (val fpu:Boolean, val ic_sets:Int, val
           ic_ways:Int, val dc_sets:Int, val dc_ways:Int)
  extends Module {
  val core = Module(new Core(fpu))
  val icache = Module(new Cache(ic_sets,ic_ways)
  val dcache = Module(new Cache(dc_sets,dc_ways))
  ...
}
class Core (val fpu:Boolean) {...}
class Cache(val sets:Int, val ways:Int) extends Module
  {...}
```

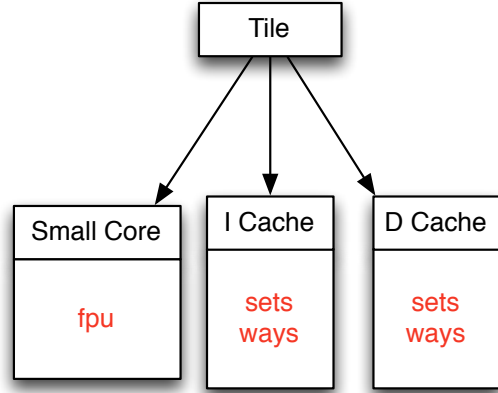


Figure 4: For a few number of parameters, the simplest scheme is to pass them directly via constructor arguments.

No source code changes are necessary to explore our parameter space, and all searchable parameters are exposed at the top. In addition, adding a new parameter, because this example is simple, requires very few changes to our source code.

3.2 Disjoint Parameter Sets

In this next design, we are designing a chip which could instantiate different cores, each with its own set of parameters. If we apply our simple solution, the number of arguments to `Tile`'s constructor would be huge as it must contain all parameters for all possible cores (which would likely be much greater than two cores!).

One could propose that a better solution would

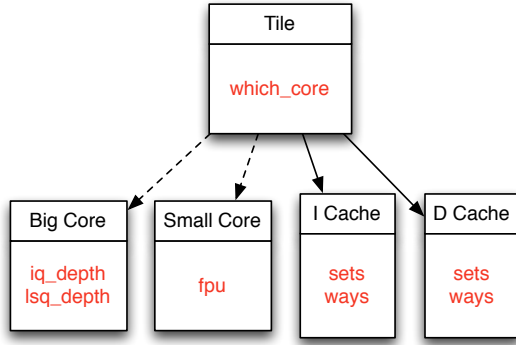


Figure 5: For disjoint parameter sets, we can group sets of parameters into configuration objects to pass as constructor arguments.

be to group parameters into configuration objects. For example, we could group all `BigCore` parameters into a `BigCoreConfig` case class, and all `SmallCore` parameters into a `SmallCoreConfig` class, both of which extend `CoreConfig`. In addition, we have our caches and `Tile` accept a `CacheConfig` and `TileConfig`, respectively, within their constructors.

```
abstract class CoreConfig {}
case class BigCoreConfig(iq_depth: Int, lsq_depth: Int)
  extends CoreConfig
case class SmallCoreConfig(fpu: Boolean) extends
  CoreConfig
case class CacheConfig(sets: Int, ways: Int)
case class TileConfig(cc: CoreConfig, icc: CacheConfig,
  dcc: CacheConfig)

class Tile (val tc: TileConfig) extends Module {
  val core = tc.cc match {
    case bcc: BigCoreConfig => Module(new
      BigCore(tc.bcc))
    case scc: SmallCoreConfig => Module(new
      SmallCore(tc.scc))
  }
  val icache = Module(new Cache(tc.icc)
  val dcache = Module(new Cache(tc.dcc))
  ...
}
```

3.3 Location-Independent Parameters

The subtle reason why nested configuration objects are extremely brittle is the structure of the nested configuration objects encodes the module hierarchy. Given a new design described in Figure 6, we assume that `BigCore`'s IQ and LSQ, as well as the icache and dcache, instantiate a `Memory` module. This `Memory` module contains a `width` parameter, and in order for the design to function correctly, all of `Memory`

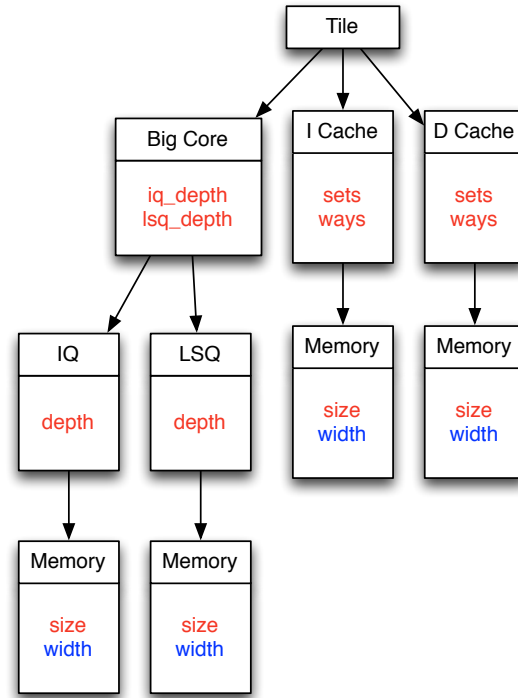


Figure 6: For location-independent parameters, every module has a parameter dictionary, which they can copy and alter before passing to each child module.

widths must be set to the same value. To ensure this requirement, the follow code might be written:

```
case class MemConfig(size: Int, banks: Int, width: Int)
case class CacheConfig(sets: Int, ways: Int, mc: MemConfig)
case class QueueConfig(depth: Int, mc: MemConfig)
case class BigCoreConfig(iqc: QueueConfig,
  lsqc: QueueConfig, mc: MemConfig)

case class TileConfig(cc: CoreConfig, icc: CacheConfig,
  dcc: CacheConfig)

class Tile (val tc: TileConfig) extends Module {
  val core = tc.cc match {
    case bcc: BigCoreConfig => Module(new
      BigCore(tc.bcc))
    case scc: SmallCoreConfig => Module(new
      SmallCore(tc.scc))
  }
  val icache = Module(new Cache(tc.icc)
  val dcache = Module(new Cache(tc.dcc))

  require(tc.dcc.mc.width == tc.icc.mc.width)
  require(tc.bcc.iqc.mc.width == tc.bcc.lsqc.mc.width)
  require(tc.dcc.mc.width == tc.bcc.lsqc.mc.width)
  ...
}
```


The series of require statements is extremely brittle, as any change in our design's hierarchy requires massive rewrites of all of these statements. Omitting the require statements is not a viable option; these statements are necessary to enforce this fundamental design requirement.

This flaw in configuration objects leads us towards the first functionality of our custom parameterization solution, namely a copy/alter dictionary of type Parameters. We use this key-value structure (map or dictionary) to store a module's parameters.

To parameterize the design in Figure 6, we implicitly pass the Parameters object and, if an alter is needed, provide a PartialFunction to the Module factory.

```
class DefaultConfig() extends ChiselConfig {
  val top:World.TopDefs = {
    (pname,site,here) => pname match {
      case IQ_depth => 10
      case LSQ_depth => 10
      case Ic_sets => 128
      case Ic_ways => 2
      case Dc_sets => 512
      case Dc_ways => 4
      case Width => 64
      // since any module querying Width should return
      // 64, the name should NOT be unique to modules
    }
  }
}

class Tile (params : Parameters) extends Module {
  implicit val p = p
  val core = Module(new Core())
  val ic_sets = p(Ic_sets)
  val ic_ways = p(Ic_ways)
  val icache = Module(new Cache(p.alterPartial({case
    Sets => ic_sets; case Ways => ic_ways}))
  // we can rename Ic_sets to Sets, effectively
  // isolating Cache's query keys from any design
  // hierarchy dependence
  val dc_sets = p(Dc_sets)
  val dc_ways = p(Dc_ways)
  val dcache = Module(new Cache, p.alterPartial({case
    Sets => dc_sets; case Ways => dc_ways}))
  // similarly we rename Dc_sets to Sets and Dc_ways to
  // Ways
}

class Core (implicit val p : Parameters) extends Module
{
  val iqdepth = p(IQ_depth)
  val iq = Module(new Queue (p.alterPartial({case Depth
    => iqdepth})))
  val lsqdepth = p(LSQ_depth)
  val lsq = Module(new Queue (p.alterPartial({case
    Depth => lsqdepth}))
  ...
}

class Queue (implicit val p : Parameters) extends
Module {
  val depth = p(Depth)
  val mem = Module(new Memory (p.alterPartial({case
    Size => depth})))
  ...
}
```

```
}
class Cache (implicit val p: Parameters) extends Module
{
  val sets = p(Sets)
  val ways = p(Ways)
  val mem = Module(new Memory(p.alterPartial({case Size
    => sets*ways})))
}
class Memory (implicit val p: Parameters) extends Module
{
  val size = p(Size)
  val width = p(Width)
}
```

Although this parameterization method is reasonably verbose, it scales well with adding parameters, requires no source changes, and allows a single parameter, such as Width, to change all leaf modules.

3.4 Location-Specific Parameters

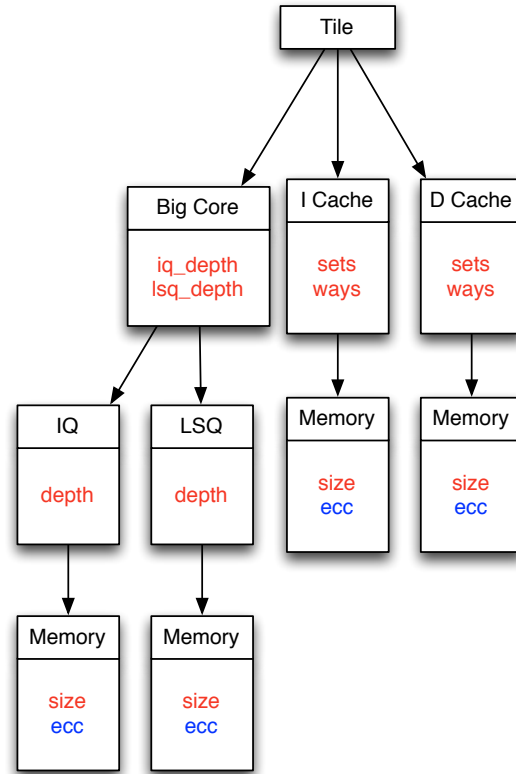


Figure 7: For location-dependent parameters, we can use the site mechanism to customize these parameters at the top level.

As we saw in the previous section, copying and altering a Parameters object can be verbose. If we wanted to add an ECC parameter to our Memory module, which depends on where the Memory is instan-

tiated, we would change source code in multiple parents to rename each parameter (e.g. ECC_icache => ECC).

In the example depicted in Figure 7, we instead use the site functionality of our Parameters object to obtain location-specific information, and tailor the value we return to that location-specific value. After adding the location-specific information, we drastically reduce the amount of code changes necessary:

```
class DefaultConfig() extends ChiselConfig {
  val top:World.TopDefs = {
    (pname,site,here) => pname match {
      case Depth => site(Queue_type) match {
        case 'iq' => 20
        case 'lsq' => 10
      }
      case Sets => site(Cache_type) match {
        case 'i' => 128
        case 'd' => 512
      }
      case Ways => site(Cache_type) match {
        case 'i' => 2
        case 'd' => 4
      }
      case Width => 64
      // since any module querying Width should return
      // 64, the name should NOT be unique to modules
      case ECC => site(Location) match {
        'incore' => false
        'incache' => true
      }
    }
  }
}

class Tile (val params: Parameters) extends Module {
  implicit val p = params
  val core = Module(new Core()(p.alterPartial({Location
    => 'incore'})))
  // we can give core and its child modules a location
  // identifier

  val cacheparams = p.alterPartial({Location =>
    'incache'})
  // we can give both caches and all their child
  // modules a location identifier
  val icache = Module(new ICache(cacheparams))
  val dcache = Module(new DCache(cacheparams))
}

class Core (implicit val p : Parameters) extends Module {
  {
    val iq = Module(new IQ())
    val lsq = Module(new LSQ())
    ...
  }
}

class IQ (implicit val p: Parameters) extends Module {
  val depth = p(Depth)
  val mem = Module(new Memory()(p.alterPartial({Size =>
    depth})))
  // in some cases, using copy/alter is preferred
  // instead of \code{site} (see Design Heuristics
  // for more details)
  ...
}

class LSQ (implicit val p: Parameters) extends Module {
```

```
  val depth = p(Depth)
  val mem = Module(new Memory()(p.alterPartial({Size =>
    depth})))
  ...
}

class ICache (implicit val p: Parameters) extends
  Module {
  val sets = p(Sets)
  val ways = p(Ways)
  val mem = Module(new Memory ()(p.alterPartial({Size
    => sets*ways})))
}

class DCache (implicit val p: Parameters) extends
  Module {
  val sets = p(Sets)
  val ways = p(Ways)
  val mem = Module(new Memory ()(p.alterPartial({Size
    => sets*ways})))
}

class Memory(implicit val p: Parameters) extends Module
{
  val size = p(Size)
  val ecc = p(ECC)
}
```

3.5 Derivative Parameters

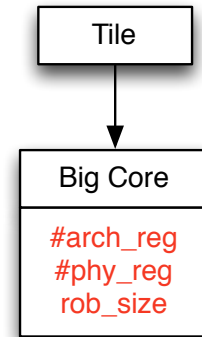


Figure 8: To derive a parameter from another top-level parameter, we can use the here functionality to avoid duplicating a parameter value.

In Figure 8, we always want our ROB to be four-thirds the size of the difference between the number physical registers and the number of architectural registers. If we express this in MyConfig.top, it could look like the following:

```
case object NUM_arch_reg extends Field[Int]
case object NUM_phy_reg extends Field[Int]
case object ROB_size extends Field[Int]
class DefaultConfig() extends ChiselConfig {
  val top:World.TopDefs = {
    (pname,site,here) => pname match {
      case NUM_arch_reg => 32
      case NUM_phy_reg => 64
```



```

    case ROB_size => 4*(64-32)/3
  }
}

```

However, if we later increase the number of physical registers, we need to remember to update the value in the derivation of the ROB size. To avoid this potential error, one should use the here functionality to query the same group of parameters:

```

class DefaultConfig() extends ChiselConfig {
  val top:World.TopDefs = {
    (pname,site,here) => pname match {
      case NUM_arch_reg => 32
      case NUM_phy_reg => 64
      case ROB_size => 4*(here(NUM_phy_reg) -
        here(NUM_arch_reg))/3
    }
  }
}

```

3.6 Renaming Parameters

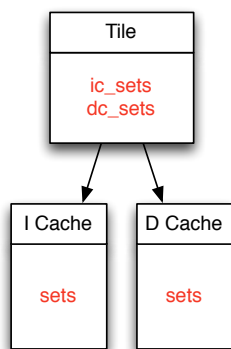


Figure 9: To rename or programmatically alter a parameter based on the previous value, one can use the up mechanism to query the parent’s Parameters object.

In Figure 9, both cache modules query for a sets parameter. However, Tile has ic_sets and dc_sets as parameters. To rename the parameters, we can read the parent value and alter the child’s Parameters object:

```

class Tile (params : Parameters) extends Module {
  implicit val p = params
  val ic_sets = params(Ic_sets)
  val ic = Module(new Cache()(p.alterPartial({case Sets
    => ic_sets})))
  val dc_sets = params(Dc_sets)
  val dc = Module(new Cache()(p.alterPartial({case Sets
    => dc_sets})))
  ...
}

```

Alternatively, we can use the ‘up’ mechanism within the Parameters.alter method to query the parent module’s Parameter object:

```

class Tile (val params : Parameters) extends Module {
  val ic_params = params.alter(
    (pname,site,here,up) => pname match {
      case Sets => up(Ic_sets)
    }
  )
  val ic = Module(new Cache(ic_params))
  ...
}

```

In general one should never use the up mechanism as it is more verbose. However, it can be useful if the parent is making significant changes to a child’s Parameters object, as all changes can be contained the Parameter.alter method because one has access to all three central mechanisms (up, site, and here).

4 External Interface

So far, this document has only described mechanisms to manipulate parameters at a top-level class (ChiselConfig). However, to actually generate multiple C++ or Verilog designs, we need to manually change these parameters.

One would prefer to express design constraints (parameter ranges, dependencies, constraints) and leave the actual instantiation of a specific design separate from the expression of the valid design space.

With that motivation, the CDE library has an additional feature based around the concept of “Knobs,” or parameters that are created specifically to explore a design space. This section will describe Knobs and their uses, the Dump Object, adding constraints to parameters/Knobs.

4.1 Knobs

A generator has some parameters that are fixed, and others that dictate the specific design point being generated. These generator-level parameters, called Knobs, have an additional key-value mapping to allow external programs and users to easily overwrite their values.

Knobs can only be instantiated within a ChiselConfig subclass’s topDefinitions:

```

package example
class MyConfig extends ChiselConfig {
  val topDefinitions:World.TopDefs = {
    (pname,site,here) => pname match {
      case NTiles => Knob('NTILES')
    }
  }
}

```

```

    case .... => .... // other non-generator
                        parameters go here
  }
}
override val knobValues:Any=>Any = {
  case 'NTILES' => 1 // generator parameter assignment
}
}

```

When the query `NTiles` matches within `topDefinitions`, the `Knob('NTILES')` is returned. Internally, the CDE library will lookup `'NTILES'` within `MyConfig.knobValues` and return 1. As described in the example in Section 2.5, to execute a generator with this specific config, provide it as the first argument:

```
sbt run MyConfig...
```

Suppose we wanted to instantiate a new design that had two tiles: simply use Scala's class inheritance and overwrite the `knobValues` method:

```

package example
class MyConfig2 extends MyConfig {
  override val knobValues:Any=>Any = {
    case 'NTILES' => 2 // will generate new design with
                        2 tiles
  }
}

```

Notice that both classes can exist in the source code, so both designs can be instantiated from the commandline. For the new design with two tiles, simply call:

```
sbt run MyConfig2 ...
```

4.2 Dump

Downstream from Chisel, other tools might need to know specific parameter/Knob assignments. If so, just pass the Knob/value to the `Dump` object, which will write the name and value to a file, then return the Knob/value:

```

package example
class MyConfig extends ChiselConfig {
  val topDefinitions:World.TopDefs = {
    (pname,site,here) => pname match {
      case Width => Dump('Width',64) // will return 64.
                                   Requires naming the parameter as the 1st
                                   argument
      case NTiles => Dump(Knob('NTILES')) // will
                                   return Knob('NTILES'), no name needed
    }
  }
}
override val knobValues:Any=>Any = {
  case 'NTILES' => 1 // generator parameter assignment
}
}

```

The name and value of each dumped parameter can be written to a `*.knb` file.

4.3 Constraints

Now that external programs/users can easily overwrite a configuration's `knobValue` method, we have provided a mechanism for defining legal ranges for Knobs. Within a `ChiselConfig`, one can overwrite another method called `topConstraints`:

```

package example
class MyConfig extends ChiselConfig {
  val topDefinitions:World.TopDefs = {
    (pname,site,here) => pname match {
      case NTiles => Knob('NTILES')
    }
  }
  override val topConstraints:List[ViewSym=>Ex[Boolean]]
    = List( { ex => ex(NTiles) > 0 },
            { ex => ex(NTiles) <= 4 })
  override val knobValues:Any=>Any = {
    case 'NTILES' => 1 // generator parameter assignment
  }
}

```

Now, if someone tried to instantiate our design with the following configuration and command, it would fail:

```

package example
class BadConfig extends ChiselConfig {
  override val knobValues:Any=>Any = {
    case 'NTILES' => 5 // would violate our constraint,
                       throws an error
  }
}

// throws 'Constraint failed' error
sbt run BadConfig ...

```

Constraints can be declared anywhere in the design, not just at the top level, by calling a `Parameter's` `constrain` method:

```

package example
class MyConfig extends ChiselConfig {
  val topDefinitions:World.TopDefs = {
    (pname,site,here) => pname match {
      case NTiles => Knob('NTILES')
    }
  }
  override val knobValues:Any=>Any = {
    case 'NTILES' => 1 // generator parameter assignment
  }
}
class Tile (val params: Parameters) extends Module {
  params.constrain( ex => ex(NTiles) > 0 )
  params.constrain( ex => ex(NTiles) <= 4 )
}
object Run {
  def main (args: Array[String]): Unit = {
    val configName = args(0)
    val config =
      Class.forName(configName).newInstance.asInstanceOf[Config]
    val world = config.toInstance
    val paramsFromConfig: Parameters =
      Parameters.root(world)
    chiselMain.run(args.drop(1), () => new

```

```

        Tile(paramsFromConfig))
    }
}

sbt runMain MyConfig ...

```

If a designer would like to see a design's constraints, parameters, knobs, and configuration parameters, these can all be dumped to files. See the code below for an example:

```

val pdFile = createOutputFile(configClassName +
    ".prm")
pdFile.write(ParameterDump.getDump)
pdFile.close
val v = createOutputFile(configClassName + ".knb")
v.write(world.getKnobs)
v.close
val d = new java.io.FileOutputStream(Driver.targetDir
    + "/" + configClassName + ".cfg")
d.write(paramsFromConfig(ConfigString))
d.close
val w = createOutputFile(configClassName + ".cst")
w.write(world.getConstraints)
w.close

```

5 Design Heuristics

TODO

References

- [1] Bachrach, J., Vo, H., Richards, B., Lee, Y., Waterman, A., Avižienis, Wawrzynek, J., Asanović
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- [2] Odersky, M., Spoon, L., Venners, B. *Programming in Scala* by Artima.
- [3] Payne, A., Wampler, D. *Programming Scala* by O'Reilly books.