CSE340 Spring 2017 Project 2

Due: **March 3rd, 2017** by 11:59pm MST

1 Introduction

In this project, you will write a C or C++ program that reads a description of a context free grammar, then, depending on the command line argument passed to the program, performs one of the following tasks: 1) for each terminal and non-terminal in the grammar determine the number of grammar rules in which the symbol appears, 2) determine *useless* symbols in the grammar and remove them, 3) calculate `FIRST` sets, 4) calculate `FOLLOW` sets and 5) determine if the grammar has a predictive parser.

2 Input Format

We specify the input format using a context-free grammar:

$$
\begin{aligned}
\text{Input} &\rightarrow \text{Terminal-list Non-terminal-list Rule-list DOUBLEHASH} \\
\text{Terminal-list} &\rightarrow \text{Id-list HASH} \\
\text{Non-terminal-list} &\rightarrow \text{Id-list HASH} \\
\text{Rule-list} &\rightarrow \text{Rule Rule-list} \mid \text{Rule} \\
\text{Id-list} &\rightarrow \text{ID Id-list} \mid \text{ID} \\
\text{Rule} &\rightarrow \text{ID ARROW Right-hand-side HASH} \\
\text{Right-hand-side} &\rightarrow \text{Id-list} \mid \epsilon
\end{aligned}
$$

The tokens used in the above grammar description are defined by the following regular expressions:

```
ID        = letter (letter + digit)*
HASH      = #
DOUBLEHASH = ##
ARROW     = ->
```

Where `digit` is the digits from 0 through 9 and `letter` is the upper and lower case letters a through z and A through Z. Tokens are case-sensitive. Tokens are space separated and there is at least one whitespace character between any two successive tokens.

3 Semantics

The first section of the input (before the first `#`) lists all terminals of the grammar. The second section of the input lists all the non-terminals of the grammar. The first non-terminal in this list is the start symbol of the grammar. The following sections each represent a grammar rule.

Each grammar rule starts with a non-terminal symbol (the left-hand side of the rule) followed by `ARROW`, then followed by a sequence of zero or more terminals and non-terminals, which represent the right-hand side of the rule. If the sequence of terminals and non-terminals in the right-hand side is empty, then it represents a rule of the form $A \rightarrow \epsilon$.

You can assume that the symbols used in the right-hand side of grammar rules are either from the non-terminals or from the terminals listed at the beginning of the input. Also the left-hand side is guaranteed to be from the list of non-terminals.

Note that the convention of using upper-case letters for non-terminals and lower-case letters for terminals is not used here, we explicitly list the terminals and non-terminals in the beginning.

3.1 Example

Here is an example input:

```
ID COMMA colon #
decl idList1 idList #
decl -> idList colon ID #
idList -> ID idList1 #
idList1 -> #
idList1 -> COMMA ID idList1 #
##
```

The first section lists all terminals of the grammar:

$$\text{Terminals} = \{\text{ID}, \text{COMMA}, \text{colon}\}$$

The second section lists all non-terminals of the grammar:

$$\text{Non-Terminals} = \{\text{decl}, \text{idList1}, \text{idList}\}$$

The rest of the input (terminated by the double-hash) specifies the grammar rules:

$$\text{decl} \rightarrow \text{idList} \ \text{colon} \ \text{ID}$$
$$\text{idList} \rightarrow \text{ID} \ \text{idList1}$$
$$\text{idList1} \rightarrow \epsilon$$
$$\text{idList1} \rightarrow \text{COMMA} \ \text{ID} \ \text{idList1}$$

Note that even though the example shows that each rule is on a line by itself, a rule can be split into multiple lines, or even multiple rules can be on the same line, according to the formal specification. The following input describes the same grammar as the above example:

```
ID COMMA colon #  decl idList1 idList #
decl -> idList colon ID # idList -> ID idList1 #
idList1 -> # idList1
-> COMMA ID idList1 #      ##
```

4 Output Specifications

Your program should read the input grammar from standard input, and read the requested task number from the first command line argument (we provide code to read the task number) then calculate the requested output based on the task number and print the results in the specified format for each task to standard output (stdout). The following specifies the exact requirements for each task number.

4.1 Task 1

For each terminal and non-terminal, in the order they are listed in the input, determine the number of productions (grammar rules) that include that symbol and output one line in the following format:

`<symbol>: <number>`

Where `<symbol>` should be replaced by the terminal or non-terminal and `<number>` should be replaced by the calculated number for that symbol.

Example: The expected output for the example input grammar given in section 3.1 for task 1 is:

```
ID: 3
COMMA: 1
colon: 1
decl: 1
idList1: 3
idList: 2
```

4.2 Task 2: Find useless symbols

Determine *useless* symbols in the grammar and remove them. Then output each rule of the modified grammar on a single line in the following format:

`<LHS> -> <RHS>`

Where `<LHS>` should be replaced by the left-hand side of the grammar rule and `<RHS>` should be replaced by the right-hand side of the grammar rule. If the grammar rule is of form $A \rightarrow \epsilon$, use `#` to represent the epsilon. Note that this is different from the input format. Also note that the order of grammar rules that are not removed from the original input grammar must be preserved.

Definition: Symbol A is *not* useless if there is a derivation starting from S (the start symbol) in which A appears and the derivation leads to a string of terminals w or the empty string ($w \in T^*$):

$$S \stackrel{*}{\Rightarrow} ...A... \stackrel{*}{\Rightarrow} w$$

Example 1: The expected output for the example input grammar given in section 3.1 for task 2 is:

```
decl -> idList colon ID
idList -> ID idList1
idList1 -> #
idList1 -> COMMA ID idList1
```

Note that none of the symbols were useless.

Example 2: Given the following input:

```
a b c #
S A B C #
S -> A B #
S -> C #
C -> c #
S -> a #
A -> a A #
B -> b #
##
```

The expected output for task 2 is:

```
S -> C
C -> c
S -> a
```

Note that A and B are useless symbols and the modified grammar has only three rules.

4.3 Task 3: Calculate `FIRST` Sets

For each of the non-terminals of the input grammar, in the order that they appear in the non-terminal section of the input, compute the `FIRST` set for that non-terminal and output one line in the following format:

```
FIRST(<symbol>) = { <set_items> }
```

Where `<symbol>` should be replaced by the non-terminal and `<set_items>` should be replaced by the comma-separated list of elements of the set ordered in the following manner:

- If ϵ belongs in the set, represent it as `#`

- If ϵ belongs in the set, it should be listed before any other elements

- All other elements of the set should be sorted in the order in which they appear in the terminal list (first section of the input)

Example: The expected output for the example input grammar given in section 3.1 for task 3 is:

```
FIRST(decl) = { ID }
FIRST(idList1) = { #, COMMA }
FIRST(iDList) = { ID }
```

4.4 Task 4: Calculate `FOLLOW` Sets

For each of the non-terminals of the input grammar, in the order that they appear in the non-terminals section of the input, compute the `FOLLOW` set for that non-terminal and output one line in the following format:

```
FOLLOW(<symbol>) = { <set_items> }
```

Where `<symbol>` should be replaced by the non-terminal and `<set_items>` should be replaced by the comma-separated list of elements of the set ordered in the following manner:

- If EOF belongs in the set, represent it as `$`

- If EOF belongs in the set, it should be listed before any other elements

- All other elements of the set should be sorted in the order in which they appear in the terminal list (first section of the input)

Example: The expected output for the example input grammar given in section 3.1 for task 4 is:

```
FOLLOW(decl) = { $ }
FOLLOW(idList1) = { colon }
FOLLOW(idList) = { colon }
```

4.5 Task 5: Determine if the grammar has a predictive parser

Determine if the grammar has a predictive parser and output either `YES` or `NO` accordingly.

Example: The expected output for the example input grammar given in section 3.1 for task 5 is:

```
YES
```

5 Implementation

5.1 Lexer

A lexer that can recognize `ID`, `ARROW`, `HASH` and `DOUBLEHASH` tokens is provided for this project. You are free to use it if you like.

5.2   Reading command-line argument

As mentioned in the introduction, your program must read the grammar from stdin and the task number from command line arguments. The following piece of code shows how to read the first command line argument and perform a task based on the value of that argument. Use this code as a starting point for your main function.

```c
/* NOTE: You should get the full version of this code as part of the project
   material, do not copy/paste from this document. */

#include <stdio.h>
#include <stdlib.h>

int main (int argc, char* argv[])
{
    int task;

    if (argc < 2) {
        printf("Error: missing argument\n");
        return 1;
    }

    task = atoi(argv[1]);

    switch (task) {
        case 1:
            // TODO: perform task 1.
            break;

        // ...

        default:
            printf("Error: unrecognized task number %d\n", task);
            break;
    }
    return 0;
}
```

6   Evaluation

Your submission will be graded on passing the automated test cases. The test cases (there will be multiple test cases in each category, each with equal weight) will be broken down in the following way (out of 100 points):

- Task 1: 10 points

- Task 2: 30 points

- Task 3: 30 points

- Task 4: 25 points

- Task 5: 5 points

Submit your code on the course submission website: `https://cse340.fulton.asu.edu/pass/`