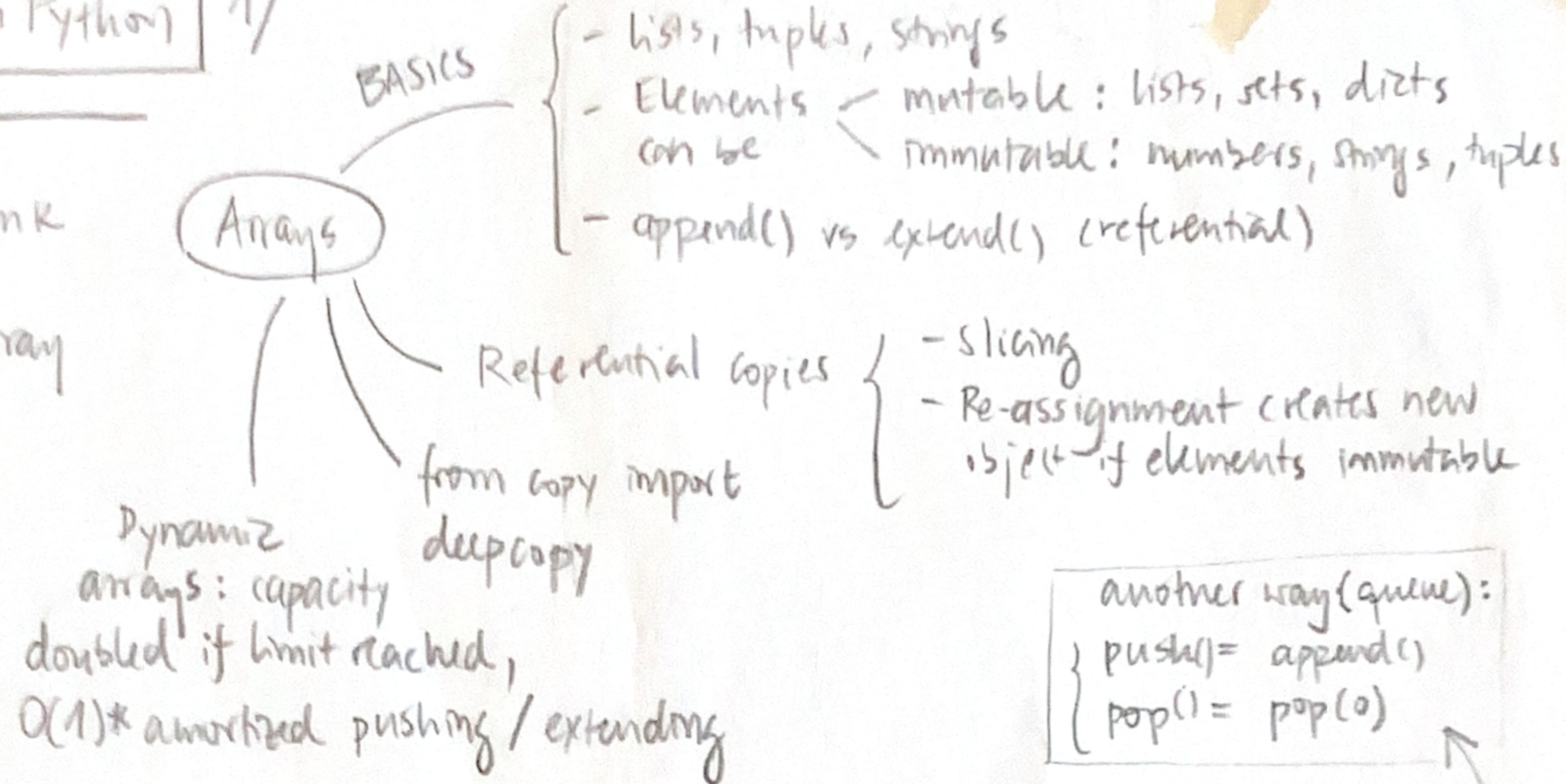
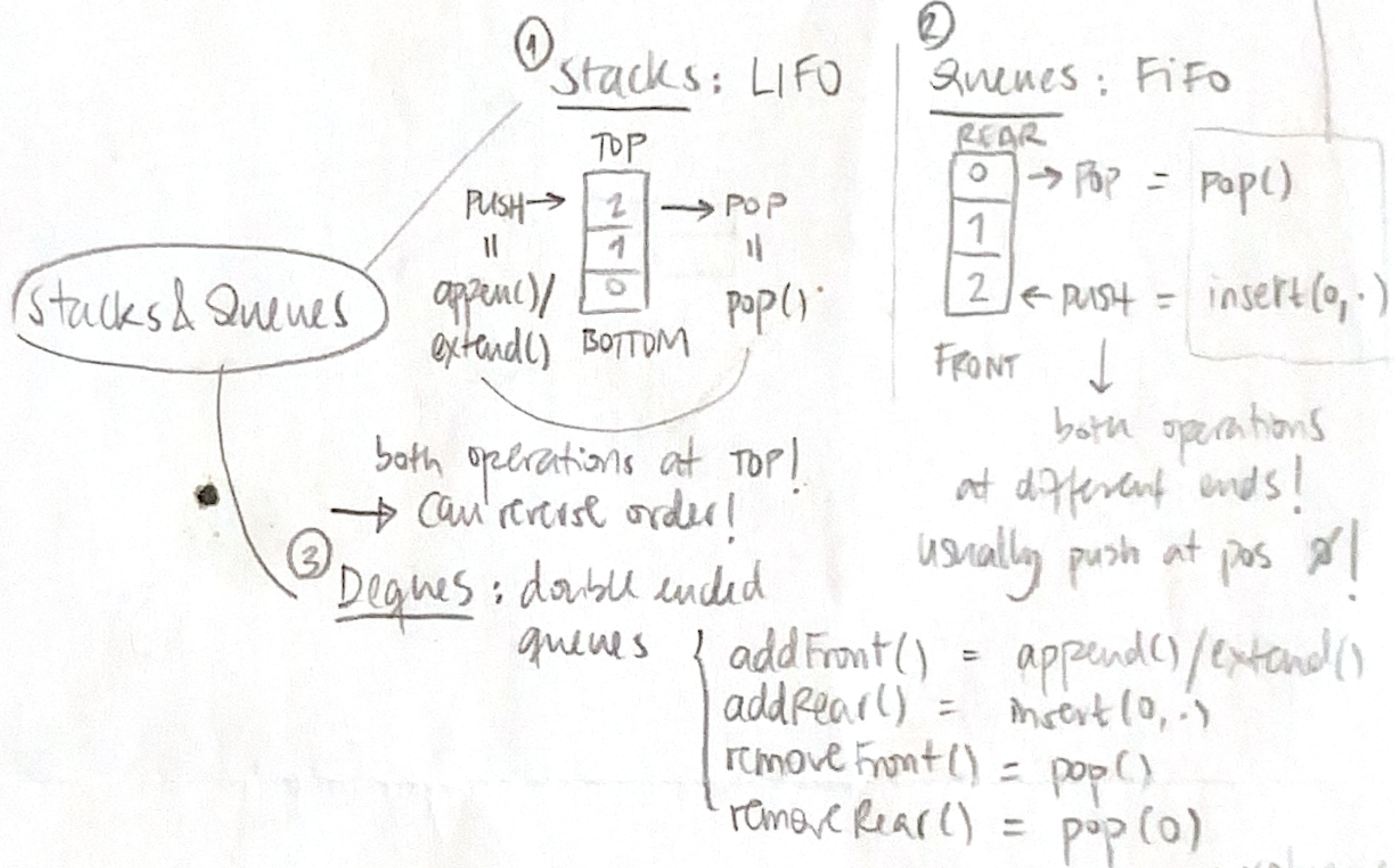


Data Structures & Algorithms in Python

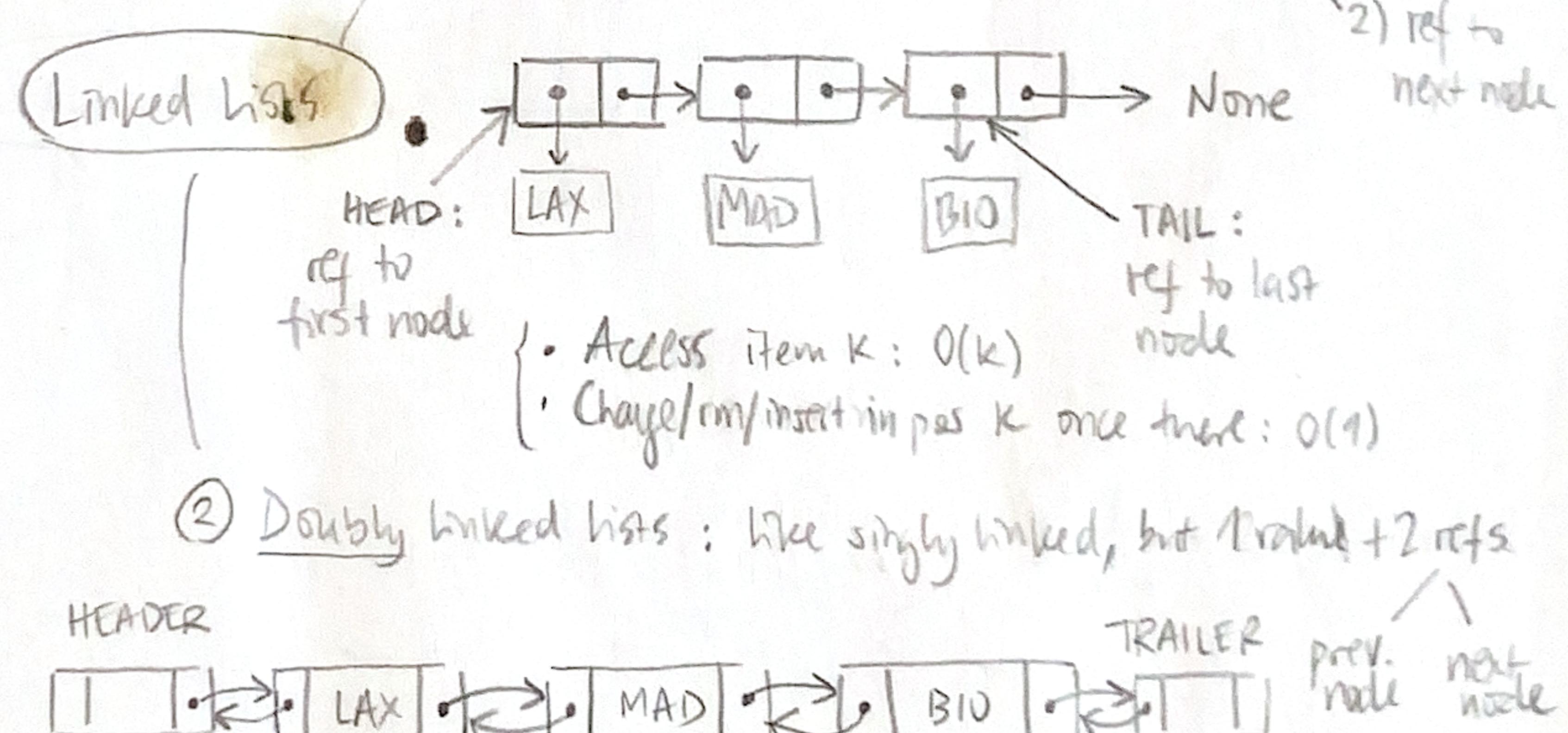
1. Given 2 strings, check if anagrams
2. Given an array, find pairs that sum K
3. Find missing element in 2 arrays
4. Largest continuous sum in an array
5. Reverse words in a sentence
6. String compression
7. Unique characters in string



0. Implement Stack & Queue
1. Check if parenthesis balanced
2. Implement a Queue using 2 Stacks



0. Implement singly/ doubly linked lists
1. Check if cycle in a singly linked list
2. Reverse a singly linked list
3. Given a list head and a value n, return the nth node before the tail



0. Boozs: factorial, \sum_i^n , sum of number digits
1. Split a stream of chars in words from a vocabulary
2. Memorize factorial: with function / with class
3. Reverse a string
4. All possible permutations of a String Δ
5. Fibonacci sequence: recursively, dynamically, iteratively
6. Given an amount n and a list of coins, what's the fewest amount of coins necessary to reach n Δ

Recursion

Advantage vs. singly linked: we can traverse from tail/trailer to head

→ Recursion: when a function calls itself - a way to avoid loops

- { Base case: when recursion stops
- Recursive case: when function is called again

Memoization: caching/storing expensive solutions to re-use, e.g. in a dict

Dynamic programming: when memorization used with recursion. Exploding solutions become feasible

- 2 ways of memoizing a function
- 1. with class wrapping technique
 - 2. function using a dict
- careful: the rec. function calls the dyn. func.

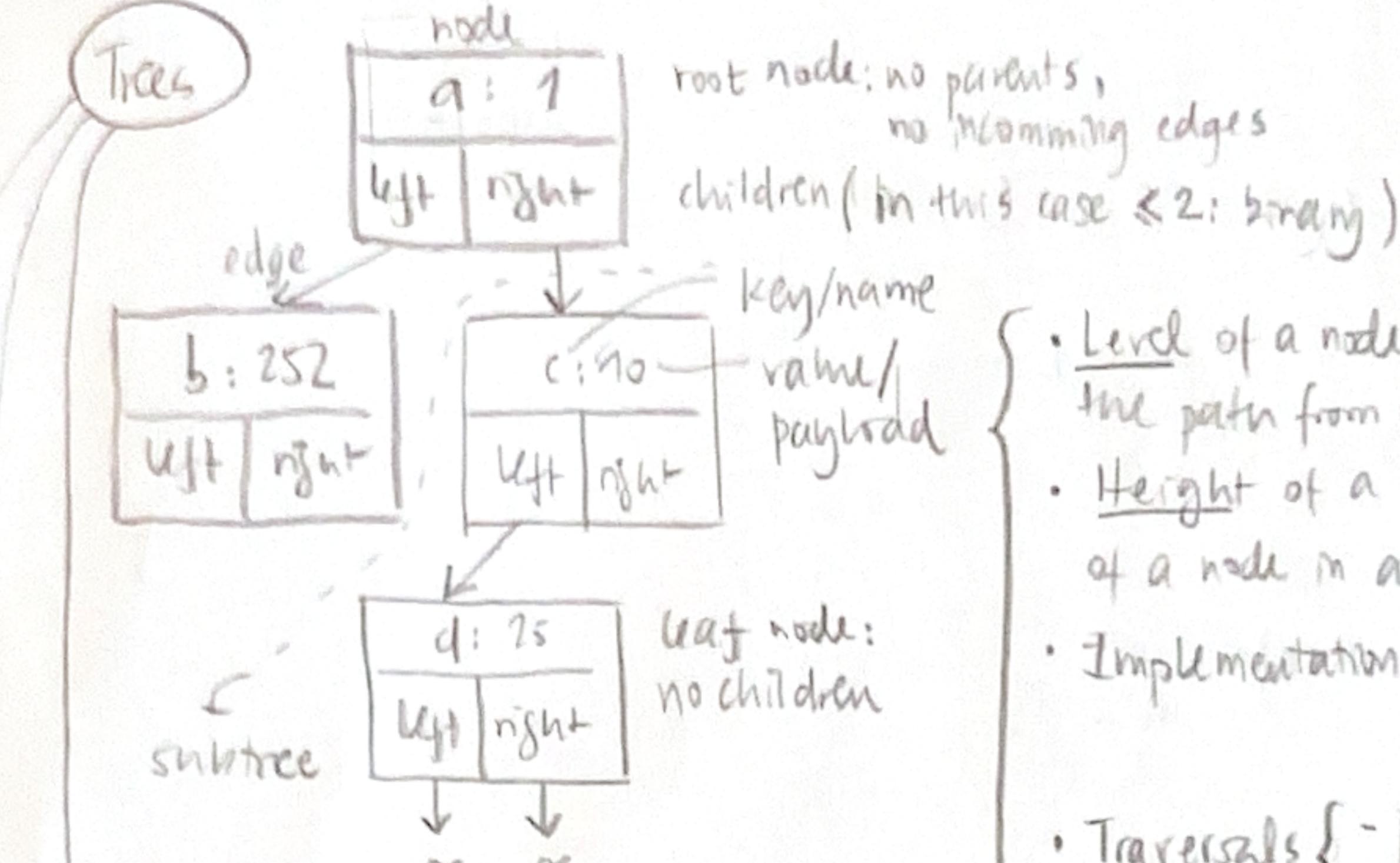
0. Tree implementation with lists / classes

1. Implement level-order traversal = breadth first.

Tries
(next page)

2. Binary heap implementation

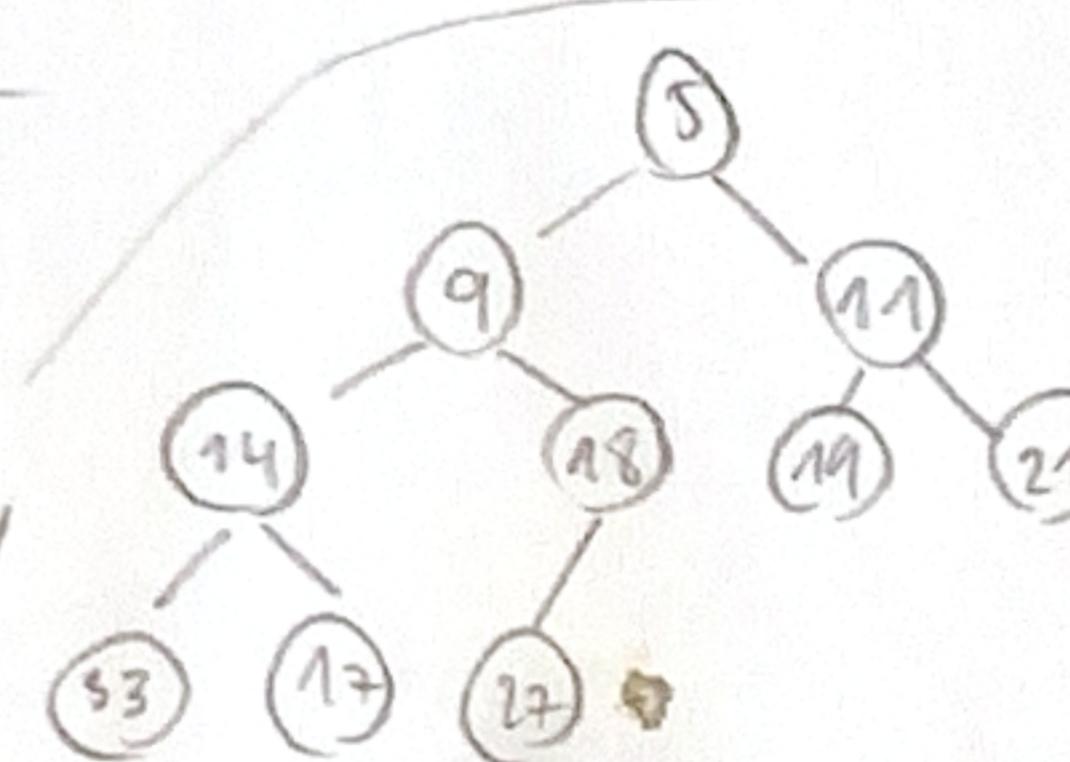
values increase/decrease



Binary heaps = Priority queues

represented by arrays, but in reality complete trees:
all nodes filled in, in last level all pushed to left

elements ordered according to relevance
 $O(\log n)$ insert/pop



p 0 1 2 3 4 5 6 7 8 9 10

| 0 | 5 | 9 | 11 | 14 | 18 | 19 | 21 | 53 | 17 | 27

Typical we care: we insert new values, which are ordered and have always the min/max ready.

- Insert items at start
- Percolate up / swap to keep heap property: swap if parent larger/smaller
- Remove min/max value = root
- Replace root with last item & percolate/swallow down: swap if children larger/smaller

To build a heap from a list: percolate down from node $i = \text{len}(\text{list})//2 : 0$, i.e., only non-leaves are percolated

3. Binary search tree implementation

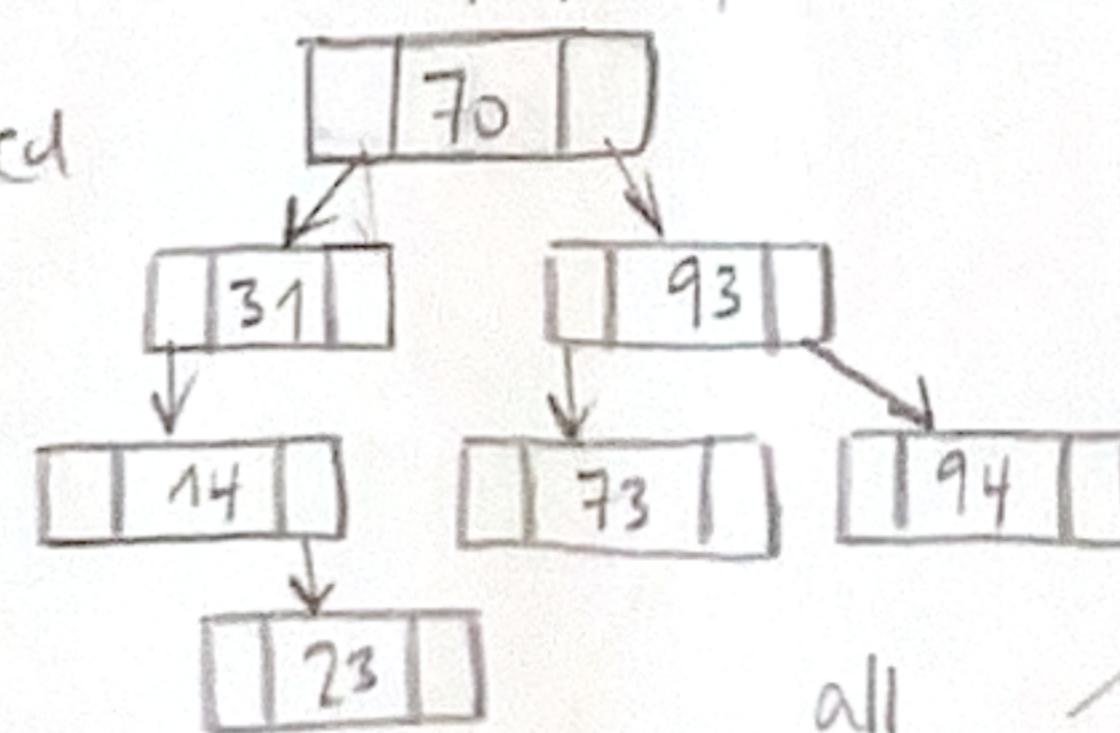
1. Check if a tree is a BST

2. Level-order print of a BST = breadth-first: FIFO list/queue used

3. Trim a BST for values (min, max)

Binary Search Trees There're 2 main key-value mappings < hash tables > BST

→ BST property: keys that are smaller than the parent are on the left subtree



an inorder traversal gives sorted values! Very important property!

RECALL: a tree can be implemented only with the Node class!

Typical we care: search values for a key in $O(\log n)$

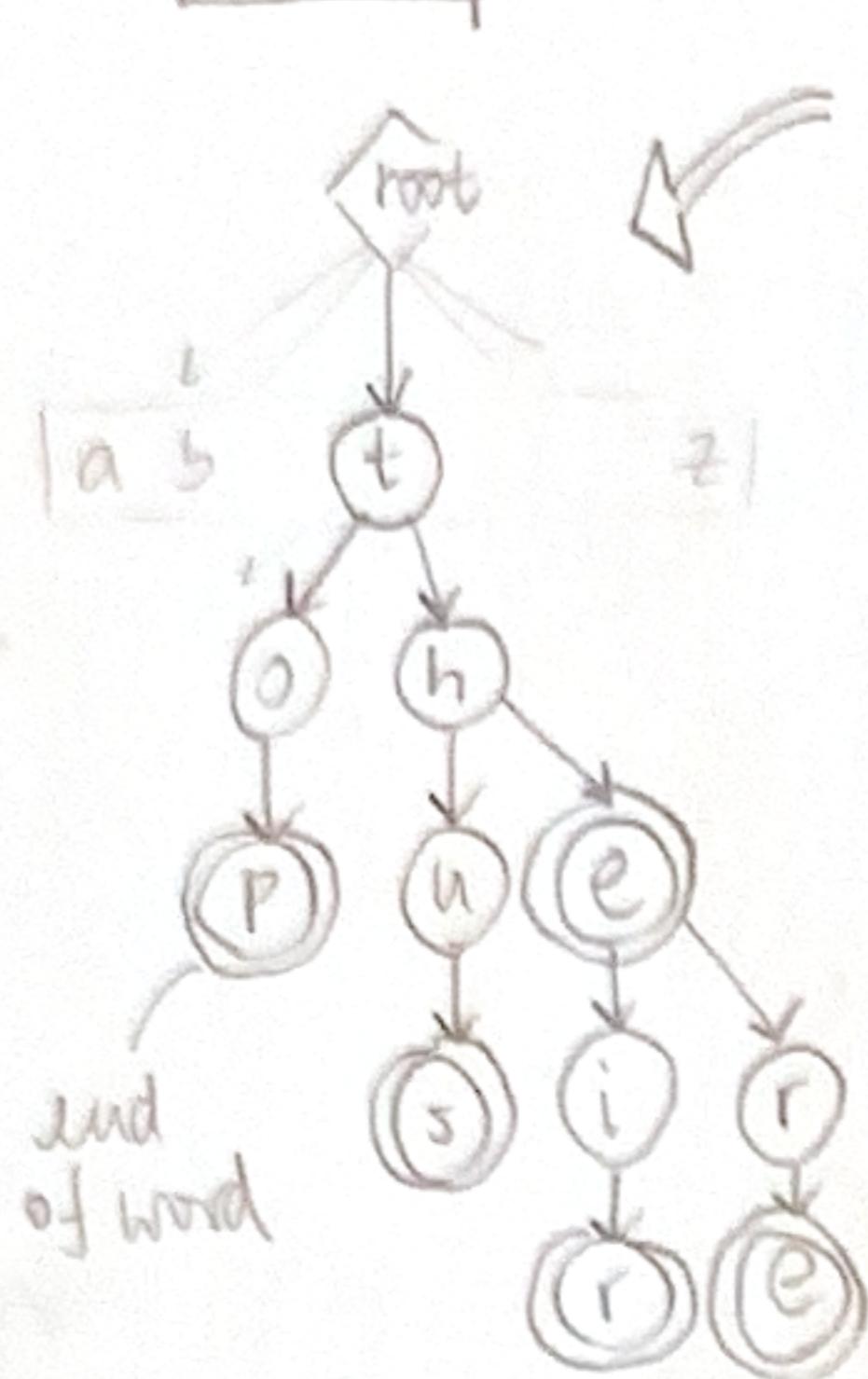
Important operations

- Insert a node: put()
- Retrieve a value given a key: get()
- Delete a node: remove()
 - 1) (case 1: node is leaf)
 - 2) (case 2: node has only 1 child)
 - 3) (case 3: - n - 2 children)

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0. Implement a trie
1. Count words in a trie
2. Print words in a trie
3. Sort words in a list
4. Check if a word can be a combination of words in a dictionary.

Tries



tree/graph-like structure well suited for lexicographical operations: order, check, ...

- Example: words in tree: [top, thus, the, their, there]
- Each node can have 26 children: 26 = number of characters
- Each inserted word is marked with end-of-word = True flag in the node of its last character.
- A unique word has a unique path
- Words with same prefixes/suffixes share parts of paths

Methods

- insert(word), $O(n)$, $n = \# \text{ chars}$
- search(word), $O(n)$, $n = \# \text{ chars}$
- delete(word), $O(n)$, $n = \# \text{ chars}$

for all 3, 3 sub-cases are defined

Subsets

0. Find the largest palindrome in a string of digits.

→ **Permutations**: count all possible ways; order matters

$$P(n, k) = \frac{n!}{(n-k)!}$$

{ A, B, C, D, E, F, G, H ($n=8$) : people
gold, silver, bronze ($k=3$) : medals

→ all possible medal permutations?

$$P(n=8, k=3) = \frac{8!}{5!} = 336$$

→ **Combinations w/o replacement**

$$C(n, k) = \frac{P(n, k)}{k!} = \frac{n!}{k!(n-k)!}$$

: count all possible unordered groups: order unimportant

$$C(n=8, k=3) = \frac{8!}{3!5!} = 56$$

→ **Combinations with replacement**

$$C_R(n, k) = C(n+k-1, k) = \frac{(n+k-1)!}{k!(n-1)!}$$

. Example: You have $n=5$ ice-creams: A, B, C, D, E
You can take $k=3$ scoops and can repeat flavour: 1, 2, 3

How many varieties do you have?

$$C_R(n=5, k=3) = \frac{7!}{3!4!} = 35$$

Search

0. Implement sequential search

→ **Sequential Search**

Visit all items in a collection linearly until we find our queried item; $O(n)$ worst case.
For unordered collections.

0. Implement binary search

→ **Binary Search**

For ordered collections. Iteratively/recursively check the midpoint of the proper half; worst case $O(\log n)$

0	1	2	3	4	5	6	7	8	9
17	20	26	31	44	54	55	65	77	93

related to binary search trees, but no tree!

Query: $q = 54$

$$\begin{aligned} 44 < 54 &\rightarrow \left[\frac{5+9}{2} \right] = 7 \\ \text{Start: } n/2 &= 4 \text{ (index)} \end{aligned}$$

Hash tables! (next page)

0. Implement a hash table

Hash tables

Tables / arrays which map keys to values using a hash function. O(1) for insert, delete, search!

Index	Key	Value
0		
1		
2		
...		
n		

index
hash(key)

keys,
aka.~items

values,
aka.~slots

table size

- Keys need to be hashable
- Common hash functions:
 - Remainder: $h(\text{key}) = \text{key} \% \text{table_size}$
 - Folding: split item and sum
 - Mid-square: square item and pick middle digits
 - String ordinals: split into chars and get ordinals
- Load factor: $\lambda = \frac{\text{num items}}{\text{table size}}$
- Collisions: when the slot of a new key is already taken;
there are several resolution approaches:
 - Open addressing: find next free slot
 - Linear probing: move sequentially until next free found $(i+1) \% \text{table_size}$
 - Quadratic probing: jump in quadratic steps to avoid clusters characteristic from
 - double hashing: apply new hash
 - Other approaches
 - Rehashing: create a new hash table; expensive
 - Chaining: allow storing a collection of key-value pairs per slot in a linked list.
- Python:
 - dictionaries are hash tables
 - make a class hashable with `__hash__` and `__eq__`