

# QCI402 - Mathematical Foundations for Quantum Computing

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The underpinnings of all scientific advancements is the ability to express natural phenomena with the art of Mathematics; this is no different for the subject of Quantum Computing. While the boundaries of quantum computing have been pushed beyond limits in theoretical terms on university blackboards, it has become of great interest to realize the theoretical computational power with the advances of hardware and technology.

However, these notes mainly concerns itself with the mathematical underpinnings of quantum computing that the course surrounds itself with. Mathematical Foundations for Quantum Computing takes a scaffolding approach designed to efficiently convey the required theoretical understanding of mathematics in order to able to learn quantum computing. As of writing, we are basing the notes on verison one of the textbook published in March 2025. In this text, we will primarily be using dirac notation for the expression of vectors, operators, and their interactions.

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# 1 Summation and Product Notations

This section primarily focuses on the common notations applied across mathematics to denote and shorten addition and product notation.

## 1.1 Summation over a single Variable

The sigma notation is defined as follows

$$\sum_{i=1}^n f(i)$$

where we use sigma  $\sum$  to represent the sum of a series. For example, the sum of all numbers in a series beginning with  $m$  and ending at index  $n$  is written as:

$$\sum_{i=m}^n a_i = a_m + a_{m+1} + a_{m+2} + \cdots + a_{n-1} + a_n$$

Sums can also be infinite, commonly seen when Sigma looks as follows:  $\sum_{i=m}^{\infty}$ . Infinite sums are either convergent or divergent. A few of the most common converging infinite sums are as follows:

$$\sum_{i=0}^{\infty} \frac{1}{2^i} = 1 + \frac{1}{2} + \frac{1}{4} + \cdots = 2$$

$$\sum_{i=0}^{\infty} \frac{1}{i^2} = \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \cdots = \frac{\pi^2}{6}$$

The first example is an infinite geometric series, and the sum of the first  $n$  terms is given by:

$$S_n = \sum_{i=0}^n \frac{1}{2^i} = \frac{1 - \frac{1}{2^n}}{1 - \frac{1}{2}}$$

As  $n \rightarrow \infty$ ,  $\frac{1}{2^n} \rightarrow 0$ . Consequently,  $S_n \rightarrow \frac{1}{1 - \frac{1}{2}} = 2$ . A rigorous proof of the second example requires extensive calculus and is not immediately obvious. While any mathematical symbol can be used for the index of a summation, it is more practical to use something other than  $i$  as in the context of complex numbers,  $i$  commonly denotes the complex number  $\sqrt{-1}$ . moreover, sume can also be specified using descriptions. For example,

$$\sum_{p \in P} f(p) \quad P \in \mathbb{N}'$$

where  $\mathbb{N}'$  is the set of all prime numbers. Summations can also contain parameters other than the index, which results in functions of those parameters. For example the discrete Fourier transform (DFT) is given by

$$\tilde{x}_k = \frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} x_n e^{-\frac{2\pi i}{N} kn}, \quad k = 0, 1, \dots, N-1$$

where  $x_n$  represents the  $N$  values index by  $n$  and  $\tilde{x}_k$  are the Fourier coefficients. Here,  $i$  is the imaginary numebr and  $N$  is a positive integer representing the dimension fo the DFT, of which we will cover in greater depth in Chapter 3. The following are some useful summation forumae commonly encountered in quantum computing:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i=1}^n i^3 = \left( \frac{n(n+1)}{2} \right)^2$$

$$\sum_{i=0}^n (a_0 + id) = (n+1) \left( a_0 + \frac{nd}{2} \right) \quad (\text{arithmetic series})$$

$$\sum_{i=0}^n a^i = \frac{1 - a^{n+1}}{1 - a} \quad (\text{geometric series})$$

$$(a+b)^n = \sum_{i=0}^n \binom{n}{i} a^{n-i} b^i \quad (\text{binomial theorem})$$

$$\frac{1}{1-x} = \sum_{n=0}^{\infty} x^n = 1 + x + x^2 + x^3 + \cdots \quad (|x| < 1)$$

$$\frac{1}{(1-x)^2} = \sum_{n=1}^{\infty} nx^{n-1} = 1 + 2x + 3x^2 + 4x^3 + \cdots \quad (|x| < 1)$$

$$\ln(1+x) = \sum_{n=1}^{\infty} \frac{(-1)^{n+1}}{n} x^n = x - \frac{x^2}{2} + \frac{x^3}{3} - \cdots \quad (|x| < 1)$$

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \cdots$$

$$\sin x = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n+1)!} x^{2n+1} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \cdots$$

$$\cos x = \sum_{n=0}^{\infty} \frac{(-1)^n}{(2n)!} x^{2n} = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \cdots$$

Below are also a list of the common summations rules and manipulations:

$$\sum_{i=m}^n a_i = \sum_{j=m}^n a_j \quad (\text{change of index variable})$$

$$\sum_{i=s}^t f(i) = \sum_{n=s}^t f(n) \quad (\text{change of index variable})$$

$$\sum_{n=s}^t f(n) = \sum_{n=s}^j f(n) + \sum_{n=j+1}^t f(n) \quad (\text{splitting a sum})$$

$$\sum_{n=s}^t f(n) = \sum_{n=0}^{t-s} f(t-n) \quad (\text{reverse order})$$

$$\sum_{n=s}^t f(n) = \sum_{n=s+p}^{t+p} f(n-p) \quad (\text{index shift})$$

$$\sum_{n=s}^t a \cdot f(n) = a \cdot \sum_{n=s}^t f(n) \quad (\text{distributivity})$$

$$\sum_{n=s}^t f(n) \pm \sum_{n=s}^t g(n) = \sum_{n=s}^t (f(n) \pm g(n)) \quad (\text{commutativity})$$

## 1.2 Products and other Notations

Similar to the  $\sum$  notation for addition, the  $\prod$  (Pi) symbol is also more commonly used to denote the product of a series of terms. In this

$$\prod_{i=m}^n a_i = a_m \cdot a_{m+1} \cdot a_{m+2} \cdot \dots \cdot a_{n-1} \cdot a_n$$

for example, the factorial of  $n$  is expressed as

$$\prod_{i=0}^n i = n!$$

and the relationship between  $\sum$  and  $\prod$ , which are

$$b^{\sum_{n=s}^t f(n)} = \prod_{n=s}^t b^{f(n)}$$

$$\sum_{n=s}^t \log_b f(n) = \log_b \prod_{n=s}^t f(n)$$

It is worth noting that in quantum computing and linear algebra, there are a few special notations such as the modulo-2 sum (bitwise XOR), or in other contexts the direct sum of linear spaces, represented by  $\oplus$ , and the tensor product represented by  $\otimes$ .

## 1.3 Summation over Multiple Variables

The double summation over a rectangular array is given by

$$\begin{aligned} \sum_{i=1, j=1}^{n_1, n_2} a_{i,j} &= \sum_{i=1}^{n_1} \sum_{j=1}^{n_2} a_{i,j} = \sum_{j=1}^{n_2} \sum_{i=1}^{n_1} a_{i,j} \\ &= a_{1,1} + a_{1,2} + a_{1,3} + a_{1,4} + \dots + a_{1,n_2} \\ &\quad + a_{2,1} + a_{2,2} + a_{2,3} + a_{2,4} + \dots + a_{2,n_2} \\ &\quad + a_{3,1} + a_{3,2} + a_{3,3} + a_{3,4} + \dots + a_{3,n_2} \\ &\quad + a_{4,1} + a_{4,2} + a_{4,3} + a_{4,4} + \dots + a_{4,n_2} \\ &\quad + \dots \\ &\quad + a_{n_1,1} + a_{n_1,2} + a_{n_1,3} + a_{n_1,4} + \dots + a_{n_1,n_2} \end{aligned}$$

Here,  $\sum_{i=1}^{n_1} \sum_{j=1}^{n_2}$  represents summing over each row first and then summing the results, while  $\sum_{j=1}^{n_2} \sum_{i=1}^{n_1}$  will represent summing over the columns and then summing those results. The term  $\sum_{i=1, j=1}^{n_1, n_2} a_{i,j}$  represents the summation over the rectangular array, irrespective of the order. The product of two sums can be expanded into a double sum as follows:

$$\begin{aligned} \left( \sum_{i=1}^m a_i \right) \left( \sum_{j=1}^n b_j \right) &= (a_1 + a_2 + \dots + a_m)(b_1 + b_2 + \dots + b_n) \\ &= a_1 b_1 + a_1 b_2 + a_1 b_3 + a_1 b_4 + \dots + a_1 b_n \\ &\quad + a_2 b_1 + a_2 b_2 + a_2 b_3 + a_2 b_4 + \dots + a_2 b_n \\ &\quad + a_3 b_1 + a_3 b_2 + a_3 b_3 + a_3 b_4 + \dots + a_3 b_n \\ &\quad + \dots \\ &\quad + a_m b_1 + a_m b_2 + a_m b_3 + a_m b_4 + \dots + a_m b_n \\ &= \sum_{i=1}^m \sum_{j=1}^n a_i b_j = \sum_{i=1}^m a_i \sum_{j=1}^n b_j \end{aligned}$$

which is actually rather intuitive given how the expansion of the standard expansion of the term  $(a+b)^2$  plays out, a more elementary application of the distributive property which the above equation generalizes over. For a triangular matrix, in this case the lower triangular matrix, the sum is given by

$$\begin{aligned}
\sum_{1 \leq j \leq n} a_{i,j} &= \sum_{i=1}^n \sum_{j=1}^i a_{i,j} = \sum_{j=1}^n \sum_{i=j}^n a_{i,j} = \sum_{j=0}^{n-1} \sum_{j=1}^{n-j} a_{i+j,i} \\
&= a_{1,1} \\
&\quad + a_{2,1} + a_{2,2} \\
&\quad + a_{3,1} + a_{3,2} + a_{3,3} \\
&\quad + a_{4,1} + a_{4,2} + a_{4,3} + a_{4,4} \\
&\quad + \dots \\
&\quad + a_{n,1} + a_{n,2} + a_{n,3} + a_{n,4} + \dots + a_{n,n}
\end{aligned}$$

where the term  $\sum_{1 \leq j \leq n} a_{i,j}$  denotes the summation over all elements in a lower triangular array including the diagonal. The first notation variation will sum up each row to the  $i$ th element then aggregate while the second notation sums each column starting from the  $j$ th element downwards then aggregate the sums. The final expression will sum along the diagonal where  $j = 0$  represents the main diagonal and  $j = n - 1$  is the first off-diagonal, which is a single term.

**Example.** Say we would like to expand the product of  $(1 + x_i)$  from 1 to  $n$ . We have

$$\prod_{i=1}^n (1 + x_i) = 1 + \sum_{k=1}^n \left( \sum_{1 \leq i_1 < \dots < i_k \leq n} \prod_{j=1}^k x_{i_j} \right)$$

This formula represents the *multinomial expansion* of a product. When you expand the equation by hand, you get the product

$$\prod_{i=1}^n (1 + x_i) = (1 + x_1)(1 + x_2) \cdots (1 + x_n)$$

If we break this down, we see that the outer summation  $\sum_{k=1}^n$  will go through each possible summation size in terms of the variables in question, and that the inner summation  $\sum_{1 \leq i_1 < \dots < i_k \leq n}$  will iterate through each possible unique product of the variables. while ensuring that they are unique. Not sure how this works, but if all  $x_i$  are the same, then we see that the equation actually simplifies to a subset of the binomial theorem

$$(1 + x)^n = \sum_{k=0}^n \binom{n}{k} x^k$$

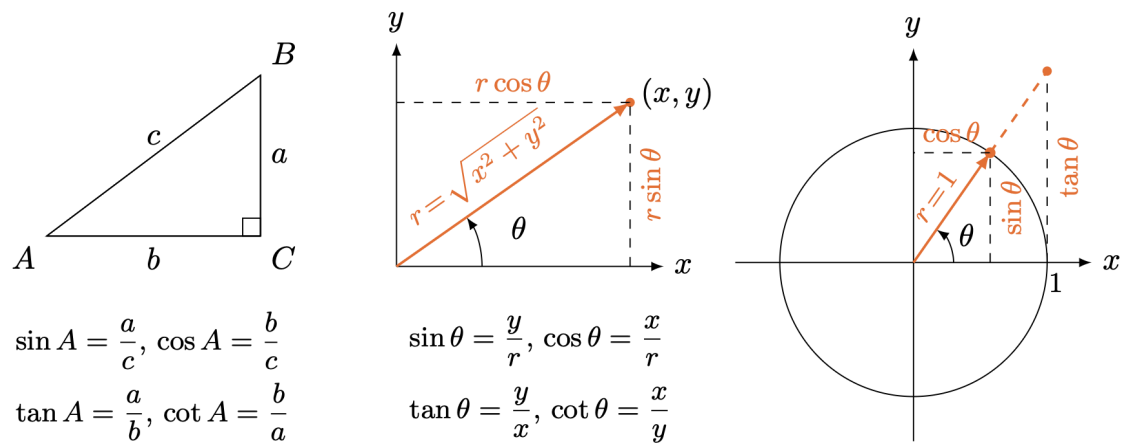
where  $\binom{n}{k}$  is the binomial coefficient representing the number of ways to choose  $k$  elements from a set of  $n$  distinct elements.

# 2 Trigonometry

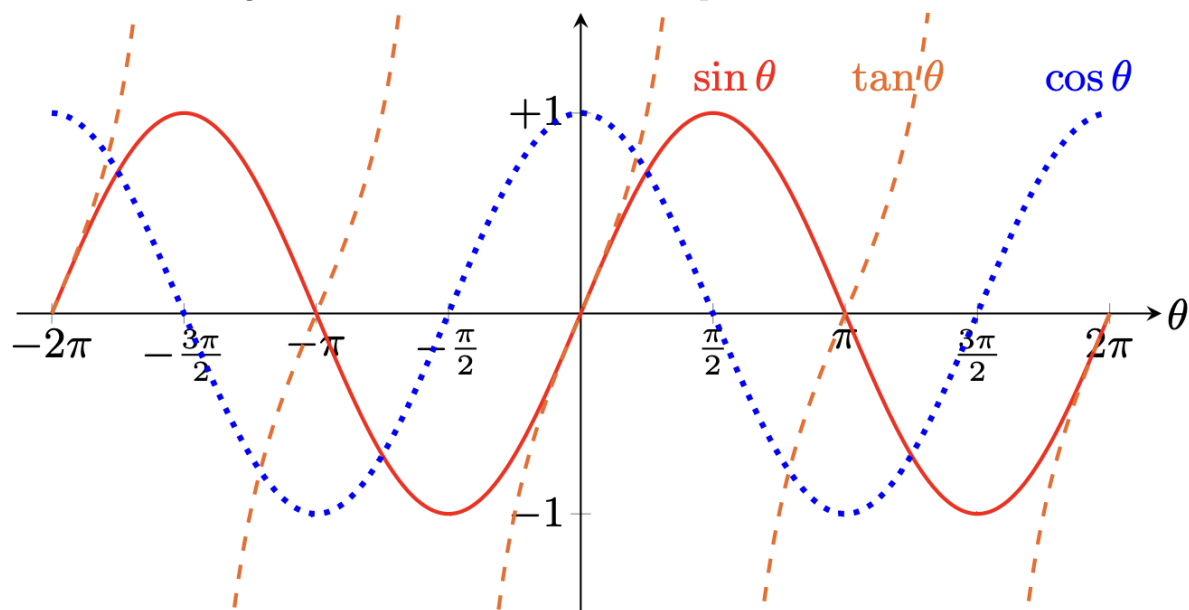
You can't escape this.

## 2.1 Definitions

I mean, where do I start? The basic trigonometric functions are defined as the ratios between the angles of a right triangle. I will not show how these ratios remain the same given the same angle, nor will I go into great mathematical detail of how to prove these items. However, we still have to go over this. Don't ask me why.



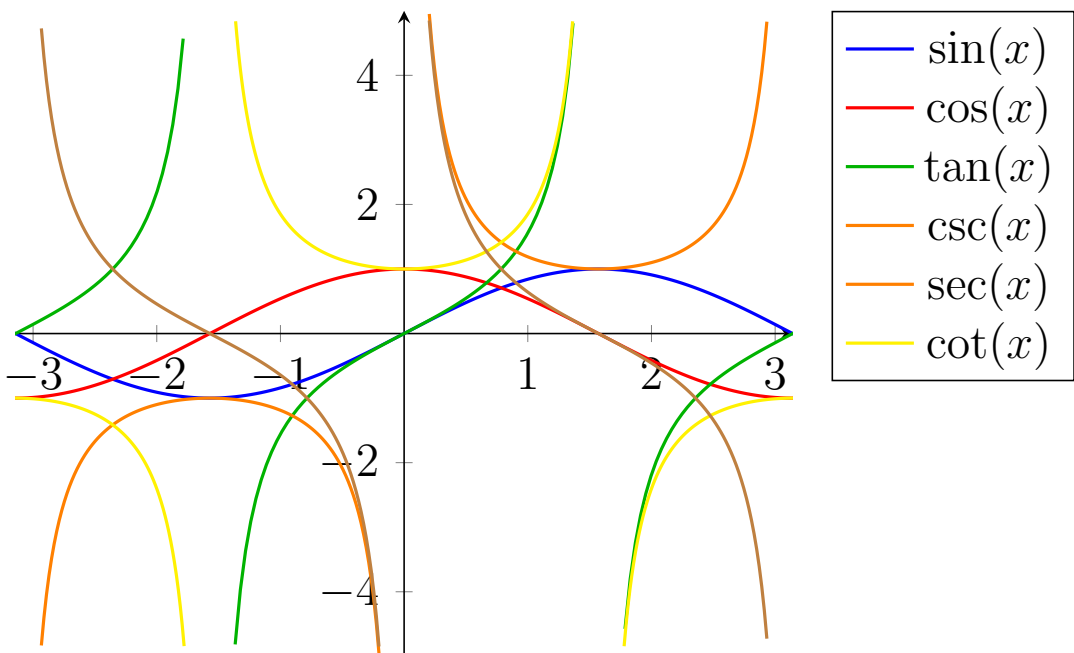
The functions of trigonometric functions can also be plotted out as follows:



## 2.2 Basic Properties and Inverse Functions

	$\sin \theta$	$\cos \theta$	$\tan \theta$	$\csc \theta$	$\sec \theta$	$\cot \theta$
Definition	$y/r$	$x/r$	$y/x$	$r/y$	$r/x$	$x/y$
Period	$2\pi$	$2\pi$	$\pi$	$2\pi$	$2\pi$	$\pi$
Range	$[-1, 1]$	$[-1, 1]$	$(-\infty, \infty)$	$(-\infty, -1] \cup [1, \infty)$	$(-\infty, \infty)$	$(-\infty, \infty)$
Zeros	$n\pi$	$(n + \frac{1}{2})\pi$	$n\pi$			$(n + \frac{1}{2})\pi$
Poles			$(n + \frac{1}{2})\pi$	$n\pi$	$(n + \frac{1}{2})\pi$	$n\pi$

Note:  $n$  is an integer.



We can also see that there are certain useful symmetric properties of the trigonometric functions

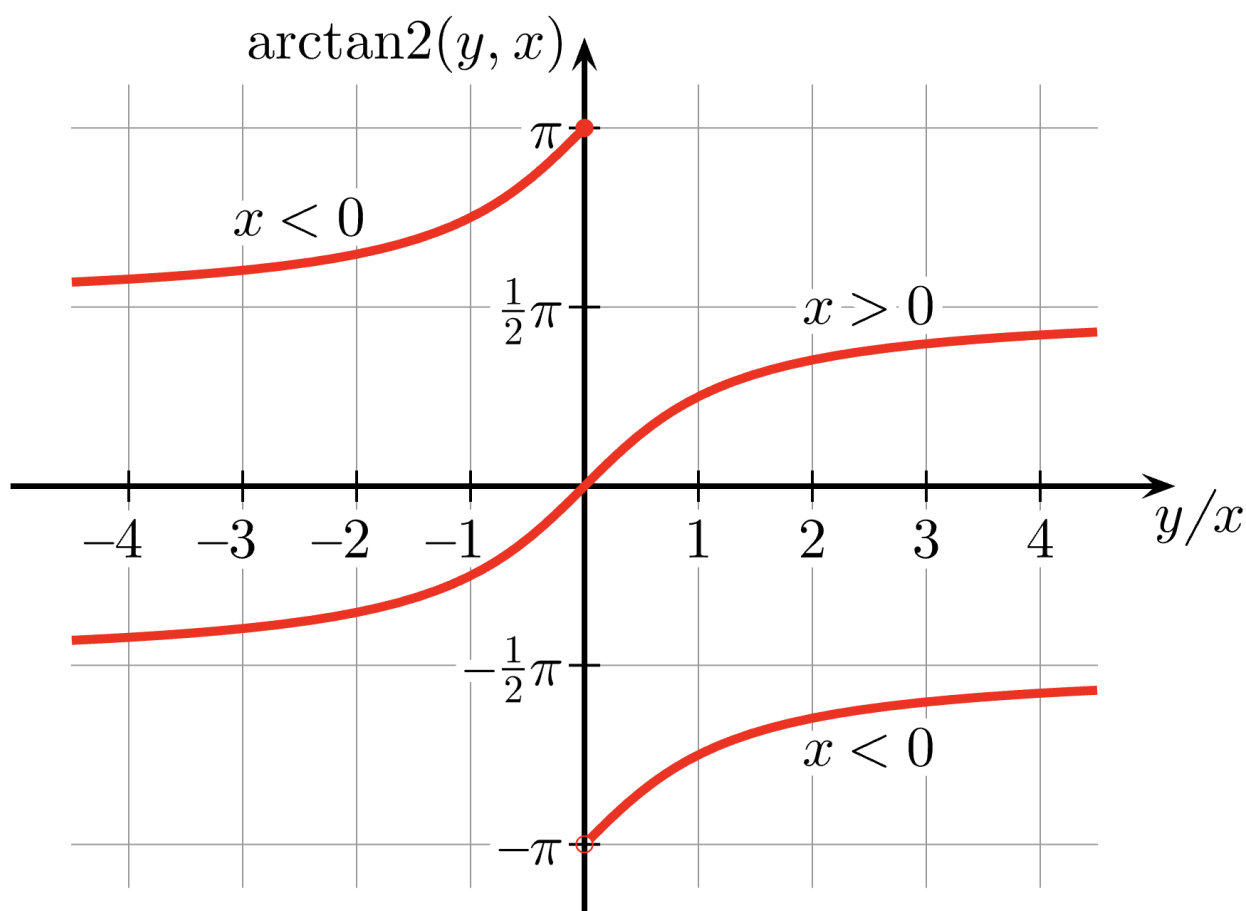
$$\begin{array}{lll}
 \sin(-\theta) = -\sin(\theta) & \sin(\pi - \theta) = \sin(\theta) & \sin(\pi + \theta) = -\sin(\theta) \\
 \cos(-\theta) = \cos(\theta) & \cos(\pi - \theta) = -\cos(\theta) & \cos(\pi + \theta) = -\cos(\theta) \\
 \tan(-\theta) = -\tan(\theta) & \tan(\pi - \theta) = -\tan(\theta) & \tan(\pi + \theta) = \tan(\theta)
 \end{array}$$

There are also some common inverse functions associated with the functions.

Function	sin	cos	tan	csc	sec	cot
Inverse	$\sin^{-1}$ arcsin	$\cos^{-1}$ arccos	$\tan^{-1}$ arctan	$\csc^{-1}$	$\sec^{-1}$	$\cot^{-1}$
Domain	$[-1, 1]$	$[-1, 1]$	$(-\infty, \infty)$	$(-\infty, -1] \cup [1, \infty)$	$(-\infty, \infty)$	$(-\infty, \infty)$
Range	$[-\frac{\pi}{2}, \frac{\pi}{2}]$	$[0, \pi]$	$(-\frac{\pi}{2}, \frac{\pi}{2})$	$[-\frac{\pi}{2}, \frac{\pi}{2}] \setminus \{0\}$	$[0, \pi] \setminus \{\frac{\pi}{2}\}$	$(-\frac{\pi}{2}, \frac{\pi}{2})$

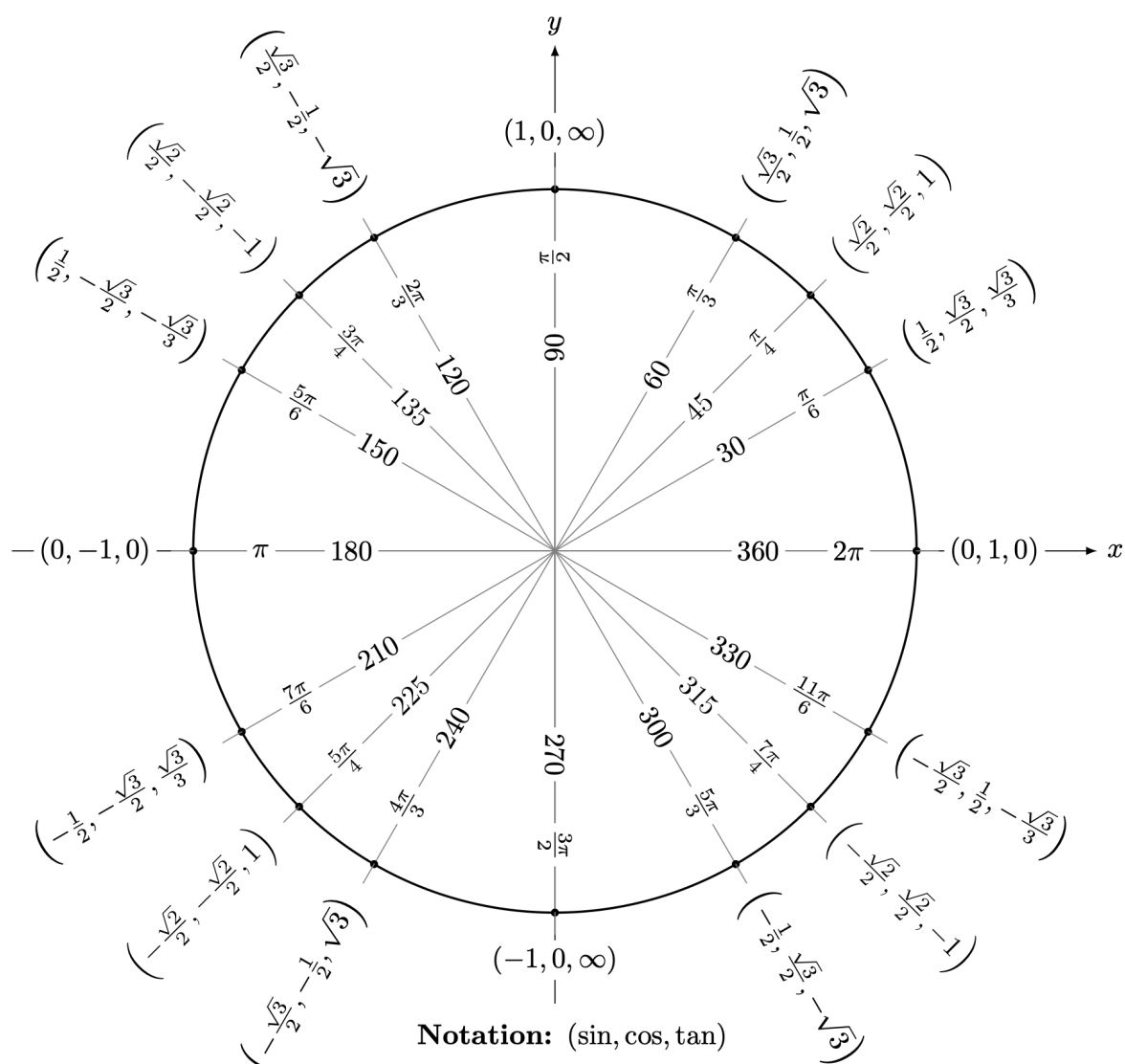
But also some interesting extensions of the commonly known inverse functions, with the example being arctan2, a function that effectively doubles the domain of the function while preserving its properties for the purpose of, say, converting values from cartesian to spherical coordinates for the azimuthal angle  $\phi$ .

$$\arctan2(y, x) = \begin{cases} \arctan\left(\frac{y}{x}\right) & \text{if } x > 0, \\ \arctan\left(\frac{y}{x}\right) + \pi & \text{if } x < 0 \text{ and } y \geq 0, \\ \arctan\left(\frac{y}{x}\right) - \pi & \text{if } x < 0 \text{ and } y < 0, \\ +\frac{\pi}{2} & \text{if } x = 0 \text{ and } y > 0, \\ -\frac{\pi}{2} & \text{if } x = 0 \text{ and } y < 0, \\ 0 & \text{if } x = 0 \text{ and } y = 0. \end{cases}$$



## 2.3 Special Angles and Function Values

There are a few special angles that are worth remembering for the trigonometric functions mentioned above, given by the wheel below.





## 2.4 Trigonometric Identities

### Reciprocal and Quotient Identities

$$\csc \theta = \frac{1}{\sin \theta} \qquad \sec \theta = \frac{1}{\cos \theta} \qquad \cot \theta = \frac{1}{\tan \theta}$$

$$\tan \theta = \frac{\sin \theta}{\cos \theta}, \qquad \cot \theta = \frac{\cos \theta}{\sin \theta}$$

### Cofunction Identities

$$\begin{aligned} \sin\left(\frac{\pi}{2} - \theta\right) &= \cos \theta, & \cos\left(\frac{\pi}{2} - \theta\right) &= \sin \theta \\ \tan\left(\frac{\pi}{2} - \theta\right) &= \cot \theta, & \cot\left(\frac{\pi}{2} - \theta\right) &= \tan \theta \\ \sec\left(\frac{\pi}{2} - \theta\right) &= \csc \theta, & \csc\left(\frac{\pi}{2} - \theta\right) &= \sec \theta \end{aligned}$$

### Pythagorean Identities

$$\begin{aligned} \sin^2 \theta + \cos^2 \theta &= 1 \\ 1 + \tan^2 \theta &= \sec^2 \theta \\ 1 + \cot^2 \theta &= \csc^2 \theta \end{aligned}$$

### Even–Odd Symmetry

$$\begin{aligned} \sin(-\theta) &= -\sin \theta, & \cos(-\theta) &= \cos \theta, & \tan(-\theta) &= -\tan \theta \\ \csc(-\theta) &= -\csc \theta, & \sec(-\theta) &= \sec \theta, & \cot(-\theta) &= -\cot \theta \end{aligned}$$

### Sum and Difference Formulas

$$\begin{aligned} \sin(\alpha \pm \beta) &= \sin \alpha \cos \beta \pm \cos \alpha \sin \beta \\ \cos(\alpha \pm \beta) &= \cos \alpha \cos \beta \mp \sin \alpha \sin \beta \\ \tan(\alpha \pm \beta) &= \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta} \end{aligned}$$

### Double Angle Formulas

$$\begin{aligned} \sin(2\theta) &= 2 \sin \theta \cos \theta \\ \cos(2\theta) &= \cos^2 \theta - \sin^2 \theta \\ &= 2 \cos^2 \theta - 1 \\ &= 1 - 2 \sin^2 \theta \\ \tan(2\theta) &= \frac{2 \tan \theta}{1 - \tan^2 \theta} \end{aligned}$$

### Half Angle Formulas

$$\begin{aligned} \sin^2\left(\frac{\theta}{2}\right) &= \frac{1 - \cos \theta}{2} \\ \cos^2\left(\frac{\theta}{2}\right) &= \frac{1 + \cos \theta}{2} \\ \tan\left(\frac{\theta}{2}\right) &= \frac{\sin \theta}{1 + \cos \theta} = \frac{1 - \cos \theta}{\sin \theta} \end{aligned}$$

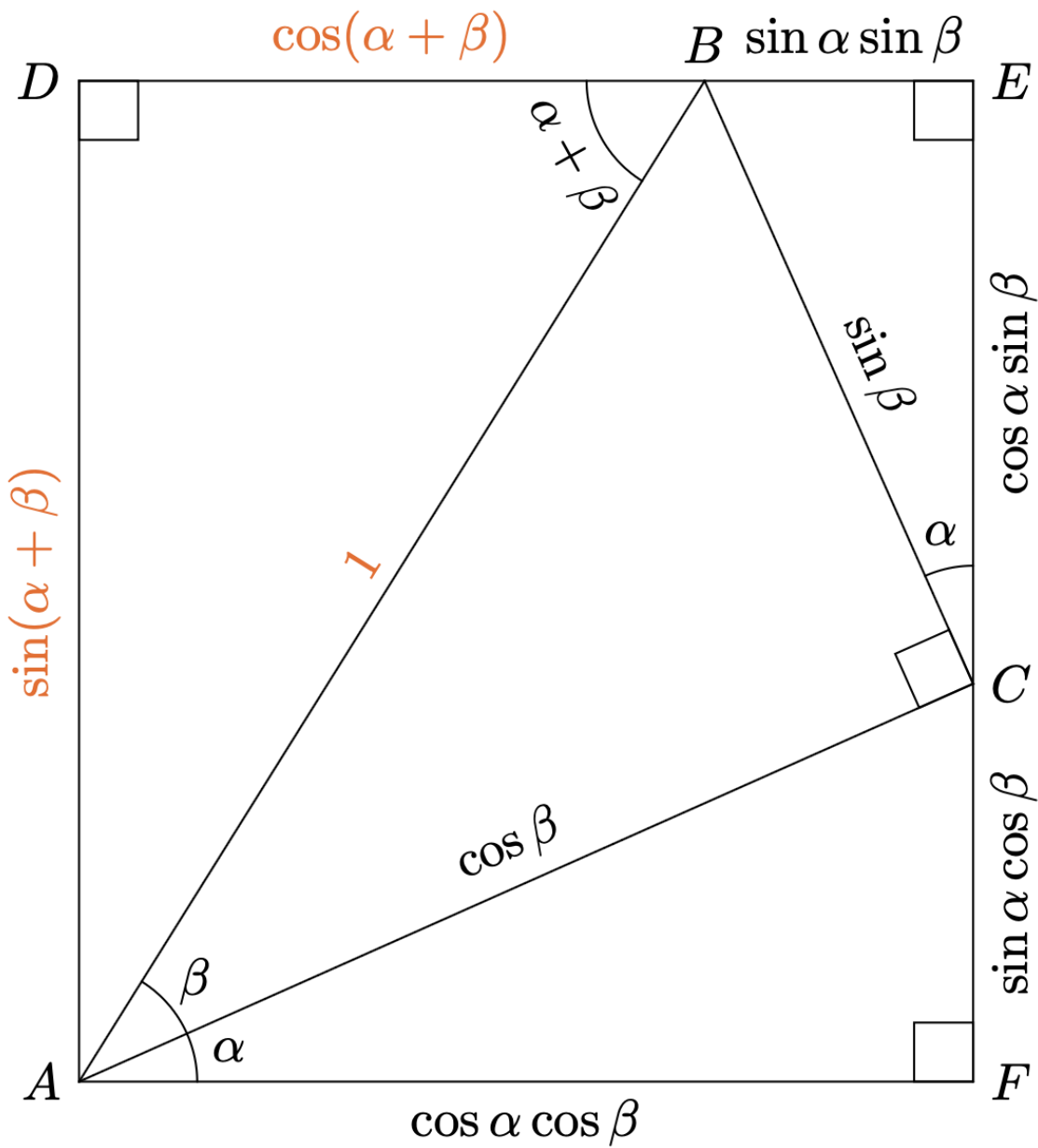
Product-to-Sum Identities

$$\begin{aligned}\sin \alpha \sin \beta &= \frac{1}{2} \left[ \cos(\alpha - \beta) - \cos(\alpha + \beta) \right] \\ \cos \alpha \cos \beta &= \frac{1}{2} \left[ \cos(\alpha - \beta) + \cos(\alpha + \beta) \right] \\ \sin \alpha \cos \beta &= \frac{1}{2} \left[ \sin(\alpha + \beta) + \sin(\alpha - \beta) \right]\end{aligned}$$

Sum-to-Product Identities

$$\begin{aligned}\sin \alpha + \sin \beta &= 2 \sin \left( \frac{\alpha + \beta}{2} \right) \cos \left( \frac{\alpha - \beta}{2} \right) \\ \sin \alpha - \sin \beta &= 2 \cos \left( \frac{\alpha + \beta}{2} \right) \sin \left( \frac{\alpha - \beta}{2} \right) \\ \cos \alpha + \cos \beta &= 2 \cos \left( \frac{\alpha + \beta}{2} \right) \cos \left( \frac{\alpha - \beta}{2} \right) \\ \cos \alpha - \cos \beta &= -2 \sin \left( \frac{\alpha + \beta}{2} \right) \sin \left( \frac{\alpha - \beta}{2} \right)\end{aligned}$$

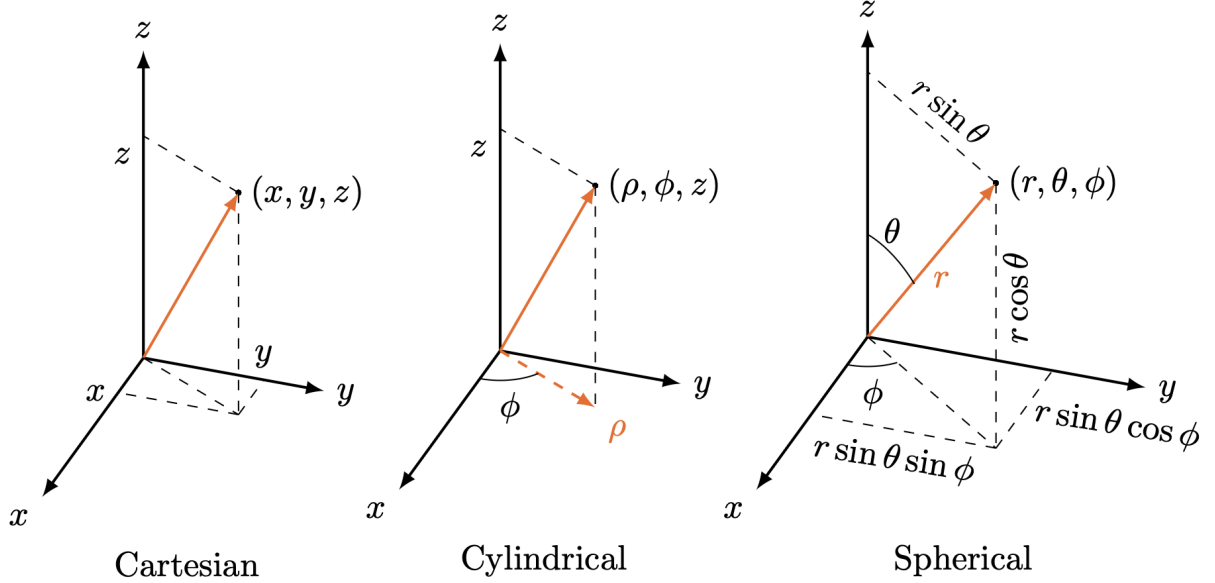
A rather nice photo to sum the seccion up is by relating the angles to each other using the following image



2.5 The Spherical Coordinate System

While the expansion of the cartesian coordinate system into three dimensions is the logical linear expansion to take, some interesting basis begin to form if we consider the angles as a

unit of measurement in three dimensions



The cylindrical coordinate system elevates the polar coordinate system into three dimensions with the addition of the  $z$  axis, while the spherical system is a more organic translation of the polar coordinate system into three dimensions. While the former is best suited for describing not only cylinder-like structures yet also helices, while the latter is best suited for rotations in three-dimensional space, which is particularly potent in the field of quantum computing.

For example, the spherical coordinate system is commonly used to represent qubit states on the Bloch sphere, employing a radius ( $r$ ), a polar angle ( $\theta$ ), and an azimuthal angle ( $\phi$ ) to represent a point in three-dimensional space. The azimuthal angle  $\phi$  is measured in the  $xy$ -plane from the positive  $x$ -axis with common values ranging from  $(-\pi, \pi]$  or  $(0, 2\pi]$ . The polar angle is commonly measured from the positive  $z$  axis towards the  $xy$ -plane, with values ranging from  $[0, \pi]$ . Note that  $r \in \mathbb{R}$ , meaning that we can cover the other half of the range simply by flipping the sign around.

Conversion is relatively simple, with conversion to and from spherical to cartesian being as follows

$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \phi &= \arctan2(y, x) \\ z &= r \cos \theta & \theta &= \arccos \frac{z}{r} \end{aligned}$$

where  $\arctan2$  was previously defined as a optimal inverse mapping onto the range of  $[-\pi, \pi]$ . We also have a few definitions

**Definition 2.1.** The Law of Sines is defined as

$$\frac{\sin A}{a} = \frac{\sin B}{b} = \frac{\sin C}{c}$$

**Definition 2.2.** The Law of Cosines is defined as

$$a^2 = b^2 + c^2 - 2bc \cos A$$

which can be rewritten as

$$\cos A = \frac{b^2 + c^2 - 2bc}{a^2}$$

**Definition 2.3.** The Law of Tangents is defined as

$$\frac{a-b}{a+b} = \frac{\tan \frac{1}{2}(A-B)}{\tan \frac{1}{2}(A+B)}$$

### 3 Complex Numbers

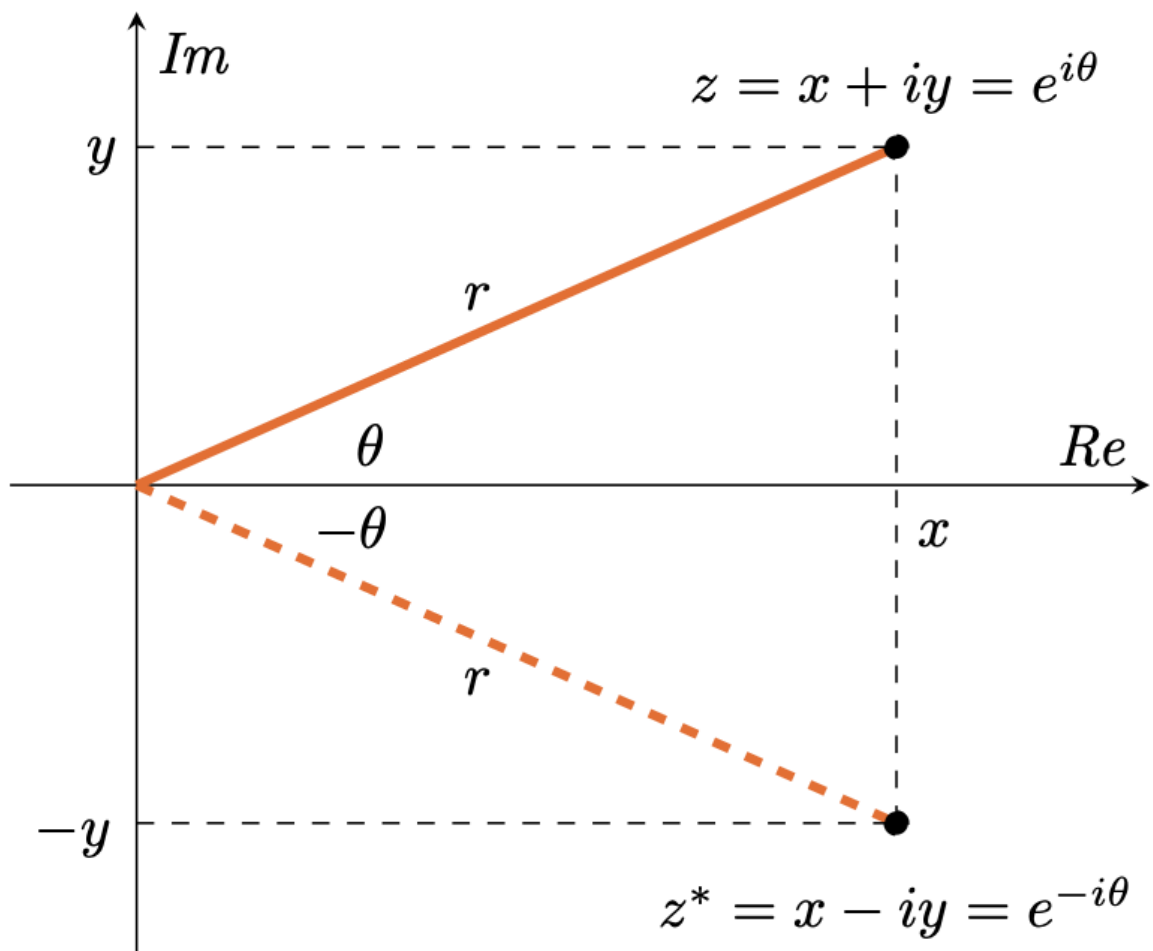
We consider numbers to be complex when they compose of a real and imaginary part, and they are not only fundamental to a complete understanding of algebra and mathematics as a whole, but also form the backbone of quantum mechanics, and, by extension, quantum computing. Mastering complex numbers is like Rosie mastering the rivet gun, so we have to study it.

#### 3.1 Cartesian Form

**Definition 3.1.** A complex number  $z$  is defined as

$$z = x + iy, \quad x, y \in \mathbb{R}, \quad i^2 = -1$$

This is called the **cartesian form** of the complex number  $z$  and corresponds to a point in the two-dimensional complex plane. We commonly refer to  $i$  as the imaginary unit. It may seem ironic that we need imaginary numbers in quantum computing, or that we really need the imaginary number. Take it as you may.



Complex numbers not motivated by quantum computing. In the numbers system, we have the real numbers  $\mathbb{N}$ , the integer numbers  $\mathbb{Z}$ , the real numbers  $\mathbb{R}$ , and the complex numbers  $\mathbb{C}$ . the set incursions go this way  $\mathbb{N} \subset \mathbb{Z} \subset \mathbb{R} \subset \mathbb{C}$ , all supersets of the preceeding set. The set of all complex numbers  $\mathbb{C}$  is closed over all algebraic operations, which include addition, subtraction, multiplication, division, power, and root, and is considered the superset of all numbers.

**Definition 3.2.** The basic components of a complex number are defined as follows.

$$\Re(z) = x, \Im(z) = y$$

which are the real and imaginary components of a complex number  $z$ . Of course the complex number itself has a few interesting properties, such as  $i^2 = -1$ ,  $i^3 = -i$ , and  $i^4 = 1$ . The complex conjugate of a complex number  $z$  is defined as

$$z^* = x - iy$$

defined as inverting the sign of the imaginary component. We can express the modulus (vector length) and argument (angle with respect to the real axis) (which are  $r$  and  $\theta$  in polar coordinates), as follows:

$$r = |z| = \sqrt{zz^*} = \sqrt{x^2 + y^2}$$

A very convenient property derived from algebra is that  $zz^* = x^2 + y^2$ .

$$\theta = \arg(z)$$

For the angle, we note that

$$\tan \frac{y}{x} \implies \arctan2(x, y)$$

where  $\arctan2$  has been defined in the previous section.

**Example.** Given  $z = 1 + \sqrt{3}i$ , we have

$$z^* = 1 - \sqrt{3}i.$$

$$|z| = \sqrt{1^2 + (\sqrt{3})^2} = 2.$$

$$zz^* = (1 + \sqrt{3}i)(1 - \sqrt{3}i) = 1 - (\sqrt{3}i)^2 = 1 + 3 = 4 = |z|^2.$$

$$\theta = \arctan\left(\frac{\sqrt{3}}{1}\right) = \frac{\pi}{3}.$$

## 3.2 Exponential Form

Now it is worth noting that while we commonly write  $z = x + iy$  to represent a complex number, we like to use the following definitions of the complex number in polar form to represent a complex number itself, defined as  $z = r(\cos \theta + i \sin \theta)$ . However, multiplication and its inverse operation, division, becomes unnecessarily difficult given the presence of another notation, namely **exponential form**.

**Definition 3.3.** The exponential/euler forms of the complex numbers can be thought of as a circular form of the function  $z = x + iy$ . In polar coordinates, we can rewrite this number as

$$z = r \cos \theta + i \sin \theta, \quad r \in \mathbb{R}$$

Conversely, the conversion between cartesian and polar are

$$x = r \cos \theta \quad y = r \sin \theta$$

The formula for  $z$  above can be rewritten as

$$z = re^{i\theta}$$

**Theorem 3.4.** Euler's formula states that for any complex number  $z = r \cos \theta + i \sin \theta$ , we have:

$$e^{i\theta} = \cos \theta + i \sin \theta$$

*Proof.* Euler's formula can be proven using the Taylor series expansion for the functions:

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \dots$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \dots$$

if we replace  $x$  in the formula for  $e^x$  with  $e^{ix}$ , we then have

$$\begin{aligned} e^{ix} &= 1 + (ix) + \frac{(ix)^2}{2!} + \frac{(ix)^3}{3!} + \frac{(ix)^4}{4!} + \frac{(ix)^5}{5!} + \frac{(ix)^6}{6!} + \frac{(ix)^7}{7!} + \dots \\ &= 1 + ix - \frac{x^2}{2!} - i\frac{x^3}{3!} + \frac{x^4}{4!} + i\frac{x^5}{5!} - \frac{x^6}{6!} - i\frac{x^7}{7!} + \dots \\ &= \left(1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots\right) + \left(ix - i\frac{x^3}{3!} + i\frac{x^5}{5!} - i\frac{x^7}{7!} + \dots\right) \\ &= \cos \theta + i \sin \theta \end{aligned}$$

which sums up essentially what mathematicians call the most beautiful proof man has known. If this proof were a female robot, mathematicians would compose harmonic waves and produce digital flowers in L<sup>A</sup>T<sub>E</sub>X to please Euler’s genius. □

As noted before, we know that the set of all algebraic operations is well defined and closed on the set of all complex numbers  $\mathbb{C}$ . Addition will be easier in cartesian form, while multiplication will be considerably simpler in exponential form. Conversion between the two is also not difficult:

	Cartesian Form	Exponential Form
	$z = x + iy$	$z = re^{i\theta}$
Conjugate	$z^* = x - iy$	$z^* = re^{-i\theta}$
Modulus	$ z  = \sqrt{zz^*} = \sqrt{x^2 + y^2}$	$ z  = r$
Conversion	$x = r \cos \theta$	$r = \sqrt{x^2 + y^2}$
	$y = r \sin \theta$	$\theta = \arctan2(y, x)$

### 3.3 Basic Operations

As we touched upon earlier, the set of all complex numbers are closed on operations of addition and subtraction:

$$z_1 + z_2 = (x_1 + x_2) + i(y_1 + y_2).$$

$$z_1 - z_2 = (x_1 - x_2) + i(y_1 - y_2).$$

As well on multiplication and division. For

$$z_1 = r_1e^{i\theta_1} \qquad z_2 = r_2e^{i\theta_2}$$

We have:

$$z_1 \cdot z_2 = r_1r_2e^{i(\theta_1+\theta_2)}.$$

$$\frac{z_1}{z_2} = \frac{r_1e^{i\theta_1}}{r_2e^{i\theta_2}} = \frac{r_1}{r_2}e^{i(\theta_1-\theta_2)}.$$

We also have the following properties for the conjugates of complex numbers

$$\begin{aligned} |z| &= |z^*| \\ (z_1 \pm z_2)^* &= z_1^* \pm z_2^* \\ (z_1 \cdot z_2)^* &= z_1^* \cdot z_2^* \\ (z_1/z_2)^* &= z_1^*/z_2^* \\ (z^x)^* &= (z^*)^x \quad x \in \mathbb{R} \\ (x^z)^* &= x^{z^*} \quad x \in \mathbb{R} \end{aligned}$$

Where the last two are not immediately obvious. To prove that  $(z^x)^* = (z^*)^x$ , it is useful to write out  $z$  using the complex notation  $re^{i\theta}$ , and the last property is best proven using the identity  $a^b = e^{b \ln a}$ . As for powers and roots of complex numbers, we have

**Theorem 3.5.** De Moivre's theorem states that

$$(\cos \theta + i \sin \theta)^s = \cos s\theta + i \sin s\theta$$

which is conveniently derived from the fact that  $z^s = r^s e^{is\theta}$ .

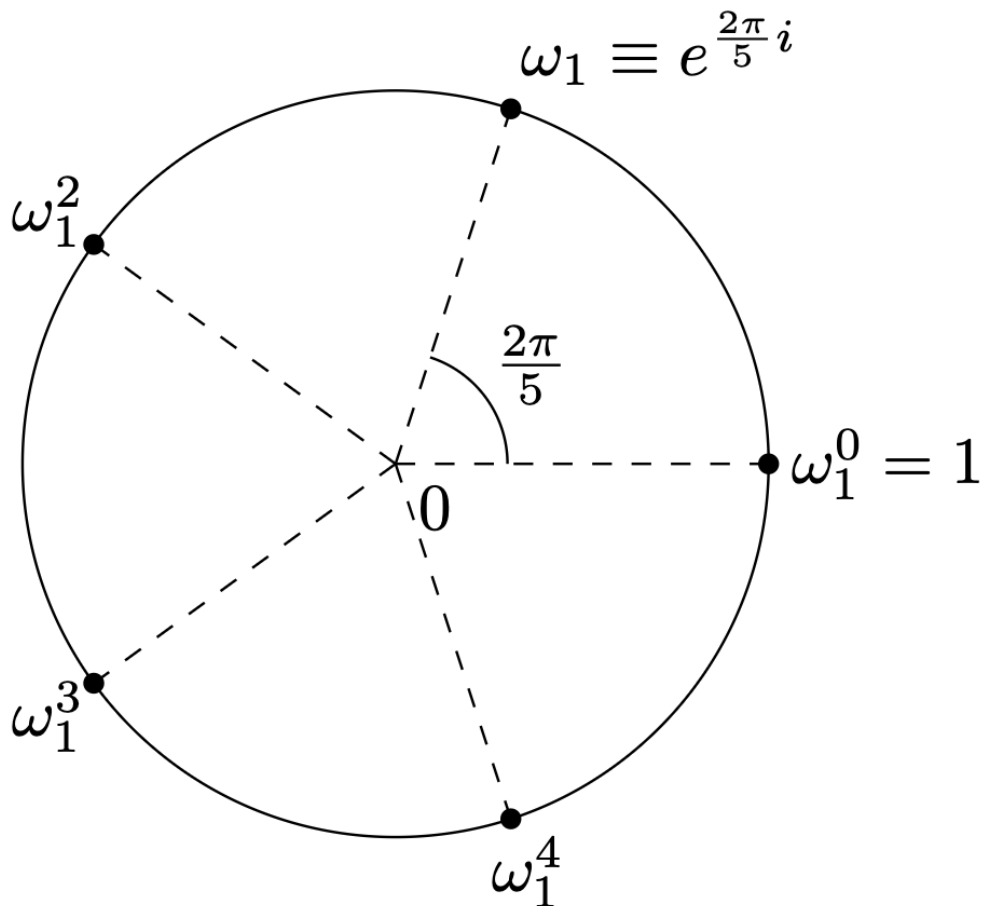
In particular, an application of this theorem is where we describe the roots of unity. Any root of unity can be described by a power of the first root of unity,

$$\omega_1 = e^{\frac{2\pi i}{n}}$$

the  $n$ -th roots of unity ( $n \in \mathbb{N}$ ) are given by

$$\omega_k = e^{\frac{2\pi i}{n}k} = \omega_1^k \quad k = 1, 2, \dots, n-1$$

which essentially says that there are  $n$  roots to the complex polynomial. For  $n = 5$ , we have the 5 roots of unity given by



In general, we say that there are  $n$  values of  $k$  that satisfy the equation  $\omega_1^n = e^{(\frac{2\pi i}{n})^n} = 1$ . From this, we can generalize what we know into the summations over  $\omega_k$ , which is any  $k$ -th root of unity except for  $\omega = 1$ .

$$\sum_{k=0}^{n-1} \omega_k = \sum_{k=0}^{n-1} \omega_1^k = 0$$

This formula can be conveniently proven by applying the formula for summing a geometric sequence to the summation. From this, we can conveniently derive a useful mathematical condition, being

**Example.** The DFT Orthonormality condition depends on two parameters  $k$  and  $l$ , and is stated as follows

$$\frac{1}{N} \sum_{n=0}^{N-1} e^{-\frac{2\pi i}{N}kn} e^{\frac{2\pi i}{N}ln} = \delta_{k-l \bmod N}$$

where  $\delta_{k-l \pmod N} = 1$  if and only if  $k \equiv l \pmod N$ , else 0. It is saying that when  $k$  is congruent to  $l$ , equivalent to  $k - l = mN$ , where the difference between  $k$  and  $l$  is divisible by some integer  $m$ . The  $\delta_{k-l \pmod N}$  term is a Kronecker delta of  $k \equiv l \pmod N$ , where the result is 1 if  $k \equiv l \pmod N$  holds and 0 in the case of  $k \not\equiv l \pmod N$ . If we define  $\omega = e^{i\frac{2\pi}{N}}$  as a primitive  $N$ th root of unity (satisfying  $\omega^N = 1$ ), we have the derivation

$$\begin{aligned} \frac{1}{N} \sum_{n=0}^{N-1} e^{-\frac{2\pi i}{N}kn} e^{\frac{2\pi i}{N}ln} &= \frac{1}{N} \sum_{n=0}^{N-1} \omega^{-kn} \omega^{ln} \\ &= \frac{1}{N} \sum_{n=0}^{N-1} \omega^{n(l-k)} \\ &= \begin{cases} \frac{1}{N} \sum_{n=0}^{N-1} 1 = 1, & \text{if } l \equiv k \pmod N, \\ \frac{1}{N} \frac{1 - \omega^{(l-k)N}}{1 - \omega^{(l-k)}} = 0, & \text{if } l \not\equiv k \pmod N \end{cases} \\ &= \delta_{k-l \pmod N}, \end{aligned}$$

where we used the fact that  $\omega^{nN} = (\omega^N)^n = 1^n = 1$  for  $n \in \mathbb{N}$ .

### 3.4 Advanced Operations

It's probably best to illustrate more advanced operations on complex numbers with the help of some examples

**Example.** Evaluating  $\sqrt{i}$  or  $\sqrt{\sqrt{1}}$  gives:

$$\sqrt{i} = \left(e^{\frac{\pi i}{2}}\right)^{\frac{1}{2}} = e^{\frac{\pi i}{4}} = \cos \frac{\pi}{4} + i \sin \frac{\pi}{4} = \frac{1}{\sqrt{2}}(1 + i)$$

The inverse is given by

$$\left(\frac{1}{\sqrt{2}}(1 + i)\right)^2 = \frac{1}{2}(1 + 2i + i^2) = \frac{(2i)}{2} = i$$

**Example.** Evaluating

$$\left(\frac{1}{2} + \frac{\sqrt{3}}{2}i\right)^{50}$$

gives

$$\left(\frac{1}{2} + \frac{\sqrt{3}}{2}i\right)^{50} = \left(e^{\frac{\pi i}{3}}\right)^{50} = e^{\frac{50\pi i}{3}} = e^{(16+\frac{2}{3})\pi i} = e^{\frac{2\pi i}{3}} = -\frac{1}{2} + \frac{\sqrt{3}}{2}i.$$

**Example.** Evaluating

$$2^{3+4i}$$

gives us

$$2^{3+4i} = 2^3 \cdot 2^{4i} = 8 \cdot e^{4\ln(2)i} = 8 \cos(4 \ln 2) + i 8 \sin(4 \ln 2).$$

**Example.** Evaluating

$$\cos(3 + 4i)$$

gives us

$$\begin{aligned} \cos(3 + 4i) &= \frac{1}{2}(e^{i(3+4i)} + e^{-i(3+4i)}) \\ &= \frac{1}{2}(e^{-4+3i} + e^{4-3i}) \\ &= \frac{1}{2}e^{-4}(\cos 3 + i \sin 3) + \frac{1}{2}e^4(\cos 3 - i \sin 3) \\ &= \frac{1}{2}(e^{-4} + e^4) \cos 3 + i \frac{1}{2}(e^{-4} - e^4) \sin 3 \end{aligned}$$



**Example.** If we have the equation  $z^5 = \frac{1}{2} + \frac{\sqrt{3}}{2}i$ , solving for  $z$  gives

$$z_k = e^{\frac{\pi i}{15}} e^{\frac{2k\pi i}{5}} \quad k = 0, 1, 2, 3, 4.$$

It is worth noting that there exists a way to express the trigonometric functions  $\sin$  and  $\cos$  as a function of euler's number. We know that

$$e^{i\theta} = \cos \theta + i \sin \theta \quad e^{-i\theta} = \cos \theta - i \sin \theta$$

from this, we can derive that

$$\begin{aligned} e^{i\theta} + e^{-i\theta} &= \cos \theta + i \sin \theta + \cos \theta - i \sin \theta \\ e^{i\theta} + e^{-i\theta} &= \cos \theta + \cos \theta \\ \cos \theta &= \frac{1}{2} (e^{i\theta} + e^{-i\theta}) \end{aligned}$$

and that

$$\begin{aligned} e^{i\theta} - e^{-i\theta} &= \cos \theta + i \sin \theta - (\cos \theta - i \sin \theta) \\ e^{i\theta} - e^{-i\theta} &= i \sin \theta - (-i \sin \theta) \\ \sin \theta &= \frac{1}{2i} (e^{i\theta} - e^{-i\theta}) \end{aligned}$$

Another way of expressing this is by saying

$$\begin{aligned} \cos x &= \Re(e^{ix}) = \frac{e^{ix} + e^{-ix}}{2} \\ \sin x &= \Im(e^{ix}) = \frac{e^{ix} - e^{-ix}}{2i} \end{aligned}$$

One final yet very important item to rememebr throughout the curriculum is that powers for complex numbers are **rotations**.

## 4 Sets, Groups, and Functions

This chapter is mainly going to fly over the various mathematical concepts that make up the backbone of many mathematical fields, including those relevant to quantum computing.

### 4.1 Sets

The concepts of sets are fundamental to many areas of mathematics. A set is a well-defined collection of distinct objects, which is also an object in its own right.

**Definition 4.1** (Set). A set is an (unordered) collection of objects, which are said to be elements or members of the set.

Let  $A = \{2, 4, 6, 8\}$  be the set of even numbers less than 10. The elements of this set are 2, 4, 6, and 8. We can also write sets using set-builder notation, such as  $B = \{x \mid x \text{ is a vowel in the English alphabet}\} = \{a, e, i, o, u\}$ .

**Definition 4.2** (Tuple). A tuple (or sequence) is an ordered list of elements.

Consider the tuple  $T = (3, 1, 4, 1, 5)$ . This is an ordered sequence where the first element is 3, the second is 1, the third is 4, the fourth is 1, and the fifth is 5. Note that the order matters and repetition is allowed, so  $(3, 1, 4, 1, 5) \neq (1, 1, 3, 4, 5)$ .

**Definition 4.3** (Cardinality). The cardinality of a set  $A$ , denoted  $|A|$ , is the number of elements in  $A$ .

For the set  $A = \{2, 4, 6, 8\}$ , the cardinality is  $|A| = 4$  since there are 4 elements. For the set  $B = \{x, y, z\}$ , we have  $|B| = 3$ . The empty set has cardinality  $|\emptyset| = 0$ .

We can categorize sets based on their cardinality into finite, countably infinite, or uncountably infinite. A set is said to be countably infinite if it can be bijectively mapped to the set of natural numbers  $\mathbb{N}$  such as the set of integers  $\mathbb{Z}$ , while the set is said to be uncountably infinite if there is no one-to-one to the set of natural numbers, such as the set of real numbers  $\mathbb{R}$ .

**Example.** Finite set:  $A = \{1, 2, 3, 4, 5\}$  has  $|A| = 5$ . Countably infinite set: The set of integers  $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$  can be put in one-to-one correspondence with the natural numbers  $\mathbb{N}$ , so  $|\mathbb{Z}| = |\mathbb{N}| = \aleph_0$ . Uncountably infinite set: The set of real numbers  $\mathbb{R}$  cannot be put in one-to-one correspondence with  $\mathbb{N}$ , so  $|\mathbb{R}| > \aleph_0$ .

**Definition 4.4** (Subset and Superset). We call  $B$  a subset of  $A$ , denoted  $B \subseteq A$  if  $\forall b \in B, b \in A$ . In this case,  $A$  is a superset of  $B$ , denoted  $B \supseteq A$ . If  $B$  is a subset of  $A$  but not equal to  $A$ , then  $B$  is called a proper subset of  $A$ , denoted  $B \subset A$ .

Let  $A = \{1, 2, 3, 4, 5\}$  and  $B = \{2, 4\}$ . Then  $B \subseteq A$  since every element of  $B$  is also in  $A$ . We can also say  $A \supseteq B$ . Since  $B \neq A$ , we have  $B \subset A$  ( $B$  is a proper subset of  $A$ ). Additionally,  $A \subseteq A$  since every set is a subset of itself.

**Definition 4.5** (Union). The union of two sets  $A$  and  $B$ , denoted by  $A \cup B$ , is the set containing all the elements in  $A$ ,  $B$ , or both.

Let  $A = \{1, 3, 5, 7\}$  and  $B = \{2, 3, 6, 7, 8\}$ . Then  $A \cup B = \{1, 2, 3, 5, 6, 7, 8\}$ , which contains all elements that appear in either set  $A$  or set  $B$  (or both).

**Definition 4.6** (Intersection). The intersection of two sets  $A$  and  $B$ , denoted by  $A \cap B$ , is the set containing all the elements in both  $A$  and  $B$ .

Using the same sets  $A = \{1, 3, 5, 7\}$  and  $B = \{2, 3, 6, 7, 8\}$ , we have  $A \cap B = \{3, 7\}$ , which contains only the elements that appear in both sets.

**Corollary 4.7** (Disjoint Sets). Two or more sets are said to be disjoint if they have no elements in common, that is, their intersection is the empty set:  $A \cap B = \emptyset$

Let  $A = \{1, 3, 5\}$  and  $B = \{2, 4, 6\}$ . These sets are disjoint because  $A \cap B = \emptyset$  - they share no common elements.

**Definition 4.8** (Difference). The difference of two sets  $A$  and  $B$ , denoted by  $A - B$  or  $A \setminus B$ , is the set containing all the elements in  $A$  but not in  $B$ .

Let  $A = \{1, 2, 3, 4, 5\}$  and  $B = \{3, 4, 6, 7\}$ . Then  $A - B = \{1, 2, 5\}$ , which contains the elements in  $A$  that are not in  $B$ . Similarly,  $B - A = \{6, 7\}$ .

**Definition 4.9** (Universal Set). The universal set, denoted by  $U$ , is the set that contains all elements under consideration, usually in relation to a particular problem or discussion. Every other set in that context is a subset of the universal set  $U$ .

In a problem about students at a university, the universal set might be  $U =$  all students at the university. If we're discussing card games, the universal set could be  $U =$  all 52 cards in a standard deck.

**Definition 4.10** (Complement). The complement of a set  $A$ , denoted by  $\bar{A}$  or  $A^c$ , is the set of all elements in the universal set  $U$  that are not in set  $A$ .

Let  $U = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  be the universal set and  $A = \{2, 4, 6, 8, 10\}$ . Then the complement of  $A$  is  $\bar{A} = \{1, 3, 5, 7, 9\}$ , containing all elements in  $U$  that are not in  $A$ .

**Definition 4.11** (Cartesian Product). The cartesian product of two sets  $A$  and  $B$ , denoted  $A \times B$ , is the set of all ordered pairs  $(a, b)$ ,  $a \in A, b \in B$ .

Let  $A = \{1, 2\}$  and  $B = \{x, y, z\}$ . Then  $A \times B = \{(1, x), (1, y), (1, z), (2, x), (2, y), (2, z)\}$ . The Cartesian product contains all possible ordered pairs where the first element comes from  $A$  and the second from  $B$ . Note that  $|A \times B| = |A| \cdot |B| = 2 \cdot 3 = 6$ .

**Definition 4.12** (Set Partitions). A partition of a set  $A$  is a collection of disjoint subsets of  $A$  such that every element in  $A$  is included in exactly one subset. These subsets are called blocks of the partition.

Let  $A = \{1, 2, 3, 4, 5, 6\}$ . One partition of  $A$  could be  $\{A_1, A_2, A_3\}$  where  $A_1 = \{1, 2\}$ ,  $A_2 = \{3, 5\}$ , and  $A_3 = \{4, 6\}$ . These subsets are disjoint ( $A_i \cap A_j = \emptyset$  for  $i \neq j$ ) and their union gives the original set ( $A_1 \cup A_2 \cup A_3 = A$ ).

**Theorem 4.13** (Set Partition). If  $\{A_1, A_2, \dots, A_n\}$  is a partition of set  $A$ , and  $B \subseteq A$ , then  $\{A_1 \cap B, A_2 \cap B, \dots, A_n \cap B\}$  is a partition of set  $B$ .

Let  $A = \{1, 2, 3, 4, 5, 6\}$  with partition  $\{A_1, A_2, A_3\}$  where  $A_1 = \{1, 2\}$ ,  $A_2 = \{3, 5\}$ , and  $A_3 = \{4, 6\}$ . If  $B = \{2, 3, 4, 5\} \subseteq A$ , then  $\{A_1 \cap B, A_2 \cap B, A_3 \cap B\} = \{\{2\}, \{3, 5\}, \{4\}\}$  forms a partition of  $B$ .

**Definition 4.14** (Power Set). The power set of  $A$  is the set of all subsets of  $A$ , denoted by  $\mathcal{P}(A)$  is the set of all subsets of  $A$ .

Let  $A = \{1, 2\}$ . The power set  $\mathcal{P}(A)$  contains all possible subsets of  $A$ :  $\mathcal{P}(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$ . Note that  $|\mathcal{P}(A)| = 2^{|A|} = 2^2 = 4$ . For any set with  $n$  elements, its power set has  $2^n$  elements. For set operations in general, there are a few laws worth noting.

Identity Laws:

$$A \cup \emptyset = A$$

$$A \cap U = A.$$

Domination Laws:

$$A \cup U = U$$

$$A \cap \emptyset = \emptyset.$$

Idempotent Laws:

$$A \cup A = A$$

$$A \cap A = A.$$

Absorption Laws:

$$A \cup (A \cap B) = A$$

$$A \cap (A \cup B) = A.$$

Complement Laws:

$$A \cup \bar{A} = U$$

$$A \cap \bar{A} = \emptyset$$

$$\overline{\bar{A}} = A.$$

Commutative Laws:

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A.$$

Associative Laws:

$$A \cup (B \cup C) = (A \cup B) \cup C$$

$$A \cap (B \cap C) = (A \cap B) \cap C.$$

Distributive Laws:

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C).$$

De Morgan's Laws:

$$\overline{A \cap B} = \bar{A} \cup \bar{B}$$

$$\overline{A \cup B} = \bar{A} \cap \bar{B}.$$

**Definition 4.15** (Totally Ordered Set). A totally ordered set is a set in which every pair of elements is comparable; for any two elements  $a$  and  $b$ , either  $a \leq b$  or  $b \leq a$  holds.

Examples include the real numbers and the integers with their usual orders.

**Definition 4.16** (Partially Ordered Set (Poset)). A partially ordered set (poset) is a set equipped with a relation  $\leq$  that is reflexive, antisymmetric, and transitive; in a poset, not every pair of elements must be comparable, so some pairs may be incomparable. For example, the power set  $\mathcal{P}(\mathbb{Z})$  ordered by inclusion ( $\subseteq$ ) contains subsets such as  $\{1, 2\}$  and  $\{2, 3\}$  that are not comparable, and the positive integers ordered by divisibility have incomparable primes like 5 and 7.

**Definition 4.17** (Supremum). Let  $S$  be a nonempty subset of a partially ordered set  $P$ . An element  $u \in P$  is the supremum of  $S$ , denoted by  $u = \sup S$ , if:

1. **Upper bound:** every  $s \in S$  satisfies  $s \leq u$ .
2. **Least upper bound:** if  $v$  is any upper bound of  $S$ , then  $u \leq v$ .

For instance,  $\sup\{-1, -2, -3, \dots\} = -1$ , and  $\sup\{\sin x : x \in [0, \pi]\} = 1$ .

**Definition 4.18** (Infimum). Let  $S$  be a nonempty subset of a partially ordered set  $P$ . An element  $\ell \in P$  is the infimum of  $S$ , denoted by  $\ell = \inf S$ , if:

1. **Lower bound:** every  $s \in S$  satisfies  $\ell \leq s$ .
2. **Greatest lower bound:** if  $v$  is any lower bound of  $S$ , then  $v \leq \ell$ .

For example,  $\inf\{e^{-x} : x > 0\} = 0$ , even though 0 is not an element of the set.

## 4.2 Groups

Groups, rings, and fields lay the foundation in mathematics that build upon sets with additional operations in algebra. These structures are ubiquitous in mathematics and physics and lay the foundation for quantum computing as well.

A group is a set equipped with a single binary operation that exhibits certain properties, much like addition and multiplication. A ring expands on this by incorporating two operations, typically referred to as multiplication and addition. A field is a more stringent structure where the set is a group under both operations, with multiplication also being commutative, and every non-zero element having a multiplicative inverse.

**Definition 4.19.** A group is a set  $G$  which is *closed* under an operation  $\cdot$  (that is  $\forall x, y \in G, x \cdot y \in G$ ) and satisfies the following properties:

1. **Identity:**  $\exists e \in G$  where  $\forall x \in G, x \cdot e = x = e \cdot x$ . We define  $e$  to be the identity element.
2. **Inverse:**  $\forall x \in G, \exists y \in G$  such that  $x \cdot y = e = y \cdot x$ , where  $e$  is the identity element identified above.
3. **Associativity:** The operation  $\cdot$  is associative for every  $x, y, z \in G$ , i.e.

$$x \cdot (y \cdot z) = (x \cdot y) \cdot z$$

In mathematical contexts, it is common to omit the symbol  $\cdot$  or  $*$  for the group operation and simply write  $x \cdot y$  as  $xy$ . Some common examples include the group of integers  $\mathbb{Z}$  under addition, denoted by  $(\mathbb{Z}, +)$ , with  $e = 0$ , and the inverse of  $x$  being  $-x$ . Another example would be the group of integers modulo  $n$ , denoted  $\mathbb{Z}/n\mathbb{Z}$ , with closure (the sum of any two integers modulo  $n$  also forming a group under modulo  $n$ ), identity, inverse, and associativity all holding under the subspace of  $\mathbb{Z}/n\mathbb{Z}$ .

**Definition 4.20** (Abelian Group). A group is said to be *abelian* if the operation  $\cdot$  is commutative  $\forall x, y \in G$ , that is,

$$x \cdot y = y \cdot x$$

While the examples so far are all abelian groups, there are some of groups that are not abelian. For example, the set of all symmetries of an equilateral triangle, known as the dihedral group  $D_3$  is not abelian. As you can see, the table of operations is not symmetrical, rendering the group as a non-abelian group. However, the identity, inverse, and associative elements/properties are all present upon verification.

	$e$	$a$	$b$	$c$	$r$	$s$
$e$	$e$	$a$	$b$	$c$	$r$	$s$
$a$	$a$	$e$	$r$	$s$	$b$	$c$
$b$	$b$	$s$	$e$	$r$	$c$	$a$
$c$	$c$	$r$	$s$	$e$	$a$	$b$
$r$	$r$	$c$	$a$	$b$	$s$	$e$
$s$	$s$	$b$	$c$	$a$	$e$	$r$

However, in the context of quantum computing, there are a few symmetry groups worth noting. They are

1. **SO( $N$ )** : The orthogonal group in  $N$  dimensions consists of all  $N \times N$  orthogonal matrices with determinant 1, representing rotations in  $\mathbb{R}^N$ . These rotations preserve distance and the orientation of objects in question. SO(2) corresponds to a circle, and SO(3) to a sphere.
2. **SU(2)** : The special unitary group of degree 2 comprises of all  $2 \times 2$  unitary matrices with determinant 1. It is closely related to SO(3) and is commonly used to describe spins and qubit states. Each rotation in SO(3) corresponds to two points in SU(2).
3. **SU( $N$ )** : This represents the special unitary group of degree  $N$ , extending the concepts of SU(2). These groups are useful in the study of quantum entanglement in quantum computing concerning  $N$ -level quantum systems.

**Definition 4.21** (Subgroups). A subgroup  $H$  of group  $G$  is a subset of  $G$  that is a group in by itself, with the same group operation in  $G$ . You can think of this as a reduced version of  $G$  where still

1. The identity element of  $G$  is in  $H$ .
2.  $\forall h_1, h_2 \in H, \quad h_1 \cdot h_2 \in H$ .
3.  $\forall h \in H, \quad h^{-1} \in H$

**Theorem 4.22** (Lagrange's Theorem). For any finite group  $G$  and any subgroup  $H$  of  $G$ , the order of  $H$  divides the order of  $G$ , i.e.

$$|G| \equiv 0 \pmod{|H|}$$

**Definition 4.23** (Coset). Given a group  $G$  and a subgroup  $H$  of  $G$ , the **left coset** of  $H$  with representative  $g \in G$  is the set  $gH = \{gh|h \in H\}$ . Similarly, the right coset is  $Hg = \{hg|h \in H\}$ .

**Theorem 4.24** (Partition Theorem for Cosets). The collection of all left cosets of a subgroup  $H$  forms a partition of the group  $G$ . This means

1. Every element of  $G$  belongs to exactly one coset of  $H$ .
2. Cosets are disjoint and have no elements in common.

**Definition 4.25** (Normal Subgroup). A subgroup  $N$  of a group  $G$  is called a normal subgroup if it is invariant under conjugation by elements of  $G$ . This means that  $\forall n \in N, \quad g \in G, \quad gng^{-1} \in N$ . In notation,  $N \triangleleft G$  if

$$gNg^{-1} = \{gng^{-1}|n \in N\} \subseteq N \forall g \in G.$$

Now, note that this is not the same as an abelian group. While the abelian group implies that every subgroup is a normal subgroup, the latter does not imply the former.

**Example** (Quaternion Group  $Q_8$ ). Let

$$Q_8 = \{\pm 1, \pm i, \pm j, \pm k\}$$

with relations

$$i^2 = j^2 = k^2 = ijk = -1.$$

This group is non-abelian since, for instance,  $ij = k$  while  $ji = -k$ , so  $ij \neq ji$ . Nevertheless, every subgroup of  $Q_8$  is normal; the subgroups are  $\{1\}$ ,  $\{\pm 1\}$ ,  $\langle i \rangle = \{\pm 1, \pm i\}$ ,  $\langle j \rangle = \{\pm 1, \pm j\}$ ,  $\langle k \rangle = \{\pm 1, \pm k\}$ , and  $Q_8$  itself, and each is invariant under conjugation by any element of  $Q_8$ . Thus  $Q_8$  is a non-abelian group in which all subgroups are normal (a Hamiltonian group).

Also, it is worth noting that the left and right cosets of a normal subgroup  $N$  are the same, allowing the group operations on cosets to be well-defined, which brings us to quotient groups:

**Definition 4.26** (Quotient Group). Let  $G$  be a group, and  $N \triangleleft G$ . The quotient group  $G/N$  is the set of cosets of  $N \in G$  with the group operations defined by:

$$(gN)(hN)(gh)N \quad \forall g, h \in G$$

**Example.** Consider the group  $\mathbb{Z}$  of integers under addition and the subgroup  $2\mathbb{Z}$  consisting of all even integers. The quotient group  $\mathbb{Z}/2\mathbb{Z}$  is the set of cosets of  $2\mathbb{Z}$  in  $\mathbb{Z}$ :

- The coset  $0 + 2\mathbb{Z} = \{\dots, -4, -2, 0, 2, 4, \dots\}$  represents the even integers.
- The coset  $1 + 2\mathbb{Z} = \{\dots, -3, -1, 1, 3, 5, \dots\}$  represents the odd integers.

Thus,  $\mathbb{Z}/2\mathbb{Z}$  has two elements:  $0 + 2\mathbb{Z}$  and  $1 + 2\mathbb{Z}$ , corresponding to the even and odd integers, respectively. The group operation is addition modulo 2.

This quotient group is *isomorphic* to  $\mathbb{Z}_2 = \{0, 1\}$  under addition modulo 2, denoted as  $\mathbb{Z}/2\mathbb{Z} \cong \mathbb{Z}_2$ . Two groups are isomorphic if they have the same structure, meaning there is a one-to-one correspondence between the elements of the two groups that preserves the group operation.

**Definition 4.27** (Cyclic Group). A group  $G$  is cyclic if  $\exists g \in G$  such that  $\forall x \in G$ ,  $x$  can be expressed as powers (repeated operations) of  $g$ .

A good example of this would be  $D_3$ , which also has multiple generators.

**Definition 4.28** (Ring). A ring is a set  $R$  equipped with two operations  $+$  and  $\times$  satisfying the following properties:

1.  $(R, +)$  forms an abelian group.
2.  $(R, \times)$  is associative, that is  $\forall a, b, c \in R$ ,

$$a \times (b \times c) = (a \times b) \times c$$

3. The distributive properties hold, that is  $\forall a, b, c \in R$ ,

$$a \times (b + c) = (a \times b) + (a \times c)$$

$$(b + c) \times a = (b \times a) + (c \times a)$$

For example, both the set of integers  $\mathbb{Z}$  and the group of integers modulo  $n$ ,  $\mathbb{Z}/n\mathbb{Z}$ , are both rings. However, multiplication notably does not have an inverse under most groups of integers. For this to hold, we need to have fields.

**Definition 4.29** (Field). A field is a set  $F$  with two operations  $+$  and  $\times$ , where

1.  $(R, +)$  forms an abelian group.
2.  $(R - \{0\}, \times)$  forms an abelian group.
3. The distributive properties holds as in rings.

Some common examples include the set of rational numbers  $\mathbb{Q}$ , the set of real numbers  $\mathbb{R}$ , and the set of complex numbers  $\mathbb{C}$ . Interestingly,  $\mathbb{Z}/p\mathbb{Z}$  where  $p$  is a prime is also a field.

### 4.3 Functions

Functions serve as mappings from one set to another. They assign each element in the domain to exactly one element in the codomain, and are essential for describing relationships using math.

**Definition 4.30** (Function). Let  $f$  be a function from set  $A$  to set  $B$ . A *function from  $A$  to  $B$* , denoted  $f : A \rightarrow B$ , is an assignment of exactly one element of  $B$  to each element of  $A$ .

We write  $f(a) = b$  to denote the assignment of  $b$  to an element  $a$  of  $A$  by the function  $f$ .

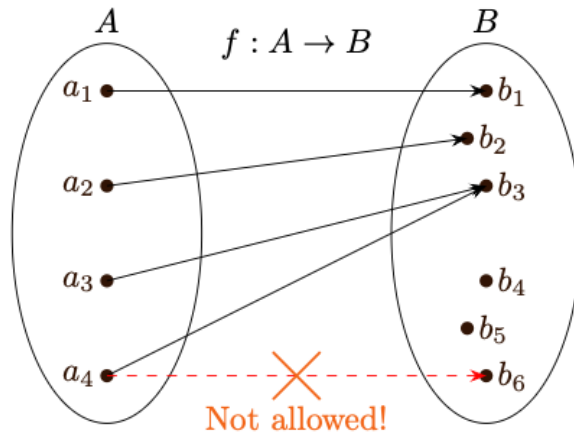
**Example.**  $f(x) = x^2$  is a function  $f : \mathbb{R} \rightarrow \mathbb{R}$ .  $r = \sqrt{x^2 + y^2}$  is a function  $f : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$ .

**Definition 4.31** (Image, Range, and Domain). Let  $f$  be a function from  $A$  to  $B$ .

- We say that  $A$  is the *domain* of  $f$  and  $B$  is the *codomain* of  $f$ .
- If  $f(a) = b$ ,  $b$  is the *image* of  $a$  and  $a$  is a *pre-image* of  $b$ .
- The *range* of  $f$  (a subset of  $B$ ) is the set of all images of elements of  $A$ .

- Let  $S$  be a subset of  $A$ . The *image* of  $S$  is a subset of  $B$  that consists of the images of the elements of  $S$ . We denote the image of  $S$  by  $f(S)$ , so that

$$f(S) = \{f(s) \mid s \in S\}$$



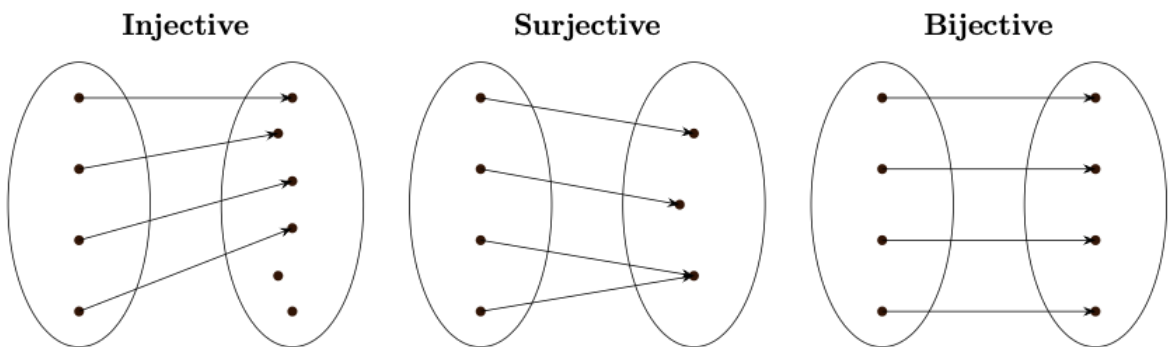
There are also a few common function types that we will define:

**Definition 4.32** (Injective). A function  $f : A \rightarrow B$  is injective if  $\forall a_1, a_2 \in A$ , we have  $f(a_1) = f(a_2) \implies a_1 = a_2$ .

**Definition 4.33** (Surjective). A function  $f : A \rightarrow B$  is surjective if the whole codomain is covered, meaning that  $\forall b \in B, \exists a \in A$  such that  $f(a) = b$ .

Sometimes we call surjective functions *onto* functions.

**Definition 4.34** (Bijective). A function is bijective if it is both *injective* and *surjective*.



Only with bijectivity established can we define the inverse of a function.

**Definition 4.35** (Inverse Function). If we take the function  $f : A \rightarrow B$ , its domain  $A$  and codomain  $B$ , we can define an inverse function  $f^{-1} : B \rightarrow A$  such that  $\forall b \in B, f^{-1}(b) = a$  if and only if  $f(a) = b$ .

Put plainly,  $\forall a \in A, f^{-1}(f(a)) = a$  and  $\forall b \in B, f(f^{-1}(b)) = b$ .

## 4.4 Common Functions and Asymptotic Behavior

This section will be a brief review of real functions ( $f : \mathbb{R} \rightarrow \mathbb{R}$ ) commonly used in quantum computing.

### Power Functions

Power functions take the form  $f(x) = x^p, \quad x \geq 0, p \in \mathbb{R}$ .

The behavior of the function varies significantly with the exponent  $p$ :

- For  $p > 0$ ,  $f(x)$  increases as  $x$  increases.  $f(x)$  exhibits a more rapid growth with a larger  $p$ .



- For  $p < 0$ ,  $f(x)$  decreases as  $x$  increases.
- When  $p = 0$ ,  $f(x) = 1$ , regardless of  $x$  (excluding  $x = 0$ ), which is a constant function.
- For  $p = 1$ ,  $f(x) = x$ , representing a linear relationship.

Key properties of power functions include the rules for exponentiation:

- Multiplying powers with the same base:  $x^a \cdot x^b = x^{a+b}$ .
- Dividing powers with the same base:  $x^a/x^b = x^{a-b}$ .

## Polynomial Functions

A polynomial function is a sum of terms  $a_i x^i$ , where  $i$  is a non-negative integer:

$$f(x) = a_0 + a_1 x + a_2 x^2 + \cdots + a_n x^n.$$

Its behavior for large  $x$  values is predominantly determined by its highest power term,  $x^n$ , where  $n$  is the degree of the polynomial.

An  $n$ -th degree polynomial has  $n$  complex roots (counting multiplicities), and according to Vieta's formulas, the sum of these roots is equal to  $-a_{n-1}/a_n$  and their product  $(-1)^n a_0/a_n$ .

**Example.** Consider the polynomial function  $f(x) = x^3 - 7x^2 + 14x - 8$ . It can be factored as:

$$f(x) = (x - 1)(x - 2)(x - 4).$$

The roots of this polynomial are  $x = 1$ ,  $x = 2$ , and  $x = 4$ , which can be found by solving the equations  $(x - 1) = 0$ ,  $(x - 2) = 0$ , and  $(x - 4) = 0$ . According to Vieta's formulas, the sum of the roots is:

$$1 + 2 + 4 = 7 = -\frac{-7}{1},$$

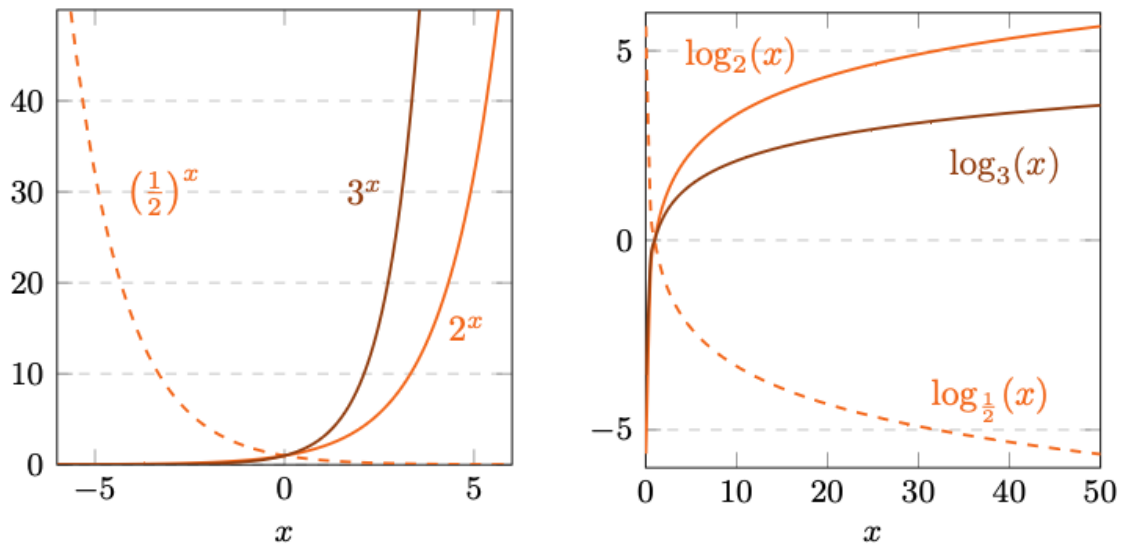
and the product is:

$$1 \cdot 2 \cdot 4 = 8 = (-1)^3 \frac{-8}{1}.$$

## Exponential Functions

Exponential functions are defined as  $f(x) = b^x$ , where  $b$  is a positive constant (called the base) and  $b \neq 1$ . The variable  $x$  is the exponent. The key characteristic is that the variable is in the exponent. Some important notes:

- Growth and Decay:
  - If  $b > 1$ ,  $f(x)$  exhibits exponential growth – increasing rapidly as  $x$  increases. Larger bases lead to faster growth.
  - If  $0 < b < 1$ ,  $f(x)$  shows exponential decay – decreasing towards zero as  $x$  increases.
- Always Positive: Exponential functions are always positive for any real-valued input  $x$ .
- Horizontal Asymptote: They approach zero for one direction of  $x$  (negative infinity for growth, positive infinity for decay).
- Base  $e$ : The natural exponential function with  $b = e$  (Euler's number,  $\approx 2.718$ ), i.e.,  $e^x$ , also denoted as  $\exp(x)$ , has special significance across mathematics.



## Logarithmic Functions

Logarithmic functions are the inverses of exponential functions. They are defined as  $f(x) = \log_b(x)$ , where  $b$  is a positive constant ( $b \neq 1$ ) and  $x > 0$ . Some key points:

- Reversing Exponentiation: If  $b^y = x$  then  $\log_b(x) = y$ .
- Growth and Behavior
  - For  $b > 1$ ,  $\log_b(x)$  increases as  $x$  increases, but very slowly.
  - For  $0 < b < 1$ ,  $\log_b(x)$  decreases as  $x$  increases.
- Vertical Asymptote: Logarithmic functions have a vertical asymptote at  $x = 0$ .
- Logarithms of 1 and the Base:  $\log_b(1) = 0$  and  $\log_b(b) = 1$ .
- The natural logarithm, written as  $\ln(x)$  has the base  $e$ .

Key Properties:

- The Product Rule:  $\log_b(xy) = \log_b(x) + \log_b(y)$
- Logarithms "Break" Exponents:  $\log_b(x^y) = y \cdot \log_b(x)$
- Changing Bases:  $\log_b(x) = \log_a(x) / \log_a(b)$

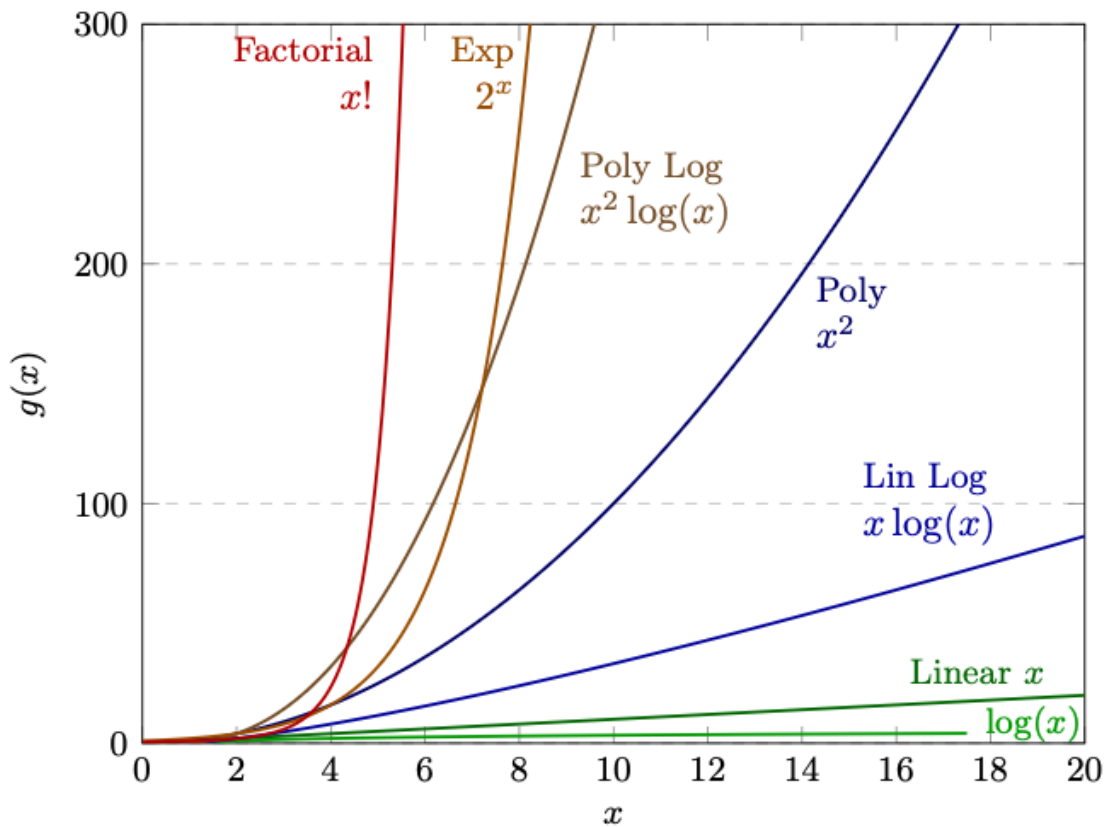
Scaling behavior, especially in the context of data structures and algorithm efficiency, is a significant topic for any computing related fields. It concerns itself with the evolution of function curves as the response becomes significantly large. We use the *Big O notation* to describe the behavior of a function  $f(x)$ ,  $x \rightarrow \infty$ . We say that  $f(x)$  is  $O(g(x))$  meaning that there is some positive constant  $c$  such that the upper bound for complexity growth does not increase faster than  $c \cdot g(x)$  for a sufficiently large  $x$ .

**Example.** For example, consider

$$f(x) = 6x^3 + 2x + 1,$$

we say that  $f(x)$  is  $O(x^3)$  as  $x \rightarrow \infty$ , as  $x^3$  is the dominant term.

Most of the common limiting functions are illustrated in the following figure:



1. Log-log:  $g(x) = \log \log(x)$

- Exhibits extremely slow growth. Algorithms within this complexity class increase their running time at a negligible rate with input size escalation.
- Applications include specialized computational geometry problems.

2. Log:  $g(x) = \log(x)$

- Denotes high efficiency. The execution time grows much slower than the input size.
- Examples include binary search in sorted arrays and operations on certain balanced tree data structures.

3. Sublinear:  $g(x) = x^p, 0 < p < 1$

- Exhibits growth slower than linear but faster than logarithmic.
- Common examples include the Grover's search algorithm in quantum computing, which has a complexity of approximately  $O(\sqrt{x})$ , and some algorithms that utilize probabilistic methods to achieve faster-than-linear performance on average.

4. Linear:  $g(x) = x$

- Indicates direct proportionality. Doubling the input size doubles the running time.
- Common examples are searching in unsorted lists and identifying max/min elements in a list.

5. Polynomial:  $g(x) = x^p, p > 1$

- The growth rate is influenced by the exponent  $p$ . Higher values lead to rapid increases in running time with input size.
- Examples: Bubble sort and insertion sort (quadratic complexity), matrix multiplication algorithms (cubic complexity or better).

6. Poly-log:  $g(x) = x^p \log(x), p \geq 1$

- Less efficient than the corresponding poly (or linear for  $p = 1$ ) but still considered scalable.

- Fast Fourier Transform (FFT) algorithms are a prime example of algorithms with linear-log complexity. Some fast sorting algorithms also approach this performance.

7. Exponential:  $g(x) = b^x, b > 1$

- Characterized by rapid growth. Algorithms in this class quickly become impractical for moderate input sizes.
- Examples: Brute-force approaches to the Traveling Salesman Problem. Currently known classical algorithms for integer factorization.

8. Factorial:  $g(x) = x!$

- Exhibits extremely rapid growth, surpassing even exponential functions in rate. Practical for only very small input sizes.
- Example: Generating all permutations of a set.

9. Hyper-exponential:  $g(x) = x^x, g(x) = b^{a^x}, g(x) = b^{x!}$ , and  $g(x) = b^{x^x}$ , etc., where  $a, b > 1$

- Exhibits growth that is even more rapid than factorial functions.
- Example: Modeling scenarios with extremely high growth rates, beyond combinatorial complexity.