Crypto HW2b

Milena Gonzalez

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Q1.

Users A and B use the Diffie-Hellman key exchange technique with a common prime q = 71 and a primitive root $\alpha = 7$.

- a) If user A has a private key $X_a = 5$, what is A's public key Y_a ? $Y_a = \alpha^{X_a} \mod q = 7^5 \mod 71 = 51$
- b) If user B has a private key $X_b = 12$, what is B's public key Y_b ? $Y_b = \alpha^{X_b} \mod q = 7^{12} \mod 71 = 4$
- c) What is the shared secret key? K = $(\alpha^{X_a})^{X_b} \mod q = 7^{5*12} \mod 71 = 30$
- d) In the Diffie-Hellman protocol, each participant selects a secret number x and sends the other participant ($\alpha \ge 1$ x mod q) for some public number α . What would happen if the participants sent each other ($x \ge 1$ mod q) instead?

The system would be insecure and can be broken easily. Since α is a public number, Eve will also know it and can easily get the key value by dividing the message by α

Q2.

A network resource X is prepared to sign a message by appending the appropriate 64- bit hash code and encrypting that hash code with X's private key as described in class (also in the textbook, Page 330).

- a) Describe the Birthday Attack where an attacker receives a valid signature for his fraudulent message?
- type of brute force attack
- exploits mathematics behind the birthday parado and probability theory
- success depends on the likelihood of collisions found between random attack attempts and a fixed degree of permutations (pigeonhole theorem)
- b) How much memory space does attacker need for an M-bit message? Attacker generates $2^{m/2}$ variations of a valid message all with essentially the same meaning, but since it is a 64 bit message, will need $2^{m/2} * 2^6$ bits of memory
- c) Assuming that attacker's computer can process 2^{20} hash/second, how long does it take at average to find pair of messages that have the same hash? $\frac{2^{m/2}*2^6}{2^{20}}$ seconds
- d) Answer (b) and (c) when 128-bit hash is used instead. $2^{m/2}*2^7, \frac{2^{m/2}*2^7}{2^{20}} {\rm seconds}$

Q3.

Use Trapdoor Oneway Function with following secrets as described in lecture notes to encrypt plaintext P = '0101 0111'. Decrypt the resulting ciphertext to obtain the plaintext P back. Show each step to get full credit.

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S = \{5, 9, 21, 45, 103, 215, 450, 946\}
a = 1019, p = 1999
public key \beta = S^*(a \mod p)
\beta = S^*\{1097, 1175, 1409, 1877, 1009, 1194, 779, 456\}
Ciphertext C = (0)(1097) + (1)(1175) + (0)(1409) + (1)(1877) + (0)(1009) + (1)(1194) + (1)(779) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1)(1194) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) + (1004) +
(1)(456) = 5481
a^{-1} \pmod{p}
    1999 = (1019)(1) + 980
                                                                                            980 = 1999 + 1019(-1)
    1019 = 980(1) + 39
                                                                                              39 = 1019 + 910(-1)
    980 = 39(25) + 5
                                                                                             5
    39 = 5(7) + 4
                                                                                             4 = 39 + 5(-7)
    5 = 4(1) + 1
                                                                                              1 = 5 + 4(1)
1 = 5 + 4(-1)
1 = 5 + (-1)[39 + 5(-7)]
1 = 39(-1) + 5(8)
1 = 39(-1) + 8[980 + 39(-25)] 1 = 980(8) + 39(-201)
1 = 980(8) + (-201)[1019 + 950(-1)]
1 = 1019(-201) + 980(209)
1 = 1019(-201) + (209)[1999 + 1019(-1)]
1 = 1019(-410) + 1999(0)
1999 - 410 = 1589 (< -a^{-1})
C*a^{-1} \mod p = 1665
Decompose:
1665 - 946 = 719
719 - 450 = 269
269 - 215 = 52
54 - 45 = 9
9 - 9 = 0
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