

# Chapter 1

## Section 1.1

N/A

## Section 1.2

N/A

# Chapter 2

## Section 2.1

N/A

## Section 2.2

### Exercise 2.2.1

We shall use induction on  $a$ , fixing  $b$  and  $c$ . First, we shall prove the base case. If  $a = 0$ ,  $a + b = 0 + b = b$  by the definition of addition of natural numbers (Definition 2.2.1). Therefore,  $(a + b) + c = b + c$ . Also by Definition 2.2.1,  $a + (b + c) = 0 + (b + c) = b + c$ . We now have  $(a + b) + c = b + c = a + (b + c)$ .

Now suppose that  $(a + b) + c = a + (b + c)$ . We then have, by Definition 2.2.1, that  $((a++) + b) + c = ((a + b)++) + c = ((a + b) + c)++ = (a + (b + c))++$ . Because of Definition 2.2.1 again,  $(a + (b + c))++ = (a++) + (b + c)$ . Now we have  $((a++) + b) + c = (a++) + (b + c)$ , which closes the induction.

### Exercise 2.2.2

We shall use induction on  $a$ . The base case is vacuously true since 0 is not positive. For the inductive step, we have to show that  $a++$  being positive implies that  $a++ = b++$  for exactly one natural number  $b$  if  $a$  being positive

implies that  $a = c++$  for exactly one natural number  $c$ . But  $a$  is equal to  $b$  if  $a++$  is equal to  $b++$ , by Axiom 2.4. This closes the induction.

### Exercise 2.2.3

- (a) By Lemma 2.2.2,  $a = a + 0$ , so  $a \geq a$ .
- (b) Because  $a \geq b$  and  $b \geq c$ ,  $a = b + n$  and  $b = c + m$  for some  $n$  and  $m$ . Therefore,  $a = b + n = (c + m) + n$ . By Proposition 2.2.5,  $(c + m) + n = c + (m + n)$ , so  $a = c + (m + n)$ .  $n + m$  is a natural number, so  $a \geq c$ .
- (c) Because  $a \geq b$  and  $b \geq a$ ,  $a = b + n$  and  $b = a + m$ . Therefore,  $a = b + n = (a + m) + n$ . By Proposition 2.2.5,  $(a + m) + n = a + (m + n)$ . Now we have  $a = a + (m + n)$ . By Lemma 2.2.2,  $a = a + 0$ , so we can use Proposition 2.2.6 to get  $0 = m + n$ . Using Corollary 2.2.9, we can deduce that  $n$  and  $m$  are both equal to 0. Therefore,  $a = b$ .
- (d) If  $a \geq b$ ,  $a = b + d$  for some  $d$ . Because of this,  $a + c = (b + d) + c$ . By Propositions 2.2.4 and 2.2.5,  $(b + d) + c = b + (d + c) = b + (c + d) = (b + c) + d$ . Therefore,  $a + c \geq b + c$ .

Conversely, if  $a + c \geq b + c$ ,  $a + c = (b + c) + d$  for some  $d$ . By Propositions 2.2.4 and 2.2.5,  $(b + c) + d = b + (c + d) + b + (d + c) = (b + d) + c$ . Using Proposition 2.2.6, we get  $a = b + d$ , and therefore,  $a \geq b$ .

- (f) If  $a < b$ ,  $b = a + c$  for some  $c$ , and  $b \neq a$ . Therefore,  $c$  is positive, because otherwise  $b = a + 0 = a$ . ( $a + 0 = a$  because of Lemma 2.2.2.)

Conversely, if  $b = a + c$  for some positive number  $c$ , we know that  $a \leq b$ . If  $a = b$ ,  $a = a + c$ . By Lemma 2.2.2,  $a = a + 0$ , so  $a + 0 = a + c$ . We can use Proposition 2.2.6 to get  $0 = c$ , but  $c$  is positive. Therefore,  $a \neq b$ , and combined with  $a \leq b$ , we get  $a < b$ .

- (e) We first prove that  $a < b$  implies that  $a++ \leq b$ . By (f),  $b = a + c$  for some positive  $c$ . By Lemma 2.2.10, exactly one  $d$  exists such that  $d++ = c$ . Therefore,  $b = a + d++$ . By Lemma 2.2.3 and Definition 2.2.1,  $b = a + d++ = (a + d)++ = (a++) + d$ . Therefore,  $a++ \leq b$ .

Conversely, if  $a++ \leq b$ , we need to prove that  $a < b$ . Since  $a++ = a + 1 \geq a$ , by (b),  $a \leq b$ . If  $a = b$ , then  $a++ \leq a$ , and we get, from (c), that  $a++ = a$ . This is a contradiction, from Axiom 2.3 if  $a = 0$  and Axiom 2.4 otherwise. Therefore,  $a < b$ .

### Exercise 2.2.6 (unfinished)

We use induction on  $n$ . When  $n = 0$ ,  $m \leq n$  implies that  $0 = m + a$ . But by Corollary 2.2.9,  $m = 0 = n$ , so  $P(n)$  being true implies that  $P(m)$  is true.

If  $a \leq n$  and  $P(n)$  together imply  $P(a)$ , we now need to prove that  $b \leq n++$  and  $P(n++)$  together imply  $P(b)$ . Since  $b \leq n++$ ,  $n++ = b + c$  for some  $c$ . If  $c$  is positive, by Lemma 2.2.10,  $n++ = b + d++$  for some  $d$ , so by Lemma 2.2.3 and Axiom 2.4,  $n = b + d$ , and  $b \leq n$ . Therefore,  $b \leq n++$  implies that  $b \leq n$  or  $n++ = b + 0 = b$ . Because of Lemmas 2.2.2 and 2.2.3,  $n++ = (n + 0)++ = n + 0++ = n + 1$ . Since  $P(n++)$  implies  $P(n)$ ,  $P(n)$  is true. But then ?

## Section 2.3

### Exercise 2.3.1

First, we prove that  $n \cdot 0 = 0$ . We induct on  $n$ . When  $n = 0$ ,  $n \cdot 0 = 0 \cdot 0 = 0$ , because of Definition 2.3.1. For the inductive hypothesis, we assume  $n \cdot 0 = 0$ . Then we need to prove that  $n++ \cdot 0 = 0$ . By Definition 2.3.1,  $n++ \cdot 0 = n \cdot 0 + 0 = n \cdot 0$ . But we already know that  $n \cdot 0 = 0$ , so  $n++ \cdot 0 = 0$ . This closes the induction.

Next, we prove that  $n \cdot m++ = nm + n$  using induction on  $n$ . When  $n = 0$ ,  $0 \cdot m++ = 0 = 0 \cdot m + 0$ . Now, if  $n \cdot m++ = nm + n$ , we need to prove that  $n++ \cdot m++ = n++ \cdot m + n++$ . We can deduce from  $n \cdot m++ = nm + n$  that  $n++ \cdot m++ = (n \cdot m++) + m++ = nm + n + m++ = nm + (n + m)++ = nm + n++ + m = (n++) \cdot m + n++$ . This closes the induction.

Now we can prove that  $nm = mn$ . We induct on  $m$ . When  $m = 0$ , we have  $n \cdot 0 = 0 = 0 \cdot n$ . For the inductive hypothesis, we assume  $nm = mn$ . Now we need to prove  $n \cdot m++ = m++ \cdot n$ . We know that  $n \cdot m++ = nm + n$  and  $m++ \cdot n = mn + n$ . But  $nm = mn$ , so  $n \cdot m++ = m++ \cdot n$ . This closes the induction.

### Exercise 2.3.4

By Proposition 3.4,  $(a + b)^2 = (a + b) \cdot (a + b) = (a + b) \cdot a + (a + b) \cdot b = a \cdot a + b \cdot a + a \cdot b + b \cdot b$ . We can rewrite  $b \cdot a + a \cdot b$  as  $a \cdot b + a \cdot b$  because of Lemma 2.3.2. We can further rewrite this as  $2ab$  because  $2(ab) = 1(ab) + ab =$

$0(ab) + ab + ab = 0 + ab + ab = ab + ab$ . We can also rewrite  $a \cdot a$  and  $b \cdot b$  as  $a^2$  and  $b^2$  respectively. Therefore,  $(a + b)^2 = a^2 + 2ab + b^2$ .

## Chapter 3

### Section 3.1

#### Exercise 3.3.1

We know that at least one of the statements “ $a = c$ ” and “ $a = d$ ” is true since  $\{a, b\} = \{c, d\}$ . Similarly, at least one of the statements “ $b = c$ ” and “ $b = d$ ” is true. If we have that  $a = c$  and  $b = d$  at the same time, then we have proved what we want. (Same for  $a = d$  and  $b = c$  at the same time.) Otherwise, then we must have that  $a = c = b$  or  $a = d = b$ . If both of these are true, then  $a = b = c = d$ , and therefore both  $a = c$  and  $b = d$ . If only one of these is true, which we will assume to be the statement that  $a = c = b$  as a similar argument holds assuming  $a = d = b$ , then  $d \notin \{a, b\} = \{a, a\}$ , while  $d \in \{c, d\} = \{a, d\}$ , so  $\{a, b\} \neq \{c, d\}$ , a contradiction. Therefore, if  $\{a, b\} = \{c, d\}$ , either both  $a = c$  and  $b = d$  or  $a = d$  and  $b = c$ .

#### Exercise 3.3.2

First, we show that these four sets exist. Axiom 3.3 tells us that  $\emptyset$  exists. The singleton set axiom (part of Axiom 3.4) and Axiom 3.1 (so that we are allowed to construct the singleton set whose element is another set) tell us that  $\{\emptyset\}$  exists. Applying them again gives that  $\{\{\emptyset\}\}$  exists. Using the pair set axiom (also part of Axiom 3.4) and Axiom 3.1 allows us to create the set  $\{\emptyset, \{\emptyset\}\}$ .

Now we show that these four sets are distinct. Since the sets  $\{\emptyset\}$ ,  $\{\{\emptyset\}\}$ , and  $\{\emptyset, \{\emptyset\}\}$  obviously all contain at least one element, if one of them is equal to  $\emptyset$ , then by Axiom 3.2, it contains no elements. But this contradicts that it has at least one element, and we have a contradiction. Now we need to prove that  $\{\emptyset\}$  is not equal to  $\{\{\emptyset\}\}$  or  $\{\emptyset, \{\emptyset\}\}$ . Both  $\{\emptyset\} \in \{\{\emptyset\}\}$  and  $\{\emptyset, \{\emptyset\}\}$  are true, but  $\{\emptyset\} \notin \{\emptyset\}$ . Therefore, by Axiom 3.2,  $\{\emptyset\}$  is not equal to any of the other sets. (We have already proved that  $\emptyset \neq \{\emptyset\}$ , so  $\{\emptyset\} \neq \emptyset$ ). The last thing that we need to prove now is that  $\{\{\emptyset\}\} \neq \{\emptyset, \{\emptyset\}\}$ . Since  $\emptyset \in \{\emptyset, \{\emptyset\}\}$

and  $\emptyset \notin \{\{\emptyset\}\}$ , by Axiom 3.2, these two sets are not equal. Therefore, all four sets are distinct.

### Exercise 3.1.3

**Proving that**  $\{a, b\} = \{a\} \cup \{b\}$

We shall prove that  $\{a, b\} = \{a\} \cup \{b\}$ . The following are equivalent:

- $c \in \{a, b\}$
- $c$  is either equal to  $a$  or  $b$
- $c \in \{a\}$  or  $c \in \{b\}$
- $c \in \{a\} \cup \{b\}$

Therefore,  $\{a, b\} = \{a\} \cup \{b\}$ .

**Proving that**  $A \cup B = B \cup A$

We now shall prove that  $A \cup B = B \cup A$ . If  $c \in A \cup B$ , then at least one of the statements  $c \in A$  and  $c \in B$  are true. By Axiom 3.5, this means that  $c \in B \cup A$ . We can replace  $A$  with  $B$  and  $B$  with  $A$  to show the other direction. Therefore,  $A \cup B = B \cup A$ .

**Proving that**  $A \cup A = A \cup \emptyset = \emptyset \cup A = A$

Because it is impossible for  $c$  to be in  $\emptyset$ , the following are equivalent:

- $c \in A$  or  $c \in A$
- $c \in A$  or  $c \in \emptyset$
- $c \in \emptyset$  or  $c \in A$
- $c \in A$

These four statements are equivalent to  $c \in A \cup A$ ,  $c \in A \cup \emptyset$ ,  $c \in \emptyset \cup A$ , and  $c \in A$ , respectively. Therefore,  $A \cup A = A \cup \emptyset = \emptyset \cup A = A$ .

### Exercise 3.1.11

The next paragraph will be written like a proof of Axiom 3.6 from Axiom 3.7: we will not define again  $A$ ,  $P(x)$ , and so on.

If  $Q(x, y)$  is the statement that  $x = y$  and  $P(x)$  are both true, then by Axiom 3.7 and the fact that for each  $x$  there exists at most one  $y$  such that  $Q(x, y)$ , there exists a set  $B = \{y : x = y \text{ and } P(x) \text{ for some } x \in A\}$ . For any  $z \in B$ , we know that  $z \in A$  and that  $P(z)$  is true. For any  $x \in A$  such that  $P(x)$  is true, we similarly know that  $x \in B$ . Therefore,  $A = B$ .

## Section 3.2

### Exercise 3.2.1

#### Proving that Axiom 3.9 implies Axiom 3.3

We can create a property  $P(x)$  that is always false regardless to the choice of  $x$ . Then the set  $\{x : P(x) \text{ is true}\}$  has no elements, as  $P(x)$  is false all the time.

#### Proving that Axiom 3.9 implies Axiom 3.4

If  $P(x)$  is the property that  $x = a$ , then  $\{x : P(x) \text{ is true}\}$  has only the element  $a$ . Similarly, if  $Q(x)$  is the property that  $x = a$  or  $x = b$ , then  $\{x : Q(x) \text{ is true}\}$  has only the elements  $a$  and  $b$ .

#### Proving that Axiom 3.9 implies Axiom 3.5

If  $P(x)$  is the property that  $x \in A$  or  $x \in B$ , then  $\{x : P(x) \text{ is true}\}$  exists, so  $A \cup B$  exists.

#### Proving that Axiom 3.9 implies Axiom 3.6

If  $Q(x)$  is the property that  $P(x)$  is true and  $x \in A$ , then  $\{x : Q(x) \text{ is true}\} = \{x \in A : P(x)\}$  exists.

#### Proving that Axiom 3.9 implies Axiom 3.7

If  $Q(y)$  is the property that there exists some  $x$  such that  $P(x, y)$  is true, then  $\{y : Q(y) \text{ is true}\} = \{y : P(x, y) \text{ is true for some } x \in A\}$  exists.

**Proving that Axiom 3.9 implies Axiom 3.8, assuming that all natural numbers are objects**

If  $P(x)$  is the property that  $x$  is a natural number, then  $\mathbf{N} = \{x : P(x) \text{ is true}\}$  exists, by Axiom 3.9.

## Section 3.3

### Exercise 3.3.2

**Proving that  $f$  and  $g$  being injective implies that  $g \circ f$  is too**

We will show that if  $x \neq x'$ , then  $(g \circ f)(x) \neq (g \circ f)(x')$ . First, since  $f$  is injective,  $f(x) \neq f(x')$ . Since  $g$  is also injective,  $g(f(x)) \neq g(f(x'))$ . Therefore,  $(g \circ f)(x) \neq (g \circ f)(x')$ .

**Proving that  $f$  and  $g$  being surjective implies that  $g \circ f$  is too**

We will show that  $g \circ f$  is surjective by showing that for any  $c \in Z$ , there exists some  $a \in X$  such that  $(g \circ f)(a) = c$ . Since  $g$  is surjective, there exists some  $b \in Y$  such that  $g(b) = c$ . Because  $f$  is also surjective, there exists some  $a \in X$  such that  $f(a) = b$ . Therefore,  $(g \circ f)(a) = c$ , and  $g \circ f$  is surjective.

### Exercise 3.3.3

**Finding when the empty function is injective**

We will prove that the empty function  $f: \emptyset \rightarrow X$  is injective for any  $X$ . Since we cannot find any  $x, x' \in \emptyset$  that are unequal, the empty function is (vacuously) injective.

**Finding when the empty function is surjective**

We will prove that the empty function  $f: \emptyset \rightarrow X$  is surjective only when  $X = \emptyset$ . The statement

“For every  $y \in X$ , there exists  $x \in \emptyset$  such that  $f(x) = y$ ”

can only be true when it is impossible that  $y \in X$ , because  $x \in \emptyset$  is impossible. But then we have that  $X = \emptyset$ , and we have proved our claim.

### Finding when the empty function is bijective

We will prove that the empty function  $f: \emptyset \rightarrow X$  is bijective only when  $X = \emptyset$ . Since  $f$  is always injective regardless of the choice of  $X$ , the empty function being bijective is equivalent to it being surjective. Therefore, since  $f$  is surjective precisely when  $X = \emptyset$ , the empty function is bijective only when  $X = \emptyset$ .

### Exercise 3.3.4

**Showing that if  $g \circ f = g \circ \tilde{f}$  and  $g$  is injective, then  $f = \tilde{f}$**

We will use proof by contradiction. If  $f \neq \tilde{f}$ , then for some  $x \in X$ , we have  $f(x) \neq \tilde{f}(x)$ . Therefore, since  $g$  is injective,  $(g \circ f)(x) \neq (g \circ \tilde{f})(x)$ , and  $g \circ f \neq g \circ \tilde{f}$ . This is a contradiction, and therefore,  $f = \tilde{f}$ . This statement is not necessarily true if  $g$  is not injective. If  $X = Y = Z = \mathbf{N}$ ,  $f(x) = 0$ ,  $\tilde{f}(x) = 1$ , and  $g(x) = 0$ , then  $f \neq \tilde{f}$ , while  $(g \circ f)(x) = 0 = (g \circ \tilde{f})(x)$ .

**Showing that if  $g \circ f = \tilde{g} \circ f$  and  $f$  is surjective, then  $g = \tilde{g}$**

We will use proof by contradiction again. If  $g \neq \tilde{g}$ , then for some  $y \in Y$ , we have  $g(y) \neq \tilde{g}(y)$ . Since  $f$  is surjective, there exists  $x \in X$  such that  $f(x) = y$ . But then  $(g \circ f)(x) \neq (\tilde{g} \circ f)(x)$ . This is a contradiction, and therefore,  $g = \tilde{g}$ . This statement is not necessarily true if  $f$  is not surjective. If  $X = Y = Z = \mathbf{N}$ ,  $f(x) = 0$ ,  $g(x) = x$ , and  $\tilde{g}(x) = 0$ , then  $g \neq \tilde{g}$ , even though  $(g \circ f)(x) = 0 = (\tilde{g} \circ f)(x)$ .

### Exercise 3.3.6

If  $f(x) = a$ , then by the definition of  $f^{-1}$ ,  $f^{-1}(f(x)) = f^{-1}(a) = x$ . If  $f^{-1}(y) = b$ , then  $f(b) = y$ , so  $f(f^{-1}(y)) = f(b) = y$ .

We can deduce that  $f^{-1}$  is bijective from Exercise 3.3.5 and the fact that the identity map  $\iota_{X \rightarrow X}$  defined as  $\iota_{X \rightarrow X}(x) = x$  for all  $x \in X$  is obviously bijective.

### Exercise 3.3.7

By Exercise 3.3.2,  $g \circ f$  is both injective and surjective, and is therefore bijective. The only thing left we have to prove is that  $(f^{-1} \circ g^{-1} \circ g \circ f)(x) = x$ .



But  $(f^{-1} \circ g^{-1} \circ g \circ f)(x) = (f^{-1} \circ f)(x) = x$ , and we have proved what we want.

## Section 3.4

### Challenge to define $f(S)$ using the axiom of specification

We can define  $f(S) := \{y \in Y : \text{there exists } x \in S \text{ such that } f(x) = y\}$ .

#### Exercise 3.4.1

First, we will show that any element of the forward image of  $V$  under  $f^{-1}$  is contained in the inverse image of  $V$  under  $f$ . If  $x$  is in the forward image of  $V$  under  $f^{-1}$ , then there exists  $v$  such that  $f^{-1}(v) = x$  and therefore  $f(x) = v$ . But then  $x$  is in the inverse image of  $V$  under  $f$ .

Next, we will show that any element  $v$  of the inverse image of  $V$  under  $f$  is contained in the forward image of  $V$  under  $f^{-1}$ . Since there exists  $x$  such that  $f(x) = v$  and therefore  $f^{-1}(v) = x$ ,  $x$  is in the forward image of  $V$  under  $f$ .

Finally, we have that the forward image of  $V$  under  $f^{-1}$  is equal to the inverse image of  $V$  under  $f$ , and that it is valid to use the notation  $f^{-1}(V)$ .

#### Exercise 3.4.6

By Axiom 3.11,  $\{0, 1\}^X$  exists. We can use Axiom 3.7 (the axiom of replacement) to create a set

$$Y = \{f^{-1}(\{1\}) : f \in \{0, 1\}^X\}.$$

Every  $S$  which is a subset of  $X$  is in  $Y$ , because defining  $f(x)$  to be 1 when  $x$  is in  $S$  and to be 0 otherwise means that  $f^{-1}(\{1\}) = S$ . Every element of  $Y$  is a subset of  $X$ , because every element of  $f^{-1}(\{1\})$  has to be in  $X$ .

#### Exercise 3.4.9

We will show that

$$\{x \in A_\beta : x \in A_\alpha \text{ for all } \alpha \in I\} = \{x \in A_{\beta'} : x \in A_\alpha \text{ for all } \alpha \in I\}.$$

We know that  $y \in \{x \in A_\beta : x \in A_\alpha \text{ for all } \alpha \in I\}$  if and only if  $y \in A_\beta$  and  $y \in A_\alpha$  for all  $\alpha \in I$ . But the latter statement implies the former, so  $y \in \{x \in A_\beta : x \in A_\alpha \text{ for all } \alpha \in I\}$  if and only if  $y \in A_\alpha$  for all  $\alpha \in I$ .

We can do something similar for  $A_{\beta'}$ . Therefore,

$$\{x \in A_\beta : x \in A_\alpha \text{ for all } \alpha \in I\} = \{x \in A_{\beta'} : x \in A_\alpha \text{ for all } \alpha \in I\}.$$

## Section 3.5

### Exercise 3.5.1

First, we will show that  $(x, y) = (x', y')$  implies that both  $x = x'$  and  $y = y'$ . Since  $(x, y) := \{\{x\}, \{x, y\}\}$ , we have to show that  $\{\{x\}, \{x, y\}\} = \{\{x'\}, \{x', y'\}\}$  implies that both  $x = x'$  and  $y = y'$ . If  $\{\{x\}, \{x, y\}\} = \{\{x'\}, \{x', y'\}\}$ , then  $\{x\} \in \{\{x'\}, \{x', y'\}\}$ . In order for this to be true, we must have  $\{x\} = \{x'\}$ , which implies  $x = x'$ , or  $\{x\} = \{x', y'\}$ , which implies  $x' = y' = x$ . But the latter case implies the former case, so we can ignore the latter one. Doing the same with  $\{x, y\}$ ,  $\{x'\}$ , and  $\{x', y'\}$  yields that  $x = x'$  and  $y = y'$ . The converse is true by Axiom 3.2.

We can show that  $X \times Y$  is a set if  $X$  and  $Y$  are both sets using Axiom 3.1 (sets are objects), Axiom 3.7 (axiom of replacement), and Axiom 3.12 (axiom of union). By Axioms 3.1 and 3.7, we can create the set  $Z = \{\{(x, y)\} : y \in Y\} : x \in X\}$ . By Axiom 3.12, we can also create the set  $\bigcup Z$ , which we can define to be  $X \times Y$ .

I won't do the additional challenge.

### Exercise 3.5.6

We will show that if  $A \times B \subseteq C \times D$  and  $A, B, C, D \neq \emptyset$ , then  $A \subseteq C$  and  $B \subseteq D$ . Since  $A$  and  $B$  are nonempty, by the single choice lemma (Lemma 3.1.5), there exist some  $a \in A$  and  $b \in B$ . Now, for any  $x \in A$ , since  $(x, b) \in A \times B \subseteq C \times D$ , we have  $x \in C$ . Similarly, for any  $y \in B$ , since  $(a, y) \in A \times B \subseteq C \times D$ , we have  $y \in D$ . Therefore,  $A \subseteq C$  and  $B \subseteq D$ .

However, we can show that if we remove the assumption that  $A, B, C$ , and  $D$  are nonempty, then this statement is false. If  $A = \{0\}$ ,  $C = \{1\}$ , and  $B = D = \emptyset$ , then  $A \times B$  is empty, as there are no elements of  $B$ . Similarly,  $C \times D = \emptyset$ , and therefore  $A \times B \subseteq C \times D$ . However,  $A \not\subseteq C$ .

**Exercise 3.5.10**

**Exercise 3.5.11**

**Exercise 3.5.13**