# Answers to Analysis I Exercises, Third Edition

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# Chapter 1

# Introduction

1.1 What is analysis?

N/A

1.2 Why do analysis?

N/A

# Chapter 2

# Starting at the beginning: the natural numbers

# 2.1 The Peano axioms

N/A

# 2.2 Addition

#### Exercise 2.2.1

We shall use induction on a, fixing b and c. First, we shall prove the base case. If a=0, a+b=0+b=b by the definition of addition of natural numbers (Definition 2.2.1). Therefore, (a+b)+c=b+c. Also by Definition 2.2.1, a+(b+c)=0+(b+c)=b+c. We now have (a+b)+c=b+c=a+(b+c). Now suppose that (a+b)+c=a+(b+c). We then have, by Definition 2.2.1, that ((a++)+b)+c=((a+b)++)+c=((a+b)+c)++=(a+(b+c))++. Because of Definition 2.2.1 again, (a+(b+c))++=(a++)+(b+c). Now we have ((a++)+b)+c=(a++)+(b+c), which closes the induction.

#### Exercise 2.2.2

We shall use induction on a. The base case is vacuously true since 0 is not positive. For the inductive step, we have to show that a++ being positive implies that a++=b++ for exactly one natural number b if a being positive

implies that a = c++ for exactly one natural number c. But a is equal to b if a++ is equal to b++, by Axiom 2.4. This closes the induction.

## Exercise 2.2.3

- (a) By Lemma 2.2.2, a = a + 0, so  $a \ge a$ .
- (b) Because  $a \ge b$  and  $b \ge c$ , a = b + n and b = c + m for some n and m. Therefore, a = b + n = (c + m) + n. By Proposition 2.2.5, (c + m) + n = c + (m + n), so a = c + (m + n). n + m is a natural number, so  $a \ge c$ .
- (c) Because  $a \ge b$  and  $b \ge a$ , a = b + n and b = a + m. Therefore, a = b + n = (a + m) + n. By Proposition 2.2.5, (a + m) + n = a + (m + n). Now we have a = a + (m + n). By Lemma 2.2.2, a = a + 0, so we can use Proposition 2.2.6 to get 0 = m + n. Using Corollary 2.2.9, we can deduce that n and m are both equal to 0. Therefore, a = b.
- (d) If  $a \ge b$ , a = b + d for some d. Because of this, a + c = (b + d) + c. By Propositions 2.2.4 and 2.2.5, (b+d)+c = b+(d+c) = b+(c+d) = (b+c)+d. Therefore, a + c > b + c.

Conversely, if  $a+c \ge b+c$ , a+c = (b+c)+d for some d. By Propositions 2.2.4 and 2.2.5, (b+c)+d=b+(c+d)+b+(d+c)=(b+d)+c. Using Proposition 2.2.6, we get a=b+d, and therefore,  $a \ge b$ .

(f) If a < b, b = a + c for some c, and  $b \ne a$ . Therefore, c is positive, because otherwise b = a + 0 = a. (a + 0 = a because of Lemma 2.2.2.)

Conversely, if b = a + c for some positive number c, we know that  $a \le b$ . If a = b, a = a + c. By Lemma 2.2.2, a = a + 0, so a + 0 = a + c. We can use Proposition 2.2.6 to get 0 = c, but c is positive. Therefore,  $a \ne b$ , and combined with  $a \le b$ , we get a < b.

(e) We first prove that a < b implies that  $a++ \le b$ . By (f), b=a+c for some positive c. By Lemma 2.2.10, exactly one d exists such that d++=c. Therefore, b=a+d++. By Lemma 2.2.3 and Definition 2.2.1, b=a+d++=(a+d)++=(a+d)++=(a+d)++=b.

Conversely, if  $a++ \le b$ , we need to prove that a < b. Since  $a++ = a+1 \ge a$ , by (b),  $a \le b$ . If a = b, then  $a++ \le a$ , and we get, from (c), that a++ = a. This is a contradiction, from Axiom 2.3 if a = 0 and Axiom 2.4 otherwise. Therefore, a < b.

# Exercise 2.2.6 (unfinished)

We use induction on n. When n = 0,  $m \le n$  implies that 0 = m + a. But by Corollary 2.2.9, m = 0 = n, so P(n) being true implies that P(m) is true.

If  $a \leq n$  and P(n) together imply P(a), we now need to prove that  $b \leq n++$  and P(n++) together imply P(b). Since  $b \leq n++$ , n++=b+c for some c. If c is positive, by Lemma 2.2.10, n++=b+d++ for some d, so by Lemma 2.2.3 and Axiom 2.4, n=b+d, and  $b \leq n$ . Therefore,  $b \leq n++$  implies that  $b \leq n$  or n++=b+0=b. Because of Lemmas 2.2.2 and 2.2.3, n++=(n+0)++=n+0++=n+1. Since P(n++) implies P(n), P(n) is true. But then ?

# 2.3 Multiplication

## Exercise 2.3.1

First, we prove that  $n \cdot 0 = 0$ . We induct on n. When n = 0,  $n \cdot 0 = 0 \cdot 0 = 0$ , because of Definition 2.3.1. For the inductive hypothesis, we assume  $n \cdot 0 = 0$ . Then we need to prove that  $n++\cdot 0=0$ . By Definition 2.3.1,  $n++\cdot 0=n\cdot 0+0=n\cdot 0$ . But we already know that  $n\cdot 0=0$ , so  $n++\cdot 0=0$ . This closes the induction.

Next, we prove that  $n \cdot m++=nm+n$  using induction on n. When  $n=0, 0 \cdot m++=0=0 \cdot m+0$ . Now, if  $n \cdot m++=nm+n$ , we need to prove that  $n++\cdot m++=n++\cdot m+n++$ . We can deduce from  $n \cdot m++=nm+n$  that  $n++\cdot m++=(n \cdot m++)+m++=nm+n+m++=nm+(n+m)++=nm+n++++m=(n++)\cdot m+n+++$ . This closes the induction.

Now we can prove that nm = mn. We induct on m. When m = 0, we have  $n \cdot 0 = 0 = 0 \cdot n$ . For the inductive hypothesis, we assume nm = mn. Now we need to prove  $n \cdot m++=m++\cdot n$ . We know that  $n \cdot m++=nm+n$  and  $m++\cdot n = mn+n$ . But nm = mn, so  $n \cdot m++=m++\cdot n$ . This closes the induction.

#### Exercise 2.3.4

By Proposition 3.4,  $(a+b)^2 = (a+b) \cdot (a+b) = (a+b) \cdot a + (a+b) \cdot b = a \cdot a + b \cdot a + a \cdot b + b \cdot b$ . We can rewrite  $b \cdot a + a \cdot b$  as  $a \cdot b + a \cdot b$  because of Lemma 2.3.2. We can further rewrite this as 2ab because 2(ab) = 1(ab) + ab = ab

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0(ab)+ab+ab=0+ab+ab=ab+ab. We can also rewrite  $a\cdot a$  and  $b\cdot b$  as  $a^2$  and  $b^2$  respectively. Therefore,  $(a+b)^2=a^2+2ab+b^2$ .

# Chapter 3

# Set theory

# 3.1 Fundamentals

#### Exercise 3.3.1

We know that at least one of the statements "a=c" and "a=d" is true since  $\{a,b\}=\{c,d\}$ . Similarly, at least one of the statements "b=c" and "b=d" is true. If we have that a=c and b=d at the same time, then we have proved what we want. (Same for a=d and b=c at the same time.) Otherwise, then we must have that a=c=b or a=d=b. If both of these are true, then a=b=c=d, and therefore both a=c and b=d. If only one of these is true, which we will assume to be the statement that a=c=b as a similar argument holds assuming a=d=b, then  $d\notin \{a,b\}=\{a,a\}$ , while  $d\in \{c,d\}=\{a,d\}$ , so  $\{a,b\}\neq \{c,d\}$ , a contradiction. Therefore, if  $\{a,b\}=\{c,d\}$ , either both a=c and b=d or a=d and b=c.

#### Exercise 3.3.2

First, we show that these four sets exist. Axiom 3.3 tells us that  $\emptyset$  exists. The singleton set axiom (part of Axiom 3.4) and Axiom 3.1 (so that we are allowed to construct the singleton set whose element is another set) tell us that  $\{\emptyset\}$  exists. Applying them again gives that  $\{\{\emptyset\}\}$  exists. Using the pair set axiom (also part of Axiom 3.4) and Axiom 3.1 allows us to create the set  $\{\emptyset, \{\emptyset\}\}$ .

Now we show that these four sets are distinct. Since the sets  $\{\emptyset\}$ ,  $\{\{\emptyset\}\}$ , and  $\{\emptyset, \{\emptyset\}\}$  obviously all contain at least one element, if one of them is equal

to  $\emptyset$ , then by Axiom 3.2, it contains no elements. But this contradicts that it has at least one element, and we have a contradiction. Now we need to prove that  $\{\emptyset\}$  is not equal to  $\{\{\emptyset\}\}$  or  $\{\emptyset, \{\emptyset\}\}$ . Both  $\{\emptyset\} \in \{\{\emptyset\}\}\}$  and  $\{\emptyset, \{\emptyset\}\}\}$  are true, but  $\{\emptyset\} \notin \{\emptyset\}$ . Therefore, by Axiom 3.2,  $\{\emptyset\}$  is not equal to any of the other sets. (We have already proved that  $\emptyset \neq \{\emptyset\}$ , so  $\{\emptyset\} \neq \emptyset$ ). The last thing that we need to prove now is that  $\{\{\emptyset\}\} \neq \{\emptyset, \{\emptyset\}\}\}$ . Since  $\emptyset \in \{\emptyset, \{\emptyset\}\}\}$  and  $\emptyset \notin \{\{\emptyset\}\}$ , by Axiom 3.2, these two sets are not equal. Therefore, all four sets are distinct.

## Exercise 3.1.3

**Proving that**  $\{a,b\} = \{a\} \cup \{b\}$  We shall prove that  $\{a,b\} = \{a\} \cup \{b\}$ . The following are equivalent:

- $c \in \{a, b\}$
- $\bullet$  c is either equal to a or b
- $c \in \{a\}$  or  $c \in \{b\}$
- $c \in \{a\} \cup \{b\}$

Therefore,  $\{a, b\} = \{a\} \cup \{b\}$ .

**Proving that**  $A \cup B = B \cup A$  We now shall prove that  $A \cup B = B \cup A$ . If  $c \in A \cup B$ , then at least one of the statements  $c \in A$  and  $c \in B$  are true. By Axiom 3.5, this means that  $c \in B \cup A$ . We can replace A with B and B with A to show the other direction. Therefore,  $A \cup B = B \cup A$ .

**Proving that**  $A \cup A = A \cup \emptyset = \emptyset \cup A = A$  Because it is impossible for c to be in  $\emptyset$ , the following are equivalent:

- $c \in A$  or  $c \in A$
- $c \in A \text{ or } c \in \emptyset$
- $c \in \emptyset$  or  $c \in A$
- $c \in A$

These four statements are equivalent to  $c \in A \cup A$ ,  $c \in A \cup \emptyset$ ,  $c \in \emptyset cup A$ , and  $c \in A$ , respectively. Therefore,  $A \cup A = A \cup \emptyset = \emptyset \cup A = A$ .

#### Exercise 3.1.11

The next paragraph will be written like a proof of Axiom 3.6 from Axiom 3.7: we will not define again A, P(x), and so on.

If Q(x,y) is the statement that x=y and P(x) are both true, then by Axiom 3.7 and the fact that for each x there exists at most one y such that Q(x,y), there exists a set  $B=\{y:x=y\text{ and }P(x)\text{ for some }x\in A\}$ . For any  $z\in B$ , we know that  $z\in A$  and that P(z) is true. For any  $x\in A$  such that P(x) is true, we similarly know that  $z\in B$ . Therefore, A=B.

# 3.2 Russell's paradox (Optional)

## Exercise 3.2.1

**Proving that Axiom 3.9 implies Axiom 3.3** We can create a property P(x) that is always false regardless to the choice of x. Then the set  $\{x : P(x) \text{ is true}\}$  has no elements, as P(x) is false all the time.

**Proving that Axiom 3.9 implies Axiom 3.4** If P(x) is the property that x = a, then  $\{x : P(x) \text{ is true}\}$  has only the element a. Similarly, if Q(x) is the property that x = a or x = b, then  $\{x : Q(x) \text{ is true}\}$  has only the elements a and b.

**Proving that Axiom 3.9 implies Axiom 3.5** If P(x) is the property that  $x \in A$  or  $x \in B$ , then  $\{x : P(x) \text{ is true}\}$  exists, so  $A \cup B$  exists.

**Proving that Axiom 3.9 implies Axiom 3.6** If Q(x) is the property that P(x) is true and  $x \in A$ , then  $\{x : Q(x) \text{ is true}\} = \{x \in A : P(x)\}$  exists.

**Proving that Axiom 3.9 implies Axiom 3.7** If Q(y) is the property that there exists some x such that P(x, y) is true, then  $\{y : Q(y) \text{ is true}\} = \{y : P(x, y) \text{ is true for some } x \in A\}$  exists.

Proving that Axiom 3.9 implies Axiom 3.8, assuming that all natural numbers are objects If P(x) is the property that x is a natural number, then  $\mathbf{N} = \{x : P(x) \text{ is true}\}$  exists, by Axiom 3.9.

# 3.3 Functions

#### Exercise 3.3.2

Proving that f and g being injective implies that  $g \circ f$  is too We will show that if  $x \neq x'$ , then  $(g \circ f)(x) \neq (g \circ f)(x')$ . First, since f is injective,  $f(x) \neq f(x')$ . Since g is also injective,  $g(f(x)) \neq g(f(x'))$ . Therefore,  $(g \circ f)(x) \neq (g \circ f)(x')$ .

**Proving that** f and g being surjective implies that  $g \circ f$  is too We will show that  $g \circ f$  is surjective by showing that for any  $c \in Z$ , there exists some  $a \in X$  such that  $(g \circ f)(a) = c$ . Since g is surjective, there exists some  $b \in Y$  such that g(b) = c. Because f is also surjective, there exists some  $a \in X$  such that f(a) = b. Therefore,  $(g \circ f)(a) = c$ , and  $g \circ f$  is surjective.

## Exercise 3.3.3

Finding when the empty function is injective We will prove that the empty function  $f: \emptyset \to X$  is injective for any X. Since we cannot find any  $x, x' \in \emptyset$  that are unequal, the empty function is (vacuously) injective.

Finding when the empty function is surjective We will prove that the empty function  $f: \emptyset \to X$  is surjective only when  $X = \emptyset$ . The statement

"For every  $y \in X$ , there exists  $x \in \emptyset$  such that f(x) = y"

can only be true when it is impossible that  $y \in X$ , because  $x \in \emptyset$  is impossible. But then we have that  $X = \emptyset$ , and we have proved our claim.

Finding when the empty function is bijective We will prove that the empty function  $f: \emptyset \to X$  is bijective only when  $X = \emptyset$ . Since f is always injective regardless of the choice of X, the empty function being bijective is equivalent to it being surjective. Therefore, since f is surjective precisely when  $X = \emptyset$ , the empty function is bijective only when  $X = \emptyset$ .

#### Exercise 3.3.4

Showing that if  $g \circ f = g \circ \tilde{f}$  and g is injective, then  $f = \tilde{f}$  We will use proof by contradiction. If  $f \neq \tilde{f}$ , then for some  $x \in X$ , we have  $f(x) \neq \tilde{f}(x)$ .

Therefore, since g is injective,  $(g \circ f)(x) \neq (g \circ \tilde{f})(x)$ , and  $g \circ f \neq g \circ \tilde{f}$ . This is a contradiction, and therefore,  $f = \tilde{f}$ . This statement is not necessarily true if g is not injective. If  $X = Y = Z = \mathbb{N}$ , f(x) = 0,  $\tilde{f}(x) = 1$ , and g(x) = 0, then  $f \neq \tilde{f}$ , while  $(g \circ f)(x) = 0 = (g \circ \tilde{f})(x)$ .

Showing that if  $g \circ f = \tilde{g} \circ f$  and f is surjective, then  $g = \tilde{g}$  We will use proof by contradiction again. If  $g \neq \tilde{g}$ , then for some  $y \in Y$ , we have  $g(y) \neq \tilde{g}(y)$ . Since f is surjective, there exists  $x \in X$  such that f(x) = y. But then  $(g \circ f)(x) \neq (\tilde{g} \circ f)(x)$ . This is a contradiction, and therefore,  $g = \tilde{g}$ . This statement is not necessarily true if f is not surjective. If  $X = Y = Z = \mathbb{N}$ , f(x) = 0, g(x) = x, and  $\tilde{g}(x) = 0$ , then  $g \neq \tilde{g}$ , even though  $(g \circ f)(x) = 0 = (\tilde{g} \circ f)(x)$ .

#### Exercise 3.3.6

If f(x) = a, then by the definition of  $f^{-1}$ ,  $f^{-1}(f(x)) = f^{-1}(a) = x$ . If  $f^{-1}(y) = b$ , then f(b) = y, so  $f(f^{-1}(y)) = f(b) = y$ .

We can deduce that  $f^{-1}$  is bijective from Exercise 3.3.5 and the fact that the identity map  $i_{X\to X}$  defined as  $i_{X\to X}(x)=x$  for all  $x\in X$  is obviously bijective.

## Exercise 3.3.7

By Exercise 3.3.2,  $g \circ f$  is both injective and surjective, and is therefore bijective. The only thing left we have to prove is that  $(f^{-1} \circ g^{-1} \circ g \circ f)(x) = x$ . But  $(f^{-1} \circ g^{-1} \circ g \circ f)(x) = (f^{-1} \circ f)(x) = x$ , and we have proved what we want.

# 3.4 Images and inverse images

# Challenge to define f(S) using the axiom of specification

We can define  $f(S) := \{ y \in Y : \text{there exists } x \in S \text{ such that } f(x) = y \}.$ 

#### Exercise 3.4.1

First, we will show that any element of the forward image of V under  $f^{-1}$  is contained in the inverse image of V under f. If x is in the forward image of V

under  $f^{-1}$ , then there exists v such that  $f^{-1}(v) = x$  and therefore f(x) = v. But then x is in the inverse image of V under f.

Next, we will show that any element v of the inverse image of V under f is contained in the forward image of V under  $f^{-1}$ . Since there exists x such that f(x) = v and therefore  $f^{-1}(v) = x$ , x is in the forward image of V under f.

Finally, we have that the forward image of V under  $f^{-1}$  is equal to the inverse image of V under f, and that it is valid to use the notation  $f^{-1}(V)$ .

#### Exercise 3.4.6

By Axiom 3.11,  $\{0,1\}^X$  exists. We can use Axiom 3.7 (the axiom of replacement) to create a set

$$Y = \{f^{-1}(\{1\}) : f \in \{0, 1\}^X\}.$$

Every S which is a subset of X is in Y, because defining f(x) to be 1 when x is in S and to be 0 otherwise means that  $f^{-1}(\{1\}) = S$ . Every element of Y is a subset of X, because every element of  $f^{-1}(\{1\})$  has to be in X.

#### Exercise 3.4.9

We will show that

$$\{x \in A_{\beta} : x \in A_{\alpha} \text{ for all } \alpha \in I\} = \{x \in A_{\beta'} : x \in A_{\alpha} \text{ for all } \alpha \in I\}.$$

We know that  $y \in \{x \in A_{\beta} : x \in A_{\alpha} \text{ for all } \alpha \in I\}$  if and only if  $y \in A_{\beta}$  and  $y \in A_{\alpha}$  for all  $\alpha \in I$ . But the latter statement implies the former, so  $y \in \{x \in A_{\beta} : x \in A_{\alpha} \text{ for all } \alpha \in I\}$  if and only if  $y \in A_{\alpha}$  for all  $\alpha \in I$ .

We can do something similar for  $A_{\beta'}$ . Therefore,

$$\{x \in A_{\beta} : x \in A_{\alpha} \text{ for all } \alpha \in I\} = \{x \in A_{\beta'} : x \in A_{\alpha} \text{ for all } \alpha \in I\}.$$

# 3.5 Cartesian products

#### Exercise 3.5.1

First, we will show that (x,y) = (x',y') implies that both x = x' and y = y'. Since  $(x,y) \coloneqq \{\{x\}, \{x,y\}\}$ , we have to show that  $\{\{x\}, \{x,y\}\}$ 

 $\{\{x'\}, \{x', y'\}\}\$  implies that both x=x' and y=y'. If  $\{\{x\}, \{x, y\}\}\}=\{\{x'\}, \{x', y'\}\}$ , then  $\{x\} \in \{\{x'\}, \{x', y'\}\}$ . In order for this to be true, we must have  $\{x\} = \{x'\}$ , which implies x=x', or  $\{x\} = \{x', y'\}$ , which implies x'=y'=x. But the latter case implies the former case, so we can ignore the latter one. Doing the same with  $\{x, y\}, \{x'\},$  and  $\{x', y'\}$  yields that x=x' and y=y'. The converse is true by Axiom 3.2.

We can show that  $X \times Y$  is a set if X and Y are both sets using Axiom 3.1 (sets are objects), Axiom 3.7 (axiom of replacement), and Axiom 3.12 (axiom of union). By Axioms 3.1 and 3.7, we can create the set  $Z = \{\{(x,y)\} : y \in Y\} : x \in X\}$ . By Axiom 3.12, we can also create the set  $\bigcup Z$ , which we can define to be  $X \times Y$ .

I won't do the additional challenge.

#### Exercise 3.5.6

We will show that if  $A \times B \subseteq C \times D$  and  $A, B, C, D \neq \emptyset$ , then  $A \subseteq C$  and  $B \subseteq D$ . Since A and B are nonempty, by the single choice lemma (Lemma 3.1.5), there exist some  $a \in A$  and  $b \in B$ . Now, for any  $x \in A$ , since  $(x, b) \in A \times B \subseteq C \times D$ , we have  $x \in C$ . Similarly, for any  $y \in B$ , since  $(a, y) \in A \times B \subseteq C \times D$ , we have  $y \in D$ . Therefore,  $A \subseteq C$  and  $B \subseteq D$ .

However, we can show that if we remove the assumption that A, B, C, and D are nonempty, then this statement is false. If  $A = \{0\}$ ,  $C = \{1\}$ , and  $B = D = \emptyset$ , then  $A \times B$  is empty, as there are no elements of B. Similarly,  $C \times D = \emptyset$ , and therefore  $A \times B \subseteq C \times D$ . However,  $A \not\subseteq C$ .

If  $A \subseteq C$  and  $B \subseteq D$ , all ordered pairs (a, b) where  $a \in A$  and  $b \in B$  are in  $C \times D$  because  $a \in C$  and  $b \in D$ . Hence,  $A \times B \subseteq C \times D$ . This statement holds even when A, B, C, and D are not all nonempty.

Assuming that  $A \times B = C \times D$  and  $A, B, C, D \neq \emptyset$ , A = C and B = D. This is because  $A \times B \subseteq C \times D$  and  $C \times D \supseteq A \times B$ , and therefore  $A \subseteq C$ ,  $A \supseteq C$ ,  $B \subseteq D$ , and  $B \supseteq D$ .

If we remove the assumption that  $A, B, C, D \neq \emptyset$ , the statement is false. The counterexample from before  $(A = \{0\}, C = \{1\}, B = D = \emptyset)$  also works here. We have that  $A \times B = C \times D = \emptyset$ , but  $A \neq C$ .

We have that A = C and B = D imply that  $A \times B = C \times D$  because  $A \subseteq C$  and  $B \subseteq D$  imply, together, that  $A \times B \subseteq C \times D$  and the same with  $\subseteq$  replaced by  $\supseteq$ .

# Exercise 3.5.13 (unfinished)

?

# 3.6 Cardinality of sets

## Exercise 3.6.2

If X has cardinality 0, there exists a bijection  $f: X \to \{i \in \mathbb{N} : 1 \le i \le 0\}$ . The set  $\{i \in \mathbb{N} : 1 \le i \le 0\}$  is obviously empty. If  $X \ne \emptyset$ , then there exists some  $x \in X$ . But then f(x) is in  $\emptyset$ , and therefore X must be empty.

If X is empty, the function  $f: X \to \emptyset$  exists (we don't need to give how to derive f(x) given x, since  $x \in X$  is impossible) and is a bijection (injectivity is because there are no elements in the domain at all, and surjectivity is from that the range is empty). Since  $\emptyset = \{i \in \mathbf{N} : 1 \le i \le 0\}$ , X has cardinality 0.

#### Exercise 3.6.3

We will induct on n. When n=0,  $\{i\in N:1\leq i\leq n\}$  is empty, so we can make M equal to any natural number, say, 0. Now, assuming that all  $f\colon\{i\in N:1\leq i\leq n\}\to \mathbf{N}$  satisfy the property that there exists some natural number M such that  $f(i)\leq M$  for all  $1\leq i\leq n$ , we will now show that for every  $g\colon\{i\in N:1\leq i\leq n++\}\to \mathbf{N}$ , there exists some natural number a such that  $g(i)\leq a$  for all  $1\leq i\leq n++$ . If  $h\colon\{i\in N:1\leq i\leq n++\}\to \mathbf{N}$  is defined so that h(x):=g(x) for all  $x\in\{i\in N:1\leq i\leq n++\}$ , there exists some b such that  $h(i)\leq b$  for all  $1\leq i\leq n$ . Defining  $a:=\begin{cases}b&\text{if }g(n++)\leq b\\g(n++)&\text{if }g(n++)>b\end{cases}$  ensures that a will always be greater than or equal to all elements of  $g(\{i\in N:1\leq i\leq n++\})$ .

#### Exercise 3.6.5

We will define  $f: A \times B \to B \times A$  by f(a, b) := (b, a), where  $a \in A$  and  $b \in B$ . Injectivity of f comes from (b, a) = (d, c) implying that b = d and a = c, further implying that (a, b) = (c, d). Surjectivity comes from the fact that for every  $(b, a) \in B \times A$  has a corresponding ordered pair (namely, (a, b)) in  $A \times B$  where f(a, b) = (b, a).

Now, we will prove that ab = ba. We will define A as  $\{i \in \mathbb{N} : 1 \le i \le a\}$  and B as  $\{j \in \mathbb{N} : 1 \le j \le b\}$ . Since A has equal cardinality with itself and the same for B (Proposition 3.6.4), A and B are both finite and have the cardinalities a and b, respectively. By Proposition 3.6.14(e),  $\#(A) \cdot \#(B) = \#(A \times B)$  and  $\#(B) \cdot \#(A) = \#(B \times A)$ . We also know that  $\#(A \times B) = \#(B \times A)$ , and therefore ab = ba.

# Exercise 3.6.7 (unfinished)

?

# Exercise 3.6.10 (unfinished)

**Attempt 1** We induct on n. (I don't think we can induct, actually, since  $\#(\bigcup_{i\in\{1,\dots,n+1\}}A_i) > n+1$  does not necessarily imply  $\#(\bigcup_{i\in\{1,\dots,n\}}A_i) > n$ ). When n=0,  $\#(\bigcup_{i\in\{1,\dots,n\}}A_i) = \#(\emptyset) = 0$ . It is obvious that  $0 \not> 0$ , so the pigeonhole principle holds vacuously when n=0. When the pigeonhole principle is true for n finite sets, now we will show that it also is true for  $A_1,\dots,A_{n++}$ . By Proposition 3.6.14(b),

$$\#(\bigcup_{i \in \{1,\dots,n++\}} A_i) \le \#(\bigcup_{i \in \{1,\dots,n\}} A_i) + \#(A_{n++}).$$

# Chapter 4

# Integers and rationals

# 4.1 The integers

#### Exercise 4.1.1

**Reflexive** If  $a, b \in \mathbb{N}$ , a + b = b + a, so by Definition 4.1.1, a - b = a - b.

Symmetric If a-b=c-d, a+d=c+b. Now, since c+b=a+d, c-d=a-b.

# Exercise 4.1.3

We know that -1 = -(1 - 0) = 0 - 1. If  $a = a_1 - a_2$ ,  $(-1) \cdot a = (0 \cdot a_1 + 1 \cdot a_2) - (0 \cdot a_2 + 1 \cdot a_1) = a_2 - a_1 = -a$ .

## Exercise 4.1.4

I will write  $x=a-b,\ y=c-d,$  and z=e-f, where  $a,b,c,d,e,f\in\mathbf{N}.$  First, we will prove x+y=y+x:

$$x + y = (a - b) + (c - d)$$

$$= (a + c) - (b + d)$$

$$y + x = (c - d) + (a - b)$$

$$= (c + a) - (d + b)$$

and therefore, because of natural number addition's commutativity, x + y = y + x.

Next, we will prove that (x + y) + z = x + (y + z):

$$\begin{split} (x+y) + z &= ((a-b) + (c-d)) + (e-f) \\ &= ((a+c) - (b+d)) + (e-f) \\ &= ((a+c) + e) - ((b+d) + f) \\ x + (y+z) &= (a-b) + ((c-d) + (e-f)) \\ &= (a-b) + ((c+e) - (d+f)) \\ &= (a+(c+e)) - (b+(d+f)) \end{split}$$

and therefore, because of natural number addition's associativity, (x+y)+z=x+(y+z).

Now, we will prove that x + 0 = 0 + x = x. We have proved above that x + y = y + x and therefore x + 0 = 0 + x, so we only need to show x + 0 = x:

$$x + 0 = (a - b) + (0 - 0)$$
  
=  $(a + 0) - (b + 0)$ .

We know that a+0=a and b+0=b, so x+0=x, and hence x+0=0+x=x.

Next, we will prove that x + (-x) = (-x) + x = 0. We have already proved that x + y = y + x and therefore x + (-x) = (-x) + x, so the only thing left is to show x + (-x) = 0.

$$x + (-x) = (a - b) + (b - a)$$

$$= (a + b) - (b + a)$$

$$= (a + b) - (a + b)$$

$$= 0 - 0$$

$$= 0.$$

Therefore, x + (-x) = (-x) + x = 0.

To show that xy = yx, we will use the definition of integer multiplication:

$$xy = (a - b)(c - d)$$

$$= (ac + bd) - (ad + bc)$$

$$yx = (c - d)(a - b)$$

$$= (ca + db) - (cb + da)$$

$$= (ac + bd) - (ad + bc)$$

$$= xy.$$

Note that (xy)z = x(yz) was proved in in the section about proving Proposition 4.1.6 on page 4.1.

For the statement x1 = 1x = x, we only have to prove that 1x = x, because we have already proved that xy = yx earlier:

$$\begin{aligned} 1x &= (1 - 0) \cdot (a - b) \\ &= (1a + 0b) - (1b + 0a) \\ &= a - b \\ &= x. \end{aligned}$$

Now we will show that x(y+z) = xy + xz:

$$\begin{split} x(y+z) &= (a-b)((c-d) + (e-f)) \\ &= (a-b)((c+e) - (d+f)) \\ &= (a(c+e) + b(d+f)) - (a(d+f) + b(c+e)) \\ &= (ac+ae+bd+bf) - (ad+af+bc+be) \\ &= ((ac+bd) - (ad+bc)) + ((ae+bf) - (af+be)) \\ &= xy + xz. \end{split}$$

We can show that (y+z)x = yx+zx using xy = yx and x(y+z) = xy+xz by simply swapping x(y+z) with (y+z)x, xy with yx, and xz with zx.

## Exercise 4.1.7

(a) If a > b, we will show that a - b is a positive natural number. Since  $a \ge b$ , there exists  $d \in \mathbb{N}$  such that a = b + d. Subtracting b from both sides gives a - b = (b + d) - b. We know that

$$(b+d) - b = (b+d) + (-b) = (d+b) + (-b) = d + (b+(-b)) = d + 0 = d.$$

Therefore, a - b = d. Since d = 0 implies a = b + 0 = b, which contradicts a > b, a - b must be a nonzero, or positive, natural number.

Now we will show that if a - b is a positive natural number, a > b. We know that  $a \ge b$ , since a = a + (-b) + b = (a - b) + b. We also know that  $a \ne b$  because otherwise, a - b would be 0. Therefore, a > b.

(b) When a > b, by (a),  $a - b \neq 0$ . Since c - c = c + (-c) = 0 (by Proposition 4.1.6 and the definition of subtraction on page 79),  $a - b + c - c \neq 0 + c - c = 0$ .

By Proposition 4.1.6 and the fact that -(-x) = x (if x = a - b, -x = b - a, and -(-x) is just a - b again),

$$a - b + c - c = a + (-b) + c + (-c)$$
  
=  $(a + (-c)) + (-b) + (-(-c)) = (a - c) - (b - c).$ 

Therefore, a + c > b + c.

- (c) Since a > b, a b is a positive natural number. Since c is also a positive natural number, by Lemma 2.3.3, (a b)c is also a positive natural number. The distributive law for natural numbers (Proposition 2.3.4) tells us that (a b)c = ac bc, and therefore ac > bc.
- (d) Since a > b, a b is a positive natural number. But this means that -a < -b, since -b = (a b) + (-a).
- (e) Since a > b and b > c, a b and b c are both positive natural numbers. By Proposition 2.2.8, (a b) + (b c) is also a positive natural number. We know that

$$(a-b) + (b-c) = a + (-b) + b + (-c) = a + 0 + (-c) = a + (-c) = a.$$

Therefore, a > c.

(f) First, we will show that at least one of the statements a > b, a < b, and a = b is true. By Lemma 4.1.5 (trichotomy of integers), a - b is either zero, positive, or negative. If it is 0,

$$a = (a - b) + b = 0 + b = b.$$

If a-b is positive, by part (a), a>b. When a-b is negative, it is the negation of some positive natural number d. Since -(-d) is equal to d itself and -(a-b)=b-a, d=b-a. Therefore, b>a. We have now proved that at least one of the statements a>b, a< b, and a=b is true.

Now, we will show that no two of these statements can be true at the same time. It is obvious from the definition of ordering on the integers (Definition 4.1.10) that a=b cannot be true at the same time as a>b or a< b. Now the only thing left is to prove that a>b and a< b cannot both be true. We will prove this using proof by contradiction. The statement that a>b

is equivalent to a-b being positive, and a < b is equivalent to b-a being positive. But by Lemma 4.1.5 (the trichotomy of integers) and the fact that b-a=-(a-b) give us a contradiction, as b-a cannot both be positive and negative. We have now proved that exactly one of the statements a > b, a < b, or a = b is true.

#### Exercise 4.1.8

If P(n) is the property that n > -1, P(0) is true, because 0 = (-1) + 1 and  $0 \neq -1$ . Also, if P(n) is true, since n++>n, by Lemma 4.1.11(e), P(n++) is true. However, P(-1) is obviously false, since -1 = -1.

# 4.2 The rationals

#### Exercise 4.2.1

Proving that equality is reflexive (for rational numbers) We know that for any  $a, b \in \mathbf{Z}$  where  $b \neq 0$ , ab = ab. Therefore, a//b = a//b.

**Proving that equality is symmetric** If for  $a, b, c, d \in \mathbf{Z}$  that satisfy  $b, d \neq 0$ , we have a//b = c//d, then ad = cb. Hence, cb = ad, and c//d = a//b.

**Proving that equality is transitive** If  $a, b, c, d, e, f \in \mathbf{Z}$  where  $b, d, f \neq 0$ , a//b = c//d, and c//d = e//f, then we must have ad = cb and cf = ed. Hence,  $ad \cdot cf = cb \cdot ed$ . By Corollary 4.1.9, we can remove c and d from both sides:

$$a \cdot f = b \cdot e$$
.

Therefore, by Definition 4.2.1, a//b = e//f, and we have proved that rational number equality is reflexive, symmetric, and transitive.

#### Exercise 4.2.3

I will assume that x = a//b, y = c//d, and z = e//f, where  $a, b, c, d, e, f \in \mathbf{Z}$  and  $b, d, f \neq 0$ .

First, we will prove that x + y = y + x:

$$x + y = (a//b) + (c//d)$$

$$= (ad + bc)//(bd)$$

$$y + x = (c//d) + (a//b)$$

$$= (cb + da)//(db)$$

$$= (ad + bc)//(bd)$$

$$= x + y.$$

The identity (x + y) + z was proved on page 84 already.

Next, we will prove that x + 0 = 0 + x = x. We have already proved that x + y = y + x, so we only need to show that x + 0 = x:

$$x + 0 = (a//b) + (0//1)$$
  
=  $(a \cdot 1 + b \cdot 0) / / (b \cdot 1)$   
=  $a//b$   
=  $x$ .

For the next identity, x + (-x) = (-x) + x = 0, we have already shown that x + y = y + x, so we only need to prove x + (-x) + 0:

$$x + (-x) = (a//b) + ((-a)//b)$$

$$= (ab + b \cdot (-a))//(b \cdot b)$$

$$= 0//(b \cdot b)$$

$$= 0.$$

The last step is valid because  $0 \cdot 1 = 0 \cdot (b \cdot b)$ . Now we prove that xy = yx:

$$xy = (a//b) \cdot (c//d)$$

$$= (ac)//(bd)$$

$$yx = (c//d) \cdot (a//b)$$

$$= (ca)//(db)$$

$$= xy.$$

We now prove that (xy)z = x(yz):

$$(xy)z = ((a//b) \cdot (c//d)) \cdot (e//f)$$

$$= ((ac)//(bd)) \cdot (e//f)$$

$$= (ace)//(bdf)$$

$$x(yz) = (a//b) \cdot ((c//d) \cdot (e//f))$$

$$= (a//b) \cdot ((ce)//(df))$$

$$= (ace)//(bdf)$$

$$= (xy)z.$$

Next, we prove x1 = 1x = x. We have already proved that xy = yx, so we only need to prove  $x \cdot 1 = x$ :

$$x \cdot 1 = (a//b) \cdot (1//1)$$
$$= (a \cdot 1)//(b \cdot 1)$$
$$= a//b$$
$$= x.$$

The next identity to prove is x(y+z) = xy + xz:

$$x(y+z) = (a//b)((c//d) + (e//f))$$

$$= (a//b)((cf + de)//(df))$$

$$= (a(cf + de))//(bdf)$$

$$= (acf + ade)//(bdf)$$

$$xy + xz = ((ac)//(bd)) + ((ae)//(bf))$$

$$= (acbf + bdae)//(b^2df)$$

$$= (acf + dae)//(bdf)$$

$$= x(y+z).$$

For the identity (y+z)x = yx + zx, we can just apply xy = yx and x(y+z) = xy + xz together.

Finally, we will prove that  $xx^{-1} = x^{-1}x = 1$ , when  $x \neq 0$ . We know that  $x^{-1}$  exists because otherwise,  $a \neq 0$ , and therefore x would be 0. Since xy = yx, we only need to show that  $xx^{-1} = 1$ :

$$xx^{-1} = (a//b)(b//a)$$
$$= (ab)//(ab)$$
$$= 1.$$

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(The last step comes from Definition 4.2.1).

#### Exercise 4.2.5

- (a) By Lemma 4.2.7, we know that exactly one of the statements "x-y=0", "x-y is positive", or "x-y is negative" is true. We now can use Definition 4.2.8 to deduce that exactly one of the statements  $x=y, \ x>y,$  or x< y is true.
- (b) If x < y, x y must be a negative rational number. Hence, y x is positive, and y > x. Similarly, we can show that y > x implies x < y.
- (c) By (b), we know that y > x and z > y. Thus, y x and z y are positive. If y x = a//b and z y = c//d (where a, b, c, and d are positive; this is possible by Definition 4.2.6),

$$z - x = z - y + y - x$$
$$= (ad + bc)/(bd)$$

which is positive, and therefore z > x. We now have proved that x < z.

(d) By (b), we know that y > x, and hence, y - x is positive. If y - x = a//b, we know that

$$y + z - (x + z) = y + z - x - z$$
  
=  $y - x$ 

and therefore y + z > x + z. Hence, x + z < y + z.

(e) Because of part (b), we know that y > x, and thus, y - x is positive. Assuming that y - x = a//b and z = c//d (where a, b, c, and d are positive; this is possible by Definition 4.2.6),

$$yz - xz = (y - x)z$$
$$= (a//b)(c//d)$$
$$= (ac)//(bd)$$

and hence, yz > xz and xz < yz.

# Exercise 4.2.6

We know that x - y is positive. Assuming that y - x = a//b and z = c//d (where a, b, c, and d are positive; this is always possible by Definition 4.2.6), we have that

$$yz - xz = (y - x)z$$
$$= (ac)//(bd)$$

and therefore, xz > yz.