

Queue Data Structures

Data Structures for Computer Professionals

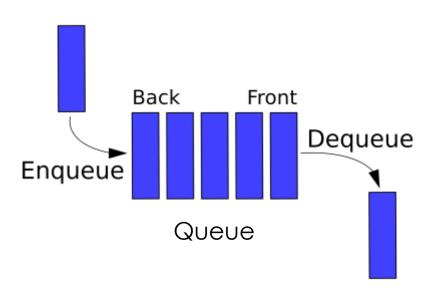
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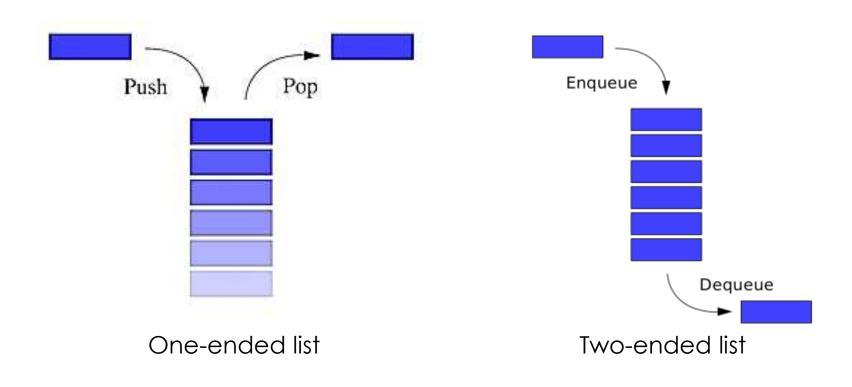
Computer Engineering, Chiang Mai University

Queue Data Structure





Queue vs Stack



Which is which?

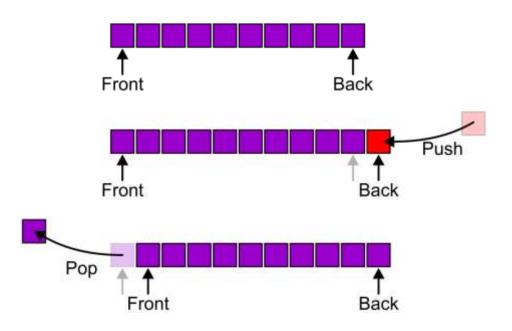
Queue Data Structure

- □ Linear ordering like Stack
- Queue is two-ended list but Stack is one-ended list
- Insert at one end of the List, remove at the other end
- Queues are "FIFO" First In, First Out
- Primary operations are Enqueue and Dequeue
- A queues ensures "fairness"
 - customers waiting on a customer hotline
 - processes waiting to run on the CPU

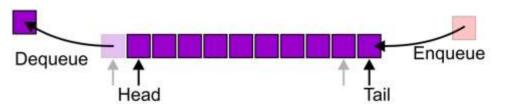
Queue API (Operations)

- Enqueue(Key): Add Key at the **end** of the queue (also called "back", "rear", or "tail")
- Key Dequeue(): Remove and return a Key from the **front** of the queue (earliest-added key)
- □ bool *IsEmpty()*: Is there any element?

FIFO



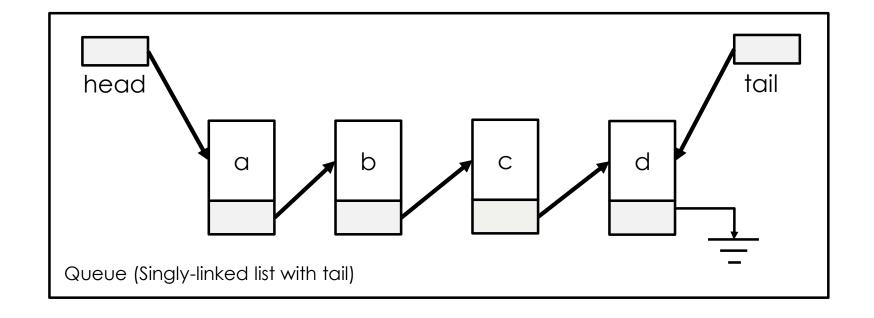
Alternative terms



A sample of CPE Applications of Queue

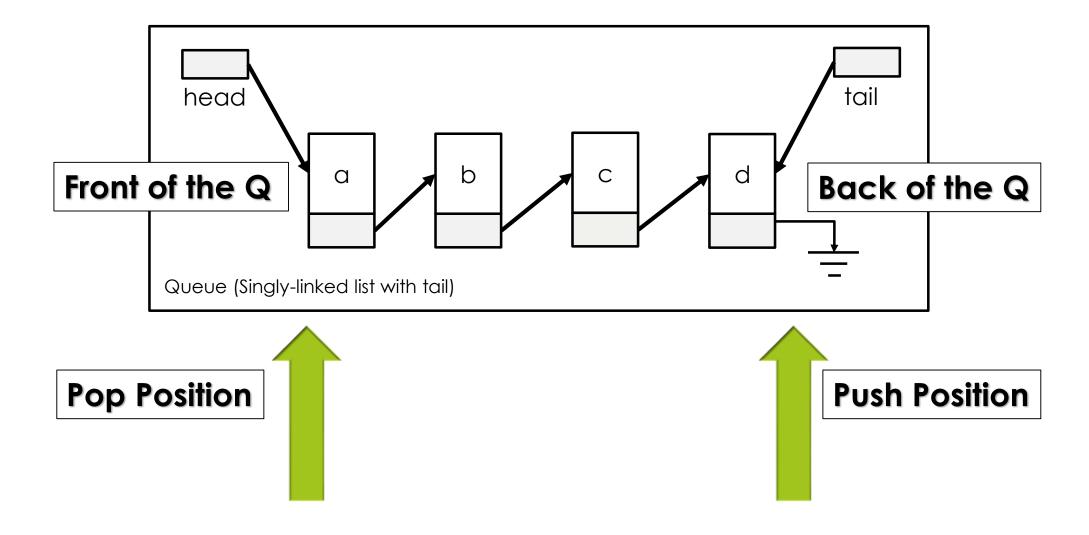
- File servers: Multiples users need to access their files on a shared file server machine are given access on a FIFO basis
- Printer Queue: Printing jobs submitted to a printer are printed in order of arrival
- Phone calls made to customer service hotlines are usually placed in a queue
- Note that Queue works best with jobs with similar priority. If the jobs are with different priority, you should use Priority Queue Data Structure instead.

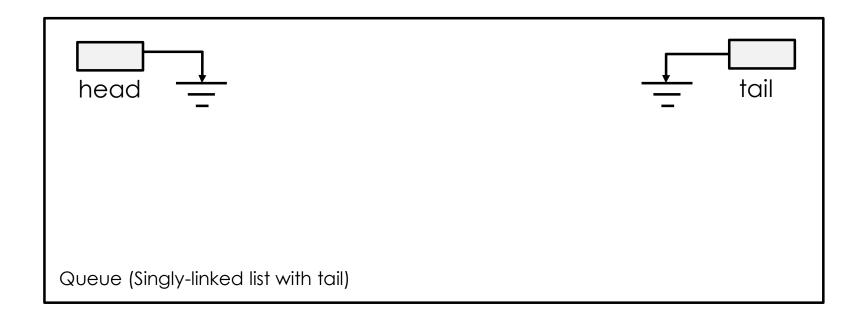
Queue Implementation with Linked List (Singly-linked list with tail)



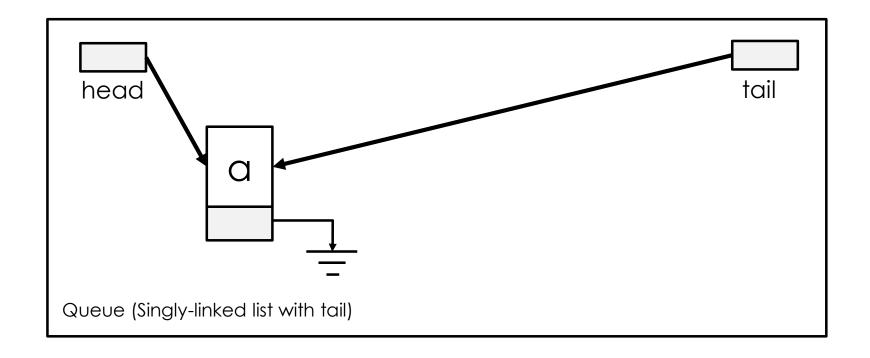
Which side should be the Push position for Enqueue? Which side should be the Pop position for Dequeue?

Queue Implementation with Linked List (Singly-linked list with tail)

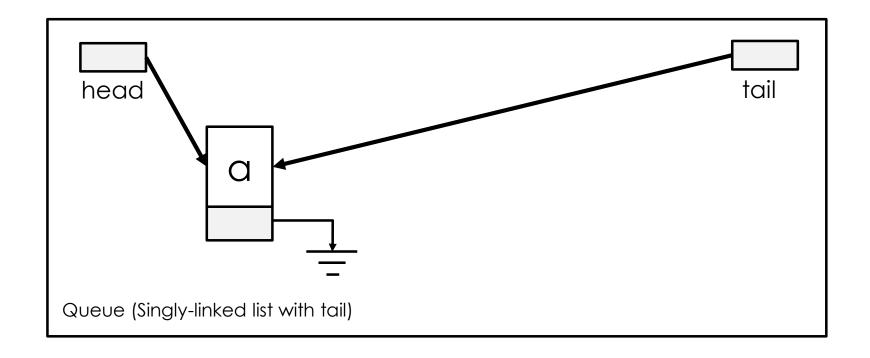




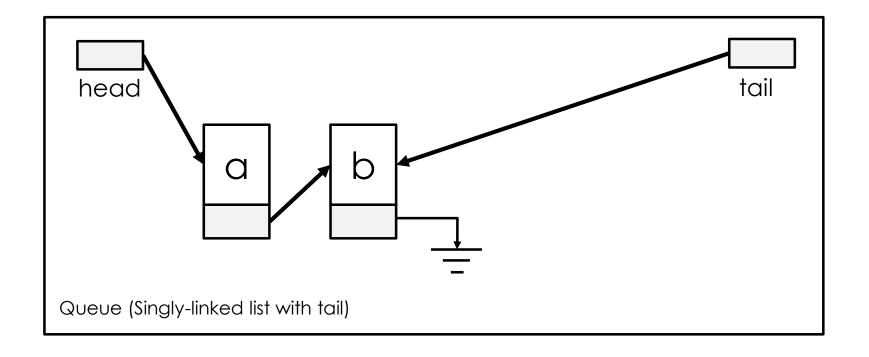
Enqueue(a)



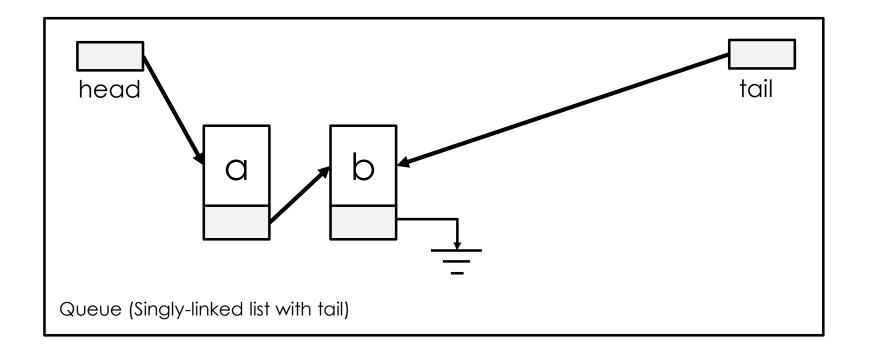
Enqueue(a)



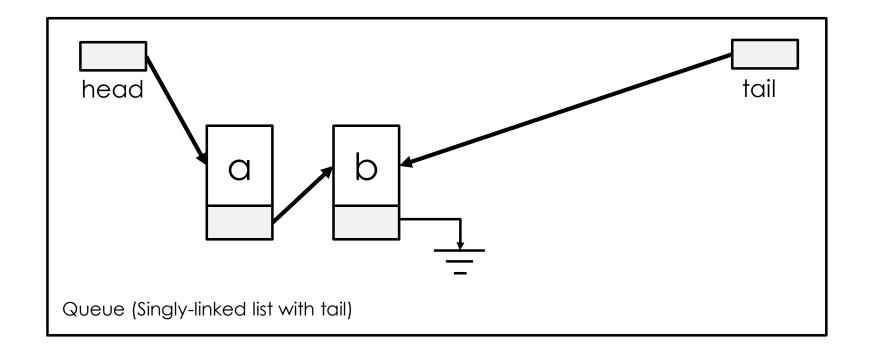
Enqueue(b)



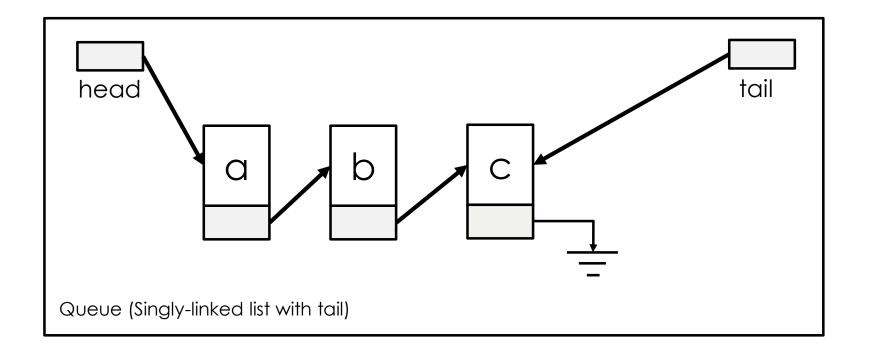
Enqueue(b)



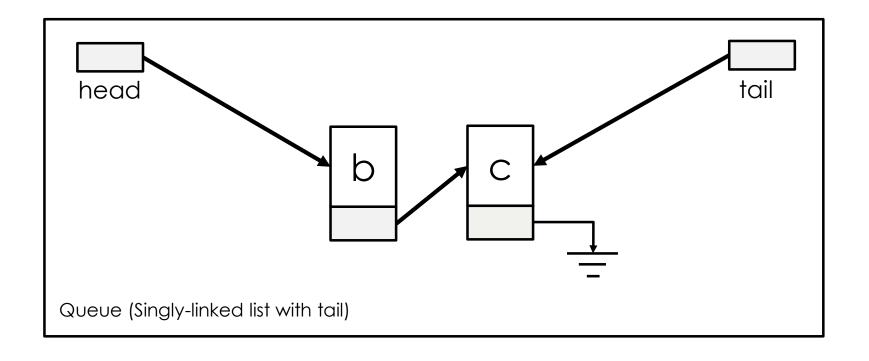
isEmpty()
$$\rightarrow$$
 false

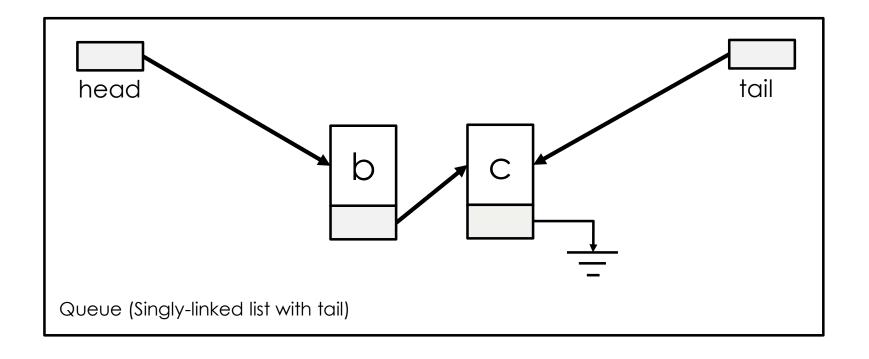


Enqueue(c)

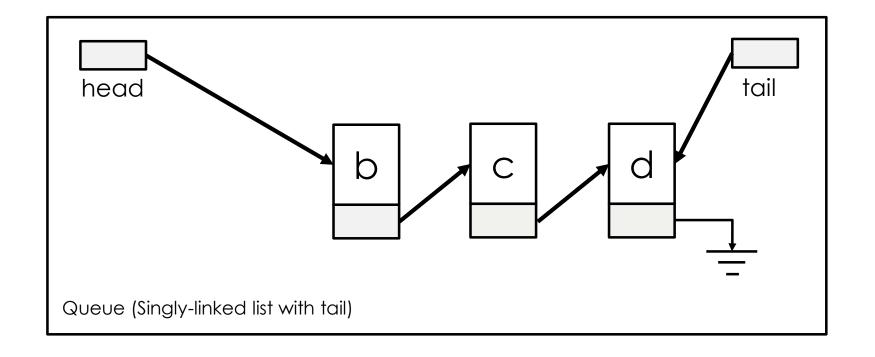


Enqueue(c)

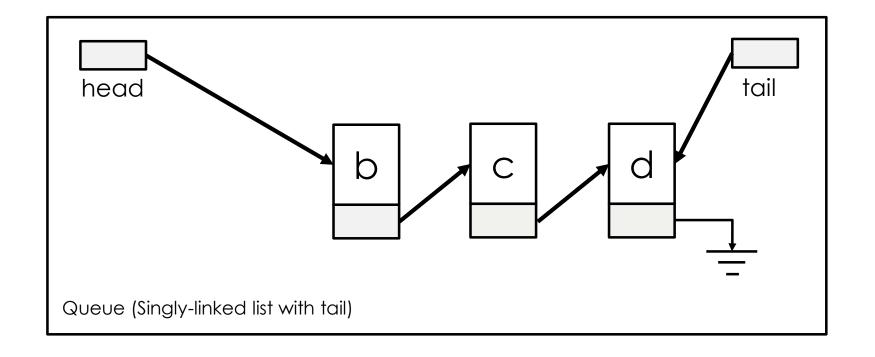




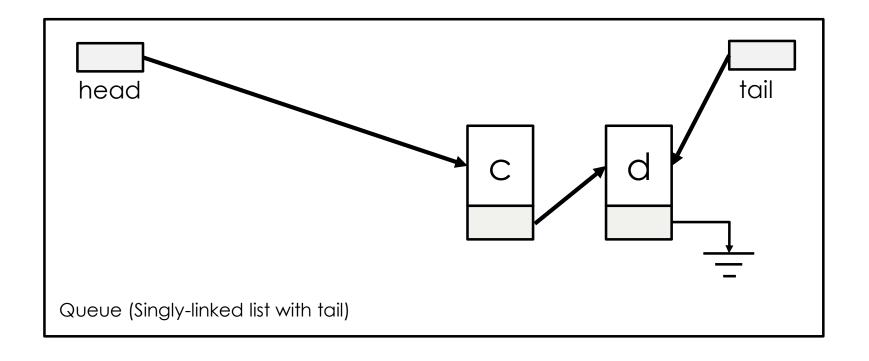
Enqueue(d)



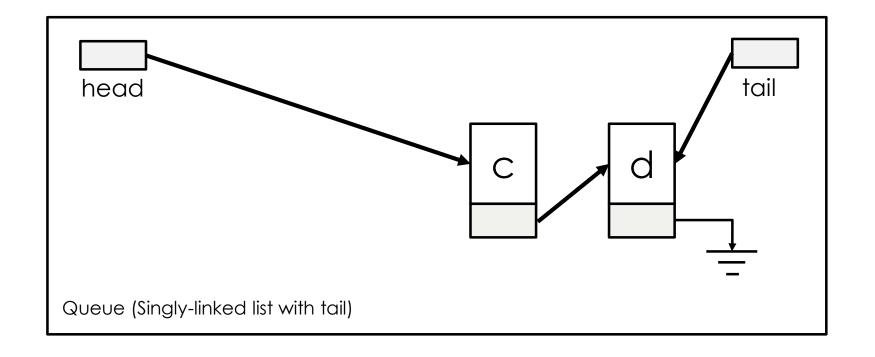
Enqueue(d)



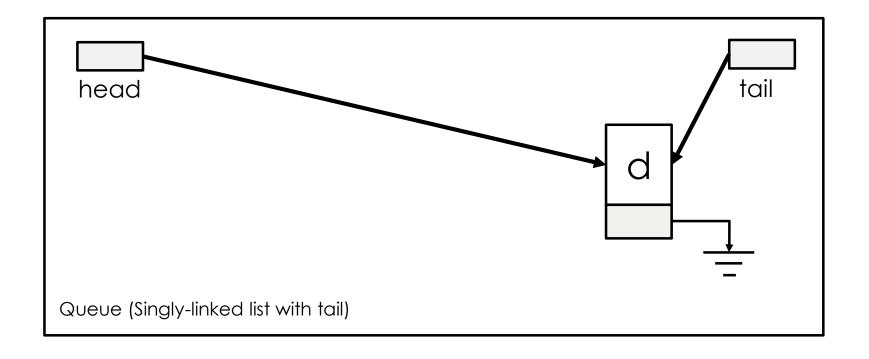
Dequeue()



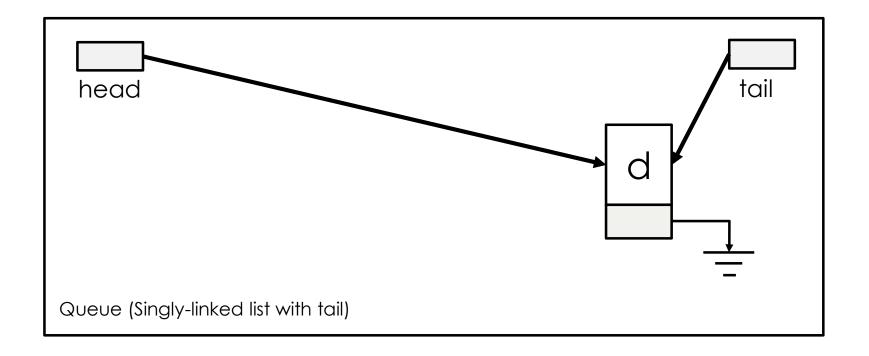
Dequeue()
$$\rightarrow$$
 b



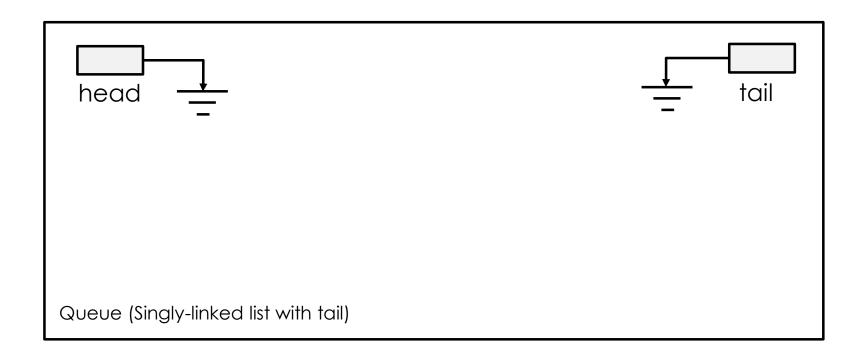
Dequeue()



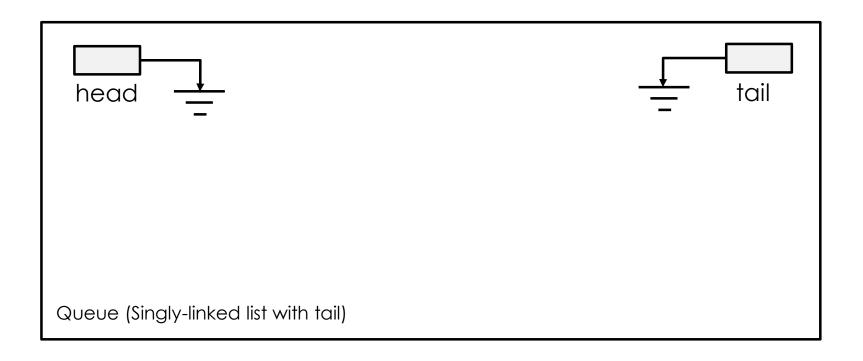
Dequeue()
$$\rightarrow$$
 c



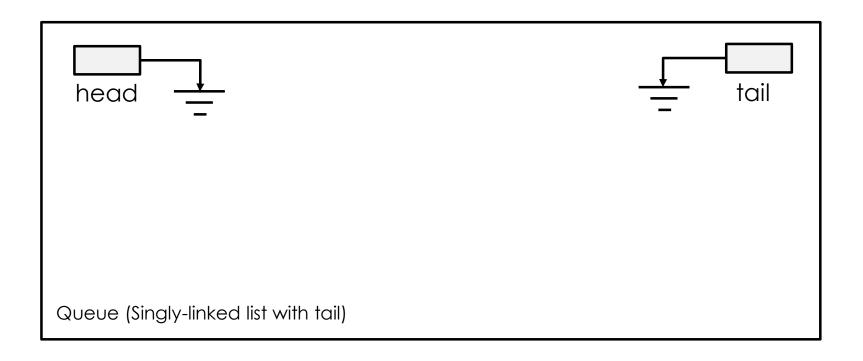
Dequeue()



Dequeue()
$$\rightarrow$$
 d

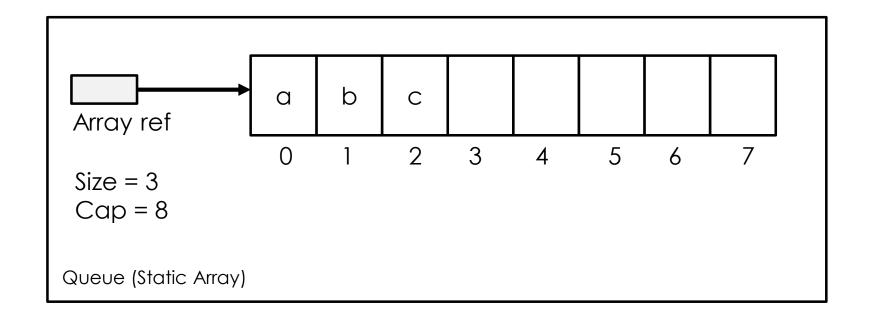


isEmpty()



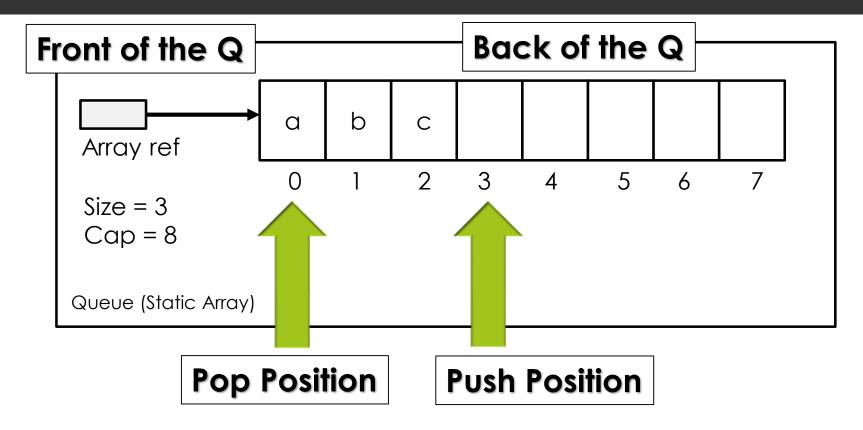
■ What is the equivalent operation(s) in the List?

| List Operations | Complexity |
|--------------------------------|---------------------------------------------|
| PushBack(Key) | O(1) |
| Key TopFront(); PopFront(); | O(1) |
| isEmpty() | O(1) |
| | PushBack(Key) Key TopFront(); PopFront(); |



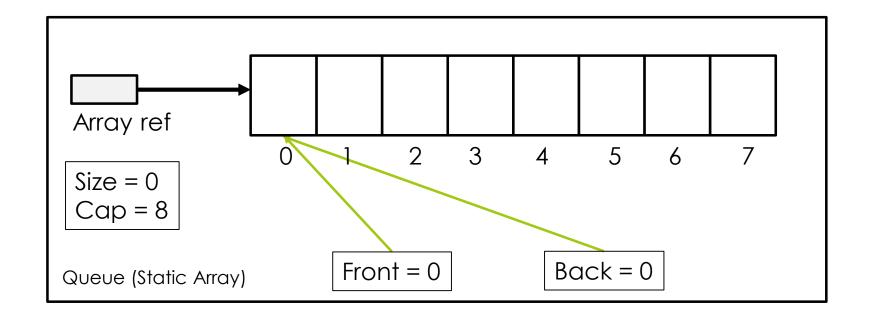
Which side should be the Push position for Enqueue? Which side should be the Pop position for Dequeue?

To prevent O(n), we should not shift anything after popping



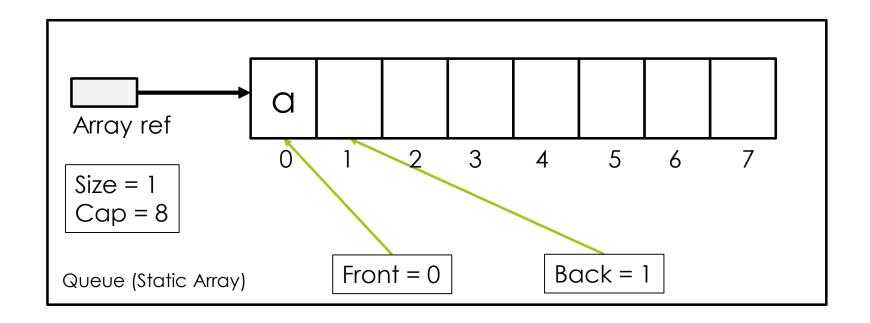
To implement Queue using Array, we need two more variables (integer type) →

- Front (indicating pop position)
- Back (indicating push position)



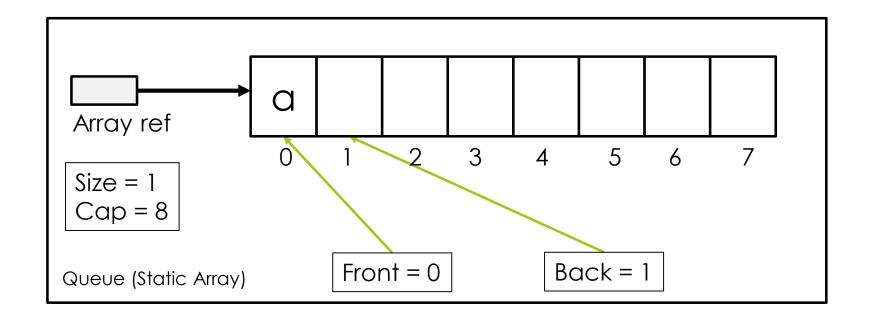
Enqueue(a)

What are the corresponding array operations?



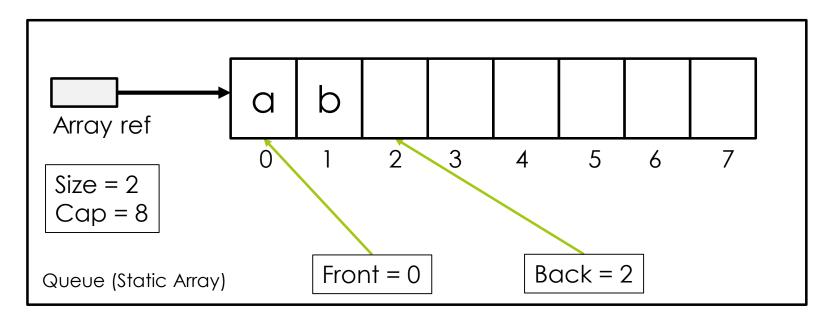
Enqueue(a)

Front can be either one block after the last item or the last item itself! In this lecture, the front variable will be one block after the last item.

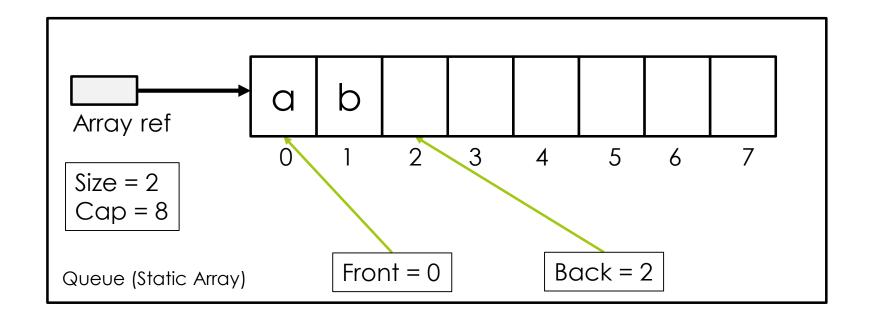


Enqueue(b)

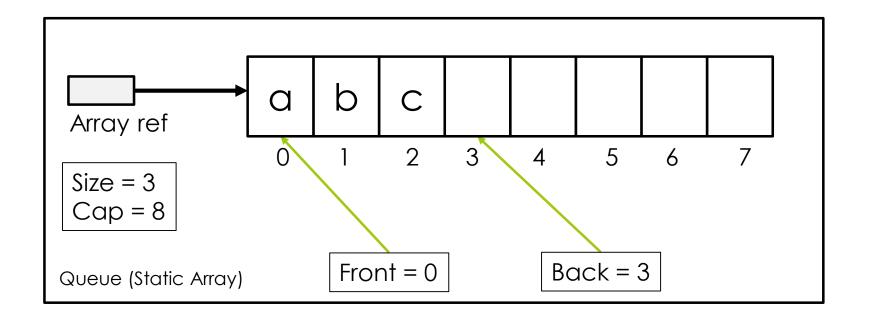
What are the corresponding Array operations?



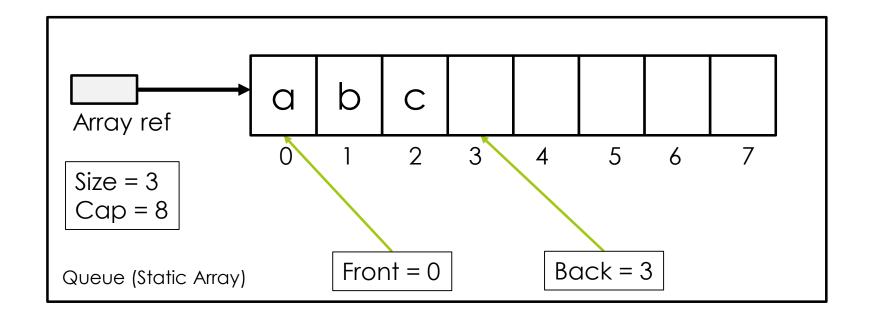
Enqueue(b)



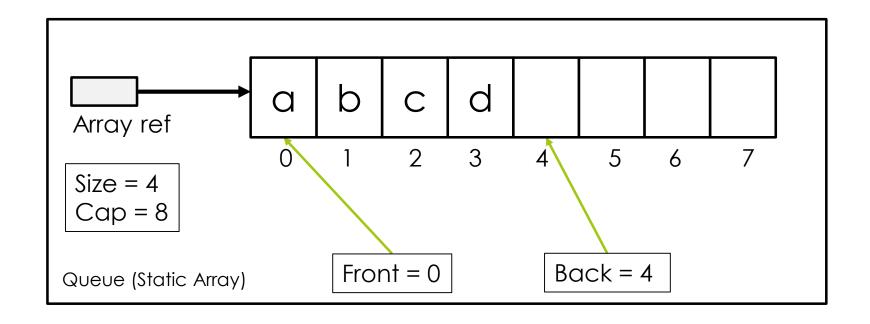
Enqueue(c)



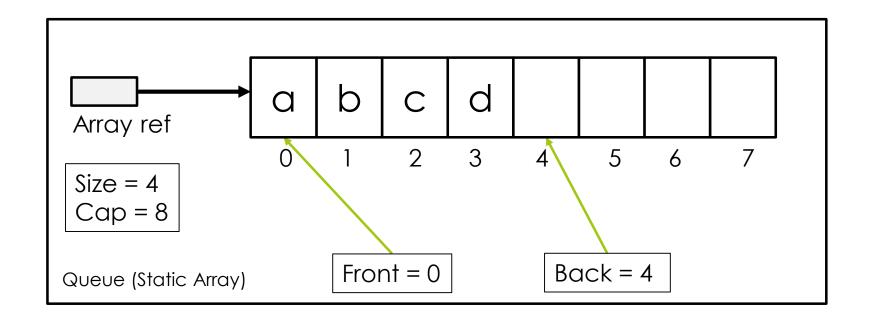
Enqueue(c)



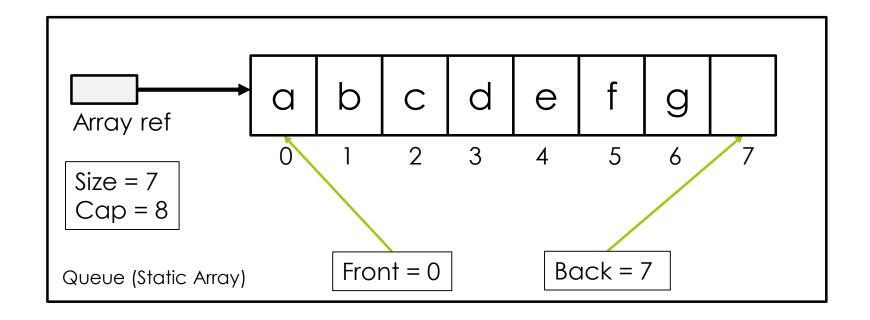
Enqueue(d)



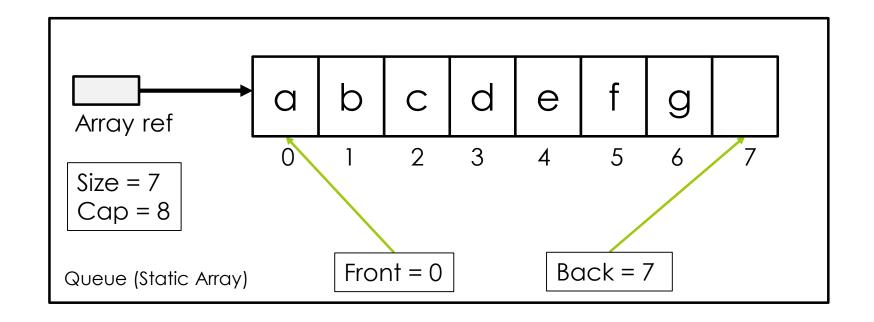
Enqueue(d)



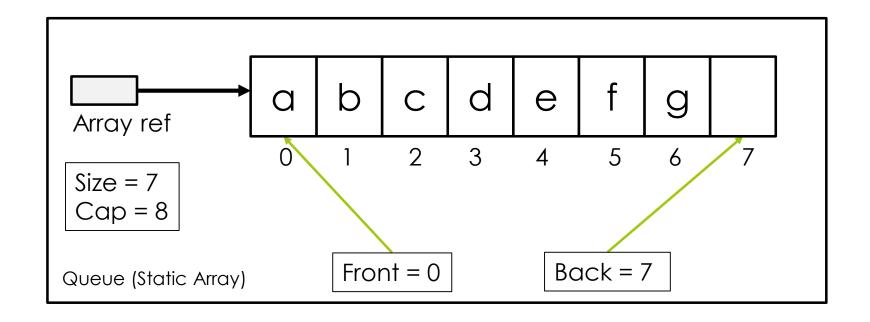
Enqueue(e); Enqueue(f); Enqueue(g);



Enqueue(e); Enqueue(f); Enqueue(g);

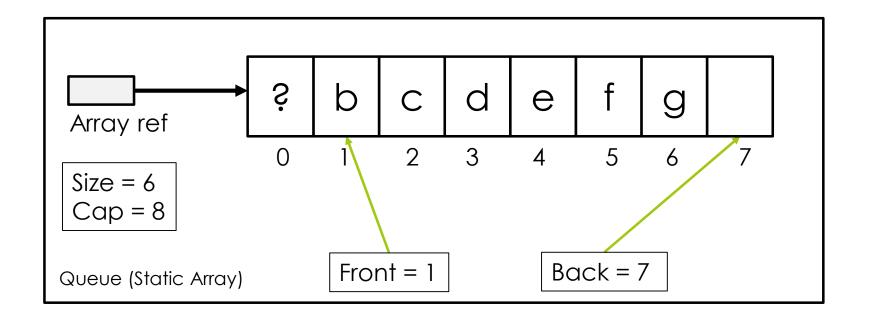


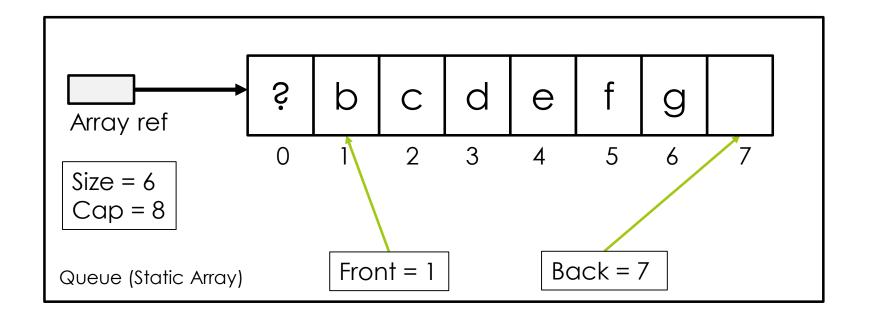
There is one more space, can we do Enqueue(h)?



Dequeue()

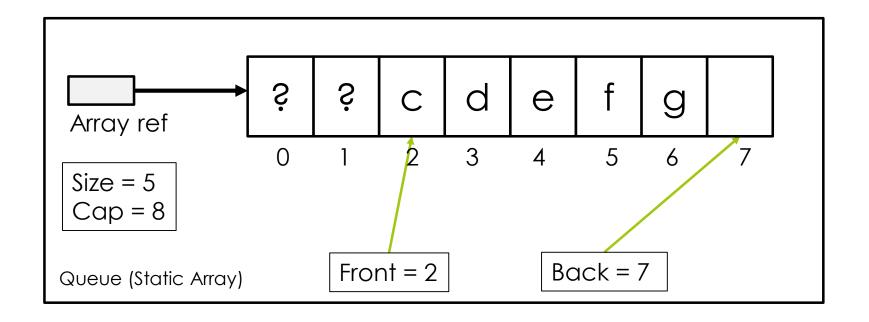
What are the corresponding Array operations?

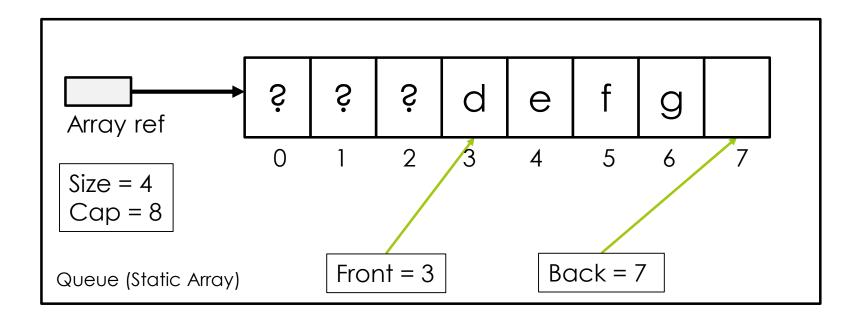


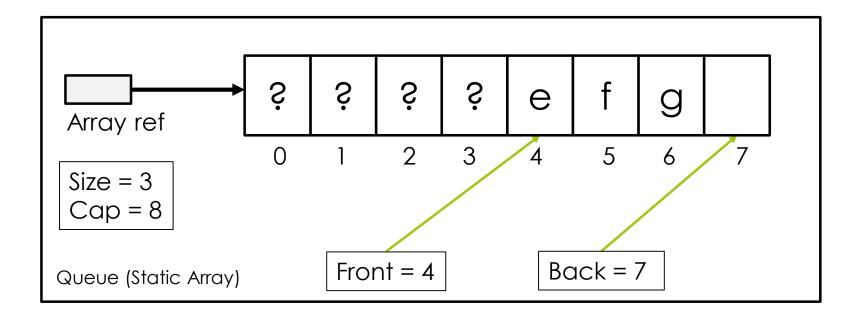


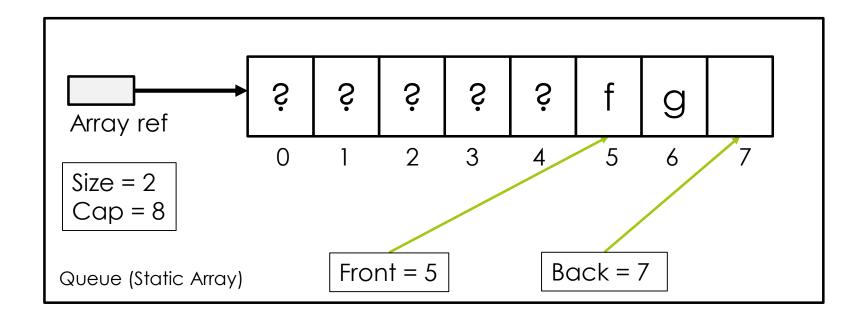
Dequeue()

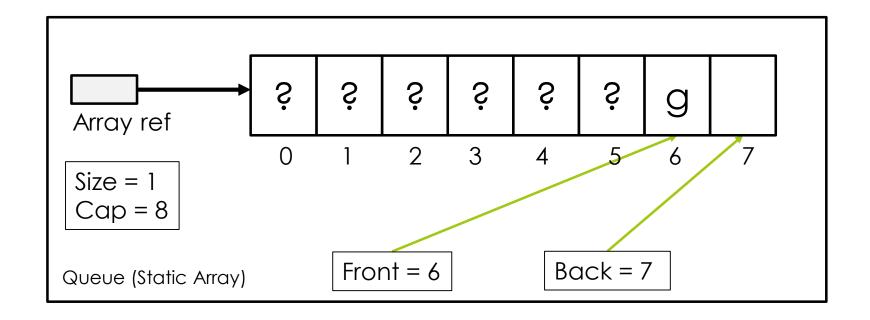
What are the corresponding Array operations?



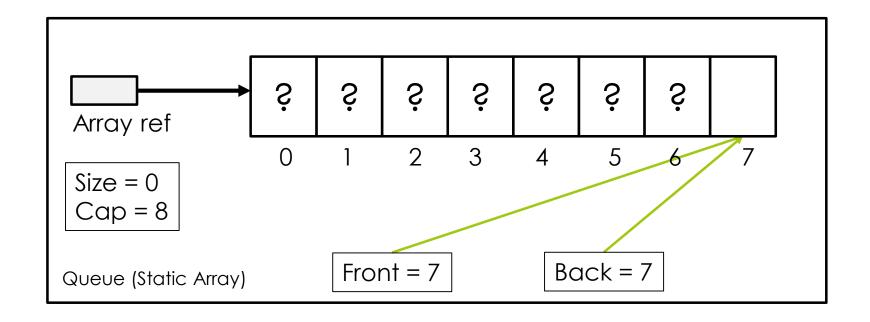


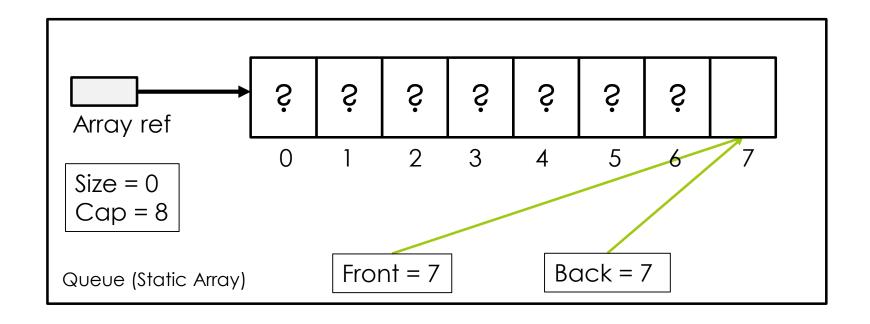




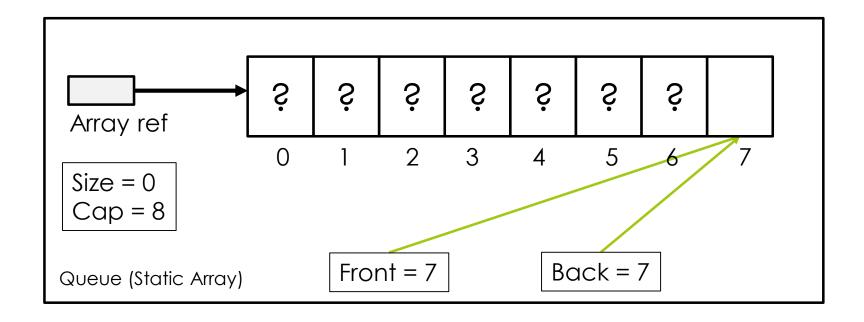


Dequeue()
$$\rightarrow$$
 f

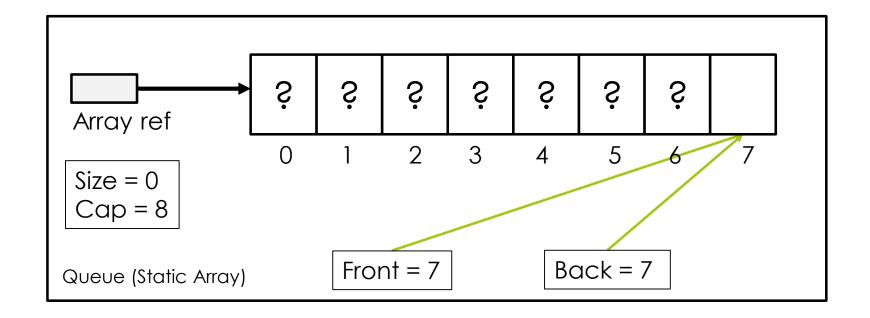




isEmpty()
$$\rightarrow$$
 true



Dequeue() → ERROR

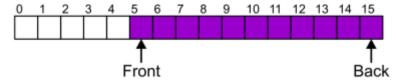


Now the pointers point to the last index, can we still do Enqueue(key)?

Full Queue?

Suppose that:

- The array capacity is 16
- We have performed 16 pushes
- We have performed 5 pops
 - The queue size is now 11



We perform one further push

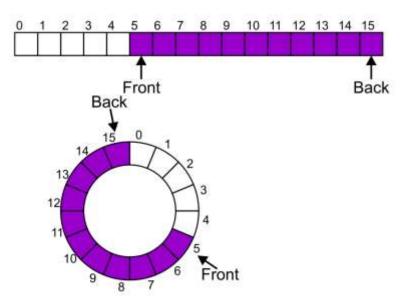
In this case, the array is not full and yet we cannot place any more objects in to the array

What should we do or just give up?

Circular Array

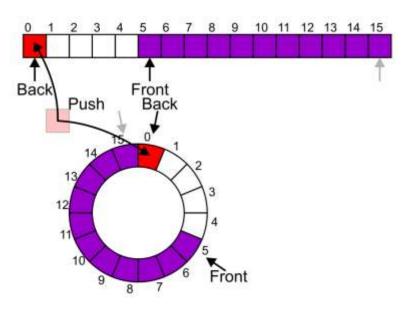
Instead of viewing the array on the range 0, ..., 15, consider the indices being cyclic:

This is referred to as a circular array



Queue as a Circular Array

How to shorten these codes to a single line without a conditional statement (if or :)?



Exceptions

As with a stack, there are a number of options which can be used if the array is filled

If the array is filled, we have five options:

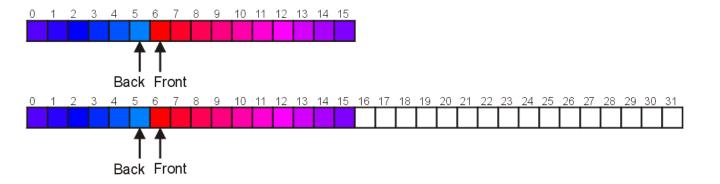
- Increase the size of the array
- Throw an exception
- Ignore the element being pushed
- Put the pushing process to "sleep" until something else pops the front of the queue

Include a member function **bool full()**

Increasing Capacity

Unfortunately, if we choose to increase the capacity, this becomes slightly more complex

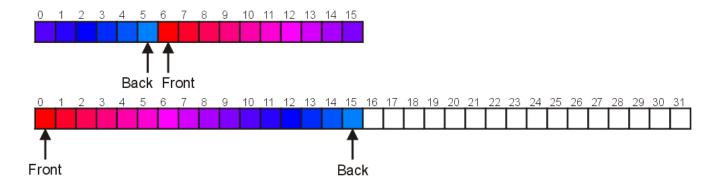
– A direct copy does not work:



Increasing Capacity

An alternate solution is normalization:

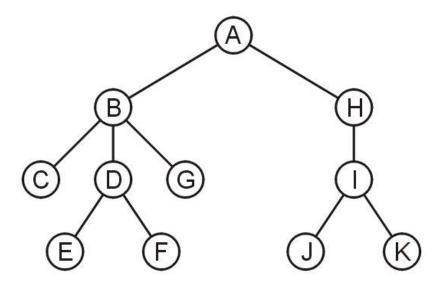
- Map the front back at position 0
- The next push would then occur in position 16



Application

Another application is performing a breadth-first traversal of a directory tree

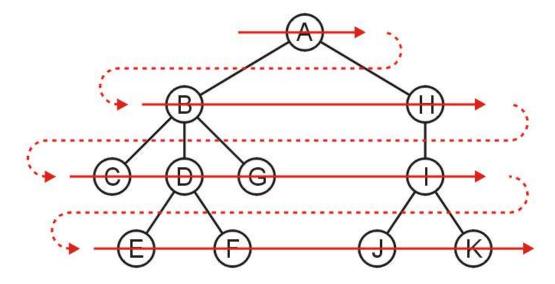
Consider searching the directory structure



We would rather search the more shallow directories first then plunge deep into searching one sub-directory and all of its contents

One such search is called a breadth-first traversal

Search all the directories at one level before descending a level



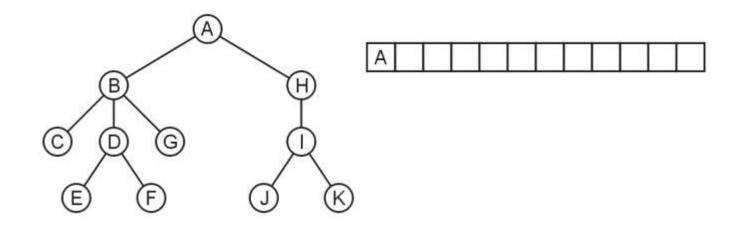
Output: **A**, **B**, **H**, **C**, **D**, **G**, **I**, **E**, **F**, **J**, **K**

The easiest implementation is:

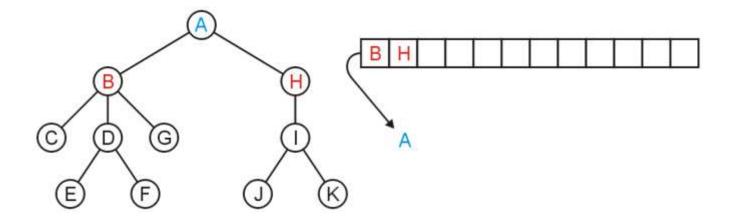
- Place the root directory into a queue
- While the queue is not empty:
 - Pop the directory at the front of the queue
 - Push all of its sub-directories into the queue

The order in which the directories come out of the queue will be in breadth-first order

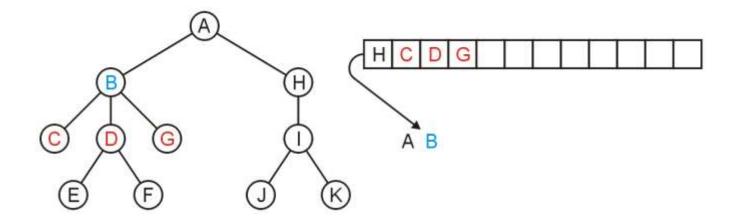
Push the root directory A



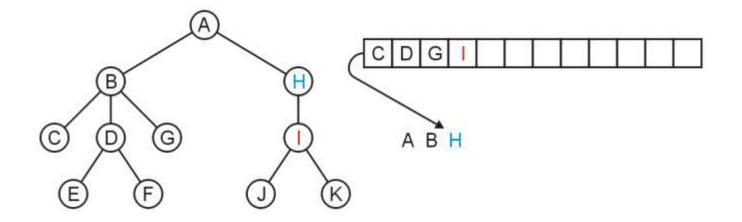
Pop A and push its two sub-directories: B and H



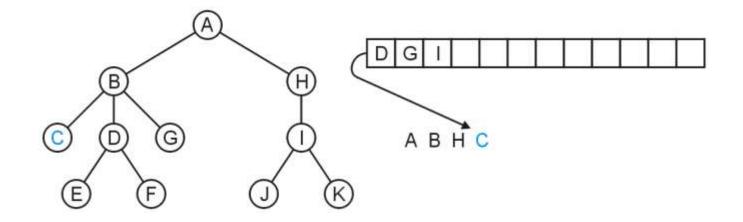
Pop B and push C, D, and G



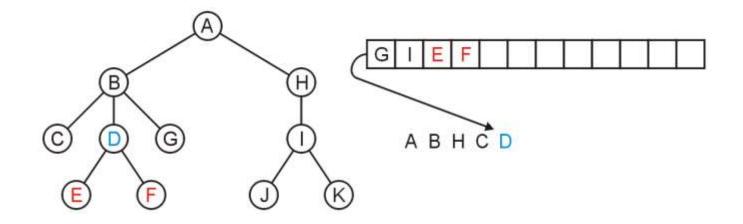
Pop H and push its one sub-directory I



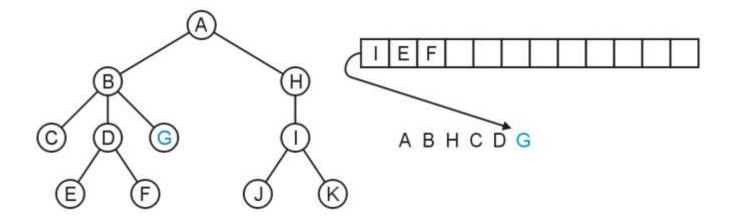
Pop C: no sub-directories



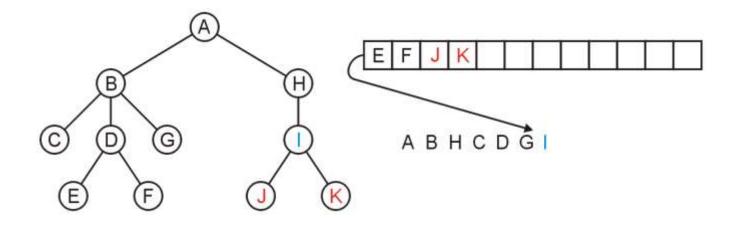
Pop D and push E and F



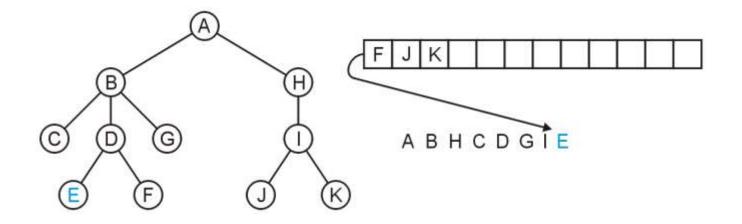
Pop G



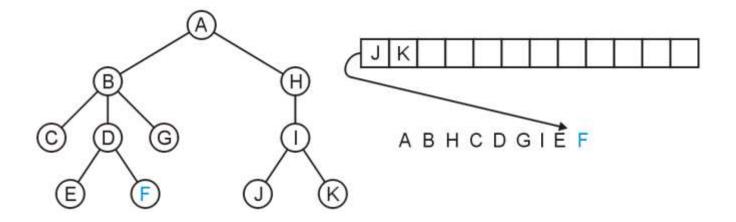
Pop I and push J and K



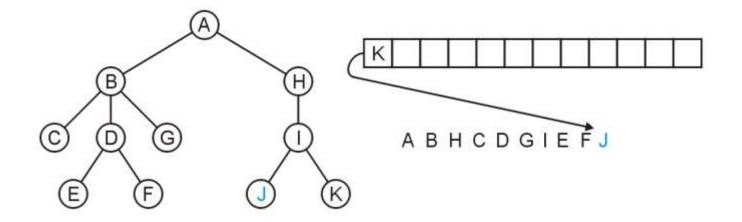
Pop E



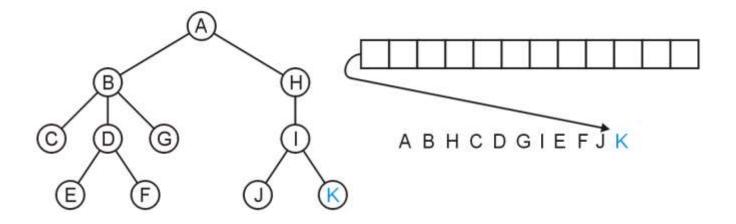
Pop F



Pop J



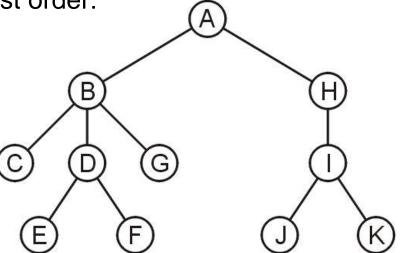
Pop K and the queue is empty



The resulting order

ABHCDGIEFJK

is in breadth-first order:



Summary

The queue is one of the most common abstract data structures

Understanding how a queue works is trivial

The implementation is only slightly more difficult than that of a stack

Applications include:

- Queuing clients in a client-server model
- Breadth-first traversals of trees