

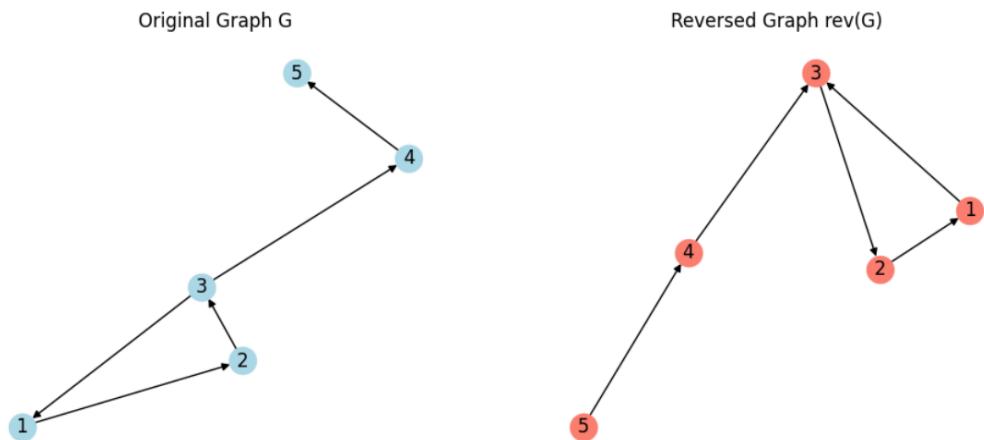
Problem1:

SCC and reversal, let G a directed graph:

1. Describe an algorithm to compute the reversal $\text{rev}(G)$ of a directed graph G in $O(V + E)$ time.

```
def reverse_graph(G):
    R = nx.DiGraph()
    R.add_nodes_from(G.nodes())
    for u, v in G.edges():
        R.add_edge(v, u)
    return R
```

1. Creates a new directed graph R with all vertices from G
2. For each edge (u, v) in G, adds edge (v, u) to R
3. Time complexity: $O(V)$ for vertices + $O(E)$ for edges = $O(V + E)$



2. Prove that for every directed graph G, the strong component graph $\text{scc}(G)$ is acyclic.

```

#build SCC graph
def scc_graph(G):
    scc = list(nx.strongly_connected_components(G))
    index = {node: i for i, comp in enumerate(scc) for node in comp}
    SG = nx.DiGraph()
    SG.add_nodes_from(range(len(scc)))

    for u, v in G.edges():
        if index[u] != index[v]:
            SG.add_edge(index[u], index[v])
    return SG

SG = scc_graph(G)

```

1. Find all strongly connected components (SCCs) using Kosaraju/Tarjan
2. Create SCC graph where each SCC is a vertex
3. Add edge between SCCs if there's any edge between their vertices
4. Acyclicity proof: If SCC graph had a cycle, all SCCs on that cycle would be mutually reachable, contradicting maximality of SCCs

3. Prove that $\text{scc}(\text{rev}(G)) = \text{rev}(\text{scc}(G))$ for every directed graph G .

```

R = reverse_graph(G)
SG = scc_graph(G)

scc_of_revG = scc_graph(R)
rev_of_sccG = reverse_graph(SG)

```

1. $\text{rev}(G)$ reverses all edges in G
2. SCCs remain the same in $\text{rev}(G)$ (strong connectivity is preserved under edge reversal)
3. Edges between SCCs in $\text{scc}(\text{rev}(G))$ are exactly the reverses of edges in $\text{scc}(G)$
4. Therefore: $\text{scc}(\text{rev}(G)) = \text{rev}(\text{scc}(G))$

4. Fix an arbitrary directed graph G . For any vertex v of G , let $S(v)$ denote the strong component of G that contains v . For all vertices u and v of G , prove that u can reach v in G if and only if $S(u)$ can reach $S(v)$ in $\text{scc}(G)$.

```

# S(v) is represented by index[v]
index = {node: i for i, comp in enumerate(scc) for node in comp}

# u can reach v in G if:
# Case 1: index[u] == index[v] (same SCC)
# Case 2: index[u] != index[v] and SCC index[u] can reach SCC index[v] in SG

```

1. Let $S(v)$ be the SCC containing vertex v
2. Forward direction: If $u \rightarrow v$ path exists in G , traverse SCCs along this path to get $S(u) \rightarrow S(v)$ path in $SCC(G)$
3. Backward direction: If $S(u) \rightarrow S(v)$ path exists in $SCC(G)$, use edges between SCCs to construct $u \rightarrow v$ path in G
4. Same SCC case: trivial reachability via SCC's internal cycles

Problem2:

```

def euler_tour(G):
    if not nx.is_strongly_connected(G):
        raise ValueError("Graph is not strongly connected")

    for v in G.nodes():
        if G.in_degree(v) != G.out_degree(v):
            raise ValueError("In-degree != out-degree → no Euler t"

    G_copy = G.copy()
    start = list(G.nodes())[0]
    stack = [start]
    circuit = []

    while stack:
        v = stack[-1]
        if G_copy.out_degree(v) > 0:
            u = next(iter(G_copy[v]))
            stack.append(u)
            G_copy.remove_edge(v, u)
        else:
            circuit.append(stack.pop())
    return circuit[::-1]

#Euler graph
E = nx.DiGraph()
E.add_edges_from([(1,2),(2,3),(3,1)])

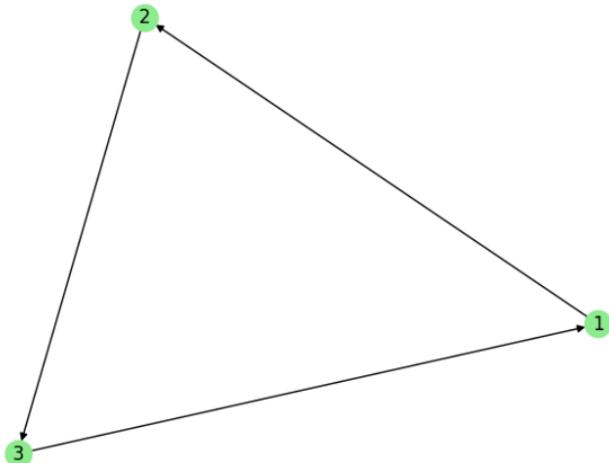
tour = euler_tour(E)
print("Euler Tour:", tour)

nx.draw(E, with_labels=True, arrows=True, node_color="lightgreen")
plt.title("Euler Graph Example")
plt.show()

```

Euler Tour: [1, 2, 3, 1]

Euler Graph Example



Problem3:

We are given:

$$A \rightarrow B$$

$$A \rightarrow C$$

$$B \rightarrow C$$

$$B \rightarrow D$$

$$C \rightarrow E$$

$$D \rightarrow E$$

$$D \rightarrow F$$

$$G \rightarrow F$$

$$G \rightarrow E$$

Topological Sort**: A linear ordering of vertices such that all edges $u \rightarrow v$ follow u before v . We compute:

- Start from A
- Use DFS-based or Kahn's Algorithm
- Multiple valid orderings exist

```

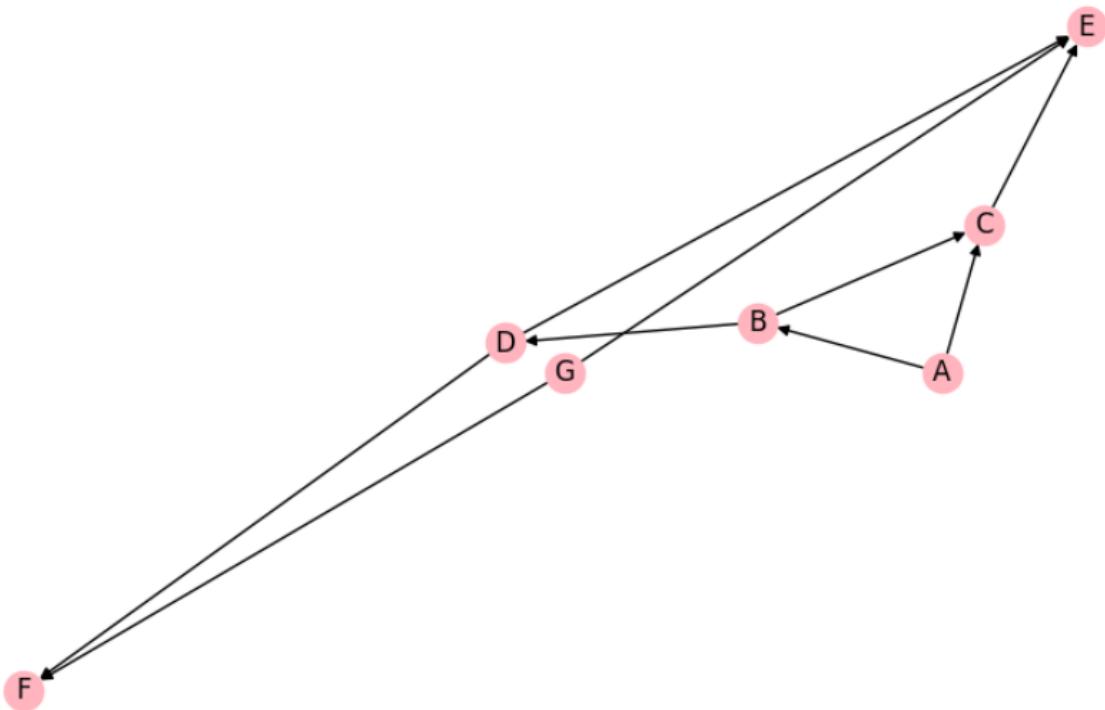
edges = [
    ("A","B"), ("A","C"),
    ("B","C"), ("B","D"),
    ("C","E"),
    ("D","E"), ("D","F"),
    ("G","F"), ("G","E")
]

G3 = nx.DiGraph()
G3.add_edges_from(edges)

plt.figure(figsize=(8,5))
nx.draw(G3, with_labels=True, node_color="lightpink", arrows=True)
plt.title("Course Prerequisite Graph")
plt.show()

# Topological sort
order = list(nx.topological_sort(G3))
print("One possible topological order:", order)

```



```
One possible topological order: ['A', 'G', 'B', 'C', 'D', 'E', 'F']
```