

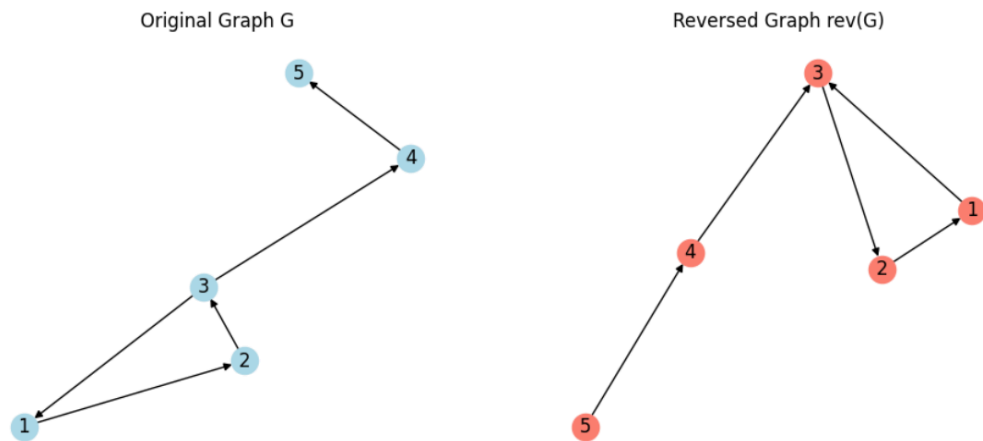
### Problem1:

SCC and reversal, let  $G$  a directed graph:

1. Describe an algorithm to compute the reversal  $\text{rev}(G)$  of a directed graph  $G$  in  $O(V + E)$  time.

```
def reverse_graph(G):  
    R = nx.DiGraph()  
    R.add_nodes_from(G.nodes())  
    for u, v in G.edges():  
        R.add_edge(v, u)  
    return R
```

1. Creates a new directed graph  $R$  with all vertices from  $G$
2. For each edge  $(u, v)$  in  $G$ , adds edge  $(v, u)$  to  $R$
3. Time complexity:  $O(V)$  for vertices +  $O(E)$  for edges =  $O(V + E)$



2. Prove that for every directed graph  $G$ , the strong component graph  $\text{scc}(G)$  is acyclic.

```
#build SCC graph
def scc_graph(G):
    scc = list(nx.strongly_connected_components(G))
    index = {node: i for i, comp in enumerate(scc) for node in comp}
    SG = nx.DiGraph()
    SG.add_nodes_from(range(len(scc)))

    for u, v in G.edges():
        if index[u] != index[v]:
            SG.add_edge(index[u], index[v])
    return SG

SG = scc_graph(G)
```

1. Find all strongly connected components (SCCs) using Kosaraju/Tarjan
  2. Create SCC graph where each SCC is a vertex
  3. Add edge between SCCs if there's any edge between their vertices
  4. Acyclicity proof: If SCC graph had a cycle, all SCCs on that cycle would be mutually reachable, contradicting maximality of SCCs
3. Prove that  $\text{scc}(\text{rev}(G)) = \text{rev}(\text{scc}(G))$  for every directed graph  $G$ .

```
R = reverse_graph(G)
SG = scc_graph(G)

scc_of_revG = scc_graph(R)
rev_of_sccG = reverse_graph(SG)
```

1.  $\text{rev}(G)$  reverses all edges in  $G$
  2. SCCs remain the same in  $\text{rev}(G)$  (strong connectivity is preserved under edge reversal)
  3. Edges between SCCs in  $\text{scc}(\text{rev}(G))$  are exactly the reverses of edges in  $\text{scc}(G)$
  4. Therefore:  $\text{scc}(\text{rev}(G)) = \text{rev}(\text{scc}(G))$
4. Fix an arbitrary directed graph  $G$ . For any vertex  $v$  of  $G$ , let  $S(v)$  denote the strong component of  $G$  that contains  $v$ . For all vertices  $u$  and  $v$  of  $G$ , prove that  $u$  can reach  $v$  in  $G$  if and only if  $S(u)$  can reach  $S(v)$  in  $\text{scc}(G)$ .

```

# S(v) is represented by index[v]
index = {node: i for i, comp in enumerate(scc) for node in comp}

# u can reach v in G if:
# Case 1: index[u] == index[v] (same SCC)
# Case 2: index[u] != index[v] and SCC index[u] can reach SCC index[v] in SG

```

1. Let  $S(v)$  be the SCC containing vertex  $v$
2. Forward direction: If  $u \rightarrow v$  path exists in  $G$ , traverse SCCs along this path to get  $S(u) \rightarrow S(v)$  path in  $\text{SCC}(G)$
3. Backward direction: If  $S(u) \rightarrow S(v)$  path exists in  $\text{SCC}(G)$ , use edges between SCCs to construct  $u \rightarrow v$  path in  $G$
4. Same SCC case: trivial reachability via SCC's internal cycles

## Problem2:

```

def euler_tour(G):
    if not nx.is_strongly_connected(G):
        raise ValueError("Graph is not strongly connected")

    for v in G.nodes():
        if G.in_degree(v) != G.out_degree(v):
            raise ValueError("In-degree != out-degree → no Euler t

    G_copy = G.copy()
    start = list(G.nodes())[0]
    stack = [start]
    circuit = []

    while stack:
        v = stack[-1]
        if G_copy.out_degree(v) > 0:
            u = next(iter(G_copy[v]))
            stack.append(u)
            G_copy.remove_edge(v, u)
        else:
            circuit.append(stack.pop())
    return circuit[::-1]

#Euler graph
E = nx.DiGraph()
E.add_edges_from([(1,2),(2,3),(3,1)])

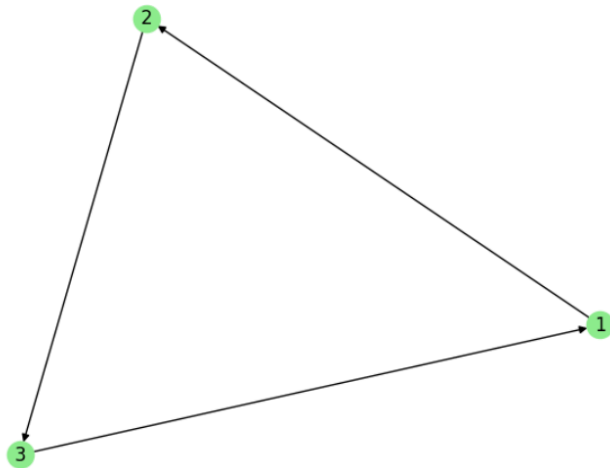
tour = euler_tour(E)
print("Euler Tour:", tour)

nx.draw(E, with_labels=True, arrows=True, node_color="lightgreen")
plt.title("Euler Graph Example")
plt.show()

```

Euler Tour: [1, 2, 3, 1]

Euler Graph Example



### Problem3:

We are given:

$A \rightarrow B$

$A \rightarrow C$

$B \rightarrow C$

$B \rightarrow D$

$C \rightarrow E$

$D \rightarrow E$

$D \rightarrow F$

$G \rightarrow F$

$G \rightarrow E$

Topological Sort\*\*: A linear ordering of vertices such that all edges  $u \rightarrow v$  follow  $u$  before  $v$ . We compute:

- Start from A
- Use DFS-based or Kahn's Algorithm
- Multiple valid orderings exist

```

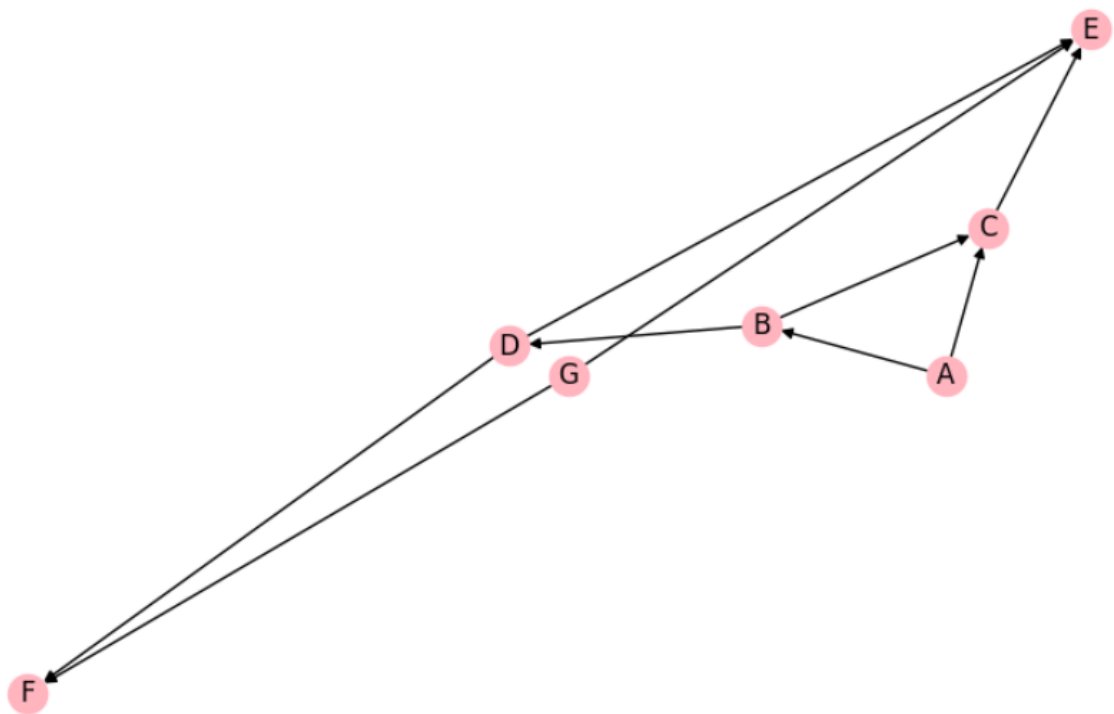
edges = [
    ("A","B"), ("A","C"),
    ("B","C"), ("B","D"),
    ("C","E"),
    ("D","E"), ("D","F"),
    ("G","F"), ("G","E")
]

G3 = nx.DiGraph()
G3.add_edges_from(edges)

plt.figure(figsize=(8,5))
nx.draw(G3, with_labels=True, node_color="lightpink", arrows=True)
plt.title("Course Prerequisite Graph")
plt.show()

# Topological sort
order = list(nx.topological_sort(G3))
print("One possible topological order:", order)

```



One possible topological order: ['A', 'G', 'B', 'C', 'D', 'E', 'F']