

# Fundamental Algorithm Techniques

## Problem Set #10

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### Problem 1: P, NP, NP-complete, NP-hard

We classify each group of problems into its corresponding complexity class. Each line below contains problems belonging to the same class.

#### 1. Class: P

Finding the maximum, linear search, shortest path in an unweighted graph, and matrix multiplication are all solvable in polynomial time.

#### 2. Class: P

Sorting a list, Dijkstra's algorithm with non-negative weights, BFS, DFS, merge sort, and quicksort all admit polynomial-time algorithms.

#### 3. Class: NP-complete

Sudoku (in its generalised  $n \times n$  form) belongs to NP, since a solution can be verified in polynomial time, and it is NP-hard via known reductions. Thus, it is NP-complete.

#### 4. Class: NP-complete

Graph 3-coloring is a classical NP-complete problem. Scheduling with conflicts can be reduced to graph coloring, and therefore belongs to the same class.

#### 5. Class: NP-complete

The Traveling Salesperson Problem (decision version), Hamiltonian Cycle, and Clique are all canonical NP-complete problems.

#### 6. Class: NP

Factoring large integers lies in NP since a factorisation can be verified efficiently. It is widely used in cryptography. It is not known whether this problem belongs to P.

#### 7. Class: Undecidable

The Halting Problem and the Busy Beaver problem are undecidable, meaning that no algorithm can solve these problems for all possible inputs.

## Problem 2: Introduction to Bayes' Theorem

A deadly disease affects 0.1% of the population:

$$P(D) = 0.001.$$

The medical test has:

- sensitivity  $P(+ | D) = 0.99$ ,
- specificity  $P(- | \neg D) = 0.99$ .

Therefore,

$$P(+ | \neg D) = 0.01, \quad P(\neg D) = 0.999.$$

Using Bayes' theorem, the probability that a patient actually has the disease given a positive test result is:

$$P(D | +) = \frac{P(+ | D) P(D)}{P(+ | D) P(D) + P(+ | \neg D) P(\neg D)}.$$

Substituting the values:

$$P(D | +) = \frac{0.99 \cdot 0.001}{0.99 \cdot 0.001 + 0.01 \cdot 0.999} = \frac{0.00099}{0.01098} \approx 0.09.$$

**Answer.** The probability that the patient actually has the disease is approximately 9%.

**Explanation.** Although the test is highly accurate, the disease itself is very rare. As a result, most positive test results come from false positives among healthy individuals. This phenomenon is known as the *base rate fallacy*.

## Problem 3: Introduction to Shannon Entropy

We consider three coins with different probabilities of landing heads. The Shannon entropy is defined as:

$$H(X) = - \sum_i p_i \log_2 p_i.$$

**Coin A: Fair coin ( $P(H) = 0.5$ )**

$$H = -(0.5 \log_2 0.5 + 0.5 \log_2 0.5) = 1 \text{ bit.}$$

A fair coin has maximum uncertainty, so each toss carries exactly one bit of information.

**Coin B: Biased coin ( $P(H) = 0.99$ )**

$$H = -(0.99 \log_2 0.99 + 0.01 \log_2 0.01) \approx 0.08 \text{ bits.}$$

Since the outcome is highly predictable, the amount of information gained from each toss is very small.

### **Coin C: Biased coin ( $P(H) = 0.01$ )**

This coin has the same entropy as Coin B:

$$H \approx 0.08 \text{ bits.}$$

Although a rare event is more surprising when it happens, such events occur so infrequently that the average information content remains low.

**Conclusion.** A fair coin is worth 1 bit of information because it has maximal uncertainty, while a strongly biased coin is worth much less, as its outcome can be predicted with high confidence.