

CSE 321 Homework #5

1-)

Firstly, For this part I can find number of sets that has equal sum zero but couldn't print them, just print how many sets there is that has zero sum. I use dynamic programming by using a 2D matrix. Initial length capacity for my matrix is 100x50. Then using by recursion and dynamic programming I fill this table.

```
26
27 arr = [2,3,-5,-8,6,-1];
28 print(Check_subset_zero(0, 0, arr, len(arr)));
```

For example, For given arr= [2,3,-5,-8,6,-1] my program output like this:

```
C:\Users\yasir\OneDrive\Masaüstü>python hw5_part1.py
4

C:\Users\yasir\OneDrive\Masaüstü>
```

It prints 5 because there is 5 possible subset of sum zero. They are

{2,3,-5}, {2,6,-8},{-5,6,-1},{3,6,-8,-1}

The complexity for this algorithm is $O(n*s)$ where s is sum of all elements and n is number of elements in given array.

2-)

In this part, I find minimum path by starting from bottom to top approach. The main idea is that start from nodes on the bottom row, After that minimum path at the i th node of k th row would be the minimum of the pathsum of its two children + the node's itself value.

Main dynamic programming usage is that:

$$Dp[i] = A[k][i] + \min(Dp[i], Dp[i+1])$$

Algorithm Works like that:

```
[
    [2],
    [5,4],
    [1,4,7],
    [8,6,9,6]
]

dp=[8,6,9,6]
```

Firstly, this is our triangle and dp will be bottom row of triangle.

```
[
    [2],
    [5,4],
    [1,4,7],
    [8,6,9,6]
]

dp=[7,10,13,6]
```

In second step our dp become like this.

I just use $Dp[i] = A[k][i] + \min(Dp[i], Dp[i+1])$ formula

For example left most element in dp become 7 because $dp[0] = 1 + \min(8,6) = 7$. And I find other dp elements like that then move from bottom to upper in my triangle.

Like this algorithm goes from bottem to top, in third step dp become like that:

```
[
    [2],
    [5,4],
    [1,4,7],
    [8,6,9,6]
]

dp=[12,14,13,6]
```

Lastly it will become :

```
[
    [2],
    [5,4],
    [1,4,7],
    [8,6,9,6]
]

dp=[14,14,13,6]
```

Here just return first element in dp it will smallest possible path. So the answer will be 14 for this triangle.

```
C:\Users\yasir\OneDrive\Masaüstü>python hw5_part2.py
14
C:\Users\yasir\OneDrive\Masaüstü>_
```

Time complexity for this algorithm is $O(n)$ where n is elements number in triangle. This is linear time because we access same element only one time in triangle.

3-)

This problem like same 0/1 knapsack problem that we have seen in class but in 0/1 knapsack problem we can select every items only once, in this problem we can select each item as much as we want.

For solving this problem dynamically I use 2 1D array instead of 2D array we use in 0/1 knapsack problem.

In algorithm, firstly I fill my dp with zeros. Then I use this dp formula: where $i \geq \text{weights}[j]$,

$$\text{dp}[i] = \max(\text{dp}[i - \text{weights}[j]] + \text{values}[j], \text{dp}[i])$$

Here weights array contains my possible items weights and values array contains those items values.

Finally the last element of my dp will be maximum value for given weights and I also hold an array selected Items.

For

weights = [5,4,2]

values= [10,4,3]

maxW = 9

My output will be like this:

```
C:\Users\yasir\OneDrive\Masaüstü>python hw5_part3.py
16
[3, 3, 1]
C:\Users\yasir\OneDrive\Masaüstü>
```

Here 16 is possible maximum choosable items total value. And [3,3,1] is selected items. It means for $W=9$ I can choose 2 times for item 3 and 1 times for item 1. Item 1 has 10 value and 5 weight, item 3 has 3 and 2 weight. So knapsack total weight will be 9 and total value will be 16 for those selected items.

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