

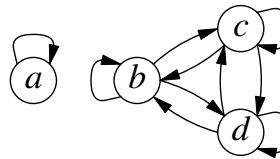
EQUIVALENCE RELATIONS

Equivalence Relation.

- A (binary) relation R on a set X is an **equivalence** relation if it is reflexive (R), symmetric (S), and transitive (T).
- We refer to these properties together, in short, as RST -properties.

Example.

- An equivalence relation on $X = \{a, b, c, d\}$. A simple way to represent this equivalence relation is to give its **equivalence classes** (defined below): $\{a\}$, $\{b, c, d\}$.



Equivalence Class $[x]$ of x in An Equivalence Relation R .

- $[x] = \{y: x \text{ related to } y \text{ in } R, \text{ i.e., } (x, y) \in R\}$. (Note: $x \in [x]$, by reflexivity of R .)
We can also say $[x] = \{y: y \text{ related to } x, \text{ i.e., } (y, x) \in R\}$, because R is symmetric.

Two Properties of Equivalence Classes.

- (P.1) If $y \in [x]$, i.e., $(x, y), (y, x) \in R$ then $[x] = [y]$. (We say x and y are **equivalent**.)
- (P.2) Two equivalence classes $[x]$ and $[y]$ are either **disjoint** ($[x] \cap [y] = \emptyset$) or **equal**.

Three Steps in Proving Property (P.1).

- (1) Suppose $y \in [x]$. For each $z \in [y]$, we have $(z, y) \in R$ and this together with $(y, x) \in R$ imply $(z, x) \in R$, by transitivity of R .
- (2) This means $z \in [x]$, and hence $[y] \subseteq [x]$.
- (3) Also, $(y, x) \in R$ implies $x \in [y]$ and hence $[x] \subseteq [y]$ as in (1)-(2). We now have $[x] = [y]$ because $[y] \subseteq [x]$ and $[x] \subseteq [y]$.

Proof of Property (P.2).

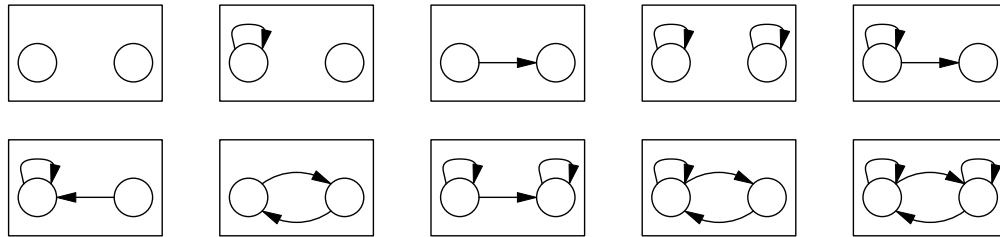
- If $[x] \cap [z] \neq \emptyset$ and $y \in [x] \cap [z]$, then $[x] = [y] = [z]$.

Disjoint Equivalence Classes Form a Partition of X .

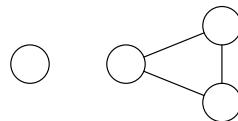
- A **partition** of X is decomposition of X into 1 or more **non-empty disjoint** subsets.
- $\#(\text{Partitions of } X) = \#(\text{equivalence relations on } X)$.

Practice Questions.

1. Shown below are the structures of all relations on $|X| = 2$. Mark those which correspond to equivalence relations.



2. Show the digraph of structures of equivalence relations on $n = 3$ items. Also, for each digraph show $\#(\text{equivalence relations})$ for that structure.
3. Show $\#(\text{equivalence relations on } n \text{ items with exactly 2 equivalence classes})$. Verify your answer for $n = 2$ and 3 from your solutions of Problems 1 and 2.
4. Shown below is the graph (undirected links and without loops to simplify the diagram) for the structure of the equivalence relation given in the previous page; show all other structures of equivalence relations on $\{a, b, c, d\}$ items. In each case, show (a) $\#(\text{equivalence classes of size } k)$ for each $k \geq 1$, and (b) $\#(\text{equivalence relations for that structure})$.



- (a) 1 equiv. class of size 1;
1 equiv. class of size 3
- (b) $\#(\text{equiv. rels.}) = 4$;

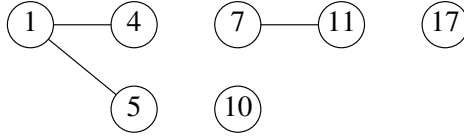
5. Complete the code below for printing the equivalence classes of an $n \times n$ equivalence relation-matrix R . Show the output for the equivalence relation with equivalence classes $\{0, 1, 3\}$, $\{2, 5\}$, $\{4, 6\}$.

```
int i = 0, //i = smallest item in an equiv. class that is found next
    j, printed[] = new int[n];
do { for (j = 0; j < n; j++) //find items in equiv. class of i
    { if (1 == R[i][j])
      { printed[j] = 1;
        System.out.println(j + " is equivalent to " + i);
      }
    }
    //find smallest item not printed yet
    for (j = 0; j < n; j++)
        if (0 == printed[j]) break; //don't use 0 == R[i][j];
    if (j < n) //there is an equiv. class with smallest item j
        .....
    else break;
} while (.....);
```

How do you avoid starting from $j = 0$ every time to find an item not printed yet?

How would you modify the above code to count $\#(\text{equivalence classes})$?

6. Let $0 \leq p \leq q$ be two fixed numbers and X a non-empty set of numbers. We define the relation $R_{p,q}$ on X by $\{(x, y): x, y \in X \text{ and } p \leq |x - y| \leq q\}$. Show the missing links in the following graph of the anti-reflexive and symmetric relation $R_{3,5}$ for $X = \{1, 4, 5, 7, 10, 11, 17\}$. Which of the transitive, non-transitive, and anti-transitive properties hold for $R_{3,5}$?



7. Consider the codes (a)-(n) below. We define *ES*-relation below two codes C and C' as follows: $(C, C') \in ES$ if the flowcharts of C and C' have the same structure when we ignore the contents of tests and assignments. First, argue that *ES*-relation satisfies *RST*-properties and hence it is an equivalence relation. Show the *ES*-equivalent classes of codes. Also, show the flowchart for each equivalence class.

- | | |
|---|---|
| (a) <code>max = first;</code>
<code>if (second > max) max = second;</code> | (b) <code>if (first > second) max = first;</code>
<code>else max = second;</code> |
| (c) <code>max = first;</code>
<code>if (second >= max) max = second;</code> | (d) <code>if (first >= second) max = first;</code>
<code>else max = second;</code> |
| (e) <code>max = first;</code>
<code>if (second > first) max = second;</code> | (f) <code>if (first > second) max = first;</code>
<code>if (first <= second) max = second;</code> |
| (g) <code>max = first;</code>
<code>if (second >= first) max = second;</code> | (h) <code>if (first >= second) max = first;</code>
<code>if (first < second) max = second;</code> |
| (i) <code>max = second;</code>
<code>if (first > max) max = first;</code> | (j) <code>if (first >= second) max = first;</code>
<code>if (first <= second) max = second;</code> |
| (k) <code>max = second;</code>
<code>if (first >= max) max = first;</code> | (l) <code>if (first < second) max = second;</code>
<code>else max = first;</code> |
| (m) <code>max = second;</code>
<code>if (first > second) max = first;</code> | (n) <code>if (first <= second) max = second;</code>
<code>else max = first;</code> |