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CS2030S Midterm Cheatsheet	cs2030s
	Definition > T-shaped blocks that convert rightwards
Compiler (Tombstone)	> Compilers (javac) convert written code to bytecode language (virtual machine language) > Compilers themselves are written in a certain language, which needs to be interpreted to the hardware instructions
	Definition Period
Interpreter (Tombstone)	> Interpreters (java) convert bytecode language to hardware instructions > Interpreters (jshell) can also convert written code directly to hardware instructions (line by line)
	Definition > Compilers (javac program.java) will create a bytecode script (.class) from the written script (.java)
Executable (Tombstone)	> Interpreters (java program.class) will run the code Definition
	> Method Name > Number of Arguments
Method Signature	> Type of Arguments > Order of Arguments
Method Dighatare	Definition
	> Models the relationship between object types using boxes and arrows
	Topics ?> Has A Relationship
Composition	<pre>?&gt; Is A Relationship ?&gt; Can Do Relationship</pre>
	Usage > Stores elements in an expanding fashion
	> <jv import="" java.util.*="" jv=""></jv>
	Types  ?> Java ArrayList
	?> Java_LinkedList
	Methods > <jv jv="" list(len)=""></jv>
	> <jv add()="" jv=""> &gt;<jv get()="" jv=""></jv></jv>
	> <jv addall(list)="" jv=""> &gt;<jv jv="" size()=""></jv></jv>
	> <uv jv="" size()=""> &gt;<jv contains(obj)="" jv=""></jv></uv>
	> <jv jv="" remove(obj)=""></jv>
Java List	> <jv clear()="" jv=""> &gt;<jv jv="" set(id,obj)=""></jv></jv>
	Usage > Can super() to parent construct
Java Super	> Can super.method() to use parent method > Can super.field to use parent class field
	Usage > Figures out which method implementation is called
	Properties
	> Class methods cannot be dynamically bound
	Process ?> Compile Time
Dynamic Binding	?> Run_Time
	Topics  ?> Compile_Time_Type
	Process
	> Identify the /Compile Time_Type of the parameters > Check the compile-time class for method
	> Identify the most specific implementation of the method to use > Store method descriptor in bytecode
Compile Time	> Then return that method's signature (name, *args) > /Type_Erasure
	Topics ?> Run Time Type
	Properties
	> Usage :- /Run_Time_Type of target, /Compile_Time_Type of arguments
	Process > Get method descriptor from /Compile Time
Run Time	> Starting from the /Run_Time_Type, go up the /Class_Heirarchy until you find the first matching method
Liskov Substitution	Definition > A subclass should not break the expectations set by the superclass
Principle	> You can do extra, but not less Definition
	> Data is immutable > Classes cannot be extended
Final Modifier	> Methods cannot be overridden  Usage
	> Understand what happens when we make method calls
	Topics ?> Stack Stack Heap Diagram
	Stack_Stack_Heap_Diagram
	>> Stack_Frame
	<pre>Process &gt; Initialise with variables / stack frames{}, including the frame for main(){}</pre>
	> Draw boxes for object instances and draw arrow to point from variable to object > If there is a new method call, create a new stack frame on top of the stack, and point arrows wherever
Stack Heap Diagram	> Once done with method, remove stack frame and clear all reference arrows  Once done with main(), remove all object frames with nothing referencing it
noup Diagram	Definition
	> Contains variables, stack frames{} for constructors / methods
	Properties > Frames can only point to the /Heap_Stack_Heap_Diagram_
Stack (Stack Page	> Variable names have no types, but rather types are based on what they point to
Stack (Stack Heap Diagram)	Examples > Block containing c1, c2, p
	Definition > Instances of objects
	Properties
	> Frames can point to each other
Heap (Stack Heap Diagram)	Examples \[ \text{(irrle frame with center and radius} \]
ay_am,	> Circle frame with center and radius Definition
	> Fixed as what the variable was defined as > Compiler will tell you what will go wrong
	Examples
Compile Time Type	> <jv c="new" circle="" colouredcircle();="" jv=""> - gives Circle</jv>

	Definition > Check what object the variable is pointing to > No compiler can tell you there will be a mistake there
	Examples
Run Time Type	> <jv c="new" circle="" colouredcircle();="" jv=""> - gives ColouredCircle  Definition</jv>
Checked Exception	> Programmer has no control > User might give an incompatible type or do something unexpected
	Usage > Systematically create specific classes for many occasions
	Properties > Cannot mix with array :- Pair <string, integer="">[2] can't work as it will convert it into Object, and it is not type safe</string,>
	Topics ?> Generic Type
	?> Generic Method ?> Type Erasure
	<pre>?&gt; Wildcard _Generics _ ?&gt; Java_Equals _Generics _ ?&gt; Type Inference _Generics _</pre>
Generics	2> Valid Generics Examples
	> <jv jv="" p="new" pair="" pair(2,3)=""> &gt;<jv integer="" pair<integer,=""> p = new Pair&lt;&gt;(2,3) JV&gt;</jv></jv>
	> <jv jv="" p="new" pair="" pair<\frac{1}{2,3}=""> &gt;<jv jv="" p="new" pair="" pair<\(2,3)=""> &gt;<jv list<?=""> L = new List<string>() JV&gt;</string></jv></jv></jv>
Valid Generics	> <uv black:="" blackstring*()="" d="New" uv=""> &gt;<uv black?,?=""> p = new Pair<integer>() JV&gt; Examples</integer></uv></uv>
Tanadid Ga	> <jv pair<=""> p = new Pair&lt;&gt;(2,3) JV&gt; &gt;<jv integer="" pair<integer,=""> p = new Pair<?,?>() JV&gt;</jv></jv>
Invalid Generics Type Inference	> <jv list<?=""> L = new List<? >() JV&gt; Definition</jv>
(Generics)	> Compiler will list out all possible types for the type parameter T, based on the bounds provided > Inference can only happen in the rightwards direction
	Purpose > Class where you can pass a type as a variable T > Compiler will make sure the types are correct
	Types
	<pre>?&gt; Raw Type ?&gt; Bounded Type ?&gt; Unbounded Type</pre>
	?> Unchecked_Type
	Topics ?> Type_Invariance
	Usage > Treats T and S as variables which can be used later
	> Cannot directly create an array of T[], can be typecasted with <jv arr="(T[])" jv="" new="" object[n]}="" t[]=""></jv>
	Examples <uv class="" pair<t,="" s=""> {   public Pair(T obj1, S obj2) {}</uv>
	Pair <string, integer=""> p = new Pair<string, integer="">()</string,></string,>
Generic Type	<pre>class Pair<t comparable<t="" extends="">, S&gt; implements Comparable<pair<t, s="">&gt; {} JV&gt;</pair<t,></t></pre>
	Usage > Once the code is compiled, all the generic types are removed, and values are typecasted > Changes all type parameters to Object if unbounded (T → Object), or specific type if bounded (T extends Comparable <t> → Comparable) ?&gt; Code Sharing</t>
	<pre>Examples <jv pair<integer,="" string=""> p = new Pair<integer, string="">(1, "one"); Integer i = p.getFirst();</integer,></jv></pre>
	JV> Becomes
	<pre><jv "one");="" i="(Integer)" integer="" p="new" p.getfirst();<="" pair="" pair(1,="" pre=""></jv></pre>
Type Erasure	JV> Usage
	> Only use it for /Java_Instance_Of
Raw Type	Disadvantages - Compiler cannot catch mistakes
	Definition > Take information out of the object
	Usage
Producer Extends	Code > <jv <?="" extends="" static="" t=""> T foo(T){} JV&gt;</jv>
	Definition > Insert information into the object
	Usage > /Return_Type > super T
Consumer Super	Code > <jv <t="" static=""> T foo(Array<? super T>){} JV&gt;</jv>