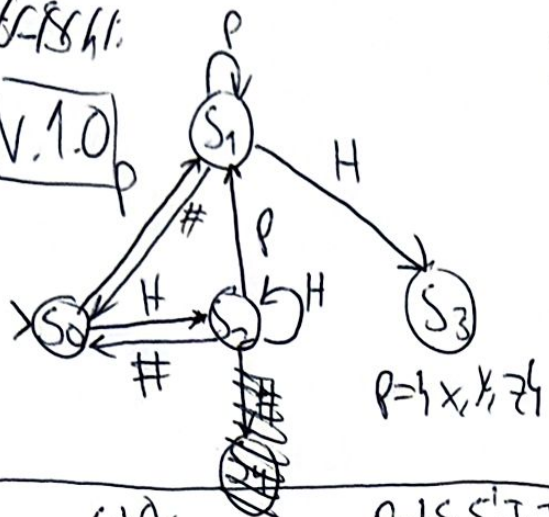


86-8841

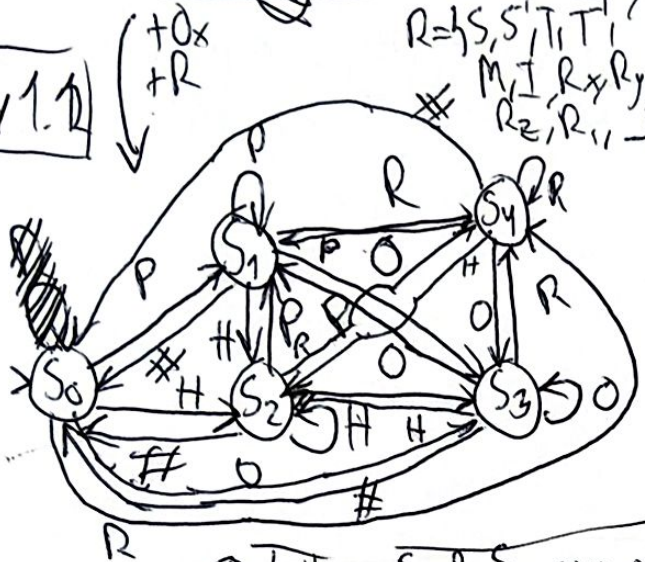
V.1.0



$$\Sigma = \{P, H, \#\}$$

- $Q = \{$
- S_0 : Initial state / New "Empty" Column
 - S_1 : Initialization Pat. found,
 - S_2 : Potential Superposition Pat.
 - S_3 : Potential Entanglement Pat.
 - ~~S_4 : Superposition Pat. found~~

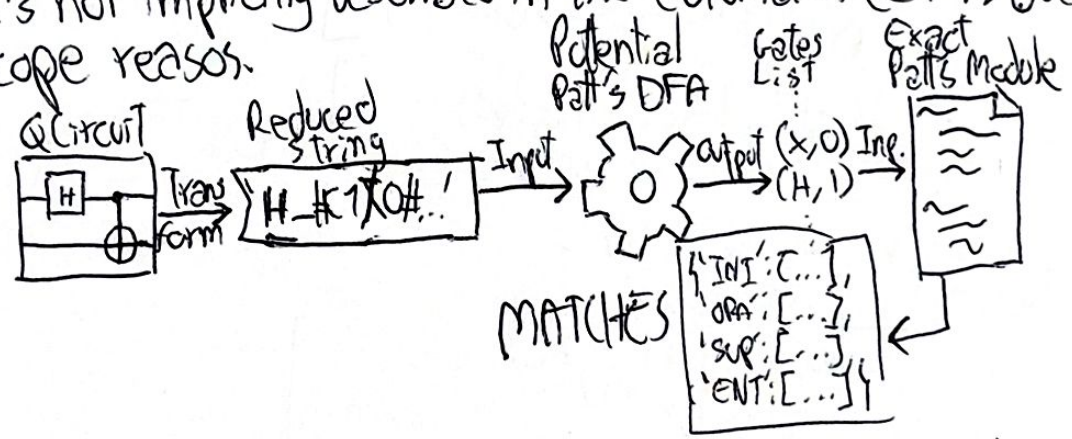
V.1.1



$$\Sigma = \{P, H, \#, R, O\}$$

- $Q = \{$
- S_0 : New Column,
 - S_1 : 'INI' potentially found, *
 - S_2 : 'SUP' potentially found, *
 - S_3 : 'CRA' found,
 - S_4 : Remaining gate found [Neutral]

* When S_1 & S_2 are reached, a method for checking whether there exists an exact match should be ran. It's not implicitly described in the automaton (DFA) due to scope reasons.



	H	P	R	#	\$	O
S_0	S_2	S_1	S_4	ERR	S_5	S_3
S_1	S_2	S_1	S_4	S_0	S_5	S_3
S_2	S_2	S_1	S_4	S_0	S_5	S_3
S_3	S_2	S_1	S_4	S_0	S_5	S_3
S_4	S_2	S_1	S_4	S_0	S_5	S_3

Tabla de Transiciones

S_5 ESTADO FINAL
 ↳ No codificado, sim. termina la ejec. de forma correcta

State Definition

(S_i) { def _s_i(self): V0

```

# print #
# State task
self.checker.check_wtf(...).
# Move gate for next state
self.gate = self._next()
# Transition
switch(self.gate):
    case 0:
        :
    case N:
    
```

!!! One gate is skipped.
switch checks i+1.

(S_i) def _s_i(self):

```

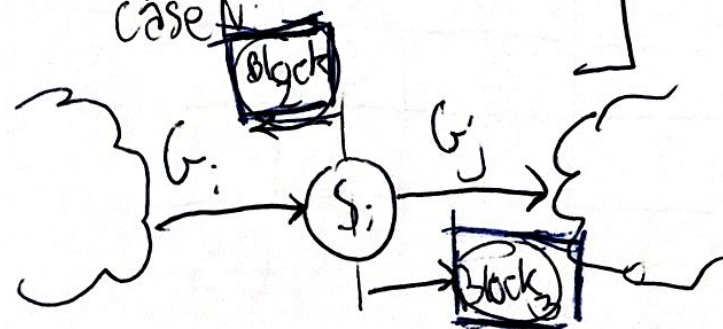
# print #
print(f' state Si → Gate {i} ...')
# state task (if needed).
self.checker.whatever(...).
# Transition
switch(self.gate):
    case 0:
        :
    case j:
        self._sj() ← return ...
    case N:
        :
    
```

Block 1 (optional) *

Block 2 (Depends on i).

Block 1 (Optional) *

Block 3 (Compulsary)



* Depends on the State task

Next Char Function

...#XYZ#_02-#... ← Input String (compressed circuit).
 self.pos. ↑ (+1) ⇒ Next method
 ⊗ self.real_pos goes apart,

```
def _next(self):
```

```
    self.pos += 1
```

```
    try:
```

```
        next_char = self.circuit[self.pos]
```

```
    except IndexError: # End of circuit assumed in error case.
```

```
        return '$'
```

```
    # specific cases.
```

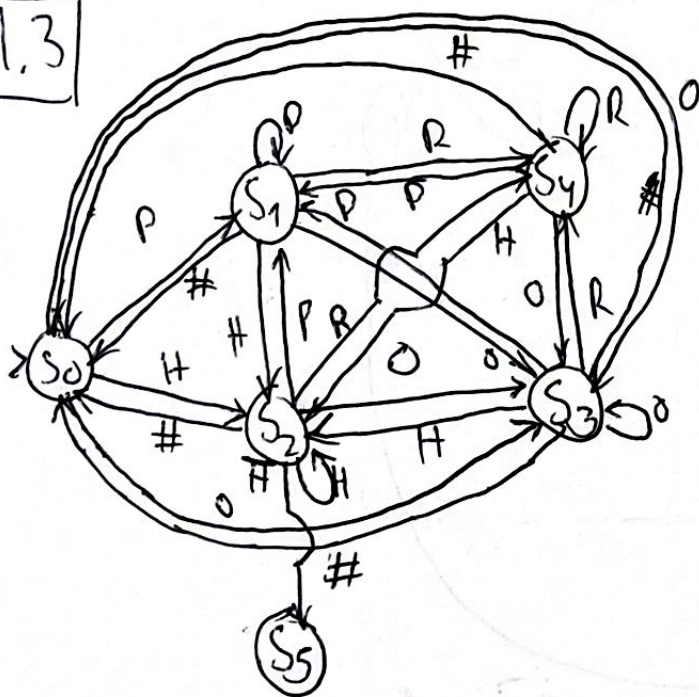
```
    if next_char == '0': # Orade length adding.
```

```
    elif next_char == '#': # Next column jump.
```

```
    ...
```

```
    return next_char.
```

V1.3



$\Sigma = \{P, H, \#, R, 0, \tau, \epsilon\}$

$Q = \{S_0: \text{New column}$

$S_1: \text{'INI' potentially found.}$

$S_2: \text{'SUR' — " — .}$

$S_3: \text{'ORA' — " — .}$

$S_4: \text{Remaining gate found [Neutral]}$

$S_5: \text{1st part of 'ENT' found.}$

$\bar{H} \equiv \text{There is not a uniform superposition.}$

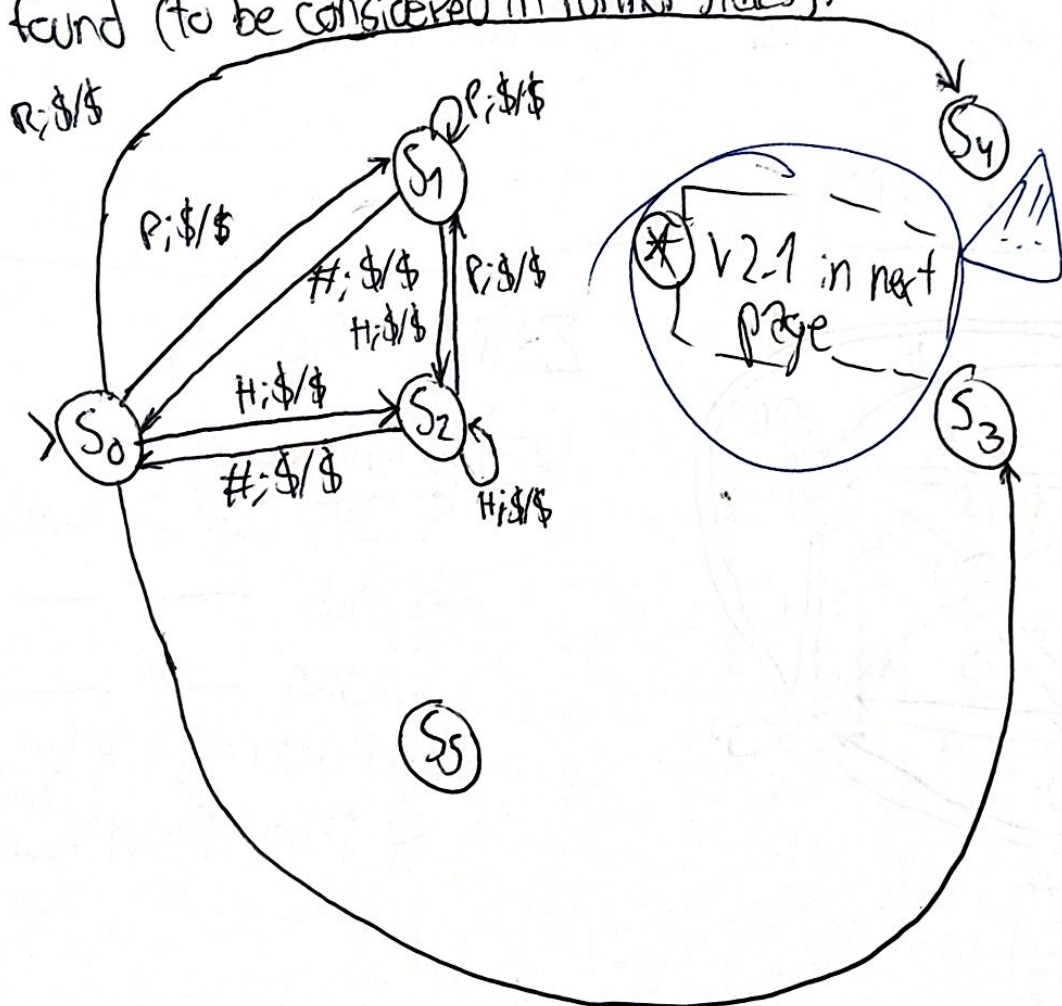
V2.1

For this new version, we will split a pattern into parts depending on its atomic behavior, such that:

- 'INI': 1 part (composed by Pauli gates).
- 'SUP': 1 part (composed by Hadamard gates).
- 'ORA': 1 part (composed by the Oracle itself).
- 'ENT': 2 parts {1) Superposition of control qubit (ENT₀).
(2) CNOT for both qubits (ENT₁).

To model the atomic-division of those patterns for its recognition, we will adapt the DFA v1.3 to a PUSHDOWN AUTOMATON (PA) v.2.1.

In the PA's stack we will store whether an atomic-part of the pattern is found (to be considered in further states).



Let's define Ψ as:

- Any symbol in Γ (stack alphabet) excepting those used in other transition from that state.
- $\Gamma / \dots / \emptyset$ means that could be found or not.

5



- $\Psi \equiv$ Any symbol in Γ (stack alphabet) excepting those used in other transitions from that state.

- $\overline{I}/\dots/\overline{S} = \overline{I}$ means any symbol except the one written under the upper line.

- Formal definition of PEPDM-PD- λ :

A pushdown automaton with λ -transitions defined as $M = (Q, \Sigma, \Gamma, q_0, z, \delta, \lambda)$
 where:

$$q_0 = S_0$$

$$z = \$$$

$$F = \{S_8\}$$

$$\Sigma = \{P, R, H, T, C, \#, O, \emptyset\} \quad \begin{matrix} \nearrow \text{End of circuit} \\ (\text{EOC}) \end{matrix}$$

$Q = \{$ S_0 : New Column (no atomic-part detected) \rightarrow No pattern in progress

S_1 : 'INI' potentially found.

S_2 : 'SUB' — " —

S_3 : 'ORA' — " —

S_4 : Remaining gate found [Neutral] (no atomic-part detected).

S_5 : Atomic Part 'ENT₀' (E) found potentially.

S_6 : Transitioning between 'ENT₀' \rightarrow 'ENT₁' (atomic-part detected).

S_7 : 'ENT' potentially found.

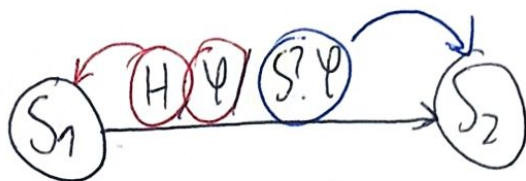
S_8 : Search has finished

$$\Gamma = \{I, S, O, E, \epsilon, \$\}$$

$$\delta = // \text{Check prev page PDA-}\lambda$$

- Code Translation for Transition

Let's take as example the following transition:



Let's now give the pseudocode for S_2 & S_1 :

```
def self._s1():
```

```
    # Inform
```

```
    ...
```

```
    # Task
```

```
    ...
```

```
    # Transition
```

```
    self.gate = self._next()
```

```
    if PAULI_R.search(self.gate) is not None:
```

```
        ...
```

```
        elif self.gate == 'H':  
            return self._s2()
```

```
def self._s2():
```

```
    # Inform
```

```
    ...
```

```
    # Task
```

```
    if checker.check_sup(gate_top):
```

```
        self.stack.pop()
```

```
        self.stack.append('S')
```

```
    # Transition
```

```
    ...
```

!!! IMPORTANT !!!

Note the transition-tasks division:

- The condition regarding the stack top is done in S_1 (left-most state).

- The stack replacing is done in S_2 (right state).

depends on transition

(q, cd, str)