

Problem 1 (2 parts, 30 points)**Storage Allocation, Arrays, and Pointers**

Part A (16 points) Assuming a **64-bit system with 64-bit memory interface and 64-bit addresses**, show how the following global variables map into static memory. Assume they are allocated starting at address 4000 and are properly aligned. **For each variable, draw a box showing its size and position** in the double word memory shown below in which byte addresses increment from left to right. **Label the box with the variable name.** Label each element of an array (e.g., M[0]). Note: int and float are 32-bits.

		4000	a	b
int a = 5;		4008	c	
int b = 25;		4016	d	
int *c = &b;		4024	e	
double d = 9.23;		4032	S[] (6 bytes)	slack (2 bytes)
double *e = &d;		4040	f	
char S[] = "Harry";		4048	g	slack (4 bytes)
char *f = &(S[3]);		4056	z	
float g = 8.1;		4064		
double z = 17.25;				

Part B (14 points) Assuming a 32-bit system, consider the following declarations:

```
int A[32][16][8] = {...};
int *q = A;
```

B.1 Complete the assignment statement below using only *q* to assign to *x* the value of A[20][10][2].

```
int x = *(q + 20*16*8+10*8+2); //an expression is ok
or int x = *(q + 2642)
```

B.2 Write the MIPS code implementation of the following assignment statement in the smallest number of instructions. A pointer to the array *A* is stored in \$3 and variables *j*, *i*, and *y* reside in \$4, \$5, and \$6, respectively. Modify only registers \$1 and \$2.

```
int y = A[0][j][i];
```

Label	Instruction	Comment
	sll \$1, \$4, 3	# j*8 (= j * Lx)
	add \$1, \$1, \$5	# j*8 + i
	sll \$1, \$1, 2	# scale by 4
	add \$1, \$1, \$3	# add A
	lw \$6, 0(\$1)	# memory read

Problem 2 (4 parts, 30 points)**Accessing Structs, Activation Frame Allocation**

Consider the following C code fragment.

```
typedef struct {
    int    age;
    float height;    // in meters
    float weight;    // in kilograms
} patient;

float Calc_and_Print_BMI(int A, float H, float W) {
    patient  Pat;
    float    BMI;

    __Pat.age = A;__ // part A

    __Pat.height = H;__ // part A

    __Pat.weight = W;__ // part A

    BMI = Compute_BMI(&Pat);
    printf("BMI: %f\n", BMI);
    return(BMI);
}

/* BMI = weight divided by height squared (units: kg/m²). */
float Compute_BMI(patient *P) {

    float height_squared = __ (P->height)*(P->height) __; // part B

    float BMI = __ P->weight __ / height_squared;    //part B

    return(BMI);
}
```

Part A (6 points) Fill in the blanks in `Calc_and_Print_BMI` with statements that assign the inputs A, H, and W to the age, height, and weight fields of `Pat`, respectively.

Part B (6 points) Fill in the blanks in `Compute_BMI` with 1) a statement that squares the height of the patient pointed to by `P` and 2) a statement that computes the body mass index (BMI) of the patient pointed to by `P` by dividing the patient's weight by `height_squared`.

Part C (12 points) Assuming a 32-bit system, give the total number of bytes allocated for the activation frame of `Calc_and_Print_BMI` and the number of bytes for the activation frame of `Compute_BMI`.

Size of activation frame of <code>Calc_and_Print_BMI</code> (in bytes):	40 bytes
Size of activation frame of <code>Compute_BMI</code> (in bytes):	24 bytes

Part D (6 points) What does the following code fragment print?

Code:

```
char Name[] = "Fred";
char *p = Name;
while (*(++p))
    printf("%c\n", *p);
```

Answer:

```
r
e
d
```

Problem 3 (1 part, 40 points)**Activation Frames**

The function Bar (below left) calls function Foo after completing code block 1. Write MIPS assembly code that properly calls Foo. Include all instructions between code block 1 and code block 2. Symbolically label all required stack entries and give their values if they are known (below right).

<pre> int Foo(int Set[], int x, int *y, int z) { <body of Foo> } int Bar() { int A[] = {40, 45, 50}; int B = 10; int i; <code block 1> A[i] = Foo(A, A[2], &B, 0); <code block 2> } </pre>	Bar's FP	9620	XXX	XXX
		9616	A[2]	50
		9612	A[1]	45
		9608	A[0]	40
		9604	B	10
	SP	9600	i	assigned in code block 1
		9596	RA	
		9592	FP	9620
		9588	Set	9608
		9584	x	50
		9580	y	9604
		9576	z	0
	SP & Foo's FP	9572	RV	

label	instruction	comment
	addi \$29, \$29, -28	# Allocate activation frame
	sw \$31, 24(\$29)	# Preserve bookkeeping info
	sw \$30, 20(\$29)	
	addi \$1, \$30, -12	# Push inputs: compute A
	sw \$1, 16(\$29)	# push A
	lw \$2, 8(\$1)	# read A[2]
	sw \$2, 12(\$29)	# push A[2]
	addi \$2, \$30, -16	# compute &B
	sw \$2, 8(\$29)	# push &B
	sw \$0, 4(\$29)	# push 0
	jal Foo	# Call Foo
	lw \$31, 24(\$29)	# Restore bookkeeping info
	lw \$30, 20(\$29)	
	lw \$3, 0(\$29)	# Read return value
	lw \$2, -20(\$30)	# Store return value in A[i]:
	sll \$2, \$2, 2	# load i, scale it, and add it
	addi \$1, \$30, -12	# to base address of A to
	addi \$1, \$1, \$2	# compute where to store return
	sw \$3, 0(\$1)	# value, and store RV there.
	addi \$29, \$29, 28	# Deallocate activation frame