

*Instructions:* This is a closed book, closed note exam. Calculators are not permitted. If you have a question, raise your hand; do not leave your seat. Please work the exam in pencil and do not separate the pages of the exam. For maximum credit, show your work.

*Good Luck!*

**Your Name (please print clearly)** \_\_\_\_\_

*This exam will be conducted according to the Georgia Tech Honor Code. I pledge to neither give nor receive unauthorized assistance on this exam and to abide by all provisions of the Honor Code.*

**Signed** \_\_\_\_\_

1	2	3	4	total
22	28	20	30	100

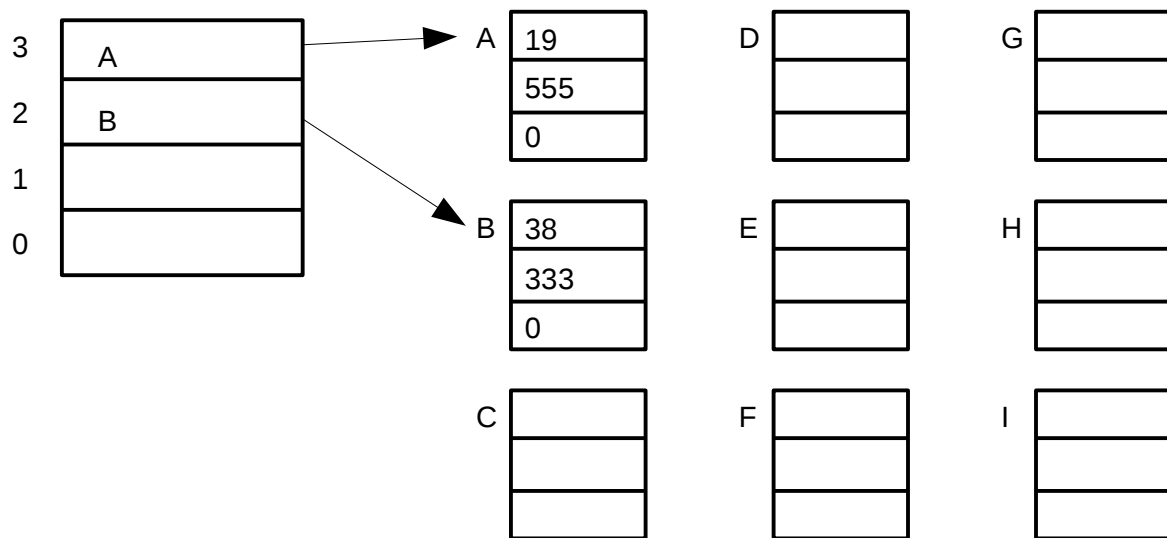


**Problem 1** (2 parts, 22 points)**Hash Tables**

**Part A** (14 points) Consider an open hash table composed of a four-bucket table, with each bucket containing a variable length list. Each list entry has three slots <key, value, next> corresponding to the three word groupings in the entries section. The hash function is  $\text{key} \bmod \text{four}$ . Inserted entries are *appended to the end* of a bucket list. The initial state of the hash table is shown. List elements as allocated by malloc are shown in the figure. The symbol to the left of each list element (A, B, C,...) is the address returned by malloc. Entries that are freed are maintained in a last-in-first-out (LIFO) free list. Assume the free list is initially empty.

Execute the access trace shown in the table below. For ease of representation, you may use the allocated blocks in any order. Show pointers both by their (symbolic) value, and with an arrow. If a value changes, cross it out and place the new value to the right. If a pointer changes, also place an x on the corresponding arrow, and draw the new arrow.

Buckets (Hash Anchor Table)



Hash Table Access Trace

#	op	key	value	#	op	key	value
1	insert	11	111	3	remove	19	n/a
2	insert	38	222	4	insert	15	777

**Part B** (8 points) Consider a different hash table that uses **7 buckets**, each containing a singly linked list of entries. The hash table contains a total of **140 entries** evenly distributed across the hash table buckets. An application performs **1000** lookups of various keys: **600** of the lookups find the key in the hash table and **400** lookups fail to find the key. The keys that are found are distributed throughout the buckets so that each position is equally likely to be where a key is found. How many key comparisons would be required for the average lookup in the hash table if each bucket list is unsorted versus sorted?

number of comparisons when each bucket list is *unsorted*:

number of comparisons when each bucket list is *sorted*:

**Problem 2** (4 parts, 28 points)**Dynamic Memory Allocation on Heap**

Consider a memory allocator (malloc and free), such as described in class. Inside the C-code for the allocator, unsigned `*heapPtr` is the address of the next word that could be allocated to the heap, and unsigned `**freePtr` is the address of the first block on the free list (and the word at the address of each block on the free list is a pointer to the next block on the free list). The allocator uses a **best fit** strategy with an **unsorted** free list, and never splits blocks.

addr	value	addr	value	addr	value	addr	value	addr	value	addr	value
<b>8000</b>	8	<b>8032</b>	20	<b>8064</b>	4	<b>8096</b>	8048	<b>8128</b>	8	<b>8160</b>	0
<b>8004</b>	1234	<b>8036</b>	0	<b>8068</b>	12	<b>8100</b>	8104	<b>8132</b>	8004	<b>8164</b>	0
<b>8008</b>	4	<b>8040</b>	43	<b>8072</b>	8036	<b>8104</b>	4	<b>8136</b>	4	<b>8168</b>	22
<b>8012</b>	16	<b>8044</b>	12	<b>8076</b>	8144	<b>8108</b>	2	<b>8140</b>	42	<b>8172</b>	7000
<b>8016</b>	8072	<b>8048</b>	8096	<b>8080</b>	8	<b>8112</b>	12	<b>8144</b>	43	<b>8176</b>	12
<b>8020</b>	8052	<b>8052</b>	12	<b>8084</b>	4	<b>8116</b>	0	<b>8148</b>	427	<b>8180</b>	41
<b>8024</b>	8132	<b>8056</b>	8	<b>8088</b>	0	<b>8120</b>	4	<b>8152</b>	8	<b>8184</b>	40
<b>8028</b>	8116	<b>8060</b>	8116	<b>8092</b>	16	<b>8124</b>	30	<b>8156</b>	0	<b>8188</b>	0

Suppose `heapPtr = 8128` and `freePtr = 8016`. Consider each part below independently.

Part A: (4 pts)	How many <b>blocks and useable bytes</b> are on the free list? blocks = ____ bytes = ____
Part B: (12 pts)	<p>B.1) What value would be returned by the call <code>malloc(10)</code>; _____</p> <p>B.2) Which (if any) values in the above map would be changed by the call in B.1?</p> <p>addr_____value_____addr_____value_____addr_____value_____no change____</p> <p>(fill in the address/value pairs above. There may be more pairs than needed)</p> <p>B.3) Fill in the values after the call in B.1: <code>heapPtr = _____ freePtr = _____</code></p>
Part C: (8 pts)	<p>C.1) What value would be returned by the call <code>malloc(22)</code>; _____</p> <p>C.2) Which (if any) values in the above map would be changed by the call in C.1?</p> <p>addr_____value_____addr_____value_____addr_____value_____no change____</p>
Part D: (4 pts)	<p>Which (if any) values in the above map would be changed by the call <code>free(8116)</code>?</p> <p>addr_____value_____addr_____value_____addr_____value_____no change____</p>

**Problem 3** (4 parts, 20 points)**Heap Management**

Consider the following three heap management strategies:

1. First fit with free list sorted by increasing size (smallest to largest).
2. First fit with free list sorted by decreasing size (largest to smallest).
3. Best fit with unsorted free list.

**Part A** (5 points) Which strategy (1, 2, or 3) has fastest average speed of malloc? \_\_\_\_\_

Why?

**Part B** (5 points) Which strategy (1, 2, or 3) has slowest average speed of malloc? \_\_\_\_\_

Why?

**Part C** (5 points) Which strategy (1, 2, or 3) has fastest average speed of free? \_\_\_\_\_

Why?

**Part D** (5 points) Which strategy (1, 2, or 3) has the worst internal fragmentation? \_\_\_\_\_

Why?

## Linked Lists

Suppose we have the following definition which is used to create singly-linked lists.

```
typedef struct Element {
    int          Num;
    struct Element *Next;
} Link;
```

**Part A** (4 points) Fill in the blank below to allocate space for a `Link` structure using `malloc` and make the variable `L1` point to the object allocated.

```
Link *L1 = _____;
```

**Part B** Complete the following recursive function which takes a pointer (called Head) to the first Link of a linked list and a pointer (called NewEnd) to a Link (which has a null Next field). The function places NewEnd on the linked list as its last element. Follow the steps specified below.

```
Link * Append(Link *Head, Link *NewEnd) {
    if (Head == NULL) return _____; /* part B.1 */
    _____; /* part B.2 */
    return _____; /* part B.3 */
}
```

**Part B.1** (5 points) Fill in what should be returned if the list is empty.

**Part B.2** (8 points) Call Append recursively to place NewEnd on the end of the rest of the list and then push Head on the result of this recursive call.

**Part B.3** (3 points) Fill in what should be returned.

**Part C** The following subroutine should free up all elements in the linked list whose first `Link` is pointed to by the input parameter `Head`. What error does it make? Write the correct code below.

```
void FreeElements(Link *Head)
{   Link *h;
    for (h = Head; h != NULL; h = h->Next)
        free(h);
}
```

**Part C.1** (4 points) What is the error? \_\_\_\_\_

**Part C.2 (6 points)** Write the correct version of FreeElements?

```
void FreeElements(Link *Head)
{
```

}