

**Problem 1** (2 parts, 24 points)**Storage Allocation, Strings, and Pointers**

**Part A** ( 16 points) **Assuming a 64-bit system with 64-bit memory interface and 64-bit addresses**, answer the following addressing questions. Assume all alignment restrictions imposed by the hardware are obeyed, and the compiler does not add additional alignment restrictions. Note: `int` and `float` are 4 bytes, and `double` is 8 bytes. For each part below, fill in the value of each expression given that the expression in the comment is true. You may find it helpful to sketch memory allocation including slack for each part. Assume variables are allocated in **global memory** in the order they are declared. **Please only write numbers in each answer box.**

**Part A1** (4 pts).

<pre>int i; // &amp;i == 1000 char c[2]; double x</pre>	<code>&amp;c[1]</code>	<b>1005</b>
	<code>&amp;x</code>	<b>1008</b>

**Part A2** (8 pts).

<pre>char *s; // &amp;s == 1000 char d = 'T'; struct {     char c;     double y;     int j;     float z; } thing;</pre>	<code>&amp;d</code>	<b>1008</b>
	<code>&amp;thing</code>	<b>1016</b>
	<code>&amp;thing.y</code>	<b>1024</b>
	<code>sizeof(thing)</code>	<b>24</b>

**Part A3** (4 pts).

<pre>int m; // &amp;m == 1000 float *q; int *p = &amp;m;</pre>	<code>&amp;q</code>	<b>1008</b>
	<code>p+1</code>	<b>1004</b>

**Part B** ( 8 points) Fill in what is printed after the following C fragment executes?

```
char *foo(char *s, char *t){
    char *z = s;
    while (*t){
        if (*t == 't' || *t == 'T')
            t++;
        else
            *s++ = *t++;
    }
    *s = *t;
    return z;
}
char a[25];
char b[] = "This cart tis there!";
printf("%s\n", foo(a,b)); //only write exactly what is printed
```

What is printed?	<b>his car is here!</b>
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**Problem 2** (4 parts, 40 points)**Accessing Inputs, Locals, Arrays**

**Part A** (16 points) Consider the following C code on a 32-bit machine:

```
typedef unsigned char color;
typedef struct {
    color r;
    color g;
    color b;
} pixel;
pixel frame[256][1024];
int i,j;
color c;
...
```

`c = frame[i][j].g` // for some `i,j`    **IMPLEMENT THIS LINE**

Assuming `i,j,frame`, and `&c` are in `$1`, `$2`, `$3`, and `$10` respectively, write MIPS code to implement the indicated line. Do not overwrite any given registers, and use additional registers beginning at `$4`. More line are provided than are necessary.

Label	Instruction	Comment
	<b>sll \$4, \$1, 10</b>	<b># 1024*i</b>
	<b>add \$4, \$4, \$2</b>	<b># 1024*i + j</b>
	<b>addi \$5, \$0, 3</b>	<b># sizeof(pixel) = 3</b>
	<b>mult \$5, \$4</b>	<b>#  pixel *(1024*i + j)</b>
	<b>mflo \$4</b>	
	<b>addi \$4, \$4, \$3</b>	<b># + frame</b>
	<b>lbu \$5, 1(\$4)</b>	<b># frame[i][j].g</b>
	<b>sb \$5, 0(\$10)</b>	<b># c = frame[i][i].g</b>

Assuming a 32-bit system, consider the following C fragment:

```
int Rotate (int Pattern[]) {
    int i = 0;
    . . .
    Pattern[i] = Pattern[i+1];    // Part C
    . . .
    return (i);                  // Part D
}
```

**Part B** (6 points) Write a MIPS code fragment to implement the beginning of Rotate's implementation: set Rotate's frame pointer and allocate its locals.

Label	Instruction	Comment
	<b>add \$30, \$29, \$0</b>	<b># set FP</b>
	<b>addi \$29, \$29, -4</b>	<b># make room for local i</b>
	<b>sw \$0, 0(\$29)</b>	<b># initialize i</b>

**Part C** (12 points) Suppose the input parameter Pattern is stored in Rotate's activation frame just above the return value slot (pointed to by the frame pointer \$30). Write a MIPS code fragment to implement the line that has the comment "PartC" in the C code above:

Pattern[i] = Pattern[i+1];

*Do not assume that the variable i is already in a register (read it from the stack).*

Label	Instruction	Comment
	<b>lw \$1, -4(\$30)</b>	<b># load i</b>
	<b>lw \$2, 4(\$30)</b>	<b># load Pattern base addr</b>
	<b>sll \$3, \$1, 2</b>	<b># scale i by  int </b>
	<b>add \$4, \$3, \$2</b>	<b># Pattern + i*4</b>
	<b>lw \$5, 4(\$4)</b>	<b># load Pattern[i+1]</b>
	<b>sw \$5, 0(\$4)</b>	<b># store in Pattern[i]</b>

**Part D** (6 points) Write a MIPS code fragment to store the return value of Rotate, deallocate its locals and return to its caller. In other words, implement the line that has the comment “Part D” in the C code above: `return(i);`

*Do not assume that the variable `i` is already in a register (read it from the stack).*

Label	Instruction	Comment
	<b>lw    \$1, -4(\$30)</b>	<b># load i</b>
	<b>sw    \$1, 0(\$30)</b>	<b># store i at RV slot</b>
	<b>addi \$29, \$29, 4</b>	<b># deallocate locals</b>
	<b>jr    \$31</b>	<b># return</b>

**Problem 3** (2 parts, 36 points)**Parameter Passing w/ Activation Frames**

The function Solver (below left) calls function Quad after completing code block 1. Write MIPS code that properly calls Quad. Include all instructions between code block 1 and code block 2. Symbolically label all required stack entries at the point just before control is transferred to Quad and give their values if they are known (below right).

```
int Quad(int x, int *y, int z[])
{
    . . .
}

int Solver(int a) {
    int b;
    int c[] = {4, 2, -1};

    <code block 1>

    c[2] = Quad(a, &b, c);

    <code block 2>
}
```

...	...	...
9604	a	6
Solver's FP 9600	XXX	XXX
9596	b	
9592	c[2]	-1
9588	c[1]	2
SP 9584	c[0]	4
9580	RA	N/A
9576	FP	9600
9572	a	6
9568	&b	9596
9564	c	9584
9560	RV	
9556		
9552		

label	instruction	comment
	<b>addi \$29, \$29, -24</b>	# Allocate activation frame
	<b>sw \$31, 20(\$29)</b>	# Preserve bookkeeping info
	<b>sw \$30, 16(\$29)</b>	
	<b>lw \$1, 4(\$30)</b>	# Push inputs: load a
	<b>sw \$1, 12(\$29)</b>	# store a
	<b>addi \$1, \$30, -4</b>	# compute &b
	<b>sw \$1, 8(\$29)</b>	# store &b
	<b>addi \$1, \$30, -16</b>	# compute c base address
	<b>sw \$1, 4(\$29)</b>	# store c base address
	<b>jal Quad</b>	# Call Quad
	<b>lw \$31, 20(\$29)</b>	# Restore bookkeeping info
	<b>lw \$30, 16(\$29)</b>	
	<b>lw \$1, 0(\$29)</b>	# Read return value
	<b>sw \$1, -8(\$30)</b>	# Store return value in c[2]
	<b>addi \$29, \$29, 24</b>	# Deallocate activation frame