

Instructions: This is a closed book, closed note exam. Calculators are not permitted. If you have a question, raise your hand and I will come to you. Please work the exam in pencil and do not separate the pages of the exam. For maximum credit, show your work.

Good Luck!

Your Name (*please print*) _____

This exam will be conducted according to the Georgia Tech Honor Code. I pledge to neither give nor receive unauthorized assistance on this exam and to abide by all provisions of the Honor Code.

Signed _____

1	2	3	4	5	6	total
20	30	25	20	30	40	165



Problem 1 (20 points)**Optimization**

Perform at least **five** standard compiler optimizations on the following C code fragment by writing the optimized version (in C) to the right. Assume **square**, **f**, and **g** are pure functions that each return an integer with no side effects to other data structures.

```
int square (int n) {  
    return (n*n);}  
  
int f (int m) { ... }  
  
int g (int c, int d) { ... }  
  
int mycode(int a, int b) {  
    int x = 100;  
    int y = 1, z=64;  
    do {  
        if (z)  
            y += x*square(z);  
        else  
            y += g(z+a*x, y+a*x);  
        printf("x:%d, y:%d\n",x,y);  
        x--;  
    } while(x>f(y*1024) + g(a, b));  
    return g(y, x);  
}
```



Briefly describe which standard compiler optimizations you applied and how they improve storage and/or execution efficiency in this code example (be specific; e.g., “replaces 2 MIPS operations with 1 operation on every loop iteration”).

- 1.
- 2.
- 3.
- 4.
- 5.

Problem 2 (2 parts, 30 points)**Linked Lists**

Suppose we have the following definition which is used to create singly linked lists.

```
typedef struct Link {
    int      ID;
    int      Value;
    struct Link *Next;
} Link;
```

Part A (6 points) Complete the following subroutine which inserts a `Link` (pointed to by the input parameter `NewLink`) into the list just after the `Link` pointed to by the input parameter `Before`. You may assume that neither input parameter is `NULL`. `Before`'s `Next` field may point to another `Link` or it may be `NULL`. `NewLink`'s `Next` field is `NULL`.

```
void SspliceIn(Link *NewLink, Link *Before){
    _____; /* part A*/
    _____; /* part A*/
}
```

Part B Complete the following recursive subroutine which takes a pointer to the head of a linked list and returns a pointer to a copy of the linked list. Follow the steps specified below.

```
Link * CopyList(Link *Head) {
    if (Head == NULL) return _____; /* part B.1 */
    _____; /* part B.2 */
    _____; /* part B.3 */
    if ( _____ ) { /* part B.4 */
        printf("Error: Insufficient space.");
        exit(1);
    }
    _____; /* part B.5 */
    _____; /* part B.5 */
    _____; /* part B.6 */
    return LinkCopy;
}
```

Part B.1 (3 points) Fill in what should be returned if the list is empty.

Part B.2 (3 points) Add a local variable called `LinkCopy` that is a pointer to a `Link` object.

Part B.3 (5 points) Allocate space for a `Link` structure using `malloc` and make `LinkCopy` point to the object allocated. Be sure to include appropriate type casting to avoid type errors.

Part B.4 (3 points) Fill in the test for whether `malloc` found enough space which controls the print statement.

Part B.5 (5 points) Copy the values of `Head`'s `ID` and `Value` fields to `LinkCopy`.

Part B.6 (5 points) Call `CopyList` recursively to copy the rest of the list and assign the result to `LinkCopy`'s `Next` field.

Problem 3 (3 parts, 25 points)**Associative Sets**

Consider a hash table that is implemented using the following struct definitions.

```
#define NUMBUCKETS 5
typedef struct Entry {
    int      Key;
    int      Value;
    struct Entry *Next;
} Entry;

typedef struct {
    Entry      *Buckets[NUMBUCKETS];
    int        Size;
} HashTable;
```

Part A (6 points) What is the value of each of these (assume a 32 bit system):

sizeof(Entry) = _____	sizeof(HashTable) = _____
-----------------------	---------------------------

Part B (10 points) Suppose the entries are maintained in a *sorted* linked in each bucket in order from small to large keys. Complete the C function `Find_Key` that *efficiently* searches the hash table for an entry corresponding to a specified key (i.e., *it should end the search as early as possible*). It should return a pointer to the matching `Entry` if `Key` is found or return `NULL` if `Key` is not found in the hash table.

```
Entry *Find_Key(HashTable *HT, int Key) {
    Entry *ThisEntry;
    int    Index;
    int    Hash(int Key); /* function prototype for hash function */
```

```
}
```

Part C (9 points) Suppose a hash table created using the structs above contains **155** entries total and the entries are evenly distributed across the **5** hash table buckets, each implemented as a *sorted* linked list of `Entry` structs. An application performs **500** lookups of various keys: **375** of the lookups find the key in the list and **125** lookups fail to find the key. The keys that are found are distributed throughout the list so that each position is equally likely to be where a key is found. What is the average number of key comparisons that would be needed for a lookup in this list implementation? (Show work. Note: you may not have to use all data provided.)

Average number of comparisons: _____

Problem 4 (3 parts, 20 points)**Garbage Collection**

Below is a snapshot of heap storage. Values that are pointers are denoted with a “\$”. The heap pointer is **\$6188**. The heap has been allocated contiguously beginning at **\$6000**, with no gaps between objects.

addr	value	addr	value	addr	value	addr	value	addr	value	addr	value
6000	8	6032	12	6064	20	6096	16	6128	12	6160	\$6004
6004	33	6036	28	6068	4	6100	\$6052	6132	\$6172	6164	0
6008	0	6040	\$6120	6072	\$6132	6104	6010	6136	\$6016	6168	16
6012	16	6044	80	6076	8	6108	5	6140	72	6172	\$6052
6016	\$6100	6048	16	6080	\$6148	6112	148	6144	20	6176	0
6020	\$6172	6052	0	6084	\$6172	6116	8	6148	6046	6180	\$6004
6024	25	6056	100	6088	4	6120	32	6152	\$6080	6184	0
6028	30	6060	0	6092	\$6080	6124	\$6080	6156	26	6188	0

Part A (9 points) Suppose the stack holds a local variable whose value is the memory address **\$6120** and register \$3 holds memory address **\$6016**. No other registers or static variables currently hold heap memory addresses. List the addresses of all objects in the heap that would be marked by a mark-and-sweep garbage collection algorithm.

Addresses of Marked Objects: _____

Part B (3 points) If a reference counting garbage collection strategy is being used, what would be the reference count of the object at address **\$6172**?

Reference count of object at \$6172 = _____

Part C (8 points) If the local variable whose value is the address **\$6120** is popped from the stack, which addresses will be reclaimed by each of the following strategies? If none, write “none.”

Reference Counting:	
Mark and Sweep:	
Old-New Space (copying):	

Problem 5 (3 parts, 30 points)**MIPS and C programming**

Part A (8 points) Suppose the instruction "jal Foo" is executed which changes the values of the following registers to:

Register	Value
\$31	3216
PC	3620

What is the address of the first instruction of the subroutine Foo and what is the address of the jal Foo instruction?

Subroutine Foo starts at address: _____

Address of jal Foo instruction: _____

Part B (12 points) Suppose variables A, B, and C are of type int and are stored in registers \$1, \$2, and \$3. Write a MIPS code fragment that computes $C = A * \min(A, B)$. Use only registers \$0, \$1, \$2, and \$3 and for maximum credit, use a minimal number of instructions and include comments.

Label	Instruction	Comment

Part C (10 points) What does the following code fragment print?

```
int A[10] = {1990, 3, 1992, 5, 1999, 1, 2001, 3, 2005, 1};
int B[10] = {1991, 2, 1993, 1, 1998, 7, 2002, 1, 2007, 6};
int i, j;
for(i=1; i<10; i=i+2)
{
    for (j=1; j<10; j=j+2)
    {
        if (A[i] == B[j])
        {
            printf("C: %d, As: %d, Bs: %d.\n", A[i], A[i-1], B[j-1]);
            break;
        }
    }
}
```

--

Problem 6 (40 points)

Activation Frames

The function `Bar` (below left) calls function `Foo` after completing code block 1. Write MIPS code that properly calls `Foo`. Include all instructions between code block 1 and code block 2. **Note that code block 1 may change the values of the local variables** (e.g., assume `i` can be any value from 0 to 2). Symbolically label all required stack entries and give their initial values if known (below right).

```
int Bar() {
    int    A[] = {9, 25, 27};
    int    i = 0;
    int    y = 5;

    (code block 1)

    A[1] = Foo(A, &y, A[i]);

    (code block 2)
}
```

Bar's FP	9900	XXX	XXX
	9896	A[2]	27
	9892	A[1]	25
	9888	A[0]	9
	9884	i	0
SP	9880	y	5
	9876		
	9872		
	9868		
	9864		
	9860		
	9856		

label	instruction	comment
		# allocate activation frame
		# preserve bookkeeping info
		# compute A
		# push A
		# compute &y
		# push &y
		# compute A[i] by loading i,
		# scaling it, and adding it to
		# the base address of A
		# push A[i]
	jal Foo	# call Foo
		# restore bookkeeping info
		# read return value
		# store return value in A[2]

MIPS Instruction Set (core)

<i>instruction</i>	<i>example</i>	<i>meaning</i>
arithmetic		
add	add \$1,\$2,\$3	$\$1 = \$2 + \$3$
subtract	sub \$1,\$2,\$3	$\$1 = \$2 - \$3$
add immediate	addi \$1,\$2,100	$\$1 = \$2 + 100$
add unsigned	addu \$1,\$2,\$3	$\$1 = \$2 + \$3$
subtract unsigned	subu \$1,\$2,\$3	$\$1 = \$2 - \$3$
add immediate unsigned	addiu \$1,\$2,100	$\$1 = \$2 + 100$
set if less than	slt \$1, \$2, \$3	if ($\$2 < \3), $\$1 = 1$ else $\$1 = 0$
set if less than immediate	slti \$1, \$2, 100	if ($\$2 < 100$), $\$1 = 1$ else $\$1 = 0$
set if less than unsigned	sltu \$1, \$2, \$3	if ($\$2 < \3), $\$1 = 1$ else $\$1 = 0$
set if < immediate unsigned	sltui \$1, \$2, 100	if ($\$2 < 100$), $\$1 = 1$ else $\$1 = 0$
multiply	mult \$2,\$3	Hi, Lo = $\$2 * \3 , 64-bit signed product
multiply unsigned	multu \$2,\$3	Hi, Lo = $\$2 * \3 , 64-bit unsigned product
divide	div \$2,\$3	Lo = $\$2 / \3 , Hi = $\$2 \bmod \3
divide unsigned	divu \$2,\$3	Lo = $\$2 / \3 , Hi = $\$2 \bmod \3 , unsigned
transfer		
move from Hi	mfhi \$1	$\$1 = \text{Hi}$
move from Lo	mflo \$1	$\$1 = \text{Lo}$
load upper immediate	lui \$1,100	$\$1 = 100 \times 2^{16}$
logic		
and	and \$1,\$2,\$3	$\$1 = \$2 \& \$3$
or	or \$1,\$2,\$3	$\$1 = \$2 \mid \$3$
and immediate	andi \$1,\$2,100	$\$1 = \$2 \& 100$
or immediate	ori \$1,\$2,100	$\$1 = \$2 \mid 100$
nor	nor \$1,\$2,\$3	$\$1 = \text{not}(\$2 \mid \$3)$
xor	xor \$1, \$2, \$3	$\$1 = \$2 \oplus \$3$
xor immediate	xori \$1, \$2, 255	$\$1 = \$2 \oplus 255$
shift		
shift left logical	sll \$1,\$2,5	$\$1 = \$2 \ll 5$ (logical)
shift left logical variable	sllv \$1,\$2,\$3	$\$1 = \$2 \ll \$3$ (logical), variable shift amt
shift right logical	srl \$1,\$2,5	$\$1 = \$2 \gg 5$ (logical)
shift right logical variable	srlv \$1,\$2,\$3	$\$1 = \$2 \gg \$3$ (logical), variable shift amt
shift right arithmetic	sra \$1,\$2,5	$\$1 = \$2 \gg 5$ (arithmetic)
shift right arithmetic variable	srav \$1,\$2,\$3	$\$1 = \$2 \gg \$3$ (arithmetic), variable shift amt
memory		
load word	lw \$1, 1000(\$2)	$\$1 = \text{memory} [\$2+1000]$
store word	sw \$1, 1000(\$2)	$\text{memory} [\$2+1000] = \1
load byte	lb \$1, 1002(\$2)	$\$1 = \text{memory} [\$2+1002]$ in least sig. byte
load byte unsigned	lbu \$1, 1002(\$2)	$\$1 = \text{memory} [\$2+1002]$ in least sig. byte
store byte	sb \$1, 1002(\$2)	$\text{memory} [\$2+1002] = \1 (byte modified only)
branch		
branch if equal	beq \$1,\$2,100	if ($\$1 = \2), $\text{PC} = \text{PC} + 4 + (100*4)$
branch if not equal	bne \$1,\$2,100	if ($\$1 \neq \2), $\text{PC} = \text{PC} + 4 + (100*4)$
jump		
jump	j 10000	$\text{PC} = 10000*4$
jump register	jr \$31	$\text{PC} = \$31$
jump and link	jal 10000	$\$31 = \text{PC} + 4$; $\text{PC} = 10000*4$