

Lab2_Part2 & Chapter3 Review

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Lab2_part2 Hints

UTF-8 Encoding Rules:

Bytes	Bit Pattern	Unicode Range
1 byte	0xxxxxxx	U+0000 to U+007F
2 bytes	110xxxxx 10xxxxxx	U+0080 to U+07FF
3 bytes	1110xxxx 10xxxxxx 10xxxxxx	U+0800 to U+FFFF
4 bytes	11110xxx 10xxxxxx 10xxxxxx 10xxxxxx	U+10000 to U+10FFFF

```
Smaller than Bound1  
JMP HANDLE_1BYTE
```

```
Smaller than Bound2  
JMP HANDLE_2BYTE
```

```
Smaller than Bound3  
JMP HANDLE_3BYTE
```

```
.....
```

```
JMP HANDLE_ERROR
```

Example: Encoding U+0800 (0x0800) to UTF-8

Step 1: Determine Encoding Size

- - U+0800 falls in the range `0x0800 - 0xFFFF` → 3-byte encoding

Step 2: Binary Representation of Code Point

- - U+0800 (hex) = `0000 1000 0000 0000` (binary)

Step 3: UTF-8 Encoding (3 Bytes)

First byte: `1110xxxx`

- Extract top 4 bits: `0000` → `11100000` (`0xE0`)

Second byte: `10xxxxxx`

- Next 6 bits: `100000` → `10100000` (`0xA0`)

Third byte: `10xxxxxx`

- Last 6 bits: `000000` → `10000000` (`0x80`)

Final UTF-8 Encoding:

- Binary: `11100000 10100000 10000000`
- Hexadecimal: `0xE0 A0 80`

1. Assume 185 and 122 are signed 8-bit decimal integers stored in sign-magnitude format. Calculate $185 + 122$. Is there overflow, underflow, or neither?

1. Understanding Sign-Magnitude Format

In 8-bit sign-magnitude format, an integer is represented as:

- 1 bit for the sign (Most Significant Bit, MSB)
 - 0 indicates a **positive** number.
 - 1 indicates a **negative** number.
- 7 bits for the magnitude, allowing values from 0 to 127.

2. Range of Representable Values

In 8-bit sign-magnitude format, the representable range is:

- Positive Range: $+0$ to $+127$
- Negative Range: -0 to -127

1. Interpretation of 185 in Sign-Magnitude Format

Sign-Magnitude Decoding:

- Binary of 185: 10111001
- Sign Bit: 1, indicating a **negative** number.
- Magnitude: 57, derived from the binary 0111001.

Final Interpretation:

185 in sign-magnitude format = -57

2. Interpretation of 122 in Sign-Magnitude Format

- Binary of 122: 01111010
- Sign Bit: 0, indicating a **positive** number.
- Magnitude: 122, directly from the binary representation.

Final Interpretation:

122 in sign-magnitude format = $+122$

Neither Overflow nor Underflow

3. Perform the Addition

$$-57 + 122 = 65$$

2. Assume 185 and 122 are signed 8-bit decimal integers stored in sign-magnitude format. Calculate $185 - 122$. Is there overflow, underflow, or neither?

Mathematical Operation:

$$-57 - 122 = -179$$

Result Validation:

- The result -179 is outside the representable range for 8-bit sign-magnitude integers, as it is less than -127.

Conclusion:

The operation results in **underflow**.

4. Assuming single precision IEEE 754 format, what decimal number is represent by this word:

1 01111101 001000000000000000000000

The IEEE 754 single-precision format consists of:

- 1 bit for the sign.
- 8 bits for the exponent.
- 23 bits for the mantissa (fraction).

1. Determine the Sign Bit

- Sign Bit: 1, indicating the number is negative.

$$\text{Sign} = (-1)^1 = -1$$

2. Decode the Exponent

- Exponent Bits: 01111101

Convert to Decimal:

$$01111101_2 = 125_{10}$$

Apply the Bias:

- IEEE 754 single-precision uses a **bias of 127**:

$$\text{Exponent} = 125 - 127 = -2$$

3. Decode the Mantissa (Fraction)

- Mantissa Bits: 00100000000000000000000

Constructing the Fraction:

The mantissa has an implicit leading 1, so:

$$1.0010000000000000000000_2$$

Convert to Decimal:

$$1 + 2^{-3} = 1 + 0.125 = 1.125$$

4. Calculate the Final Value

$$\text{Value} = \text{Sign} \times \text{Mantissa} \times 2^{\text{Exponent}}$$

$$\text{Value} = -1 \times 1.125 \times 2^{-2}$$

$$2^{-2} = \frac{1}{4} = 0.25$$

$$\text{Value} = -1.125 \times 0.25 = -0.28125$$

5. The floating-point format to be used in this problem is an 8-bit IEEE 754 normalized format with 1 sign bit, 4 exponent bits, and 3 mantissa bits. It is identical to the 32-bit and 64-bit formats in terms of the meaning of fields and special encodings. The exponent field employs an excess-7 coding. The bit fields in a number are (sign, exponent, mantissa). Assume that we use unbiased rounding to the nearest even specified in the IEEE floating point standard.

Field	Bits	Description
Sign	1	0 for positive, 1 for negative
Exponent	4	Excess-7 encoding (bias = 7)
Mantissa	3	Normalized form with an implicit leading 1

a) Encode the following numbers to the 8 bit IEEE format

1) $0.0011011_{\text{binary}}$

Normalize the Binary Number

$$0.0011011 = 1.1011 \times 2^{-3}$$

Components of the Normalized Form:

- Sign Bit: 0 (since the number is positive)
- Mantissa: The 3 bits after the leading 1:
 - $1.1011 \rightarrow 101$
- Exponent: -3 (due to shifting the binary point 3 places to the right)

Final 8-bit Representation:

0 0100 101

Calculate the Biased Exponent:

$$\text{Biased Exponent} = -3 + 7 = 4$$

Convert 4 to 4-bit Binary:

$$4 = 0100_2$$

i) Decode the following 8-bit IEEE number into their decimal value: 1 1010 101

- Binary Exponent: 1010
- Convert to Decimal:

$$1010_2 = 10_{10}$$

- Apply Excess-7 Bias:

$$\text{Exponent} = 10 - 7 = 3$$

$$\text{Value} = \text{Sign} \times \text{Mantissa} \times 2^{\text{Exponent}}$$

$$\text{Value} = -1 \times 1.625 \times 2^3$$

$$2^3 = 8$$

$$\text{Value} = -1.625 \times 8 = -13.0$$

- The IEEE 754 format uses an implicit leading 1, giving us:

$$1.101$$

Convert to Decimal:

$$1 + 0.5 + 0.125 = 1.625$$

j) Decide which number in the following pairs are greater in value (the numbers are in 8-bit IEEE 754 format):

1) 0 1000 100 and 0 1000 111

The second number is greater in value

b) Perform the computation $1.011_{\text{binary}} + 0.0011011_{\text{binary}}$ showing the correct state of the guard, round bits and sticky bits. There are three mantissa bits.

1. What are Guard, Round, and Sticky Bits?

Guard Bit (G):

- The first bit beyond the mantissa bits.
- Helps in rounding decisions.

Round Bit (R):

- The second bit beyond the mantissa bits.
- Assists in rounding by indicating whether the truncated portion is closer to 0 or 1.

Sticky Bit (S):

- Represents whether any bits beyond the round bit are non-zero.
- If any bits shifted out during normalization are 1, the sticky bit is set to 1, otherwise, it remains 0.
- Captures information about whether the discarded bits are non-zero, contributing to correct rounding.

2. Problem Recap: Adding $1.011_2 + 0.0011011_2$

1. Normalize the Numbers:

Binary	Normalized Form	Exponent	Mantissa
1.011_2	1.011×2^0	0	011
0.0011011_2	1.1011×2^{-3}	-3	101

2. Align the Exponents:

The higher exponent is 0, so shift the second number right by 3 places to align:

$$1.1011 \times 2^{-3} = 0.0011011 \times 2^0$$

$$0.0011011 \times 2^0$$

3. Bit Representation (Including Guard, Round, and Sticky Bits)

After Shifting to Align Exponents:

Mantissa (3 bits) = 001

Guard Bit (G) = 1

Round Bit (R) = 0

Sticky Bit (S) = 1

Binary Addition:

$$1.011 + 0.001 = 1.100$$

Bit Values After Addition:

Bit Type	Bit Value
Mantissa	100
Guard Bit (G)	1
Round Bit (R)	0
Sticky Bit (S)	1

$$1.101_2 = 1.625$$

G R S	Action	Easy Rule
1 1 X	Round Up	If round bit is 1, always round up
1 0 1	Round Up	Sticky bit 1? Round up!
1 0 0	Round to Nearest Even	Mantissa odd? Round up
0 X X	Do Not Round Up	Guard bit 0 means no rounding

6. Using 32-bit IEEE 754 single precision floating point with one(1) sign bit, eight (8) exponent bits and twenty three (23) mantissa bits, show the representation of $-11/16$ (-0.6875).

Field	Bits	Description
Sign Bit	1	0 for positive, 1 for negative
Exponent	8	Excess-127 encoding (bias = 127)
Mantissa	23	Normalized with an implicit leading 1

Fractional Conversion:

$$-0.6875 = -\left(\frac{11}{16}\right)$$

Convert the fractional part to binary:

$$0.6875 \times 2 = 1.375 \rightarrow 1$$

$$0.375 \times 2 = 0.75 \rightarrow 0$$

$$0.75 \times 2 = 1.5 \rightarrow 1$$

$$0.5 \times 2 = 1.0 \rightarrow 1$$

Combine the bits:

$$0.6875_{10} = 0.1011_2$$

Normalized Binary Form:

$$0.1011 = 1.011 \times 2^{-1}$$

- Mantissa: 011
- Exponent: -1

The exponent in excess-127 format:

$$\text{Biased Exponent} = -1 + 127 = 126$$

Convert 126 to 8-bit binary:

$$126 = 01111110_2$$

The normalized mantissa is 1.011, but we store only the fractional part:

$$\text{Mantissa} = 01100000000000000000000$$

The remaining bits are padded with zeros to complete 23 bits.

$$\text{Sign Bit} = 1$$

$$\text{Exponent} = 01111110$$

$$\text{Mantissa} = 01100000000000000000000$$

Combined Binary Representation:

$$1 \ 01111110 \ 01100000000000000000000$$

Feedback Survey

