

PEMBAHASAN UTS ASA 2017

```
II) Function SamaMatriks (A, B : matriks; n : integer) -> boolean {
    I, j : integer
    For I = 1 to n {
        For j = 1 to n {
            If A[i,j] != B[i,j] {
                Retrurn false
            }
        }
    }
    Return true
}
```

II) Algoritma pwncarian beruntun

```
Procedure PencarianBeruntun (input a1, a2, ..., an : integer, x : integer,
output idx : integer) {
    I : integer
    Ketemu : boolean
    I = 1
    Ketemu = false
    While (I <= n) and (!ketemu) {
        If (a[i] = x) {
            ketemu = true
        } else {
            i = I++
        }
    }
    If ketemu {
        Idx = 1
    } else {
        Idx = 0
    }
}
```

SOAL 1

- A. I) terbaik : matriks berukuran 1×1 , terburuk : semua elemen matriks
A = elemen matriks B
II) terbaik : nilai a1 adalah x, terburuk : tidak ada nilai $a[i]=x$
- B. I) terbaik : 2, terburuk : n^2
II) terbaik : 4, terburuk : $n + 4$
- C. I) kompleksitas polinomial, ii) kompleksitas linear

SOAL 2

- A. $T(n) = 2n + 120 = O(n)$
 $n \geq 1, 2n + 120 < 2n + 120n = 122n$
 $C = 122, n_0 = 1$
- B. $T(n) = 3n^3 + 6n^2 + n + 8 = O(n^3)$
 $n \geq 1, T(n) = 3n^3 + 6n^2 + n + 8 < 3n^2 + 6n^3 + 9n^3 = 18n^3$
 $C = 18, n_0 = 1$

- C. $T(n) = \log n^5 = O(\log n)$
 $n \geq 1, 5 \log n \leq 6 \log n$
 $C = 6, n_0 = 1$
- D. $T(n) = 6 \cdot 2^n + n^2 = O(2^n)$
 $n \geq 1, 6 \cdot 2^n + n^2 \leq 6 \cdot 2^n + 7 \cdot 2^n = 13 \cdot 2^n$
 $C = 13, n_0 = 1$
- E. $T(n) = 10 \log 3^n = O(n)$
 $n \geq 1, n \cdot 10 \log 3 \leq 4n$
 $C = 4, n_0 = 1$

SOAL 3

- a. $T(n) = 5n^3 + 6n^2 \log n$
 $n \geq 1, 5n^3 + 6n^2 \log n \leq 5n^3 + 6n^3 = 11n^3$
 $C = 11, n_0 = 1$
- $T(n) = O(n^3)$
 $n \geq 1, 5n^3 + 6n^2 \log n \geq 5n^3, C = 5$
 $T(n) = \Omega(n^3)$
 $T(n) = \Theta(n^3)$
- b. $T(n) = 3n^4 + 6n^3 + 18n + 2$
 $n \geq 1, 3n^4 + 6n^3 + 18n + 2 \leq 3n^4 + 6n^4 + 10n^4 + 2n^4 = 20n^4$
 $C = 20, n_0 = 1$

SOAL 4

- A. Relasi Rekurens :
- $T(n) = 0, n = 0$
 $T(n) = T(n-1) + 1, n > 0$
- B. Kompleksitas waktu algoritma
- $T(n) = T(n-1) + 1 \rightarrow T(n-1) = T(n-2) + 1$
 $= [T(n-2) + 1] + 1 = T(n-2) + 2 \rightarrow T(n-2) = T(n-3) + 1$
 \dots
 \dots
 $= T(0) + n$
- $\rightarrow T(n) + n = O(n)$