

Binary Search Tree

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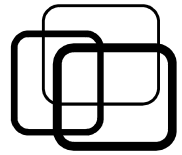
- Concepts.
- Implementation.

Contents



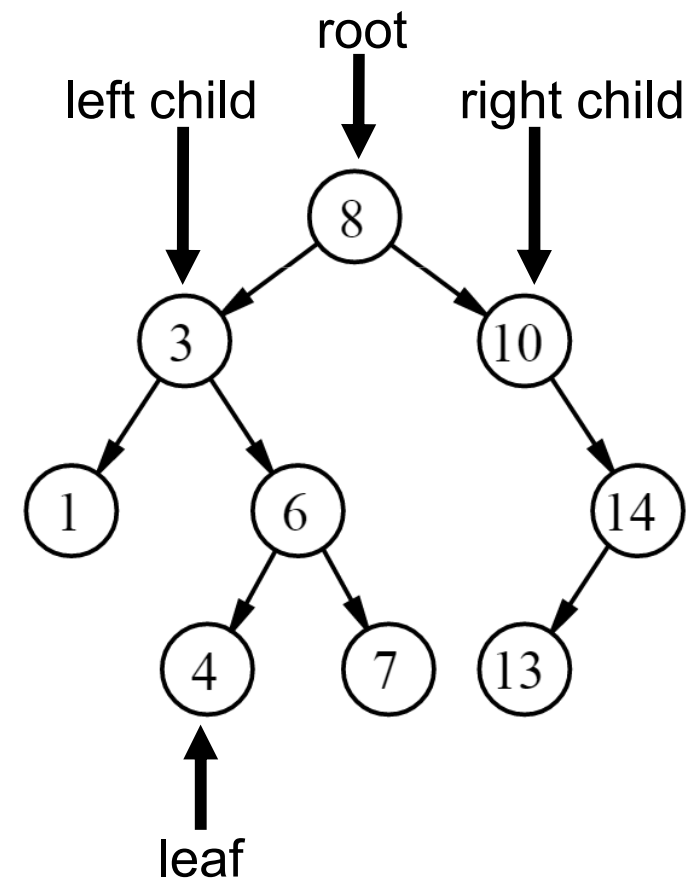
- **Concepts.**
- **Implementation.**

Concepts



■ Binary search tree:

- A binary linked data structure.
- Linear linked vs. binary linked.
- Each node has:
 - At most one parent.
 - At most two children.
 - **All left < node <= all right.**
- Root: node has no parent.
- Leaf: node has no children.



Concepts

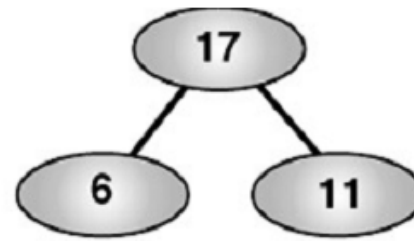


■ Example:

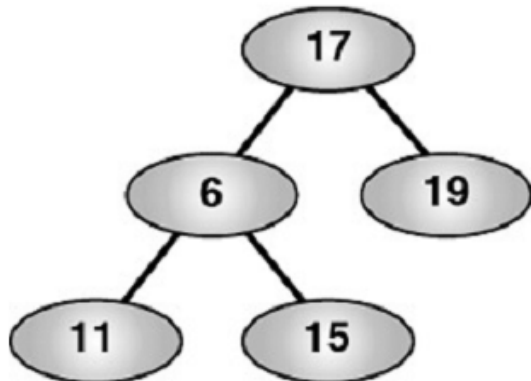
■ Which tree is binary search tree?



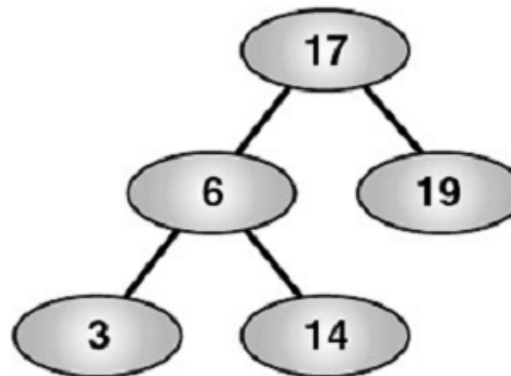
(a)



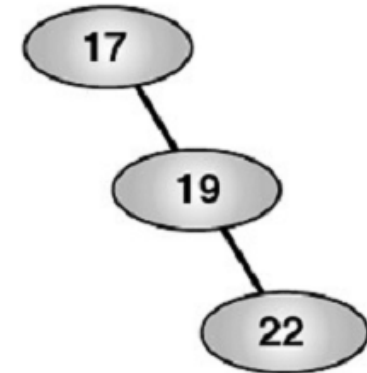
(b)



(c)



(d)

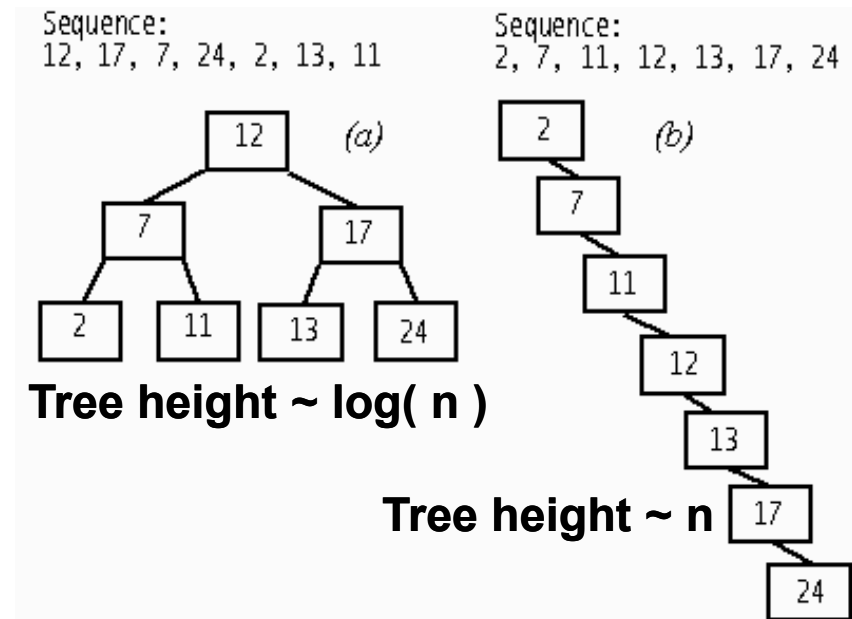
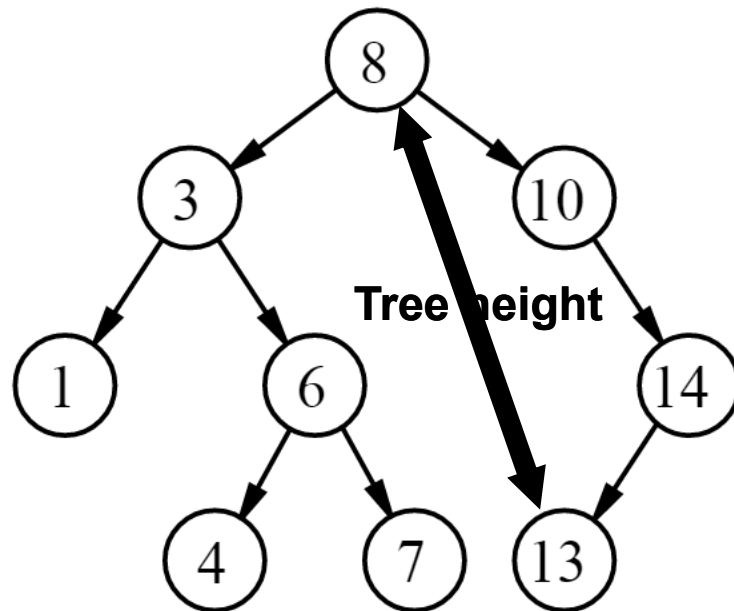


(e)



■ Tree height:

- Height of node: longest path from node to leaf.
- Tree height = height of root.
- Important factor of performance:
 - For a n -node tree: $\log(n) \leq \text{tree height} \leq n$.



Contents



- Concepts.
- **Implementation.**

Implementation

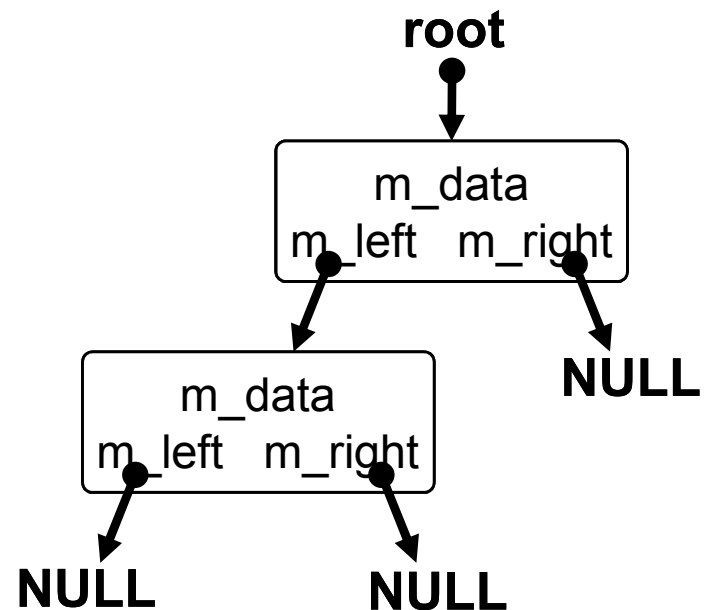


■ ADT Binary search tree:

■ Values:

```
struct TreeNode {  
    int      m_data;  
    TreeNode *m_left;  
    TreeNode *m_right;  
};
```

```
class BSTree {  
private:  
    TreeNode* m_root;  
};
```



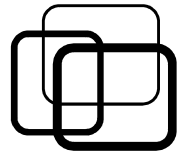


■ ADT Binary search tree:

■ Operations:

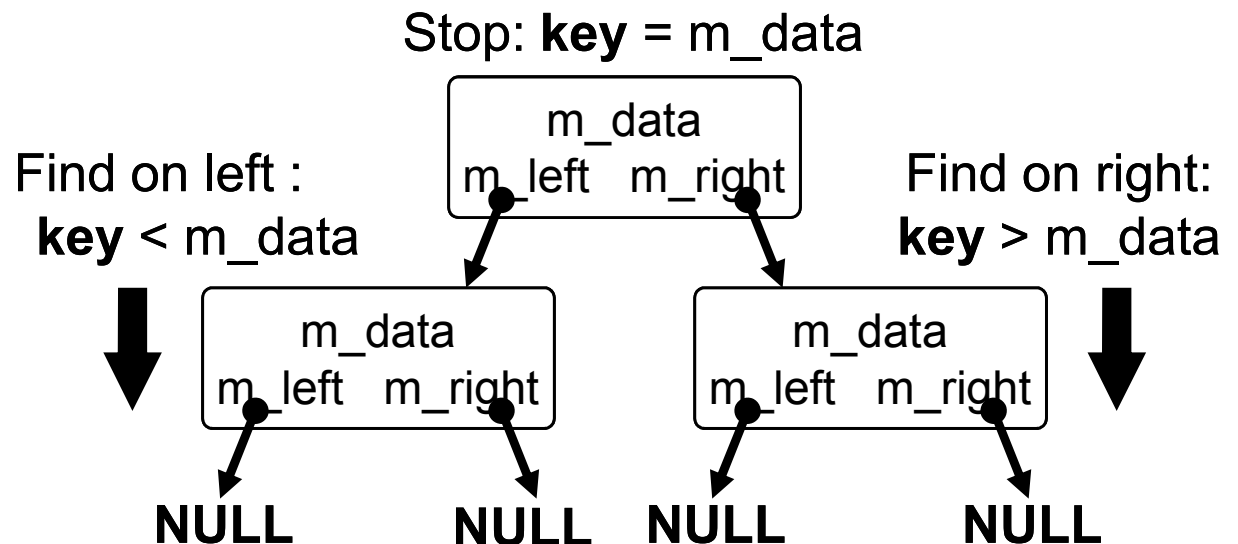
- Initialization.
- Check empty.
- Find a key.
- Add a key.
- Remove a key.
- Visit.

Implementation

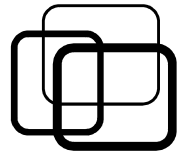


■ Find a key:

```
Find key at a node {  
    if ( key < m_data )  
        Find key on left child;  
    if ( key > m_data )  
        Find key on right child;  
  
    return current node;  
}
```



Implementation



■ Add a key (keep tree condition):

Add key at a node {

if **NULL** node

Make node with key

Stop

if (**key** < m_data)

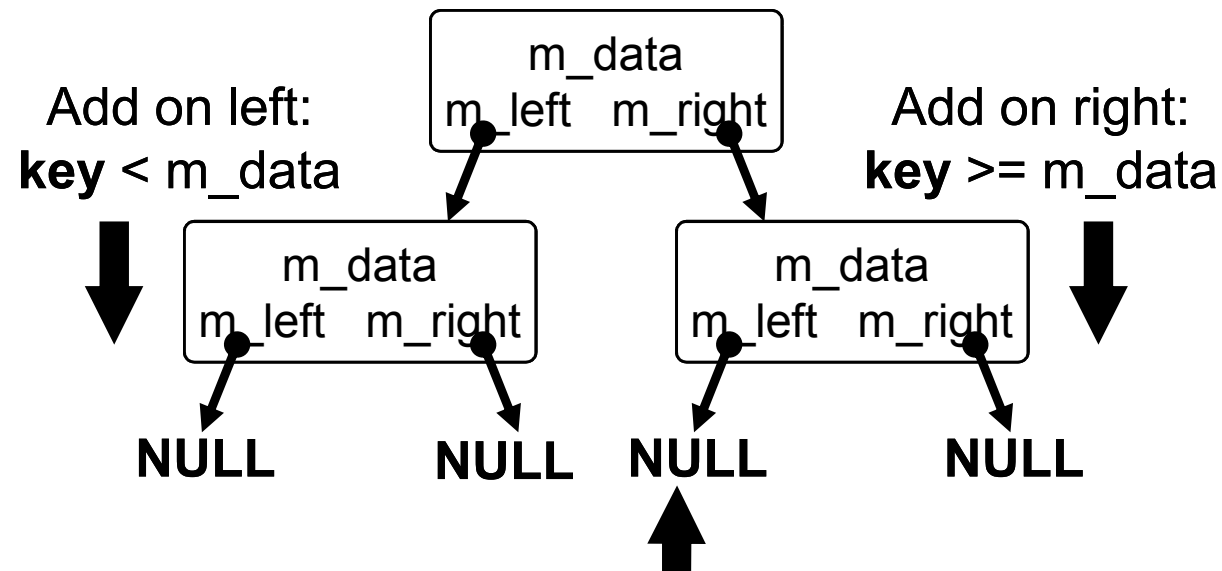
Add key on left child

else

Add key on right child

}

key is added at NULL child!
Which NULL?

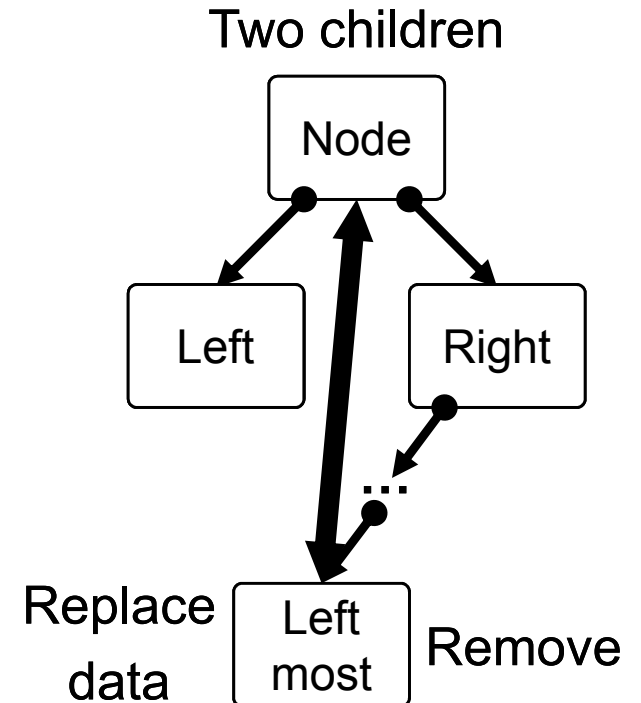
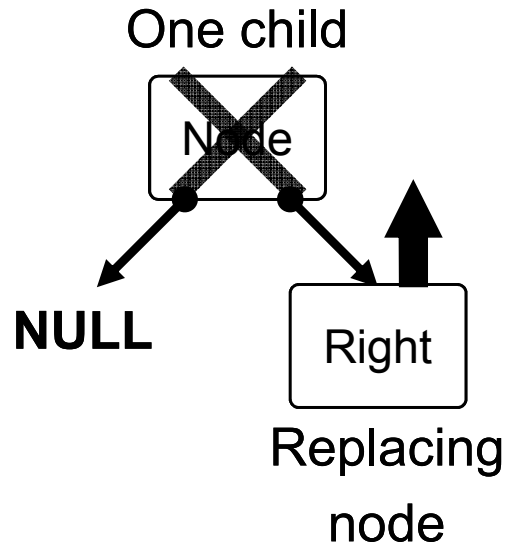
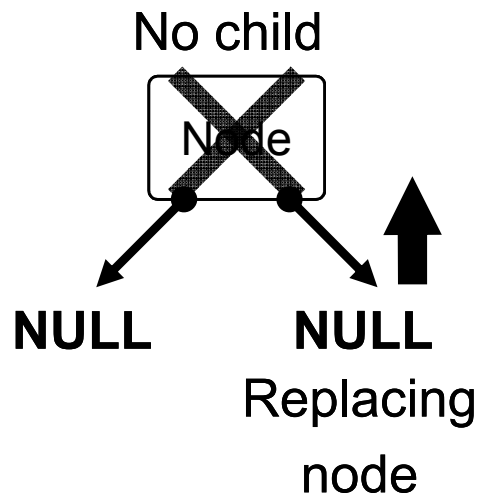


Implementation



■ Remove a key (keep tree condition):

Remove key at a node {
Find key at node
if found
Replace node
}





■ Remove a key (keep tree condition):

```
Replace a node {  
    if ( no left child )  
        replace = m_right;  
        delete current node;  
    else if ( no right child )  
        replace = m_left;  
        delete current node;  
    else  
        replace = current node;  
        leftMost = find left most on right;  
        m_data = leftMost->m_data;  
        Remove leftMost->m_data at leftMost;  
  
    current node = replace;  
}
```

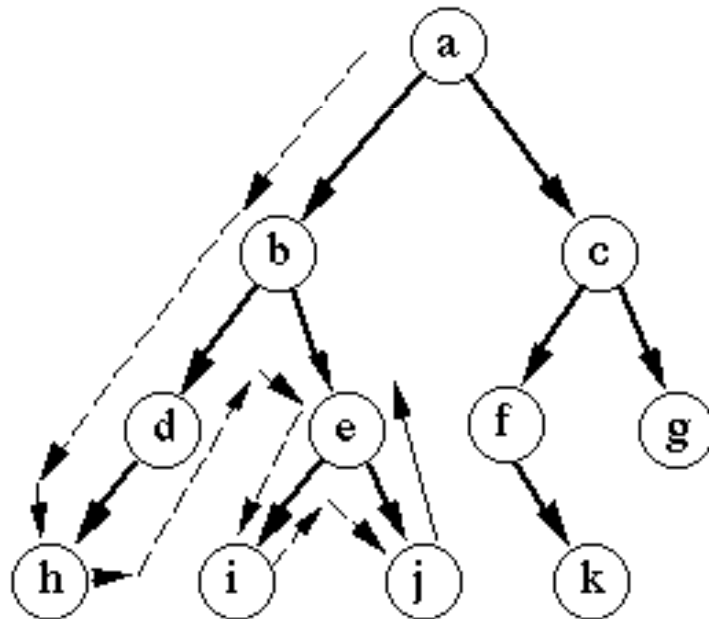
Implementation



■ Visit a node:

Depth-first visit:

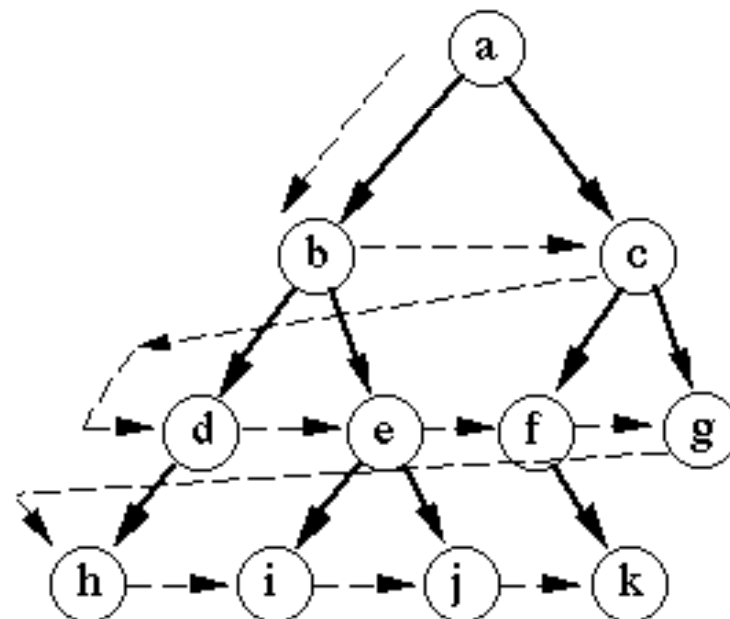
- Visit left branch first.
- Then visit right branch.



Depth-first search

Breadth-first visit:

- Visit 1st generation children first.
- Then visit 2nd generation children.



Breadth-first search

Implementation



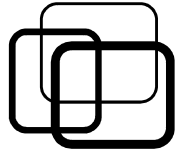
■ Visit a node:

```
Visit depth NLR {  
    Do something on node...  
    Visit depth left child  
    Visit depth right child  
}
```

```
Visit depth LNR {  
    Visit depth left child  
    Do something on node.  
    Visit depth right child  
}
```

```
Visit depth LRN {  
    Visit depth left child  
    Visit depth right child  
    Do something on node...  
}
```

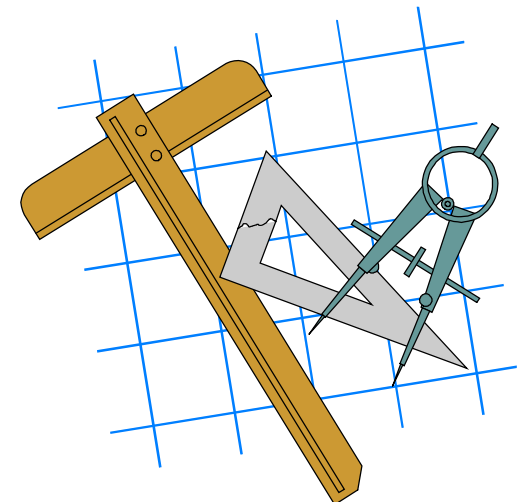
```
Visit breadth {  
    push node to QUEUE;  
    loop QUEUE is not empty {  
        node = pop QUEUE;  
  
        Do something on node...  
  
        if ( node has left child )  
            push left child to QUEUE;  
        if ( node has right child )  
            push right child to QUEUE;  
    }  
}
```



■ Practice 8.1:

Construct class **BSTree** has the following methods:

- Initialize.
- Check empty.
- Count nodes.
- Measure height.
- Find key.
- Add key.
- Remove key.
- Visit depth NLR.
- Visit breadth.





■ Practice 8.2:

Provide class **BSTree** with the following methods:

- Construct tree from an array of integers.
- Export tree to an array depth first NLR.
- Export tree to an array breadth first.

Example:

```
int a [ ] = { 5, 2, 4, 1, 3 };  
int size = 5;  
BSTree t( a, size );  
t.exportDepthFirst( a, size );  
t.exportBreadthFirst( a, size );
```

