Linked List

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Contents



- Linked List Concepts.
- Linked List Operations.
- Linked List Improvement.

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Linked List Concepts



Limitations of Array:

■ Characteristics:

Continuous memory storage.

Advantages:

- Random element access by indexing.
- > Very efficient for fix-sized storage.

■ Disadvantages:

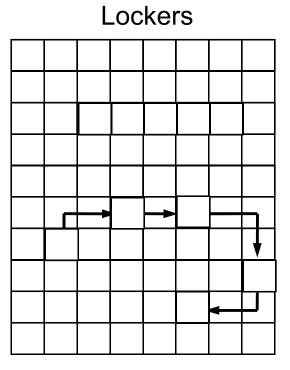
- > Resize array requires memory reallocation.
- > Add or remove element need element shifted.
- Allocate large continuous memory is difficult.

Linked List Concepts



Linked list solution:

- Locker renting problem:
 - > Need to rent lockers to store N items.
 - > Each locker keeps only 1 item.
- Continuous solution → Array.
- Dis-continuous solution:
 - > Rent N arbitrary lockers.
 - Each locker keeps:
 - > 1 items.
 - > Address of next locker.
 - Only keep address of first locker.



Linked List Concepts



Singly linked list:

- A discontinuous data structure.
- Element = Data + Link to next element.
- Last element link to NULL.
- Head: link to first element.

```
head data next NULL
```

■ C declaration:

```
struct SNode {
    int data; SNode *head;
    SNode *next; };
};
```

Contents

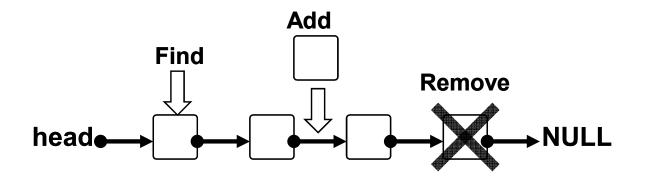


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Operations on linked list:

- Initialize.
- Check empty.
- Find element.
- Add element.
- Remove element.





- Initialize list:
 - At first, list has no element.

head NULL

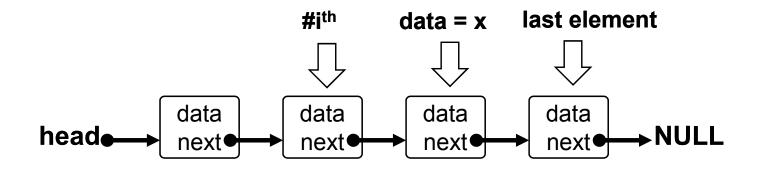
- Check empty list:
 - An empty list has no element.

head → NULL ???



■ Find element:

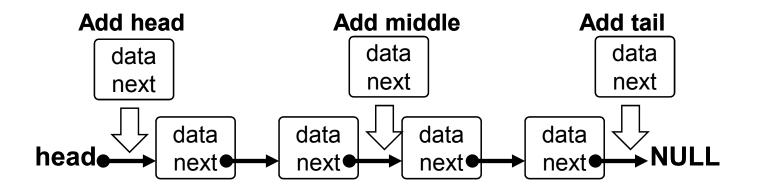
- Find #ith element.
- Find element has data x.
- Find last element.





Add element:

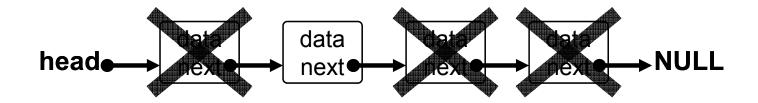
- Add head.
- Add tail.
- Add middle:
 - > After #ith element.
 - > Keep order (ascending).





Remove element:

- Remove head.
- Remove tail.
- Remove middle:
 - > #ith element.
 - > Element has data x.
- Remove all.



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Singly linked list vs. Dynamic array:

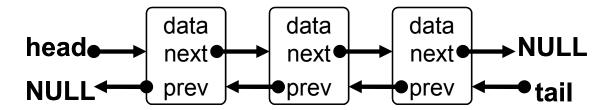
	Dynamic array	Singly linked List
Organization	Continuous	Dis-continuous
Resize	Require re-allocation Complexity: O(n)	No re-allocation Complexity: O(1)
Access element	Random access Complexity: O(1)	Sequential access Complexity: O(n)
Search	Forward/Backward Complexity: O(n)	Forward only Complexity: O (n)
Add/Remove element	Require element shifting Complexity: O(n)	No shifting Complexity: O(1)
Memory cost	No extra memory	Require extra memory Cost: 4 * n bytes

→ Efficient way to store **sequential** and **variable-size** data.



Doubly linked list:

- Element = Data + Link to next + Link to previous.
- Head: forward traverse.
- Tail: backward traverse.



C declaration:

```
struct DNode struct DList
{
  int data; DNode *head;
  DNode *next; DNode *tail;
  DNode *prev; };
};
```



Doubly linked list vs. Dynamic array:

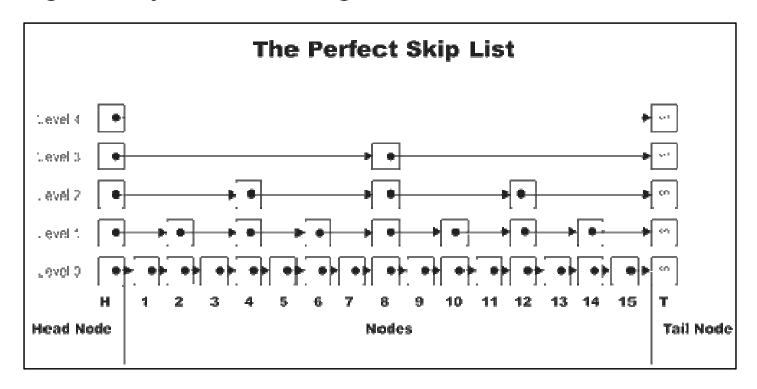
	Dynamic array	Doubly linked list
Organization	Continuous	Dis-continuous
Resize	Require re-allocation Complexity: O(n)	No re-allocation Complexity: O(1)
Access element	Random access Complexity: O(1)	Sequential access Complexity: O(n)
Search	Forward/Backward Complexity: O(n)	Forward/Backward Complexity: O (n)
Add/Remove element	Require element shifting Complexity: O(n)	No shifting Complexity: O(1)
Memory cost	No extra memory	Require extra memory Cost: 8 * n bytes

→ Good for *two-directional traverse*, *variable-size* data.



Skip list:

- Set of singly linked lists organized in layers.
- Lower layer: fine-grained nodes, slower "lane".
- Higher layer: coarse-grained nodes, faster "lane".





Skip list vs. Ordered array:

	Ordered Array	Skip List
Organization	Continuous	Dis-continuous
Resize	Require re-allocation Complexity: O(n)	No re-allocation Complexity: O(1)
Access element	Random access Complexity: O(1)	Sequential access Complexity: O(log(n))
Binary search	Forward/Backward Complexity: O(n*log(n))	Forward only Complexity: O (n*log(n))
Add/Remove element	Require element shifting Complexity: O(n)	No shifting Complexity: O(1)
Memory cost	No extra memory	Require extra memory Cost: log(n) * 4 * n bytes

→ Efficient way to store *variable-size* and *ordered* data.

Summary



Singly linked list concepts:

- Dis-continuous storage.
- Node = data + next.
- Last node points to NULL.

Singly linked list operations:

- Initialize, check empty.
- Find, add, remove.

Singly linked list improvements:

- Doubly linked list.
- Circularly linked list.

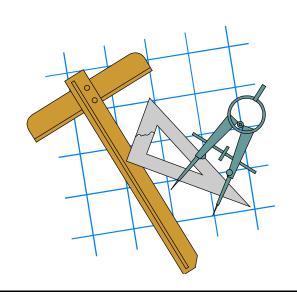


Practice



■ Practice 2.1:

Write C/C++ program to implement operations on Circularly linked lists.



Practice



■ Practice 2.2:

Write C/C++ program to implements the following operations on Singly linked list:

- Count nodes in list.
- Reverse list.
- Add new node (keep order).

