# User Manual for FFTLog-and-Beyond

Xiao Fang

#### 1 Versions

- Version 2.0: 2021/07/14. fftlogx released. Supported new forms of integrals. Upgraded for use in LSST-DESC N5K (Non-Limber method) Challenge.
- Version 1.0: 2019/11/22. Independent pure Python version and C version released.

# 2 Problem Description

The integrals we are solving are

$$F(y) = \int_0^\infty \frac{dx}{x} f(x) j_\ell^{(n)}(xy) , \qquad (1)$$

$$H(y) = \int_0^\infty \frac{dx}{x} f(x) |j_\ell(xy)|^2 , \qquad (2)$$

where f(x) is an input array,  $j_{\ell}$  is the order- $\ell$  spherical Bessel function of the first kind, the superscript <sup>(n)</sup> denotes the order of derivative. This type of integrals are numerically challenging due to the rapidly oscillatory nature of the spherical Bessel functions, especially when the input f(x) data array correspond to sampling array x over a large range (*i.e.*, over several orders of magnitude).

In weak lensing, the integral often takes form of

$$S(y) = \int_0^\infty \frac{dx}{x} f(x) \frac{j_\ell(xy)}{(xy)^2} , \qquad (3)$$

which could be solved by moving y dependence outside of the integral, and reducing to form F(y). However, we will present a numerically better approach.

# 3 Efficient Computation with FFTLog-and-Beyond

The essential idea is to expand f(x) into a series of power-laws and solve each component integral analytically. We require a logarithmic sampling of x with linear spacing in  $\ln(x)$  equal to  $\Delta_{\ln x}$ , i.e.,  $x_q = x_0 \exp(q\Delta_{\ln x})$  with  $x_0$  being the smallest value in the x array. The power-law decomposition then means

$$f(x_q) = \frac{1}{N} \sum_{m=-N/2}^{N/2} c_m x_0^{\nu} \left(\frac{x_q}{x_0}\right)^{\nu+i\eta_m} , \qquad (4)$$

where N is the sample size of the input function,  $\eta_m = 2\pi m/(N\Delta_{\ln x})$ , and  $\nu$  is the bias index. The Fourier coefficients satisfy  $c_m^* = c_{-m}$  since function f(x) is real, and are computed by discrete Fourier transforming the "biased" input function  $f(x)/x^{\nu}$  as

$$c_m = W_m \sum_{q=0}^{N-1} \frac{f(x_q)}{x_q^{\nu}} e^{-2\pi i mq/N} , \qquad (5)$$

where  $W_m$  is a window function which smooths the edges of the  $c_m$  array and takes the form of Eq. (C.1) in McEwen et al (2016, arXiv: 1603.04826). This filtering is found to reduce the ringing effects.

Each term is now analytically solvable. For n = 0, i.e., no derivative,

$$F(y) = \frac{1}{Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} \int_0^{\infty} \frac{dx}{x} x^{\nu+i\eta_m} j_{\ell}(x)$$

$$= \frac{\sqrt{\pi}}{4Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} g_{\ell}(\nu+i\eta_m) , \qquad (6)$$

where the first equality uses change of variable  $xy \to x$ . Function  $g_{\ell}(z)$  is given by

$$g_{\ell}(z) = 2^{z} \frac{\Gamma\left(\frac{\ell+z}{2}\right)}{\Gamma\left(\frac{3+\ell-z}{2}\right)} , \quad -\ell < \Re(z) < 2 , \tag{7}$$

giving the range of bias index  $-\ell < \nu < 2$ .

Finally, assuming that y is logarithmically sampled with the same linear spacing  $\Delta_{\ln y} = \Delta_{\ln x}$  in  $\ln y$ , we can write the last summation in Eq. (6) as

$$F(y_p) = \frac{\sqrt{\pi}}{4y_p^{\nu}} \text{IFFT} \left[ c_m^*(x_0 y_0)^{i\eta_m} g_\ell(\nu - i\eta_m) \right] , \qquad (8)$$

where  $y_p$   $(p = 0, 1, \dots, N - 1)$  is the p-th element in the y array. IFFT stands for the Inverse Fast Fourier Transform. In summary, this method performs two FFT operations, one in computing  $c_m$ , one in the final summation over m. Thus, the total time complexity is  $\mathcal{O}(N \log N)$ . So far we have described the principle of the FFTLog algorithm.

For n > 0, following the same procedure of power-law decomposition, we have

$$F_n(y) = \frac{1}{Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} \int_0^{\infty} \frac{dx}{x} x^{\nu + i\eta_m} j_{\ell}^{(n)}(x) . \tag{9}$$

Again, the integral for each m has an analytic solution, which can be shown with integration by parts. We write the solution in the same form with the FFTLog, i.e.,

$$F_n(y) = \frac{\sqrt{\pi}}{4Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} \tilde{g}_{\ell}(n, \nu + i\eta_m) , \qquad (10)$$

and its discrete version assuming  $\Delta_{\ln y} = \Delta_{\ln x}$ ,

$$F_n(y_p) = \frac{\sqrt{\pi}}{4y_p^{\nu}} \text{IFFT} \left[ c_m^* (x_0 y_0)^{i\eta_m} \tilde{g}_{\ell}(n, \nu - i\eta_m) \right] , \qquad (11)$$

where  $\tilde{g}_{\ell}(n,z) = 4\pi^{-1/2} \int_0^\infty dx \, x^{z-1} j_{\ell}^{(n)}(x)$ . For n=0,  $\tilde{g}_{\ell}(0,z) = g_{\ell}(z)$ , and for n=1,2, it is given by

$$\tilde{g}_{\ell}(1,z) = -2^{z-1}(z-1)\frac{\Gamma\left(\frac{\ell+z-1}{2}\right)}{\Gamma\left(\frac{4+\ell-z}{2}\right)} , \quad \begin{pmatrix} 0 < \Re(z) < 2, & \text{for } \ell = 0\\ 1 - \ell < \Re(z) < 2, & \text{for } \ell \ge 1 \end{pmatrix} , \tag{12}$$

$$\tilde{g}_{\ell}(2,z) = 2^{z-2}(z-1)(z-2)\frac{\Gamma\left(\frac{\ell+z-2}{2}\right)}{\Gamma\left(\frac{5+\ell-z}{2}\right)} , \quad \left(\begin{array}{c} -\ell < \Re(z) < 2 , & \text{for } \ell = 0, 1\\ 2 - \ell < \Re(z) < 2 , & \text{for } \ell \ge 2 \end{array}\right) .$$
(13)

We choose  $\nu=1$  for all  $\ell$ 's. With this generalized FFTLog algorithm, the integral containing one derivative of a spherical Bessel function also takes 2 FFT operations to compute.

Most generally, for  $n \in \mathcal{N}$ ,

$$\int_0^\infty dx \, x^{\alpha - 1} j_\ell^{(n)}(x) = (-1)^n \frac{\sqrt{\pi}}{4} 2^{\alpha - n} \frac{\Gamma(\alpha)}{\Gamma(\alpha - n)} \frac{\Gamma(\frac{\ell + \alpha - n}{2})}{\Gamma(\frac{3 + n + \ell - \alpha}{2})} \,, \quad \left( \begin{array}{c} -\ell < \Re(\alpha) < 2 \,, & \text{for } \ell < n \\ n - \ell < \Re(\alpha) < 2 \,, & \text{for } \ell < n \end{array} \right). \tag{14}$$

For H(y), a similar method can be developed.

$$H(y) = \frac{1}{Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} \int_0^{\infty} \frac{dx}{x} x^{\nu+i\eta_m} |j_{\ell}(x)|^2$$

$$= \frac{\sqrt{\pi}}{4Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} h_{\ell}(\nu + i\eta_m) , \qquad (15)$$

where

$$h_{\ell}(z) = \frac{4}{\sqrt{\pi}} \frac{\pi}{8} 2^{z} \frac{\Gamma\left(\ell + \frac{z}{2}\right) \Gamma(2 - z)}{\Gamma\left(2 + \ell - \frac{z}{2}\right) \left[\Gamma\left(\frac{3 - z}{2}\right)\right]^{2}} = \frac{\Gamma\left(\ell + \frac{z}{2}\right) \Gamma\left(\frac{2 - z}{2}\right)}{\Gamma\left(2 + \ell - \frac{z}{2}\right) \Gamma\left(\frac{3 - z}{2}\right)},$$
(16)

where the second equality utilizes the gamma function's duplication formula,  $\Gamma(2-z)=\frac{2^{1-z}}{\sqrt{\pi}}\Gamma\left(\frac{2-z}{2}\right)\Gamma\left(\frac{3-z}{2}\right)$ . The final result in discrete form is then

$$H(y_p) = \frac{\sqrt{\pi}}{4y_p^{\nu}} \text{IFFT} \left[ c_m^*(x_0 y_0)^{i\eta_m} h_{\ell}(\nu - i\eta_m) \right] , \qquad (17)$$

For S(y), following the power-law decomposition, we have

$$S(y) = \frac{1}{Ny^{\nu}} \sum_{m=-N/2}^{N/2} c_m x_0^{-i\eta_m} y^{-i\eta_m} \int_0^\infty \frac{dx}{x} x^{\nu+i\eta_m} \frac{j_{\ell}(x)}{x^2} . \tag{18}$$

Now the discrete result becomes

$$S(y_p) = \frac{\sqrt{\pi}}{4y_p^{\nu}} \text{IFFT} \left[ c_m^* (x_0 y_0)^{i\eta_m} g_{\ell} (\nu - 2 - i\eta_m) \right] . \tag{19}$$

Similar expressions can be derived for integrals with derivatives.

### 4 Using the Code

This repository contains independent python module and C module for computing Eqs. (1-3).

#### 4.1 fftlogx Version: C implementation + Python wrapper

This is the most recent and most recommended version. It has C code in src/ and a python wrapper in fftlogx/. The code contains all functionality presented in this Note.

To use the code, run make command at the repository, which will build the C file and generate a shared library file libfftlogx.so. The python wrapper should then be used. See test files in fftlogx/ for examples.

#### 4.2 Pure Python Version (Version 1, not recommended)

The main code is all in python/fftlog.py, along with python/test.py which provides an example of calling the module to compute the integrals with f(x) being the power spectrum in Pk\_test, and n = 0, 1, 2, respectively. We also have a Hankel function, defined as replacing  $j_{\ell}$  as  $J_n$  in Eq. (1). If you find them useful in your research, please cite Fang et al (2019).

## 4.3 C Version (Version 1, not recommended)

The code, sitting in cfftlog folder, uses FFTW library. The main code is in cfftlog.c, with auxiliary functions (e.g., extrapolation, window function,  $g_{\ell}$  functions) defined in utils.c and utils\_complex.c.

cfftlog.c provides different ways to run the computation. cfftlog provides the simplest usage. cfftlog\_ells function enables to compute the same integral with an array of  $\ell$  values. This is more optimal than calling single function cfftlog many times since FFTW plans have to be re-created and destroyed over and over again, and same for the FFTW\_COMPLEX arrays used for the FFTs.

cfftlog\_ells\_increment does the same thing as cfftlog\_ells, but increment the results (instead of refreshing). This saves memory for some applications.

test.c and test\_ells.c show examples of calling those functions.