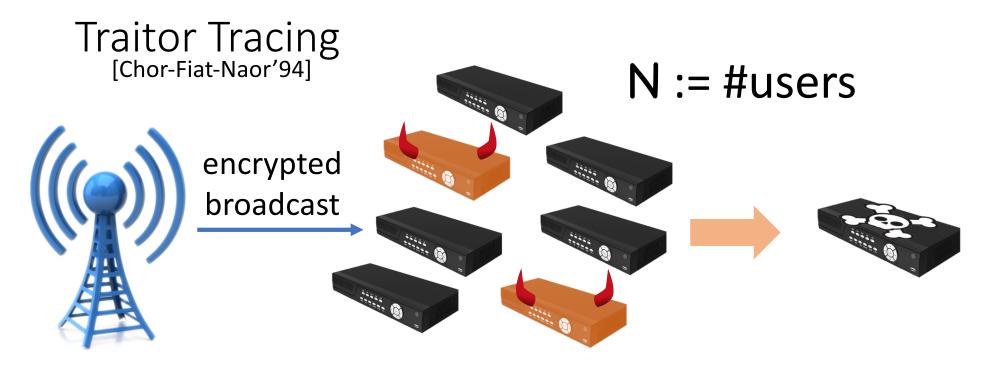
# New Techniques for Traitor Tracing: Size N<sup>1/3</sup> and More from Pairings

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#### **Requirement**

Given pirate decoder, can identify the traitor(s)

- \* Even if arbitrarily many users collude
- \* Even if decoder fails most of the time

### Main Objective?

[me'13]

"The goal is to build collusion-resistant traitor tracing where ciphertext overhead in terms of N is minimized"

Sentiment common to much of the literature

## Not the whole story...

Boneh-Naor'02:

 $PKE \longrightarrow |ctxt| = O(1)$ 

Combinatorial, uses "fingerprinting codes" [Boneh-Shaw'95]

Different views on why it doesn't "count"

#### **Problem 1:**

Only "threshold" secure

(Can only trace decoder if Pr[decrypt] ≥ 0.9)

#### **Problem 2:**

 $\Omega(N^2)$ -sized secret keys

Considered too large

## Main Objective, Take 2

[me'20]

"The goal is to build collusion-resistant traitor tracing offering the best parameter-size *trade-offs* in terms of N"

"And ideally, without the threshold limitation"

#### What's Known

$$(P, K, C) = |P| = P(N) \times poly(\lambda)$$
  
 $(S, K, C) = |SK| = K(N) \times poly(\lambda)$   
 $|C| + |C| + |C|$ 

Boneh-Sahai-Waters'06: Pairings  $(N^{1/2}, 1, N^{1/2})$ Garg-Gentry-Halevi-Raykova-Sahai-Waters'13, Boneh-Z'14: (1,1,1)Goyal-Koppula-Waters'18: LWE (1,1,1) Trivial:

PKE  $\rightarrow$  (N,1,N)

IBE  $\rightarrow$  (1,1,N)

Boneh-Naor'02:

PKE  $\rightarrow$  (N<sup>2</sup>,N<sup>2</sup>,1)

IBE  $\rightarrow$  (1,N<sup>2</sup>,1)

**Threshold** 

## Some Previously Open Questions

```
PKE, IBE,
Pairing-free groups, \rightarrow (*, N<sup>1.99</sup>, N<sup>0.99</sup>)?
or Factoring-like (even w/ threshold tracing)
```

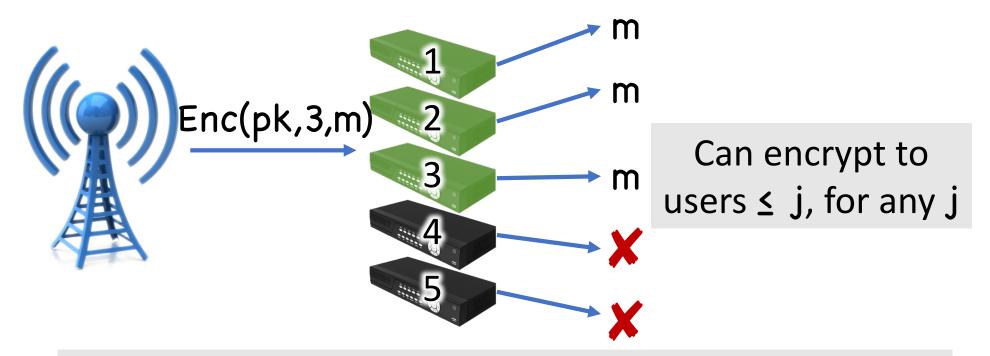
Pairings 
$$\rightarrow$$
 (\*, N<sup>1.99</sup>, N<sup>0.49</sup>)? (even w/ threshold tracing)

#### Observation

(no threshold **or** fully sublinear)

All the "best" collusion-resistant schemes in the literature follow "PLBE" framework

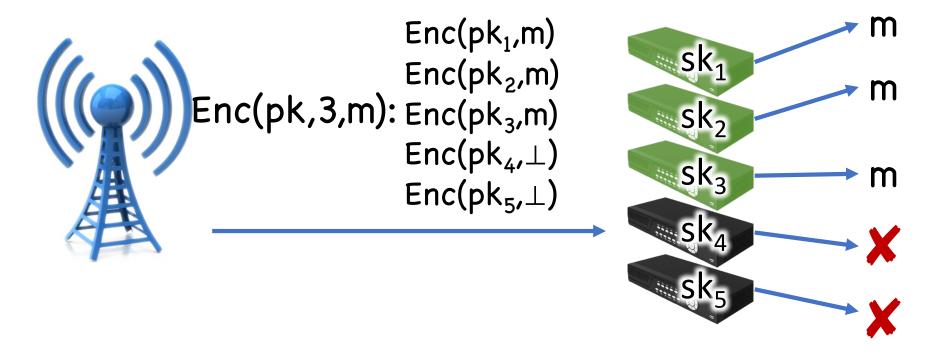
#### Private Linear Broadcast Encryption (PLBE)



Plus: User i learns nothing about j, except whether i≤j

Thm ([Boneh-Sahai-Waters'06]): PLBE → Traitor Tracing

### Trivial PLBE



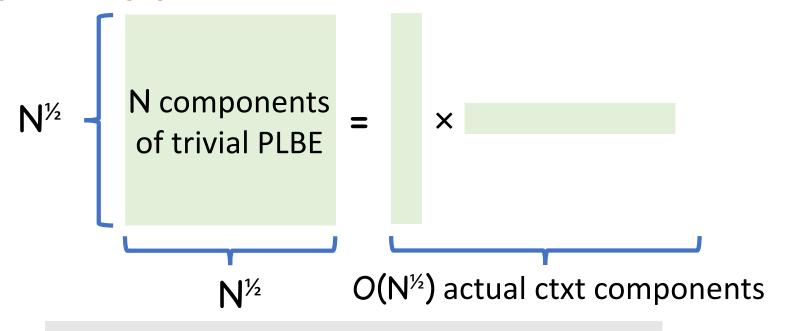
### PLBE-Based Traitor Tracing

Trivial PLBE: O(N)-sized ciphertexts

All the "best" traitor tracing schemes = improved algebraic constructions of PLBE

## The N<sup>1</sup>/<sub>2</sub> Barrier for Pairings

 $e(g^a,g^b) = e(g,g)^{ab}$   $\rightarrow$  Degree-2 functions in exponent



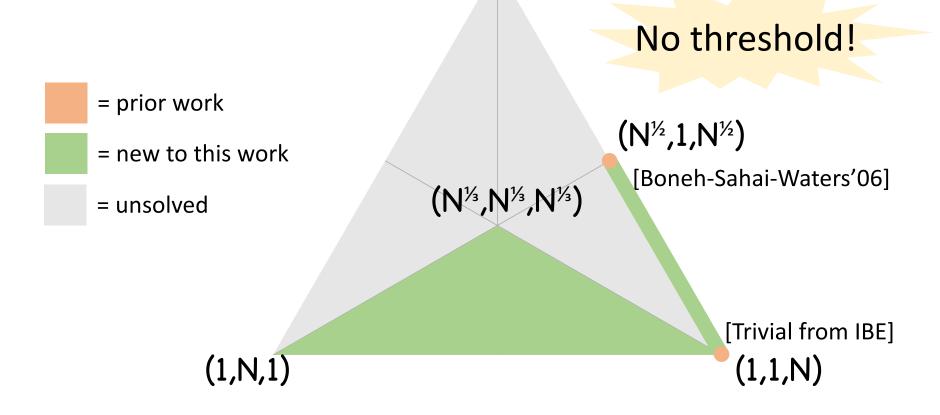
 $N^{\frac{1}{2}}$  = best known PLBE from pairings

This Work: New techniques for (collusion-resistant) traitor tracing



New parameter trade-offs from pairings and other primitives

## Parameters from Pairings P×K×C=N Simplex: (N,1,1)



#### Other Results

#### No threshold!

Pairings  $\rightarrow$  (N<sup>1-a</sup>,1,N<sup>a</sup>)  $\forall$  a  $\in$  [½,1] w/ Broadcast

Compare w/ [Boneh-Water'06]: Pairings  $\rightarrow$  (N<sup>1/2</sup>,N<sup>1/2</sup>,N<sup>1/2</sup>)

Pairings  $\rightarrow$  (N<sup>1-a</sup>,N<sup>1-a</sup>,N<sup>a</sup>)  $\forall a \in [0,1]$  w/ Broadcast

Compare w/ [Goyal-Quach-Waters-Wichs'19] : Pairings + LWE  $\rightarrow$  (N,N<sup>2</sup>,N<sup> $\epsilon$ </sup>)

#### Other Results

PKE 
$$\rightarrow$$
 (N<sup>2-a</sup>,N<sup>2-2a</sup>,N<sup>a</sup>)  $\forall$  a  $\in$  [0,1]

No threshold!

$$\mathsf{IBE} \Longrightarrow (1,\mathsf{N}^{2-2a},\mathsf{N}^a) \ \forall \, a \in [0,1]$$

$$a=0 \rightarrow |ctxt| = O(1)$$
  
 $a=\frac{2}{3} \rightarrow |sk|=|ctxt|=O(N^{\frac{2}{3}})$ 

First fully sub-linear schemes from pairingfree groups or factoring-like assumptions [Cocks'01,Döttling-Garg'17]

### Techniques

Generically remove thresholds w/o asymptotically changing (P,K,C)



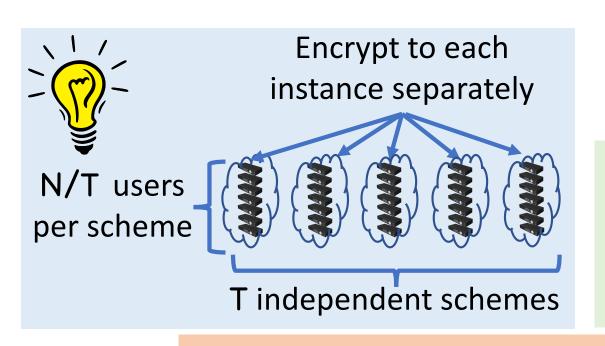
"risky" no risky (个K)

Threshold\* Broadcast > traitor tracing

New algebraic instantiations from pairings

\* Not to be confused w/ threshold tracing

## Trading off C for P,K: Generalizing Trivial PLBE



Often, using IBE techniques

Parameters;

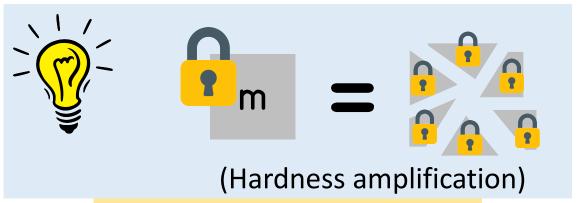
 $P(N) \rightarrow \times P(N/T)$ 

 $K(N) \rightarrow K(N/T)$ 

 $C(N) \rightarrow T \times C(N/T)$ 

Note: Factor T loss Threshold tracing

### Removing Thresholds



Key feature: #(shares) independent of N

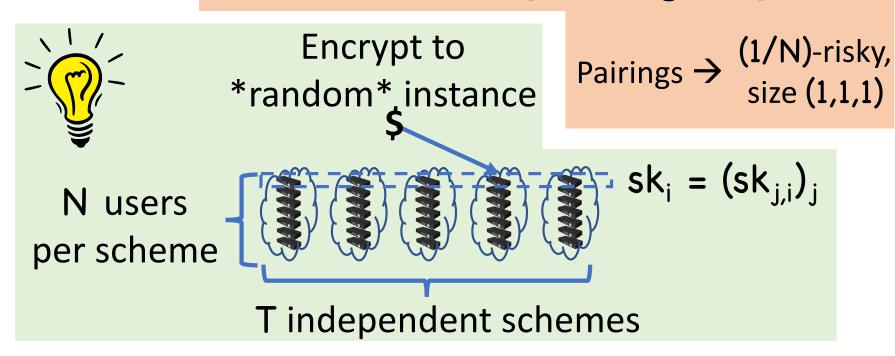
#### Parameters:

 $P(N) \rightarrow P(N)$   $K(N) \rightarrow K(N)$  $C(N) \rightarrow C(N)$ 

Already enough for PKE/IBE results

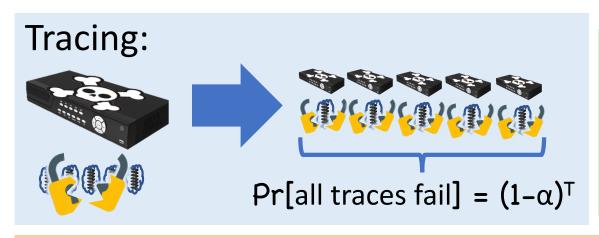
## Mitigating Risk

 $\alpha$ -Risky Tracing: Pr[false positive]  $\leq$  negletical negative  $\leq$  1- $\alpha$ 



## Mitigating Risk

IBE techniques



Parameters:

 $P(N) \rightarrow (N)$ 

 $K(N) \rightarrow \alpha^{-1} \times K(N)$ 

 $C(N) \rightarrow C(N)$ 

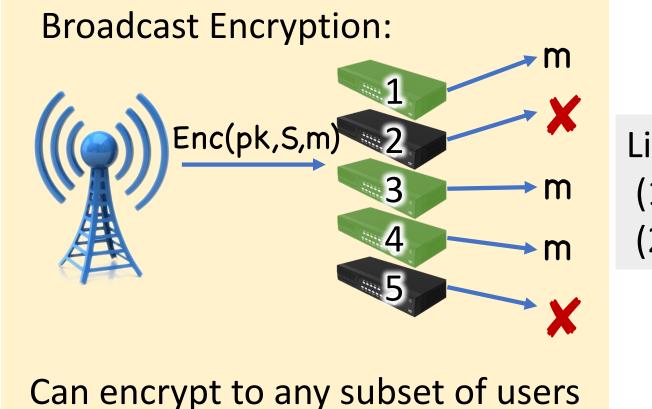
Require

 Only threshold scheme

Then apply threshold elimination

Enough for (1,N,1)

# Threshold\* Broadcast -> Traitor Tracing



Like PLBE, except:

- (1) Arbitrary S
- (2) S public

\* Not to be confused w/ threshold tracing

# Threshold\* Broadcast -> Traitor Tracing



How to encrypt to \*secret\* sets, when S is public?



Assign users (semi-)random identities (Only user/tracer knows their identity)

Problem: can "guess" user identity

Solution: generalize to threshold functionality

\* Not to be confused w/ threshold tracing

## Putting It All Together

[Attrapadung-Herranz-Laguillaumie-Libert-Panafieu-Ràfols'12]:

(N,N,1) Threshold Broadcast

Optimize for tracing app

(N<sup>1/3</sup>,N<sup>1/3</sup>,N<sup>1/3</sup>)
Tracing

Apply compilers

Combine w/ "risky" tracing

#### Lessons Learned

PLBE \*not\* inherent to traitor tracing

Thresholds no longer limitation

Risky and threshold tracing useful stepping stones