Math Reference Sheet

For semi-frequent stuff that I don't remember off the top of my head.

BEST Theorem

Counts number of Eulerian circuits in directed graphs.

$$\mathrm{ec}(G) = t_w(G) \prod_{v \in V} (\deg(v) - 1)!$$

 $t_w(G)$ denotes the number of arborescences (directed trees pointing towards the root) rooted at any arbitrary node w and is calculated via Kirchhoff's matrix tree theorem.

Kirchhoff's Matrix Tree Theorem

Counts number of spanning trees of a graph. The algorithm builds the Laplacian matrix defined as the following for a simple graph:

$$L_{i,j} := egin{cases} \deg(v_i) & ext{if } i = j \ -1 & ext{if } i
eq j ext{ and } v_i ext{ is adjacent to } v_j \ 0 & ext{otherwise} \end{cases}$$

The answer is any cofactor of this matrix (e.g. the determinant after deleting the last row and column).

- Cayley's formula: the number of spanning trees of a complete graph of size n is n^{n-2}
- ullet For multigraphs, $L_{i,j}$ equals -m, where m is the number of edges between i and j, self-loops are excluded
- For directed multigraphs, $L_{i,j}$ equals -m, where m is the number of edges from i to j, and $L_{i,i}$ equals the indegree of i minus the number of loops at i
- Removing the *i*th row and column and taking the determinant gives the number of oriented spanning trees rooted at (pointing towards) vertex *i*.

Stirling Numbers of the First Kind

Counts number of permutations of length n with k cycles.

$$dp[n+1][k] = n \cdot dp[n][k] + dp[n][k-1] \ dp[0][0] = 1 \ dp[n][0] = dp[0][k] = 0$$

Explanation: if you have n elements split into k cycles and are inserting a new element, you can either create a new cycle or insert it directly behind any of the previous n elements in an existing cycle.

The generating function for signed Stirling numbers of the first kind can be computed for a fixed n in $\mathcal{O}(n\log^2 n)$:

$$\sum_{k=0}^n s(n,k)x^k = x(x-1)\ldots(x-(n-1))$$

The unsigned Stirling numbers of the first kind are similar:

$$\sum_{k=0}^{n} \begin{bmatrix} n \\ k \end{bmatrix} = x(x+1)\dots(x+n-1)$$

The unsigned Stirling numbers of the first kind can be computed for a fixed k in $\mathcal{O}(n \log n)$:

$$egin{bmatrix} n \ k \end{bmatrix} = rac{n!}{k!} [x^n] \Biggl(\sum_{n=1}^{\infty} rac{x^n}{n} \Biggr)^k$$

Stirling Numbers of the Second Kind

Counts number of ways to partition n labeled objects into k non-empty unlabeled subsets.

$$dp[n+1][k] = k \cdot dp[n][k] + dp[n][k-1]$$

 $dp[0][0] = 1$
 $dp[n][0] = dp[0][k] = 0$

Explanation: the n+1th object is either a singleton or not. If it is a singleton, distribute the remaining n objects among k-1 groups, otherwise insert n+1 into one of the k groups and distribute the remaining n objects in k groups as well.

The generating function for Stirling numbers of the second kind can be computed for a fixed n in $\mathcal{O}(n \log n)$:

$$\sum_{k=0}^n S(n,k) x^k = \left(\sum_{i=0}^n rac{(-1)^i}{i!}
ight) \left(\sum_{i=0}^n rac{i^n}{i!}
ight) \mod x^{n+1}$$

They can also be computed for a fixed k in $\mathcal{O}(n \log n)$:

$$S(n,k)=n![x^n]rac{(e^x-1)^k}{k!}$$
 $e^x=\sum_{n=0}^{\infty}rac{x^n}{n!}$

Partition Function

Counts number of ways to partition n into non-negative integer parts. Partition of n into k parts follows the following recurrence:

$$dp[n][k] = dp[n-1][k-1] + dp[n-k][k]$$

 $dp[0][0] = 1$
 $dp[n][0] = dp[0][k] = 0$

Explanation: there are two possibilities. If we include a 1 in the partition, we simply partition the remaining of n-1 into k-1 parts. Otherwise, each part has size greater than 1, so we subtract 1 from each part and solve recursively.

p(n) denotes the number of partitions of n and its generating function can be computed in $\mathcal{O}(n \log n)$:

$$\sum_{n=0}^{\infty}p(n)x^n=\prod_{k=1}^{\infty}rac{1}{1-x^k}$$

The denominator can be computed in $\mathcal{O}(n)$ time with the pentagonal number theorem:

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

Derangements

Counts the number of permutations where $p_i
eq i$ for all i.

$$dp[n] = (n-1)(dp[n-1] + dp[n-2]) \ dp[0] = 1 \ dp[1] = 0$$

Explanation: With n elements, consider index 1 which may receive n-1 different values. There are two cases. Either index 1 swaps with another index i, so we count the number of derangements among the remaining n-2 indices. Alternatively, index 1 receives value i but index i does not receive value 1, so this is equivalent to counting the number of derangements of n-1 indices as we can renumber value 1 as value i.

The probability of getting a derangement from a random shuffle is $\frac{1}{e}$, which means the number of shuffles needed to get a derangement is effectively constant.

Catalan Numbers

Shows up in numerous counting problems, such as number of correct bracket sequences of length 2n, number of full binary trees with n+1 leaves, etc.

$$\{1,1,2,5,14,42,132,429,\ldots\}$$
 $C_n=rac{1}{n+1}inom{2n}{n}$ $C_{n+1}=\sum_{i=0}^n C_i C_{n-i}$ $C_0=1$

Proof for combinatorial definition of Catalan numbers is based on reflection argument for grid paths.