

# Proposal: Method Calling Convention

## 1 Synopsis

To encode frame conditions, the pillar of conditionals, a straightforward approach is compile statements in a new frame as a new Dafny method. Here's a challenge:

qregs are expanded into its pure data representation such as `seq<nat>` and `seq<seq<nat>>` after compilation without using heap-allocated objects. A method where the users assumed that they passed a reference to the qreg object needs to emit a return value for the post-state of the qreg passed in. The type of the qreg passed to the method could change through the execution of the body, the return value needs to be calculated based on the final type accordingly.

### 1.1 Example

The following method

```
method GHZ(q : qreg[10])
  requires { q[0 .. 10] : nor  $\mapsto$   $\otimes$  i . (0) }
  ensures { q[0 .. 10] : en01  $\mapsto$   $\sum$  j  $\in$  [0 .. 2] .  $\otimes$  k  $\in$  [0 .. 10] . ( j ) }
{
  ...
}
```

should be compiled into something like

```
method GHZ_Compiled(q_in : seq<nat>) returns (q_out : seq<seq<nat>>)
  requires { q_in[0 .. 10] :  $\llbracket$  nor  $\mapsto$   $\otimes$  i . (0)  $\rrbracket$  }
  ensures { q_out[0 .. 10] :  $\llbracket$  en01  $\mapsto$   $\sum$  j  $\in$  [0 .. 2] .  $\otimes$  k  $\in$  [0 .. 10] . ( j )  $\rrbracket$  }
{
  ...
}
```

## 2 Proposal

Methods can in fact have multiple parameter, therefore it's important to have a uniformed order to compile those input and return parameters.

### 2.1 Method Type, Reworked

To get a deterministic description of methods, the new method type is defined as followed.

```
data MethodType = MethodType
{ mtSrcParams :: [MethodElem]
, mtSrcReturns :: [MethodElem]
, mtInstantiate :: Map.Map Var Range -> [(Partition, QTy)]
}
```

```
data MethodElem
= MTyPure Var Ty
| MTyQuantum Var Exp'
deriving (Show, Eq, Ord)
```

The spotlight here is `mtInstantiate` which takes a correspondence between `qreg` variables used in definition and the `qreg` slices passed by the caller and returns partition-level typing requirements to be satisfied. The instantiator is deterministic, and therefore can determine the order of quantum states to be passed into the compiled method.

## 2.2 Input

```
(define-metafunction qafny
  call-conv : x ... -> x ...
  [(call-conv (n x ...))
   (n x' ...)]
  (where (x' ...) (call-conv x ...))
  (where n (not (is-qreg n))))]
[(call-conv (q x ...))
 (x' ... y' ...)]
  (where (x' ...) (call-conv x ...))
  (where (y' ...) (call-conv-q/env q)))]
```

where `call-conv-q/env` resolves emit symbols based on the type information encoded in the **pre**-condition. Passed qubits are compiled as emit variables corresponding to each range decided by the type information provided by the `instantiatd` instance. For convenience, those variables are put after non-quantum parameters.

## 2.3 Output

```
(define-metafunction qafny
  call-conv-ret : x ... -> x ...
  [(call-conv-ret (n x ...))
   (x' ...)]
  (where (x' ...) (call-conv-ret x ...))
  (where n (not (is-qreg n))))]
[(call-conv-ret (q x ...))
 (y' ... x' ...)]
  (where (x' ...) (call-conv-ret x ...))
  (where (y' ...) (call-conv-q/env q)))]
```

where `call-conv-q/env` resolves emit symbols based on the type information encoded in the **post**-condition.

# 3 Implementation

## 3.1 Pass qreg

How to pass a `qreg` argument? Pass the `qreg` seems suffices, but it's not enough because chances are that method requires less qubits than what the caller could provide. Therefore, the syntax should allow the caller to pass a slice of qubits, which can be of exactly the same syntax as passing a range.