Qafny Design and Bookkeeping

1 Typing

1.1 Typing with Sessions and Ranges

For example, let's say we have the full session

```
s = \{ x [0 .. 1]; y [0 .. 1]; z [0 .. 1] \}
```

For a λ application statement, let's say

```
x [0 ... 1] *= \lambda (x \Rightarrow x + 1)
```

The typing procedure should locate x[0 .. 1] the corr. session s and identify that the session s is indeed of CH type. With the CH type ensured, we then should emit the term for the oracle operator that should *only* modify the range corr. to x[0 .. 1] instead of the entire session.

Therefore, a reasonable way to do this is to have following functions

```
- | locate session with a range
findSession :: Range → Transform Session
- | locate a session with a (sub-)session
subSession :: Session → Transform Session
- | the typing of session (already done)
instance Typing Session QTy
- | sub-range based codegen
augWithSession :: Session → ? → Transform ?
```

However, I think there might be other case where you need to compare if two sessions are exactly the same.

1.1.1 Sidenote

At the same time, by viewing a session as a row, I see the connection between the subtyping between sessions w/ or w/o ranges and the open and close rows in the row polymorphism.

1.2 Type Coercion, or Casting

I need to implement a new function to perform type coercion between QTy's. When subtyping is satisfied, insert a cast to lift the subtype to its super type. This needs a support from QPreludeUntyped module.

2 QPreludeUntyped and Dafny 4.0

2.1 Ghost Functions

In the new standard, function is now function method by default, and ghost function represents function in Dafny 3.

2.2 abstract module

Dafny 4.0 introduced the notation abstract module which are modules not to be compiled. The upstream libraries seems to be adapted for that. (That's exactly what we need.) To conform to the standard, Codegen now generates a default abstract module wrapper to enable us to use those functions.

3 Compilation

3.1 separates Keyword and Body

I introduced **separates** keyword as a hint for the split-combine semantics for non-trivial guards. Let's say, if we have a guard

```
- _ leading
if (f (x [0 .. n]))
   separates x <_ .. m>
- _ trailing
if (f (x [0 .. n]))
   separates x <m .. _>
- segment
if (f (x [0 .. n]))
   separates x <i .. j>
```

The separates side condition asserts that only the sequence of states in x between i and j actually satisfies f(x[0..n]). I will need to insert assertions when implementing assertion translations.

3.2 CH Biview

CH is now expected to have two views: bitvector versus nat.

I think it's unnecessary to track views by installing a new state field. A tentative solution is to make 2 CH types, CHb and CHn, to distinguish those two views.

3.3 Biview Preference

From previous implementations, the GHZ one favors CHb and the Shor's one favors CHn. I think there's a pattern: in GHZ, the guard is of CH type and the λ oracle is *flip* (a very bit-level operation), while in Shor's, the guard is Had type and the λ oracle is modulo multiplication which is on nat.

3.4 Placeholder for Phase Calculus

The following instance is responsible for mapping a session type to its emission type which should include the emitted type for kets as well as phases.

```
instance Typing QTy [Ty]
```

Currently, this list is always a singleton and is flatten by the only1 combinator. Some lifting operation and bijective mapping will be expected when starting the phase calculus implementation.

4 Language Design

4.1 State Predicate

What would state predicates be like now?

5 Misc

5.1 Biview

Coincidentally, the choice between a sequence of bitvectors vs a sequence of nats is closely related to the idea in [Wadler 1987]

5.2 Point-free Translation of build0p2

```
x <> ("&&" <!> y)
= (<) x ("&&" <!> y)
= (flip (<)) ("&&" <!> y) $ x
= (flip (<)) (("&&" <!>) y) $ x
= (flip (<)) (("&&" <!>) y) $ x
```

References

[1] P. Wadler. "Views: a way for pattern matching to cohabit with data abstraction". In: *Proceedings of the 14th ACM SIGACT-SIGPLAN symposium on Principles of programming languages*. POPL '87. New York, NY, USA: Association for Computing Machinery, Oct. 1987, pp. 307–313. ISBN: 978-0-89791-215-0. DOI: 10.1145/41625.41653. URL: https://doi.org/10.1145/41625.41653 (visited on 02/10/2023).